

## **CprE 2810: Digital Logic**

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http://www.ece.iastate.edu/~alexs/classes/

# Synchronous Sequential Circuits Basic Design Steps

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#### **Administrative Stuff**

• Homework 10 is due on Monday Nov. 10 @ 10pm.

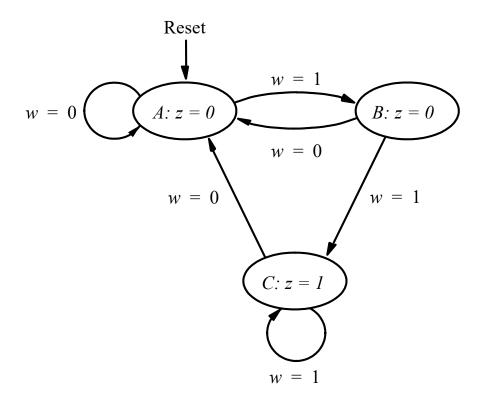
#### **Administrative Stuff**

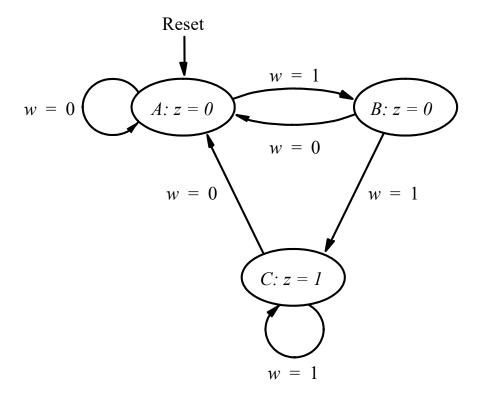
• We are starting with Chapter 6 from the textbook

# First Design Pattern: Moore Machines

### **Moore Machine:**

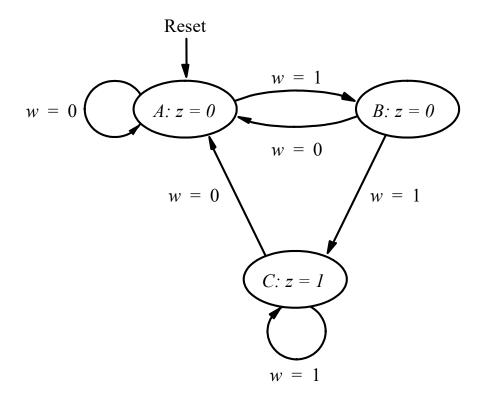
### A Type of Finite State Machine (FSM)



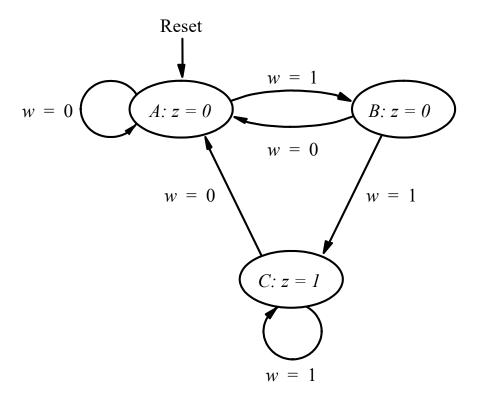


- Finite number of states (nodes).
- Discrete state transitions (edges).
- Only "in" one state at a time.
- One reset state
- Every state has an outgoing state transition for each possible input.

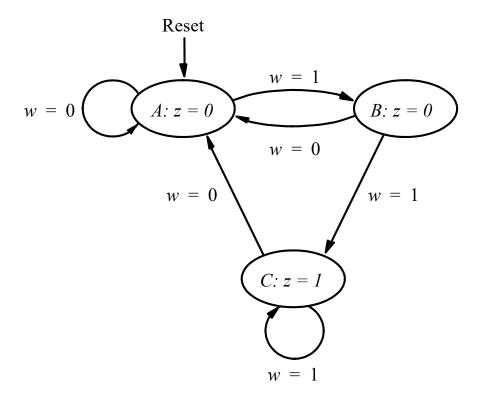
[ Figure 6.3 from the textbook ]



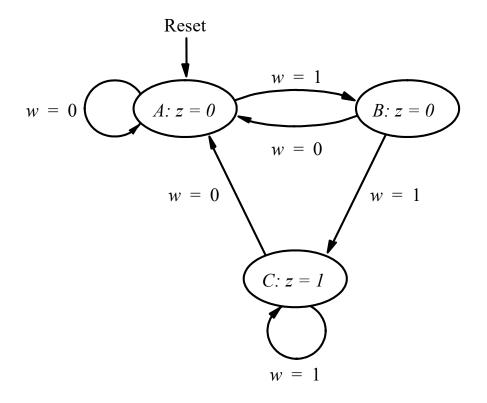
The next state depends on both the current state and the current input.



The output depends only on the current state.

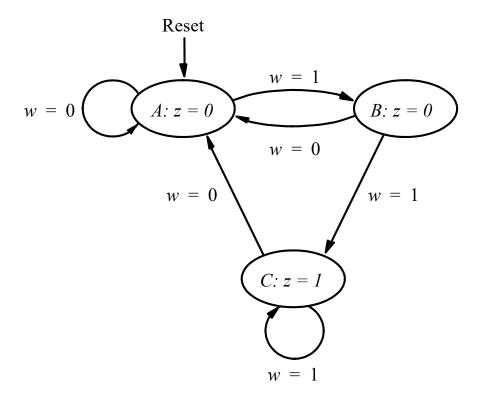


Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	t <sub>7</sub>	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0

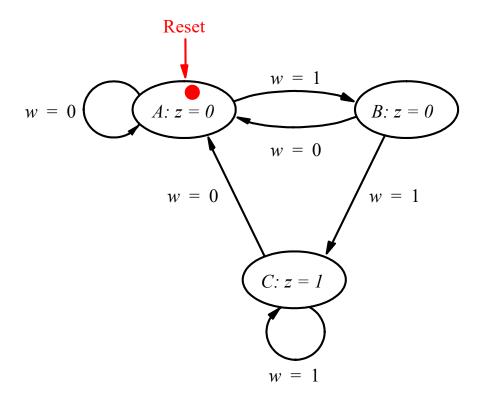


Clockey	ycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	t <sub>10</sub>
input	<i>w</i> :	0	1	0	1	1	0	1	1	1	0	1
output	<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0

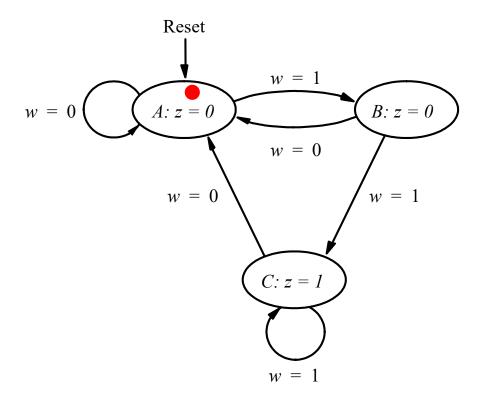
#### Let's do a simulation



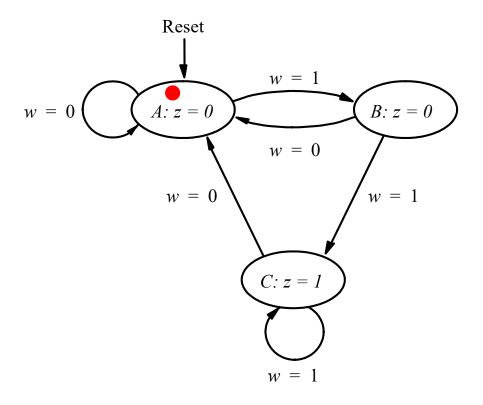
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	t <sub>7</sub>	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



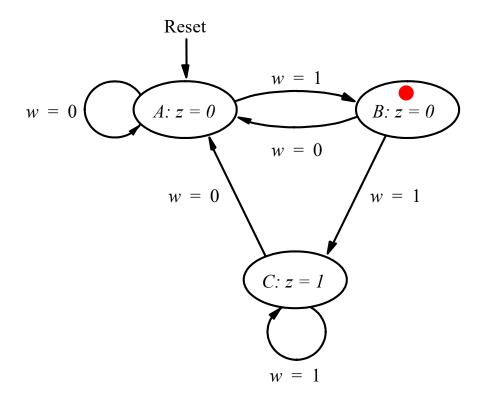
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
z:	0	0	0	0	0	1	0	0	1	1	0



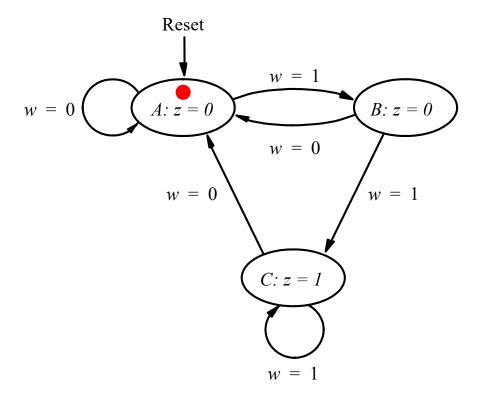
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
z:	0	0	0	0	0	1	0	0	1	1	0



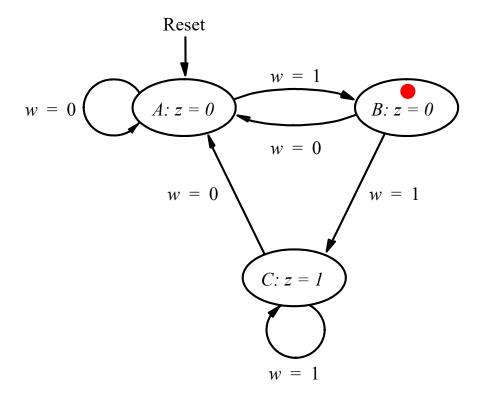
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



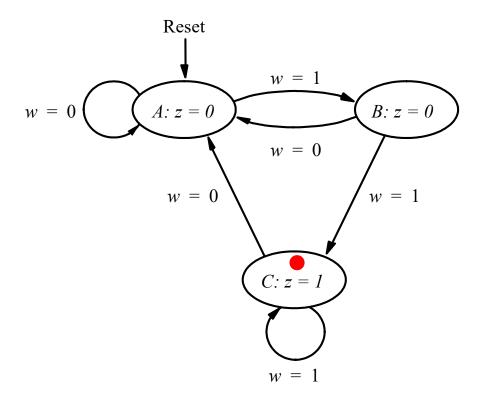
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



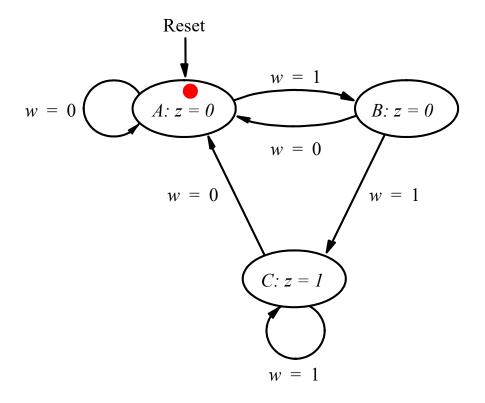
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



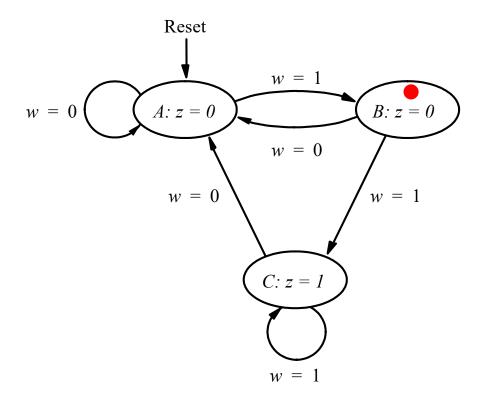
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
											1
z:	0	0	0	0	0	1	0	0	1	1	0



