

## Section 8.5

# Equivalence Relations

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Now we group properties of relations together to define new types of important relations.

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**Definition:** A relation  $R$  on a set  $A$  is an *equivalence relation* iff  $R$  is

- reflexive
- symmetric

and

- transitive
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It is easy to recognize equivalence relations using digraphs.

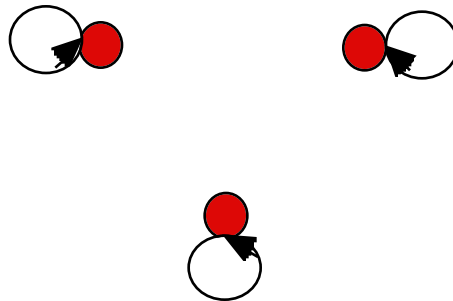
- The subset of all elements related to a particular element forms a universal relation (contains all possible arcs) on that subset. The (sub)digraph representing the subset is called a *complete* (sub)digraph. All arcs are present.

- The number of such subsets is called the *rank* of the equivalence relation

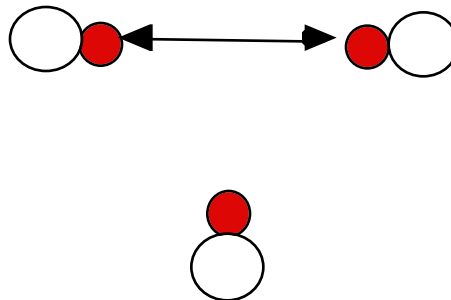
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Examples:

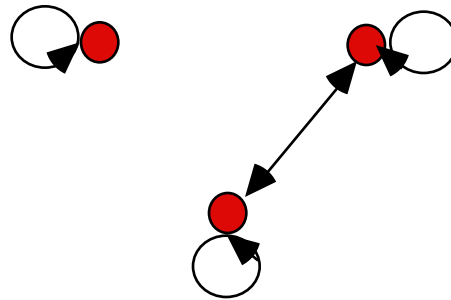
A has 3 elements:



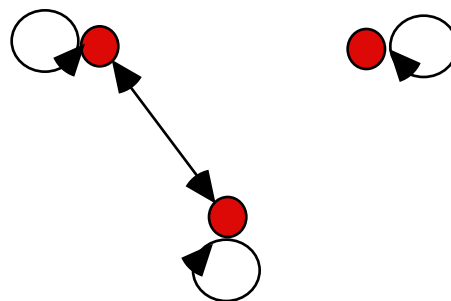
**rank = 3**



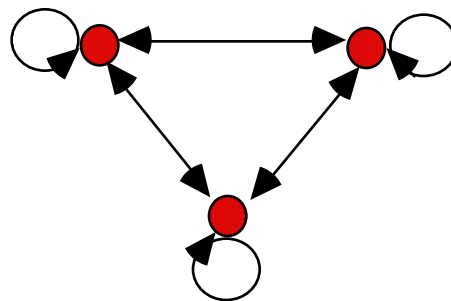
**rank = 2**



**rank = 2**



**rank = 2**



**rank = 1**



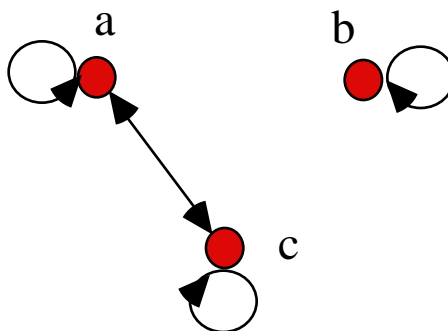
- Each of the subsets is called an *equivalence class*.
- A bracket around an element means the equivalence class in which the element lies.

$$[x] = \{y \mid \langle x, y \rangle \text{ is in } R\}$$

- The element in the bracket is called a *representative* of the equivalence class. We could have chosen any one.



Examples:



$$[a] = \{a, c\}, [c] = \{a, c\}, [b] = \{b\}.$$

$$\text{rank} = 2$$



An interesting counting problem:

Count the number of equivalence relations on a set  $A$  with  $n$  elements. Can you find a recurrence relation?

The answers are

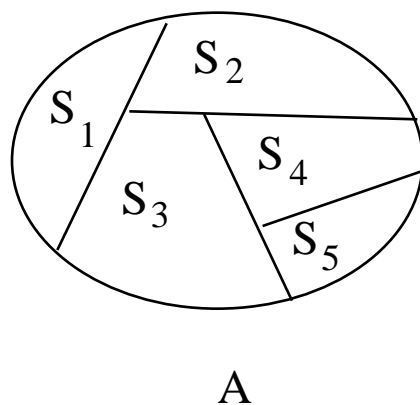
- 1 for  $n = 1$
- 3 for  $n = 2$
- 5 for  $n = 3$

How many for  $n = 4$ ?

