

PVS Quick Reference

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- o PVS
 - . Enter pvs pvs filename.pvs
 - . Exit pvs C-X C-C
 - . Abort command C-G
 - o Context
 - . change context M-X CC
 - . save context M-X SC
 - . context path M-X CP
 - . list pvs files M-X LF
 - o Window
 - . split horizontal C-X 2
 - . split vertical C-X 3
 - . switch C-X 0
 - . delete windows C-X 1
 - o Buffers (temp work)
 - . list buffers C-X C-B
 - . select a buffer C-X B
 - . kill buffer C-X K
 - o File handling
 - . create new M-X new-pvs-file
 - . save C-X C-S
 - . open C-X C-F
 - o Specification
 - . parse M-X parse
 - . typecheck M-X typecheck
 - . prove M-X prove
 - o prover
 - . exit (exit)
 - . continue :continue n
 - . try prop strategy (prop)
 - . try bddsimp strategy (bddsimp)
 - . flatten (flatten fnum)
 - . split (split fnum)
 - . go to next subgoal (postpone)
 - . add premise (lemma "lemma-name")
 - . check status of proofs M-X status-pvs-file
 - . dump prfs and pvs file M-X dump-pvs-file
 - . undump M-X undump-pvs-file
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example1_pvs.txt

EXAMPLE 1:

```
myTheory: THEORY % name of theory is myTHEORY
BEGIN
  p, q, r: bool % declaring boolean variables p, q, and r

  axiom_1: AXIOM p IMPLIES q
  axiom_2: AXIOM q IMPLIES r
  th_1: THEOREM p IMPLIES r
END myTheory
```

EXAMPLE 2:

```
addsum: THEORY
BEGIN
  n: var nat

  sum (n): RECURSIVE nat =
    (if n = 0 THEN 0 ELSE n + sum(n -1) ENDIF)
  MEASURE n

  lem1: LEMMA sum(n) = n*(n+1)/2
END mySTACK
```

TUE 01/25/05

1. MISC

PVS site? _____

2. WHAT IS PVS?

How will one use PVS?
(say in your own words)

3. SPECIFICATION LANGUAGE

- Difference with programming language? _____
- Know syntax for declaring vars and formulas.
- Difference between AXIOMS and THEOREMS? _____
- SPECIFY:
Given: $p \rightarrow q, q \rightarrow p$
To prove: $(p \text{ OR } q) \rightarrow (p \text{ AND } q)$

4. PROVER

What is sequent form? _____

how to move a formula from theory to antecedant?

what commands to use to prove?

5. DEMO

You need to know how to:

- CREATE a pvs specs file
- Use the prover to PROVE a formula given some previously proved formulas or axioms.

Exercise1: THEORY

BEGIN

$a, b, c, n, t, h, s, p, q$: bool

lem0: LEMMA $(a \wedge (a \Rightarrow b)) \Rightarrow b$

lem1: LEMMA $(c \wedge n) \Rightarrow t$

lem2: LEMMA $h \wedge \neg s$

lem3: LEMMA $(h \wedge \neg (s \vee c)) \Rightarrow p$

lem4: LEMMA $(n \wedge \neg t) \Rightarrow p$

lem5: LEMMA $\neg (p \vee q) \equiv (\neg p \wedge \neg q)$

END Exercise1

Verbose proof for `lem0`.

`lem0`:

$$\frac{}{\{1\} \quad (a \wedge (a \Rightarrow b)) \Rightarrow b}$$

`lem0`:

$$\frac{}{\{1\} \quad (a \wedge (a \Rightarrow b)) \Rightarrow b}$$

Applying disjunctive simplification to flatten sequent,

`lem0`:

$$\frac{\begin{array}{l} \{-1\} \quad a \\ \{-2\} \quad (a \Rightarrow b) \end{array}}{\{1\} \quad b}$$

Splitting conjunctions,

we get 2 subgoals:

`lem0.1`:

$$\frac{\begin{array}{l} \{-1\} \quad b \\ \{-2\} \quad a \end{array}}{\{1\} \quad b}$$

which is trivially true.

This completes the proof of `lem0.1`.

`lem0.2`:

$$\frac{\{-1\} \quad a}{\begin{array}{l} \{1\} \quad a \\ \{2\} \quad b \end{array}}$$

which is trivially true.

This completes the proof of `lem0.2`.

Q.E.D.