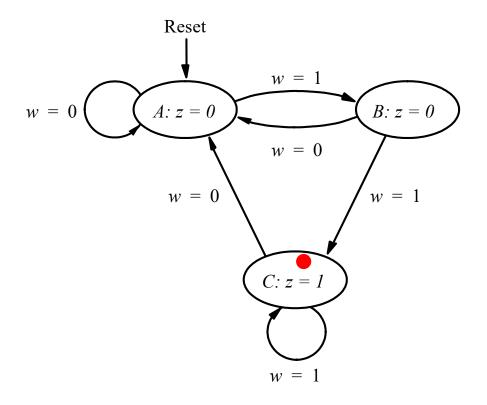
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
Clockcycle: w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



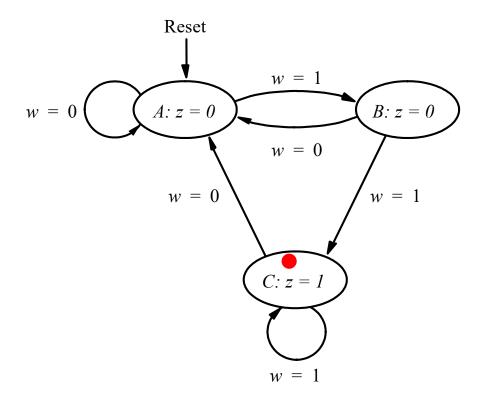
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	t <sub>6</sub>	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



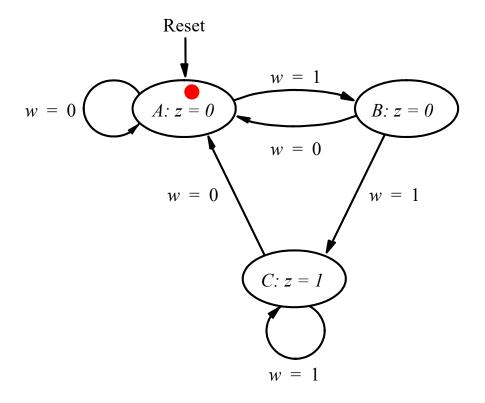
Clockcycle: w:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



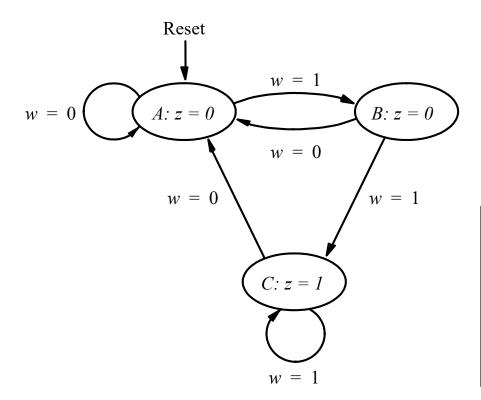
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	t <sub>10</sub>
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0



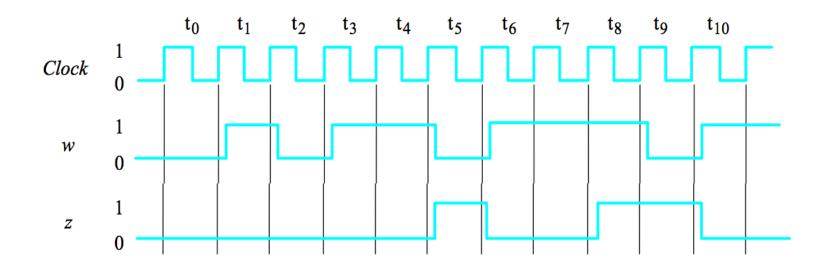
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	<b>t</b> 9	t <sub>10</sub>
w:	0	1	0	1	1	0	1	1	1	0	1
z:	0	0	0	0	0	1	0	0	1	1	0



In general, we need to start tracing from the beginning to know which state the FSM is in. It may not be clear from a short sequence of outputs.

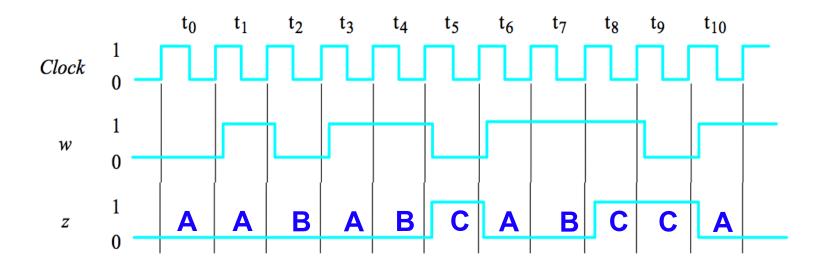
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
z:	0	0	0	0	0	1	0	0	1	1	0

## **Inferring the States**

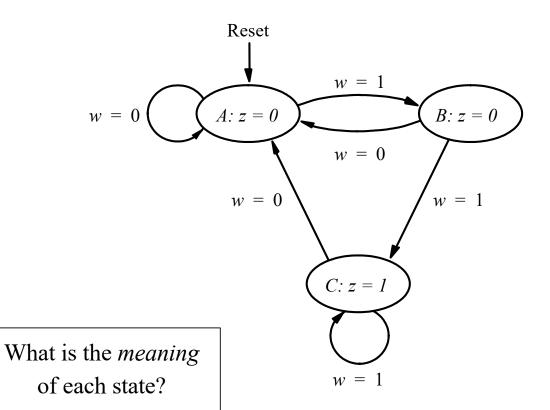


Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	t <sub>6</sub>	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0

## **Inferring the States**



Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
z:	0	0	0	0	0	1	0	0	1	1	0



Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0

### What is a State?

It is not really a memory of every past input (We might run out of space to remember it all!)

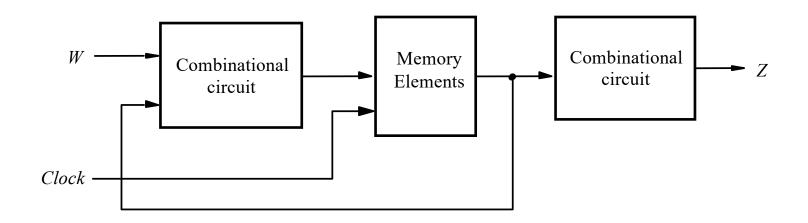
Rather, it is a characterization or snapshot of the pattern of inputs that have come before.

## **Moore Machine Implementation**

The state diagram is just an illustration to help us describe and reason about how the FSM will behave in each of its states.

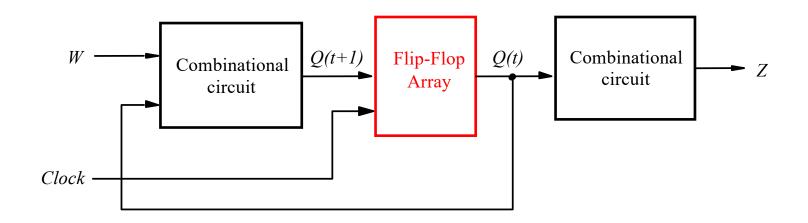
So, how do we turn it into a circuit?

## **Moore Machine Implementation**



Note: The W and Z lines need not be wires. They can be buses.

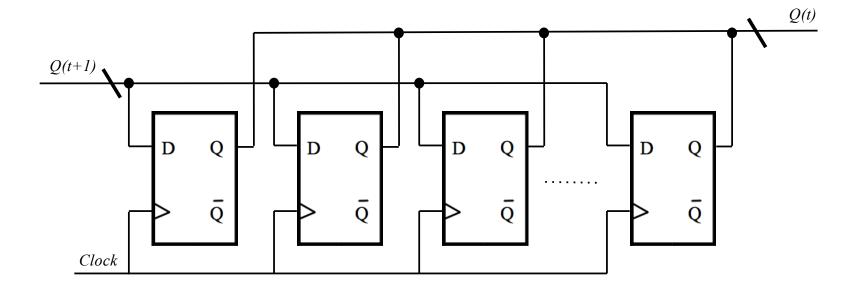
## **State Storage**



Any usable "memory" of the preceding input sequence is encoded in the flip-flop array.

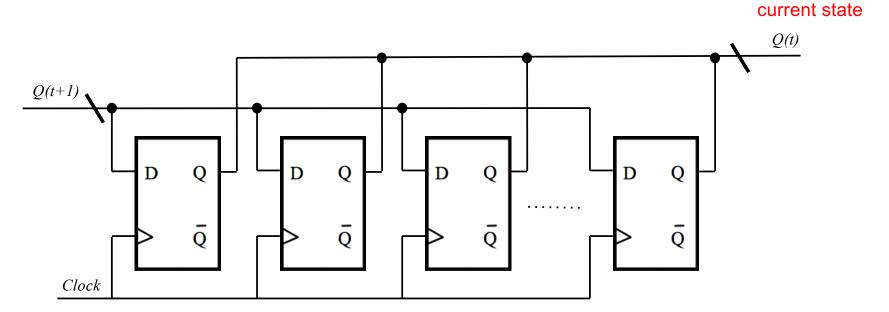
#### **FSM States**

The Flip-Flop array stores an encoding of the current state.



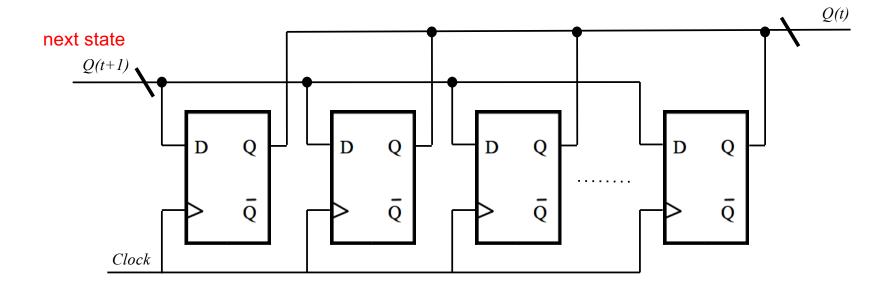
### **FSM States**

The Flip-Flop array stores an encoding of the current state.

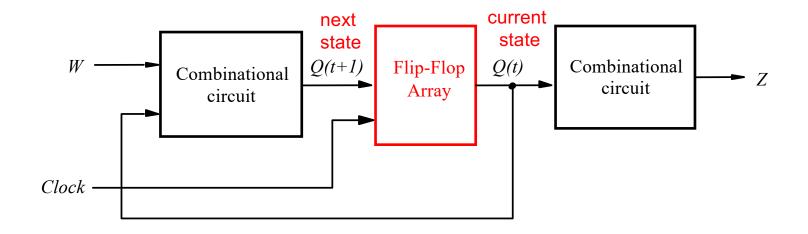


#### **FSM States**

The Flip-Flop array stores an encoding of the current state.



### **State Storage**



Any usable "memory" of the preceding input sequence is encoded in the flip-flop array.

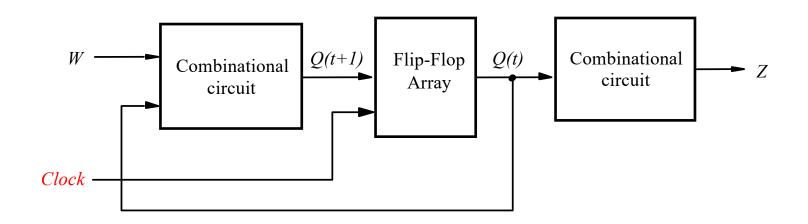
# **State Encoding**

Each of the states in our design is identified by a distinct code.