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**Definition:** Let  $S_1, S_2, \dots, S_n$  be a collection of subsets of  $A$ . Then the collection forms a *partition* of  $A$  if the subsets are nonempty, disjoint and *exhaust*  $A$ :

- $S_i \cap S_j = \emptyset$  if  $i \neq j$
- $\bigcup S_i = A$



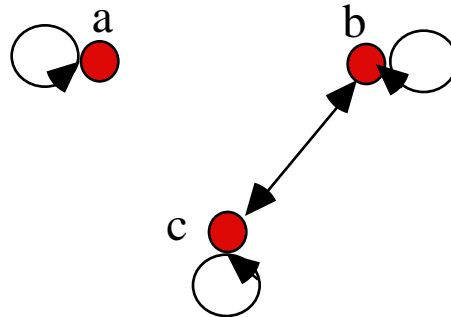
**Theorem:** The equivalence classes of an equivalence relation  $R$  *partition* the set  $A$  into disjoint nonempty subsets whose union is the entire set.

This partition is denoted  $A/R$  and called

- the *quotient set*, or
  - *the partition of  $A$  induced by  $R$* , or,
  - *$A$  modulo  $R$ .*
- 

Examples:

- $A \times A$
- $A =$
-



$$A = [a] \quad [b] = [a] \quad [c] = \{a\} \quad \{b, c\}$$

$$\text{rank} = 2$$


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**Theorem:** Let  $R$  be an equivalence relation on  $A$ . Then either

$$[a] = [b]$$

or

$$[a] \cap [b] = \emptyset$$


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**Theorem:** If  $R_1$  and  $R_2$  are equivalence relations on  $A$  then  $R_1 \cap R_2$  is an equivalence relation on  $A$ .

Proof: It suffices to show that the intersection of

- reflexive relations is reflexive,

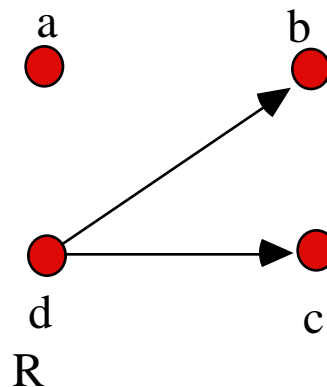
- symmetric relations is symmetric,
- and
- transitive relations is transitive.

You provide the details.

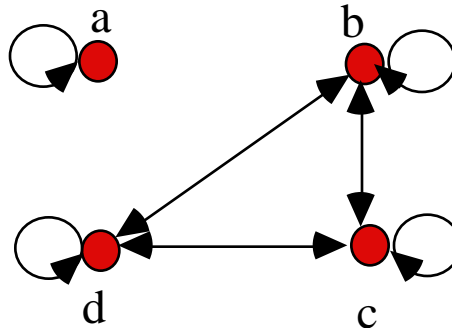
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**Definition:** Let  $R$  be a relation on  $A$ . Then the reflexive, symmetric, transitive closure of  $R$ ,  $\text{tsr}(R)$ , is an equivalence relation on  $A$ , called the *equivalence relation induced by  $R$* .

Example:







$$\begin{aligned} & \text{tsr}(R) \\ & \text{rank} = 2 \\ A = [a] \quad [b] = \{a\} \quad \{b, c, d\} \\ A/R = \{\{a\}, \{b, c, d\}\} \end{aligned}$$


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**Theorem:**  $\text{tsr}(R)$  is an equivalence relation

Proof:

We have to be careful and show that  $\text{tsr}(R)$  is still symmetric and reflexive.

- Since we only add arcs vs. deleting arcs when computing closures it must be that  $\text{tsr}(R)$  is reflexive since all loops  $\langle x, x \rangle$  on the diagraph must be present when constructing  $r(R)$ .

- If there is an arc  $\langle x, y \rangle$  then the symmetric closure of  $r(R)$  ensures there is an arc  $\langle y, x \rangle$ .

- Now argue that if we construct the transitive closure of  $sr(R)$  and we add an edge  $\langle x, z \rangle$  because there is a path from  $x$  to  $z$ , then there must also exist a path from  $z$  to  $x$  (why?) and hence we also must add an edge  $\langle z, x \rangle$ . Hence the transitive closure of  $sr(R)$  is symmetric.

Q. E. D.

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