Verbose proof for lem4.

lem4:

$$\frac{}{\{1\} \quad (n \wedge \neg t) \Rightarrow p}$$

lem4:

$$\frac{}{\{1\} \quad (n \wedge \neg t) \Rightarrow p}$$

Applying lem1

lem4:

$$\frac{\{-1\} \quad (c \wedge n) \Rightarrow t}{\{1\} \quad (n \wedge \neg t) \Rightarrow p}$$

Applying lem2

lem4:

$$\frac{\begin{array}{l} \{-1\} \quad h \wedge \neg s \\ \{-2\} \quad (c \wedge n) \Rightarrow t \end{array}}{\{1\} \quad (n \wedge \neg t) \Rightarrow p}$$

Applying lem3

lem4:

$$\frac{\begin{array}{l} \{-1\} \quad (h \wedge \neg (s \vee c)) \Rightarrow p \\ \{-2\} \quad h \wedge \neg s \\ \{-3\} \quad (c \wedge n) \Rightarrow t \end{array}}{\{1\} \quad (n \wedge \neg t) \Rightarrow p}$$

Applying disjunctive simplification to flatten sequent,

lem4:

$$\frac{\begin{array}{l} \{-1\} \quad (h \wedge \neg (s \vee c)) \Rightarrow p \\ \{-2\} \quad h \wedge \neg s \\ \{-3\} \quad (c \wedge n) \Rightarrow t \\ \{-4\} \quad n \end{array}}{\begin{array}{l} \{1\} \quad t \\ \{2\} \quad p \end{array}}$$

Applying disjunctive simplification to flatten sequent,

lem4:

{-1}	$(h \wedge \neg (s \vee c)) \Rightarrow p$
{-2}	h
{-3}	$(c \wedge n) \Rightarrow t$
{-4}	n
<hr/>	
{1}	s
{2}	t
{3}	p

Splitting conjunctions,
we get 4 subgoals:

lem4.1:

{-1}	p
{-2}	h
{-3}	$(c \wedge n) \Rightarrow t$
{-4}	n
<hr/>	
{1}	s
{2}	t
{3}	p

which is trivially true.

This completes the proof of lem4.1.

lem4.2:

{-1}	h
{-2}	$(c \wedge n) \Rightarrow t$
{-3}	n
<hr/>	
{1}	h
{2}	s
{3}	t
{4}	p

which is trivially true.

This completes the proof of lem4.2.

lem4.3:

{-1}	s
{-2}	h
{-3}	$(c \wedge n) \Rightarrow t$
{-4}	n
<hr/>	
{1}	s
{2}	t
{3}	p

which is trivially true.

This completes the proof of lem4.3.

lem4.4:

{-1}	c
{-2}	h
{-3}	$(c \wedge n) \Rightarrow t$
{-4}	n
{1}	s
{2}	t
{3}	p

Splitting conjunctions,
we get 3 subgoals:

lem4.4.1:

{-1}	t
{-2}	c
{-3}	h
{-4}	n
{1}	s
{2}	t
{3}	p

which is trivially true.

This completes the proof of lem4.4.1.

lem4.4.2:

{-1}	c
{-2}	h
{-3}	n
{1}	c
{2}	s
{3}	t
{4}	p

which is trivially true.

This completes the proof of lem4.4.2.

lem4.4.3:

{-1}	c
{-2}	h
{-3}	n
{1}	n
{2}	s
{3}	t
{4}	p

which is trivially true.

This completes the proof of lem4.4.3.

Q.E.D.

Verbose proof for `lem5`.

`lem5`:

$$\frac{}{\{1\} \quad \neg (p \vee q) \equiv (\neg p \wedge \neg q)}$$

`lem5`:

$$\frac{}{\{1\} \quad \neg (p \vee q) \equiv (\neg p \wedge \neg q)}$$

Splitting conjunctions,
we get 2 subgoals:

`lem5.1`:

$$\frac{}{\{1\} \quad \neg (p \vee q) \supset (\neg p \wedge \neg q)}$$

Applying disjunctive simplification to flatten sequent,

`lem5.1`:

$$\frac{}{\begin{array}{l} \{1\} \quad p \\ \{2\} \quad q \\ \{3\} \quad (\neg p \wedge \neg q) \end{array}}$$

Splitting conjunctions,
we get 2 subgoals:

`lem5.1.1`:

$$\frac{}{\begin{array}{l} \{-1\} \quad p \\ \{1\} \quad p \\ \{2\} \quad q \end{array}}$$

which is trivially true.

This completes the proof of `lem5.1.1`.

`lem5.1.2`:

$$\frac{}{\begin{array}{l} \{-1\} \quad q \\ \{1\} \quad p \\ \{2\} \quad q \end{array}}$$

which is trivially true.

This completes the proof of `lem5.1.2`.

`lem5.2`:

$$\frac{}{\{1\} \quad (\neg p \wedge \neg q) \supset \neg (p \vee q)}$$

Applying disjunctive simplification to flatten sequent,

lem5.2:

{-1}	$(p \vee q)$	
{1}	p	
{2}	q	

Splitting conjunctions,
we get 2 subgoals:

lem5.2.1:

{-1}	p	
{1}	p	
{2}	q	

which is trivially true.

This completes the proof of lem5.2.1.

lem5.2.2:

{-1}	q	
{1}	p	
{2}	q	

which is trivially true.

This completes the proof of lem5.2.2.

Q.E.D.