If we use 3 flip-flops, then the FSM can have up to  $2^3 = 8$  distinct states.

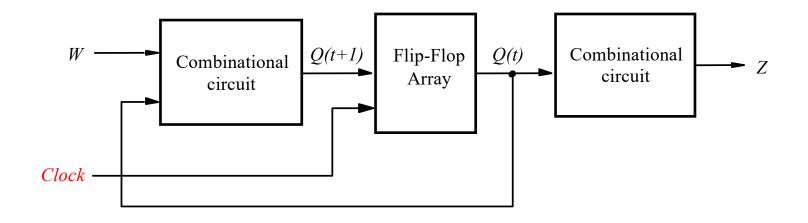
So, when the flip-flop array contains the code 011, we say that the machine is in state 011.

# **Synchronous Design**



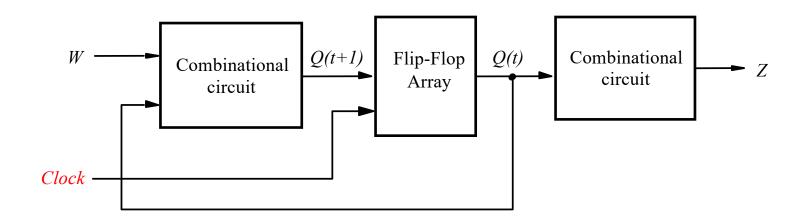
Every active clock edge causes a state transition.

# Synchronous Design



We expect the input signals to be stable before the *active clock edge* occurs.

# Synchronous Design



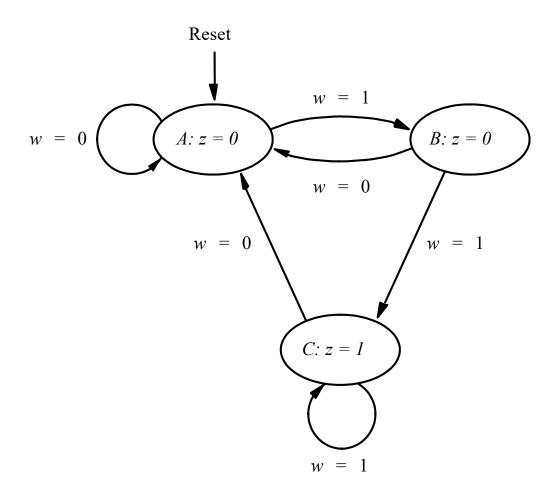
There is a whole other class of sequential circuits that are asynchronous, but we will not study them in this course.

#### Sequential Circuits: Key Ideas

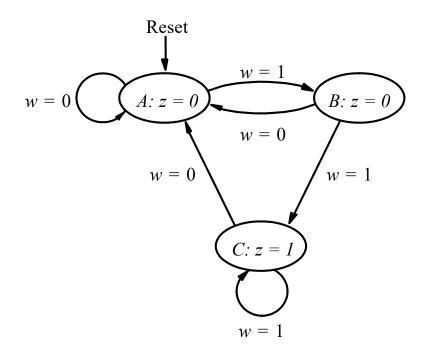
The current output depends on something about the preceding sequence of inputs (and maybe the current output).

Using *memory elements* (i.e., flip-flops), we design the circuit to remember some *relevant* information about the prior inputs.

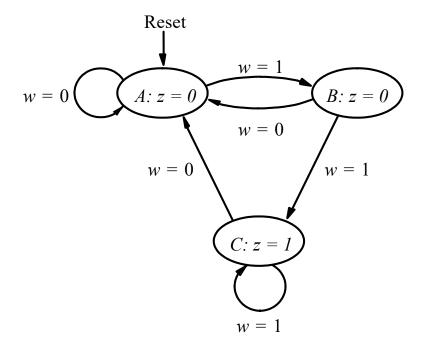
# **Moore Machine Example**



We need to find both the *next state logic* and the *output logic* implied by this machine.



Present	Next	Output	
state	w = 0	w = 1	Z
A			
В			
C			



Present	Next	Output	
state	w = 0	w = 1	$\overline{z}$
A	A	В	0
В	A	C	0
C	A	C	1

Figure 6.4 from the textbook ]

#### **How to represent the States?**

One way is to encode each state with a 2-bit binary number

 $A \sim 00$ 

B ~ 01

C ~ 10

#### How to represent the states?

One way is to encode each state with a 2-bit binary number

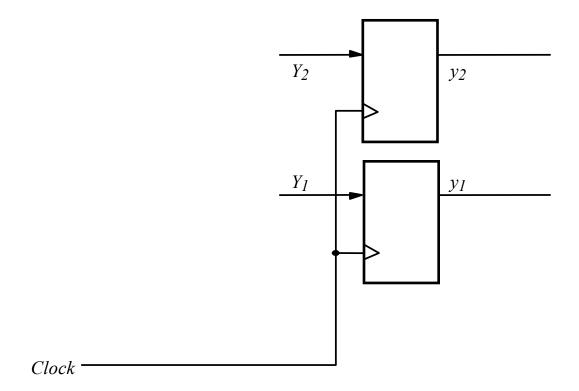
 $A \sim 00$ 

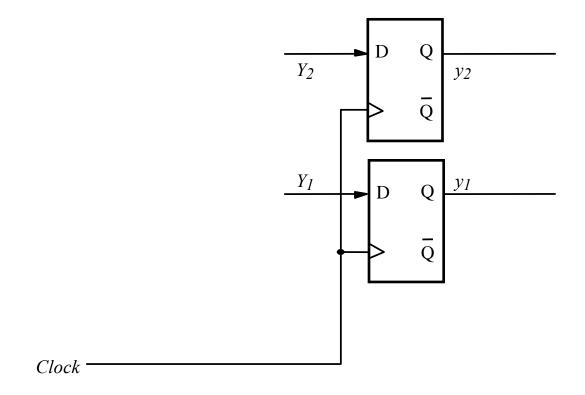
B ~ 01

C ~ 10

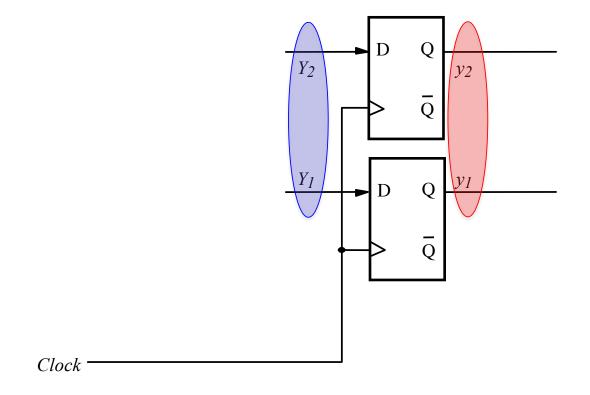
How many flip-flops do we need?

# Let's use two flip-flops to hold the machine's state





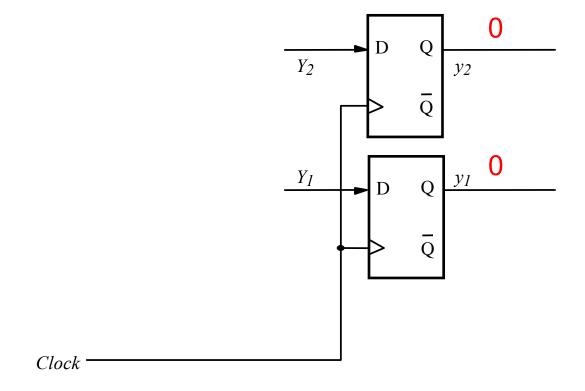
Let's pick D Flip-Flops.



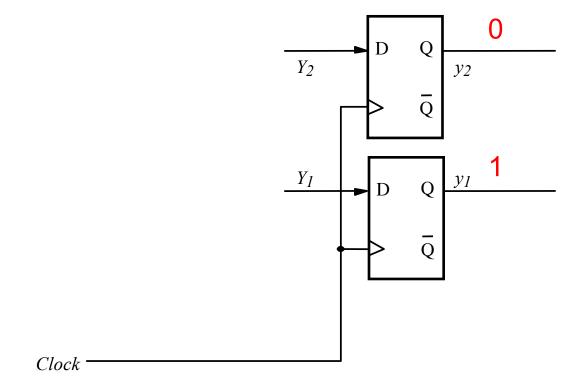
We will call  $y_1$  and  $y_2$  the present state variables.

We will call  $Y_1$  and  $Y_2$  the *next state variables*.

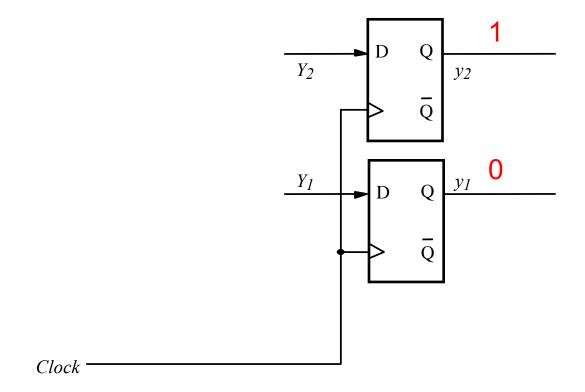
[ Figure 6.5 from the textbook ]



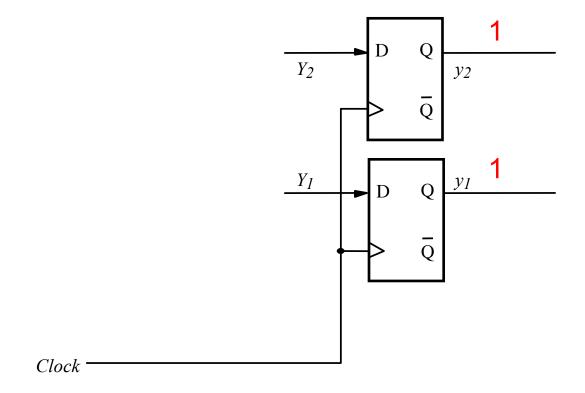
Two zeros on the output JOINTLY represent state A.



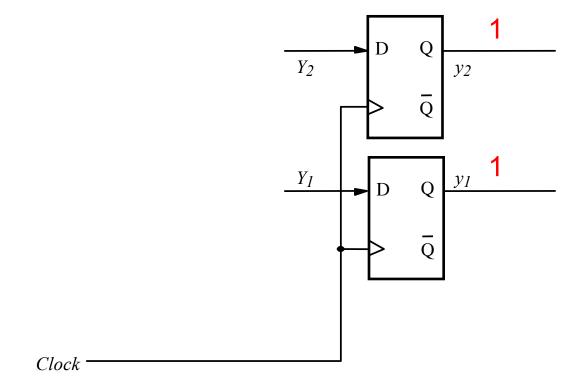
This flip-flop output pattern represents state B.



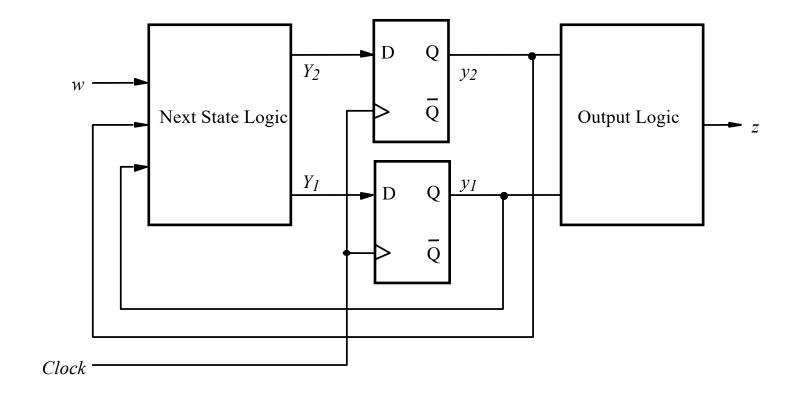
This flip-flop output pattern represents state C.



What does this flip-flop output pattern represent?



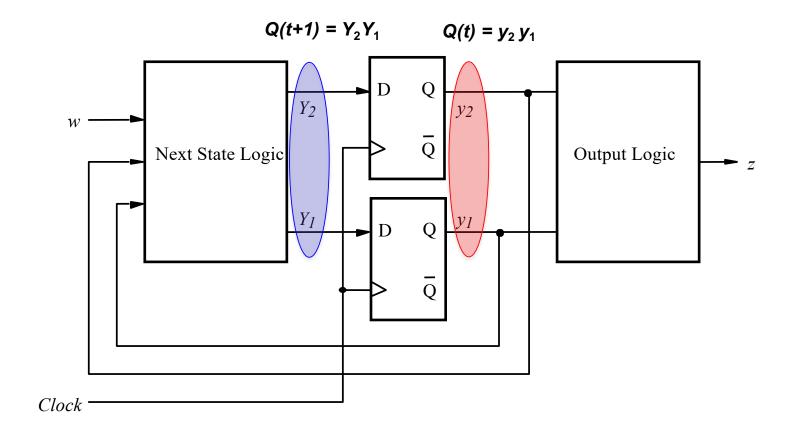
This would be state D, but we don't have one in this example. So, this is an impossible state.



We will call  $y_1$  and  $y_2$  the present state variables.

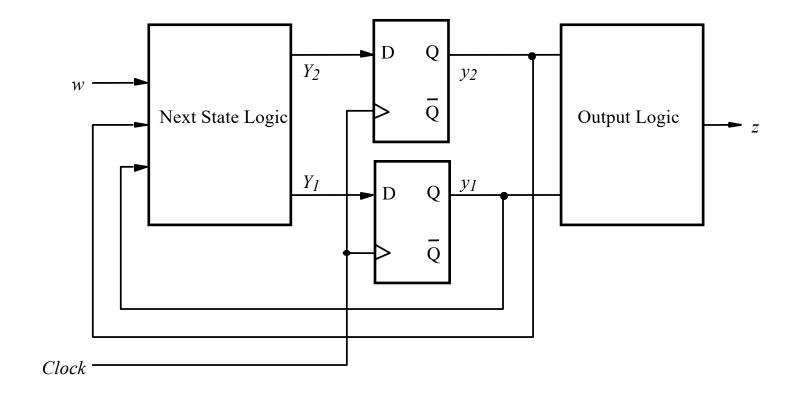
We will call  $Y_1$  and  $Y_2$  the *next state variables*.

[ Figure 6.5 from the textbook ]

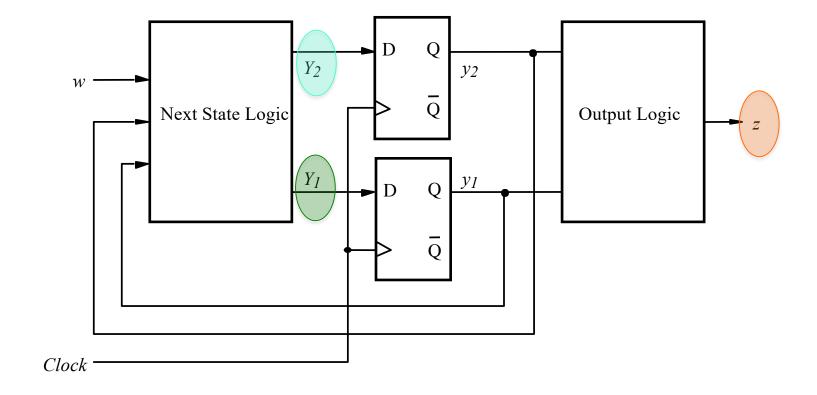


We will call  $y_1$  and  $y_2$  the present state variables.

We will call  $Y_1$  and  $Y_2$  the *next state variables*.



We need to find logic expressions for  $Y_1(w, y_1, y_2)$ ,  $Y_2(w, y_1, y_2)$ , and  $z(y_1, y_2)$ .



We need to find logic expressions for  $Y_1(w, y_1, y_2)$ ,  $Y_2(w, y_1, y_2)$ , and  $z(y_1, y_2)$ .

Present	Next	Output	
state	w = 0	w = 1	$\overline{z}$
A	A	В	0
В	A	C	0
C	A	C	1

# Suppose we encoded our states in the same order in which they were labeled:

A ~ 00

B ~ 01

C ~ 10

Present	Next	Output	
state	w = 0	w = 1	z
A	A	В	0
В	A	$\mathbf{C}$	0
C	A	C	1

	Present	Next state		
	state	w = 0 $w = 1$	Output	
			z	
A	00			
В	01			
C	10			
	11			

The finite state machine will never reach a state encoded as 11.

Figure 6.6 from the textbook ]

Present	Next	Output	
state	w = 0	w = 1	Z
A	A	В	0
В	A	$\mathbf{C}$	0
C	A	C	1

	Present	Next state		
	state	w = 0	w = 1	Output
	<i>y</i> 2 <i>y</i> 1	$Y_2Y_1$	$Y_2Y_1$	Z
4	00	00	01	0
В	01	00	10	0
С	10	00	10	1
	11	dd	dd	d

We arbitrarily chose these as our state encodings.
We could have used others.

Figure 6.6 from the textbook ]

Present	Next s		
state	w = 0 $w = 1$		Output
<i>y</i> <sub>2</sub> <i>y</i> <sub>1</sub>	$Y_2Y_1$	$y_2 y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

$Q(t) = y_2 y_1$ and	Q(t+1) =	$Y_2Y_1$
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w	<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	$Y_2$	$Y_I$
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	Z
0	0	
0	1	
1	0	
1	1	

[ Figure 6.6 from the textbook ]

Present	Next s		
state	w = 0	w = 1	Output
<i>y</i> 2 <i>y</i> 1	$Y_2Y_1$	$Y_2Y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

$Q(t) = y_2 y_1$ and $Q$	$t(t+1) = Y_2Y$	1
--------------------------	-----------------	---

w	<i>y</i> <sub>2</sub>	<i>y</i> 1	$Y_2$	$Y_{I}$
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

<i>y</i> <sub>2</sub>	<i>y</i> 1	Z
0	0	0
0	1	0
1	0	1
1	1	d

[ Figure 6.6 from the textbook ]

Present	Next s	tate	
state	w = 0	w = 1	Output
<i>y</i> 2 <i>y</i> 1	$Y_2Y_1$	$Y_2Y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

$Q(t) = y_2 y_1$ and $Q$	$t(t+1) = Y_2Y$	1
--------------------------	-----------------	---

w	<i>y</i> <sub>2</sub>	<i>y</i> 1	$Y_2$	$Y_{I}$
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	Z
0	0	0
0	1	0
1	0	1
1	1	d

[ Figure 6.6 from the textbook ]

Present	Next s	tate	
state	w = 0	w = 1	Output
<i>y</i> 2 <i>y</i> 1	$Y_2Y_1$	$y_2y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

$Q(t) = y_2y_1 \text{ and } Q(t+1) = Y_2Y_1$
--

w	<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	$Y_2$	$Y_I$
0	0	0	0	
0	0	1	0	
0	1	0	0	
0	1	1	d	
1	0	0		
1	0	1		
1	1	0		
1	1	1		

<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	Z
0	0	0
0	1	0
1	0	1
1	1	d

[ Figure 6.6 from the textbook ]

Present	Next state		
state	w = 0 $w = 1$		Output
<i>y</i> 2 <i>y</i> 1	<sup>Y</sup> 2 <sup>Y</sup> 1	$y_2y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

$Q(t) = y_2 y_1$ and $Q$	$t(t+1) = Y_2Y_1$	1
--------------------------	-------------------	---

w	<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	$Y_2$	$Y_{I}$
0	0	0	0	
0	0	1	0	
0	1	0	0	
0	1	1	d	
1	0	0	0	
1	0	1	1	
1	1	0	1	
1	1	1	d	

<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	Z
0	0	0
0	1	0
1	0	1
1	1	d

[ Figure 6.6 from the textbook ]

Present	Next s	tate	
state	w = 0	w = 1	Output
<i>y</i> 2 <i>y</i> 1	$Y_2Y_1$	$Y_2Y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

$Q(t) = y_2 y_1$ and $Q$	$t(t+1) = Y_2Y_1$	1
--------------------------	-------------------	---

w	<i>y</i> <sub>2</sub>	<i>y</i> 1	$Y_2$	$Y_I$
0	0	0	0	
0	0	1	0	
0	1	0	0	
0	1	1	d	
1	0	0	0	
1	0	1	1	
1	1	0	1	
1	1	1	d	

<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	Z
0	0	0
0	1	0
1	0	1
1	1	d

[ Figure 6.6 from the textbook ]

Present	Next s	tate	
state	w = 0	w = 1	Output
<i>y</i> 2 <i>y</i> 1	$Y_2Y_1$	$Y_2Y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

$Q(t) = y_2y_1 \text{ and } Q(t+1) = Y_2Y_1$	Q(t)	$= y_2y_1$	and Q	(t+1) =	$Y_2Y_1$
--	------	------------	-------	---------	----------

w	<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	$Y_2$	$Y_{I}$
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	d	d
1	0	0	0	
1	0	1	1	
1	1	0	1	
1	1	1	d	

<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	Z
0	0	0
0	1	0
1	0	1
1	1	d

[ Figure 6.6 from the textbook ]

Present	Next s		
state	w = 0 $w = 1$		Output
<i>y</i> 2 <i>y</i> 1	$y_2y_1$	$y_2y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

$Q(t) = y_2 y_1$ and $Q$	$t(t+1) = Y_2Y$	1
--------------------------	-----------------	---

w	<i>y</i> <sub>2</sub>	<i>y</i> 1	$Y_2$	$Y_{I}$	
0	0	0	0	0	
0	0	1	0	0	
0	1	0	0	0	
0	1	1	d	d	
1	0	0	0	1	
1	0	1	1	0	
1	1	0	1	0	
1	1	1	d	d	

<i>y</i> <sub>2</sub>	<i>y</i> 1	Z
0	0	0
0	1	0
1	0	1
1	1	d

[ Figure 6.6 from the textbook ]

Present	Next s	tate	
state	w = 0 $w = 1$		Output
<i>y</i> 2 <i>y</i> 1	$Y_2Y_1$	$Y_2Y_1$	Z
00	00	01	0
01	00	10	0
10	00	10	1
11	dd	dd	d

$Q(t) = y_2 y_1$ and $Q$	$t(t+1) = Y_2Y$	1
--------------------------	-----------------	---

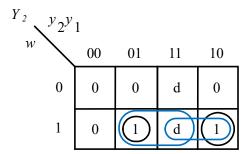
w	<i>y</i> <sub>2</sub>	<i>y</i> 1	$Y_2$	$Y_{I}$
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	d	d
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	d	d

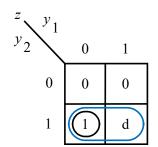
<i>y</i> <sub>2</sub>	<i>y</i> <sub>1</sub>	Z
0	0	0
0	1	0
1	0	1
1	1	d

[ Figure 6.6 from the textbook ]

Note that the textbook draws these K-Maps differently from all previous K-maps (the most significant bit indexes the rows).

$Y_I$	<i>y</i> <sub>2</sub> <i>y</i>	1			
w		00	01	11	10
	0	0	0	d	0
	1	1	0	d	0



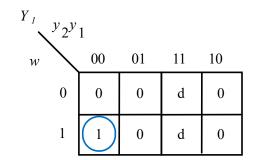


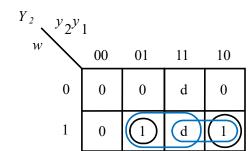
$$Q(t) = y_2y_1 \text{ and } Q(t+1) = Y_2Y_1$$

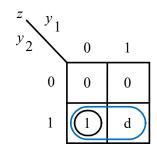
w	<i>y</i> <sub>2</sub>	<i>y</i> 1	$Y_2$	$Y_I$
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	d	d
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	d	d

<i>y</i> <sub>2</sub>	<i>y</i> 1	Z
0	0	0
0	1	0
1	0	1
1	1	d

#### Don't care conditions simplify the combinatorial logic







Ignoring don't cares

$$Y_1 = w\overline{y}_1\overline{y}_2$$

$$Y_2 = wy_1\overline{y_2} + w\overline{y_1}y_2$$

$$z = \overline{y_1}y_2$$

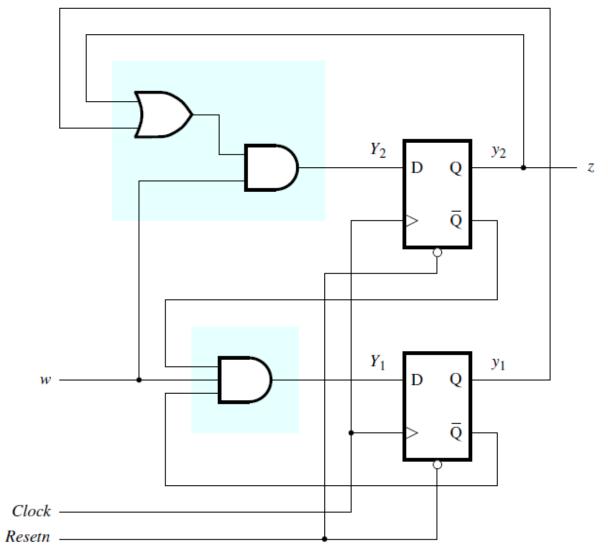
Using don't cares

$$Y_1 = w\overline{y}_1\overline{y}_2$$

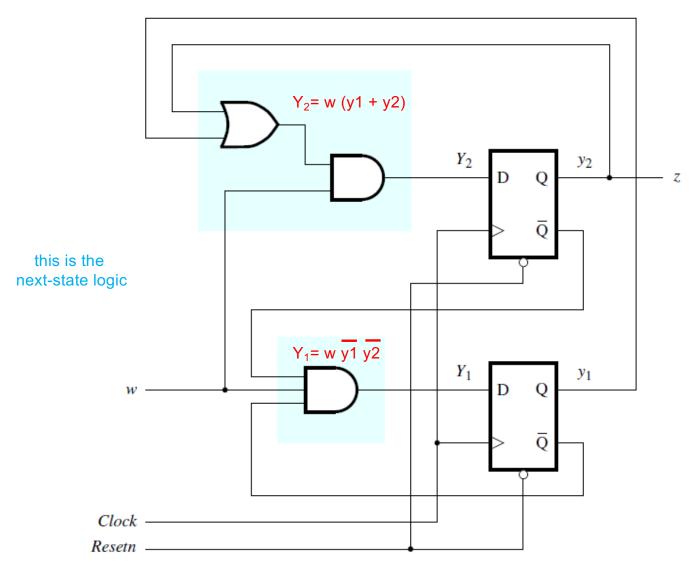
$$Y_2 = wy_1 + wy_2$$
  
=  $w(y_1 + y_2)$ 

$$z = y_2$$

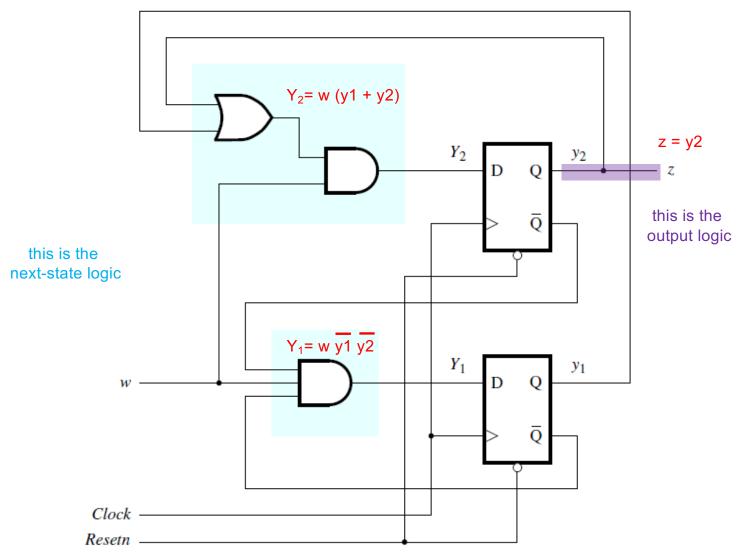
[ Figure 6.7 from the textbook ]



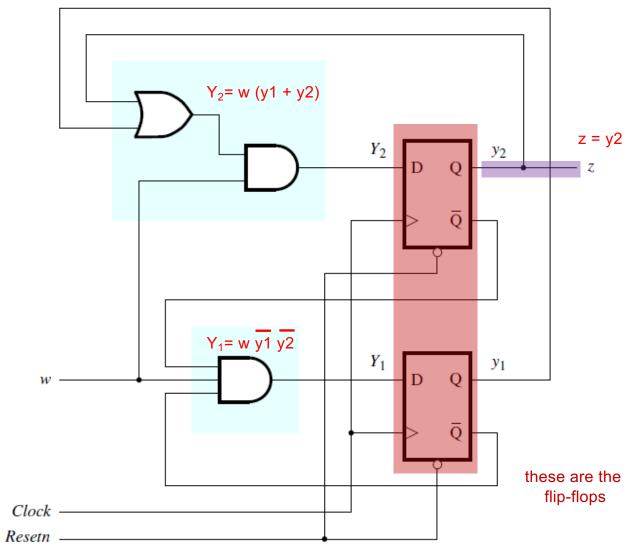
[ Figure 6.8 from the textbook ]



[ Figure 6.8 from the textbook ]

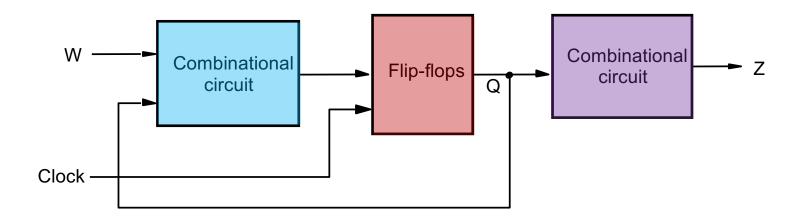


[ Figure 6.8 from the textbook ]

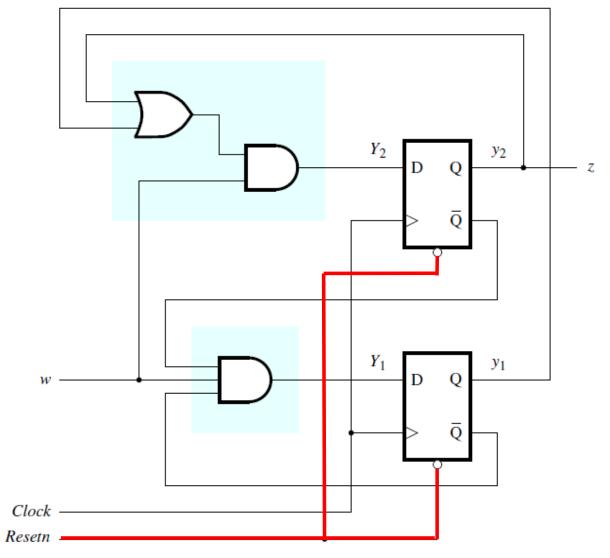


[ Figure 6.8 from the textbook ]

# **Moore Type**

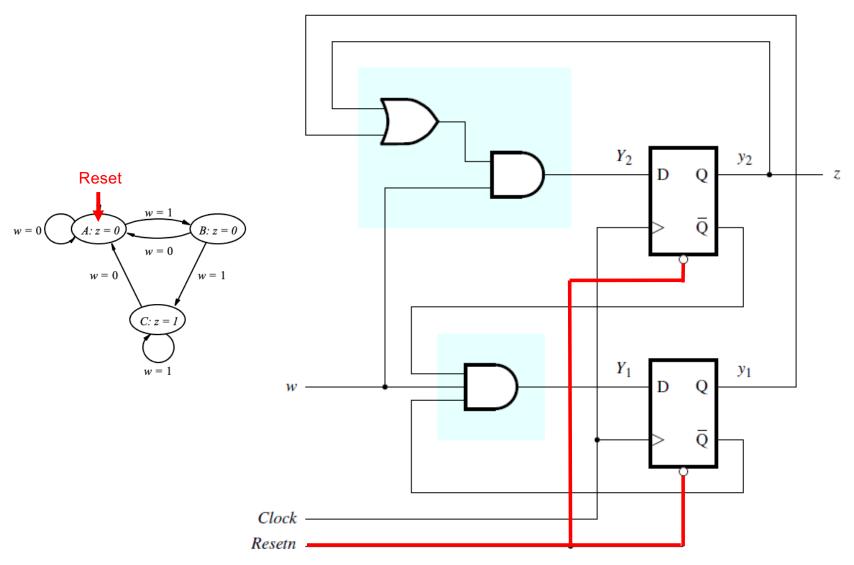


### Don't Forget to Add the Reset Line



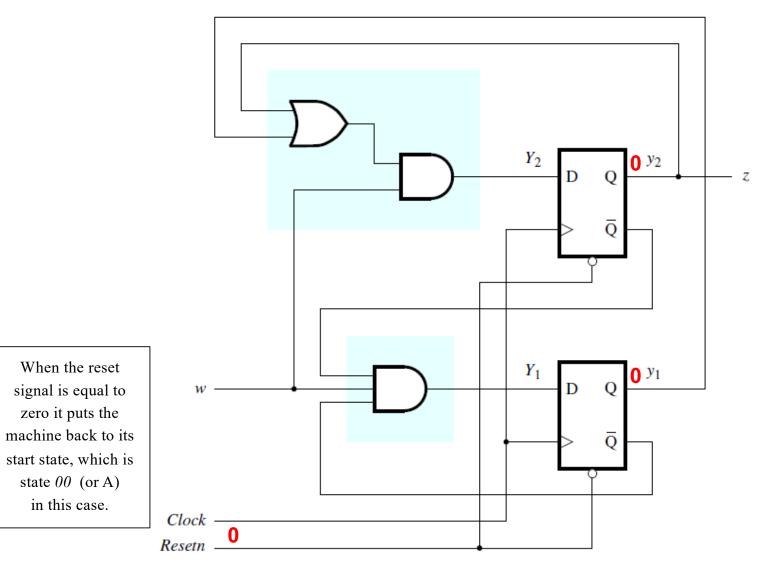
[ Figure 6.8 from the textbook ]

### Don't Forget to Add the Reset Line



[ Figure 6.8 from the textbook ]

#### State A=00



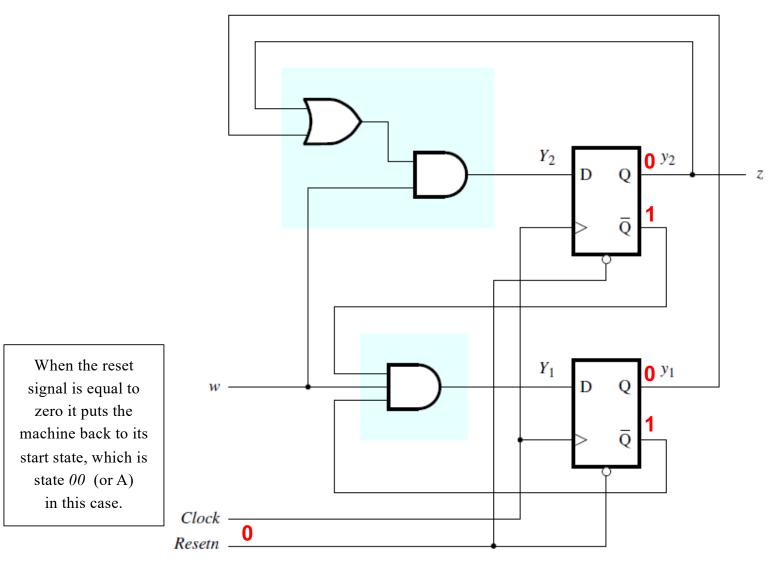
When the reset

signal is equal to zero it puts the

state  $\theta\theta$  (or A) in this case.

[ Figure 6.8 from the textbook ]

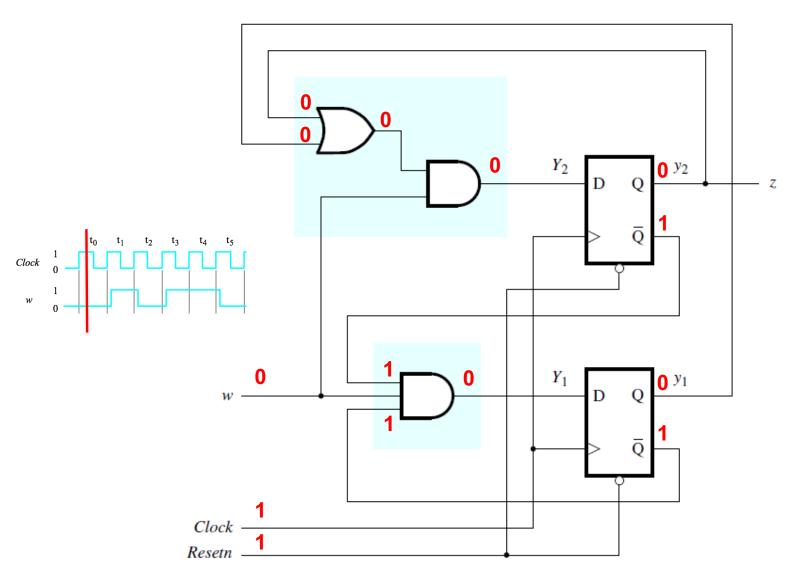
#### State A=00



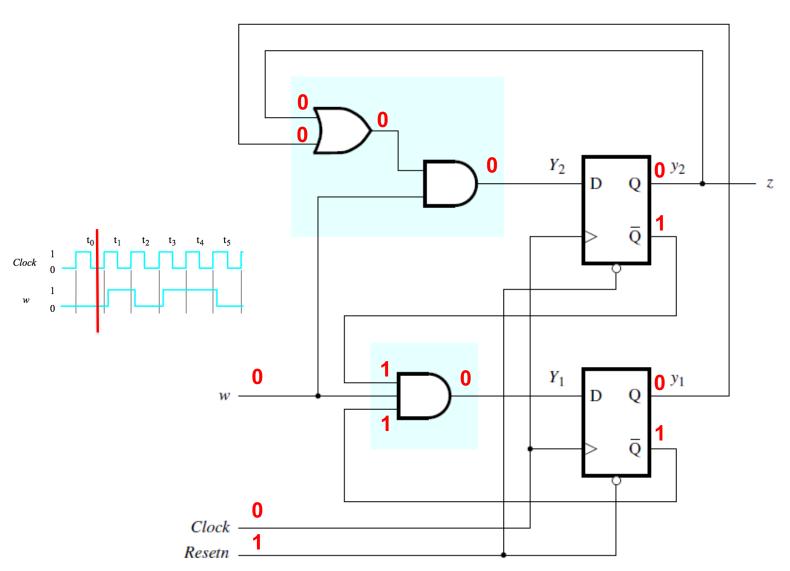
[ Figure 6.8 from the textbook ]

State =  $y_2 y_1$ 

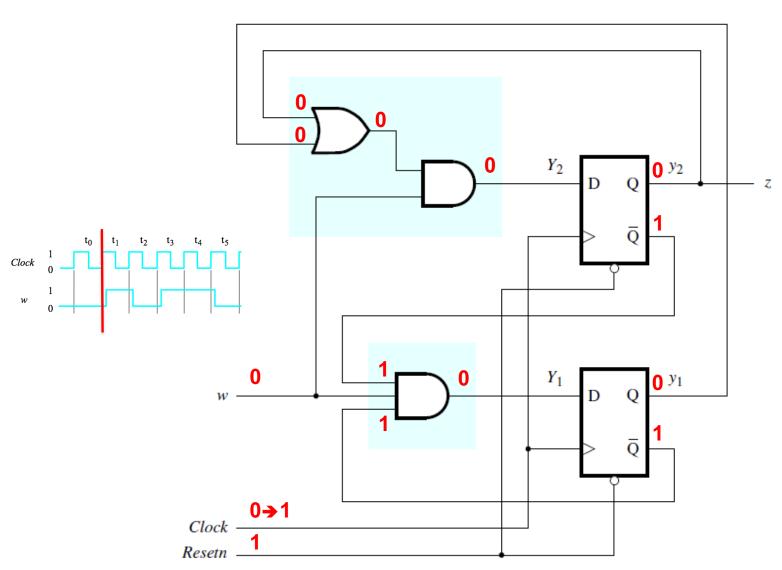
State A=00



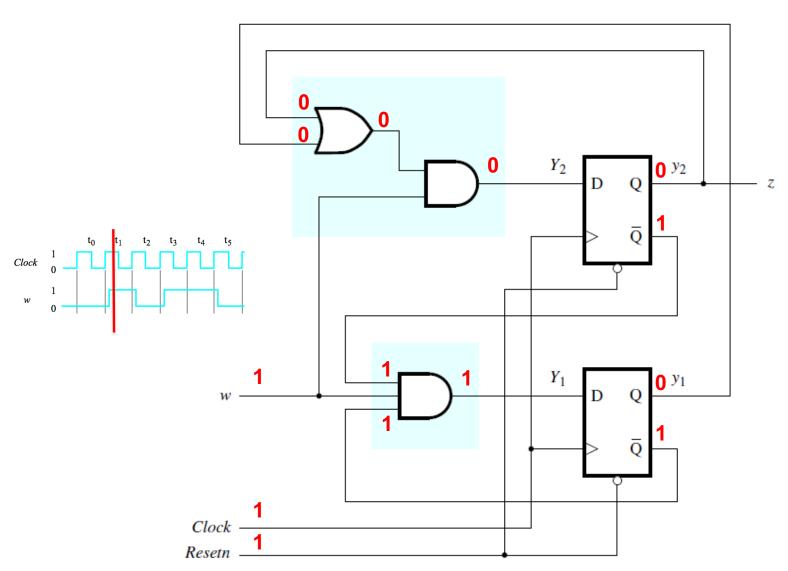
State A=00



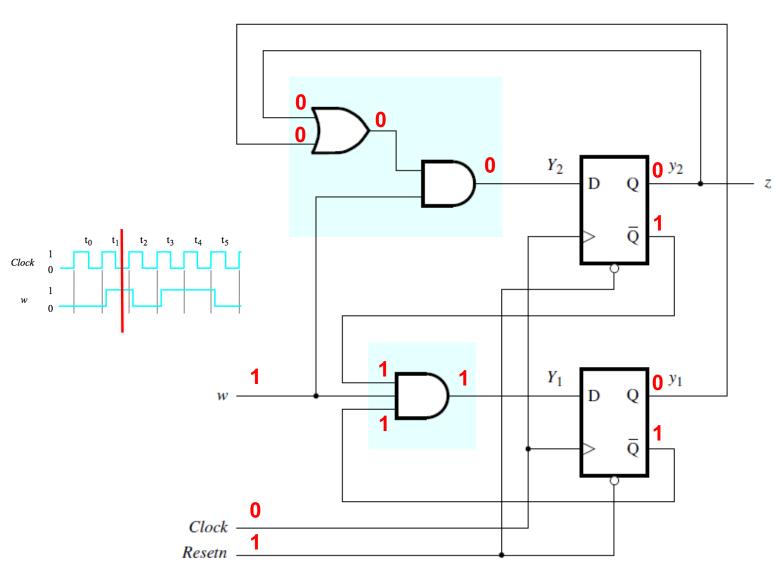
State A=00



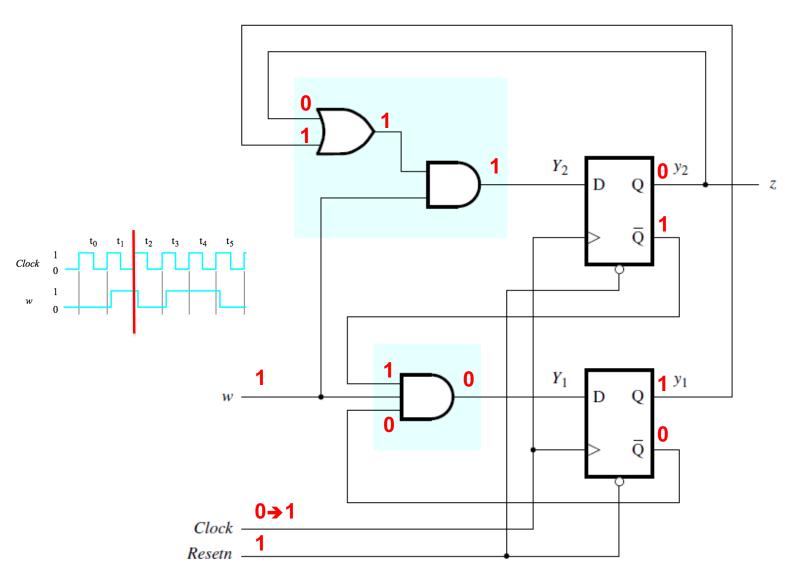
State A=00



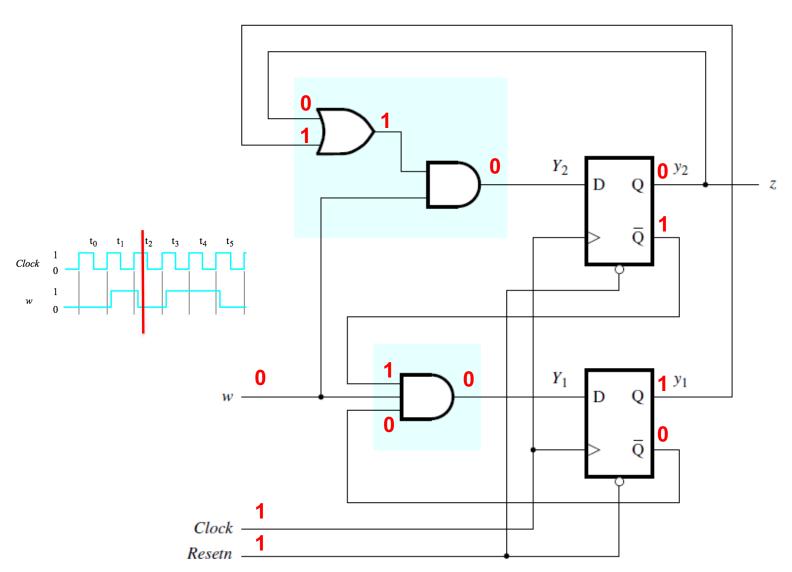
State A=00



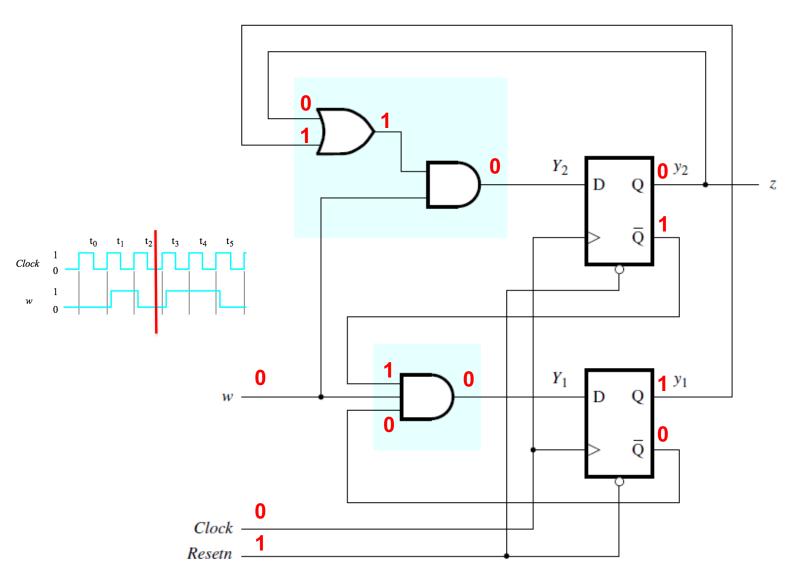
State B=01



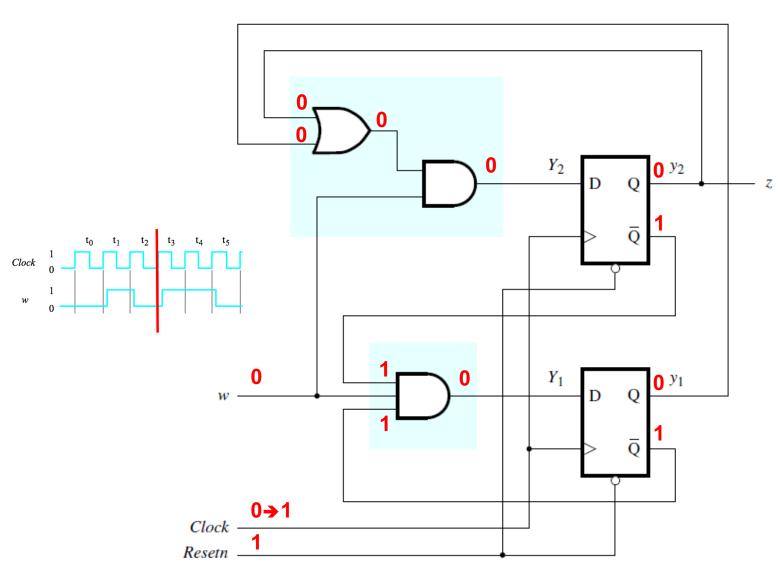
State B=01



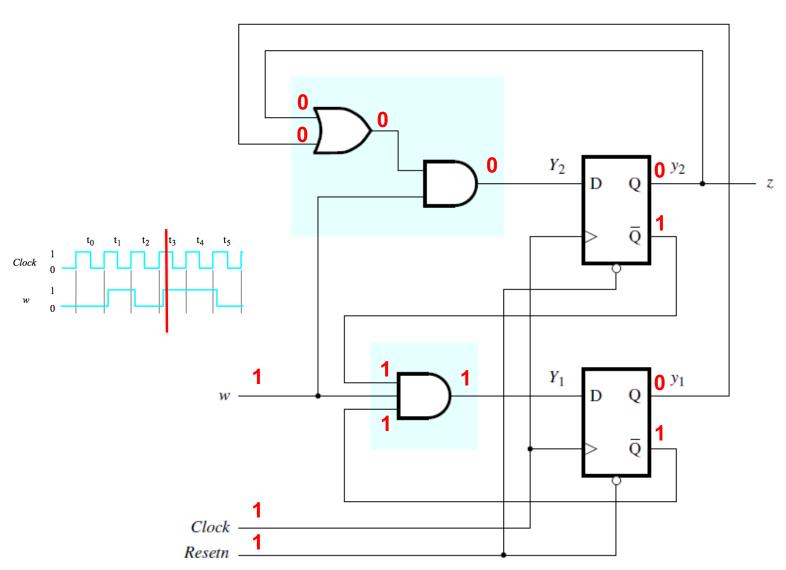
State B=01



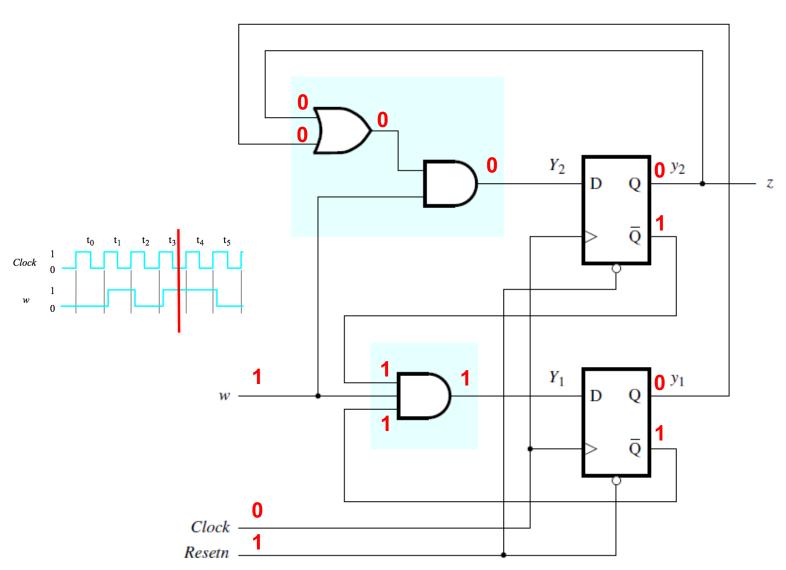
State A=00



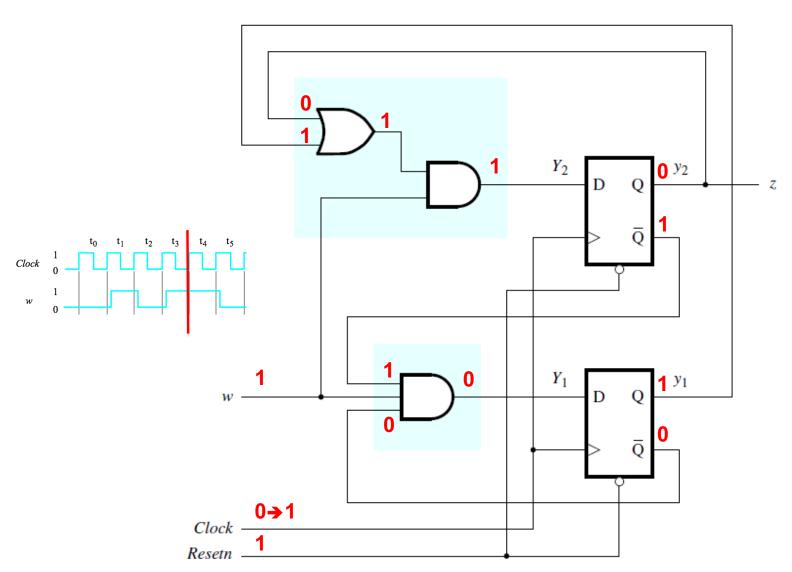
State A=00



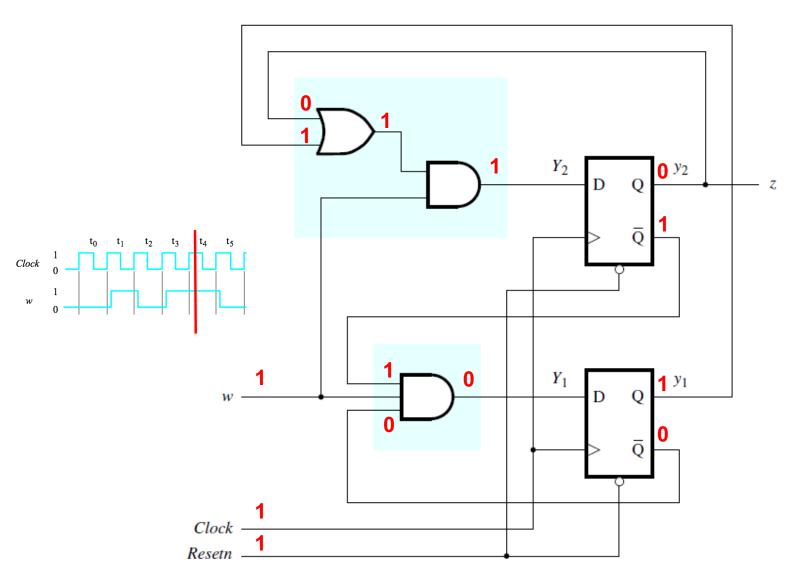
State A=00



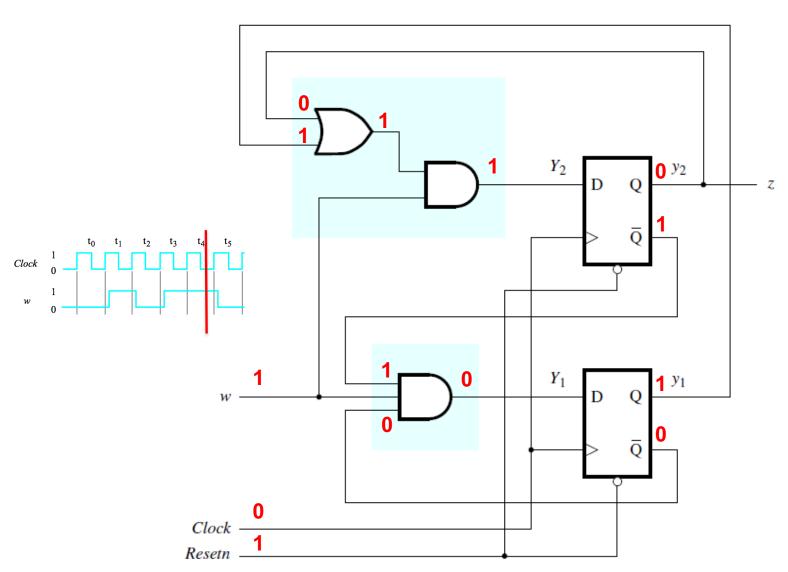
State B=01



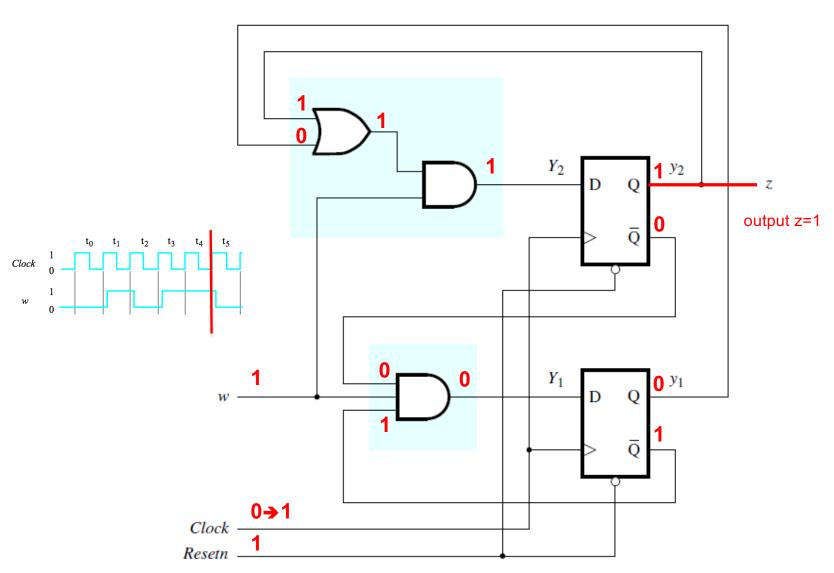
State B=01



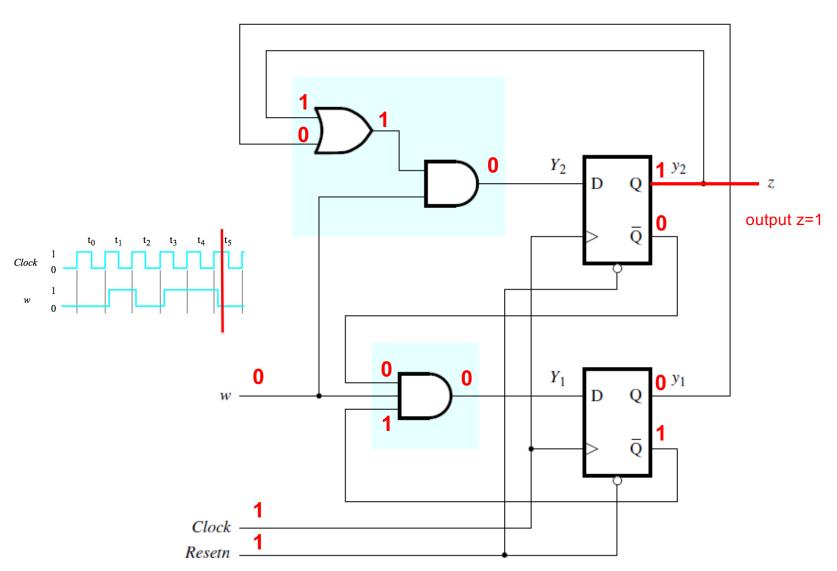
State B=01



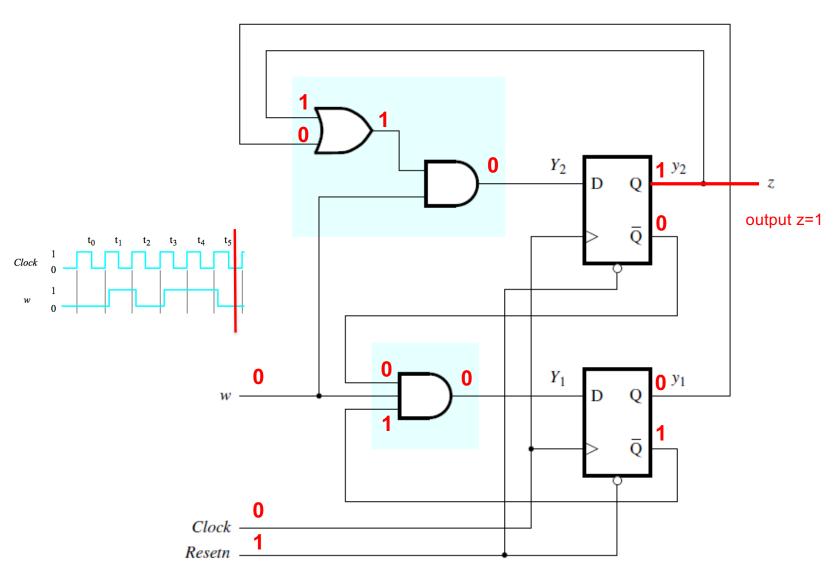
State C=10



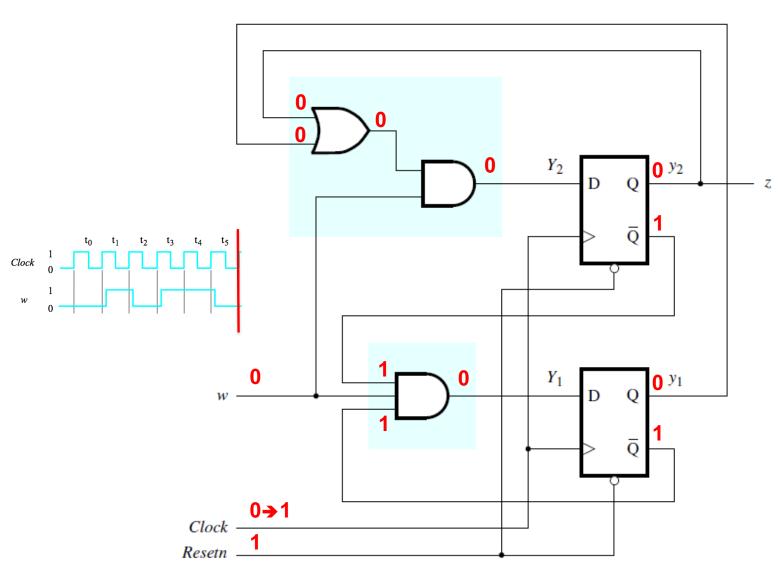
State C=10



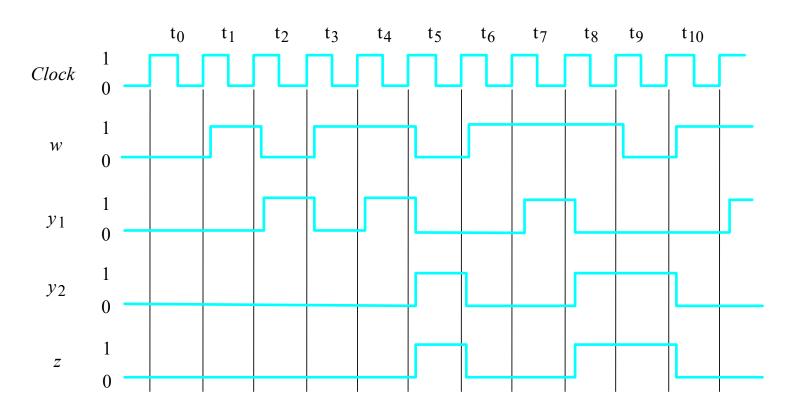
State C=10



State A=00

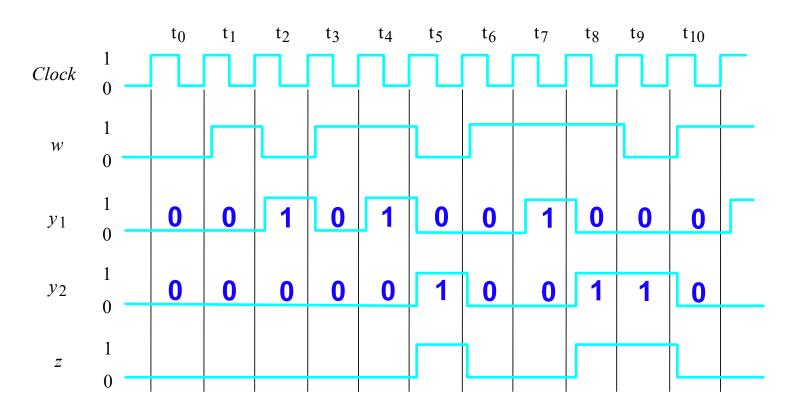


Clockcycle: w: z:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
<i>z</i> :	0	0	0	0	0	1	0	0	1	1	0

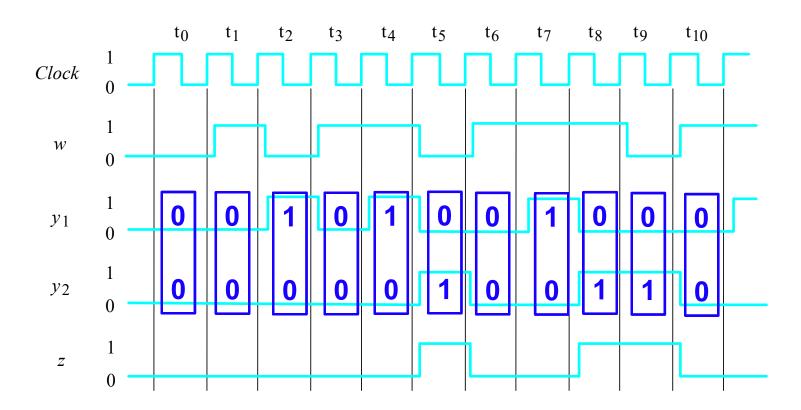


[ Figure 6.9 from the textbook ]

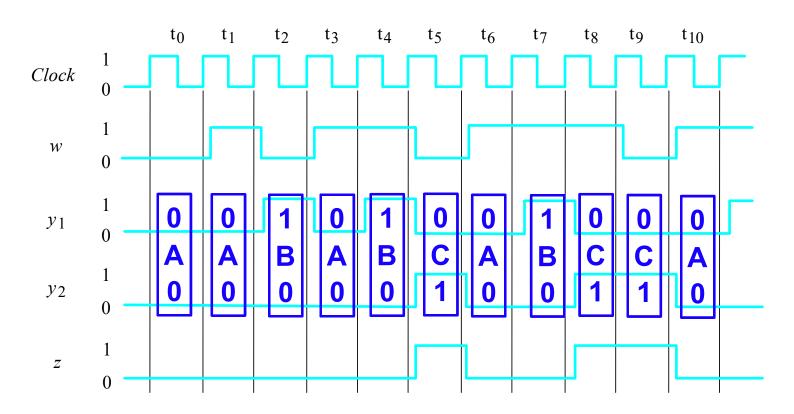
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
z:	0	0	0	0	0	1	0	0	1	1	0



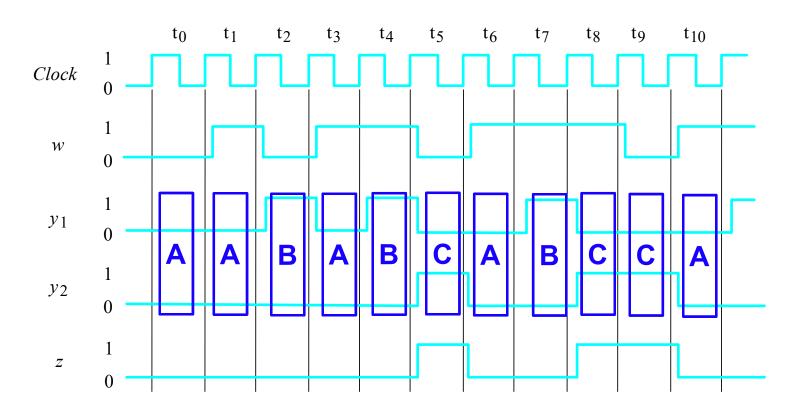
Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	t <sub>10</sub>
w:	0	1	0	1	1	0	1	1	1	0	1
z:	0	0	0	0	0	1	0	0	1	1	0



Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	t <sub>10</sub>
w:	0	1	0	1	1	0	1	1	1	0	1
z:	0	0	0	0	0	1	0	0	1	1	0



Clockcycle:	$t_0$	$t_1$	$t_2$	$t_3$	$t_4$	$t_5$	$t_6$	$t_7$	$t_8$	t <sub>9</sub>	$t_{10}$
w:	0	1	0	1	1	0	1	1	1	0	1
z:	0	0	0	0	0	1	0	0	1	1	0



## Summary: Designing a Moore Machine

- Obtain the circuit specification.
- Derive a state diagram.
- Derive the state table.
- Decide on a state encoding.
- Encode the state table.
- Derive the output logic and next-state logic.
- Draw the circuit diagram
- Add a reset signal.

**Questions?** 

