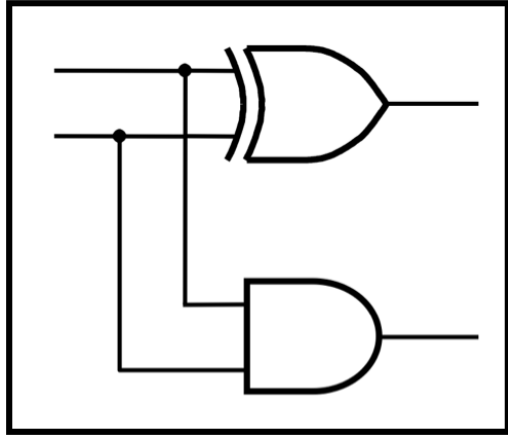




CprE 2810: Digital Logic

Instructor: Alexander Stoytchev

<http://www.ece.iastate.edu/~alexs/classes/>

Code Converters

CprE 2810: Digital Logic
Iowa State University, Ames, IA
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Administrative Stuff

- **HW 7 is out**
- **It is due on Monday (Oct 20) @ 10pm**
- **We will start with Chapter 5 on Friday.**

Quick Review

Decoders

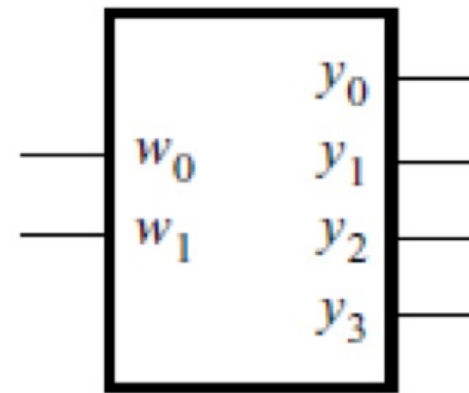
2-to-4 Decoder (Definition)

- Has two inputs: w_1 and w_0
- Has four outputs: y_0 , y_1 , y_2 , and y_3
- If $w_1=0$ and $w_0=0$, then the output y_0 is set to 1
- If $w_1=0$ and $w_0=1$, then the output y_1 is set to 1
- If $w_1=1$ and $w_0=0$, then the output y_2 is set to 1
- If $w_1=1$ and $w_0=1$, then the output y_3 is set to 1
- Only one output is set to 1. All others are set to 0.

Truth Table and Graphical Symbol for a 2-to-4 Decoder

w_1	w_0	y_0	y_1	y_2	y_3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

(a) Truth table



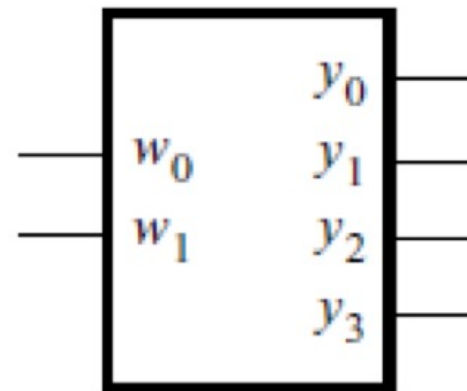
(b) Graphical symbol

Truth Table and Graphical Symbol for a 2-to-4 Decoder

w_1	w_0	y_0	y_1	y_2	y_3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

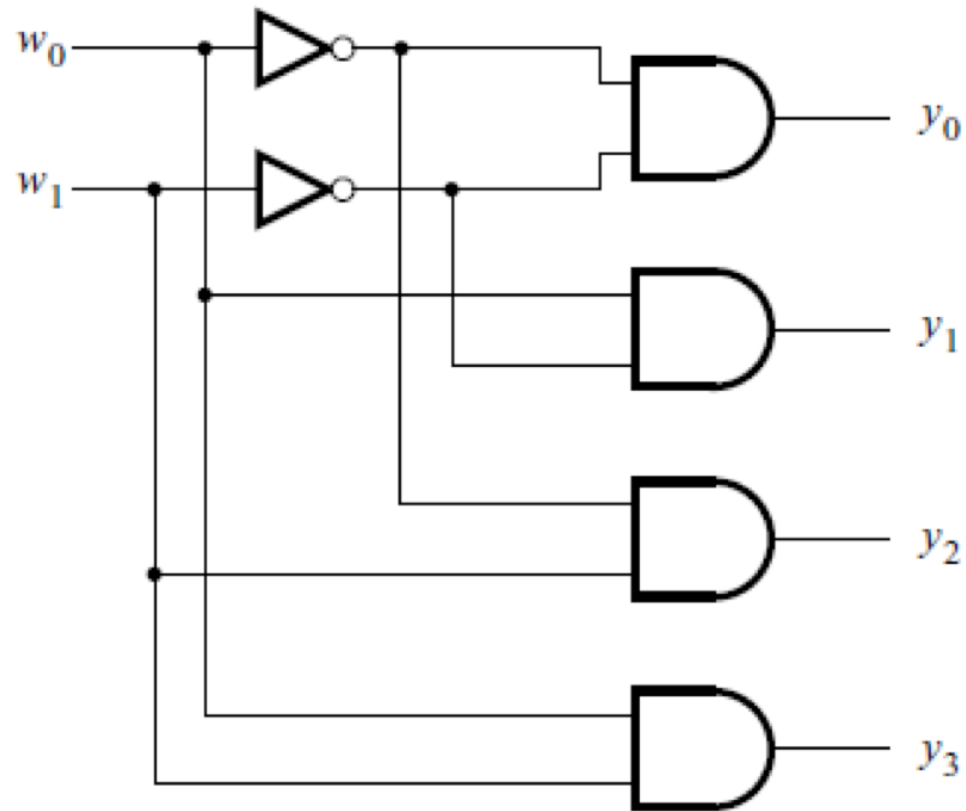
The outputs are “one-hot” encoded

(a) Truth table



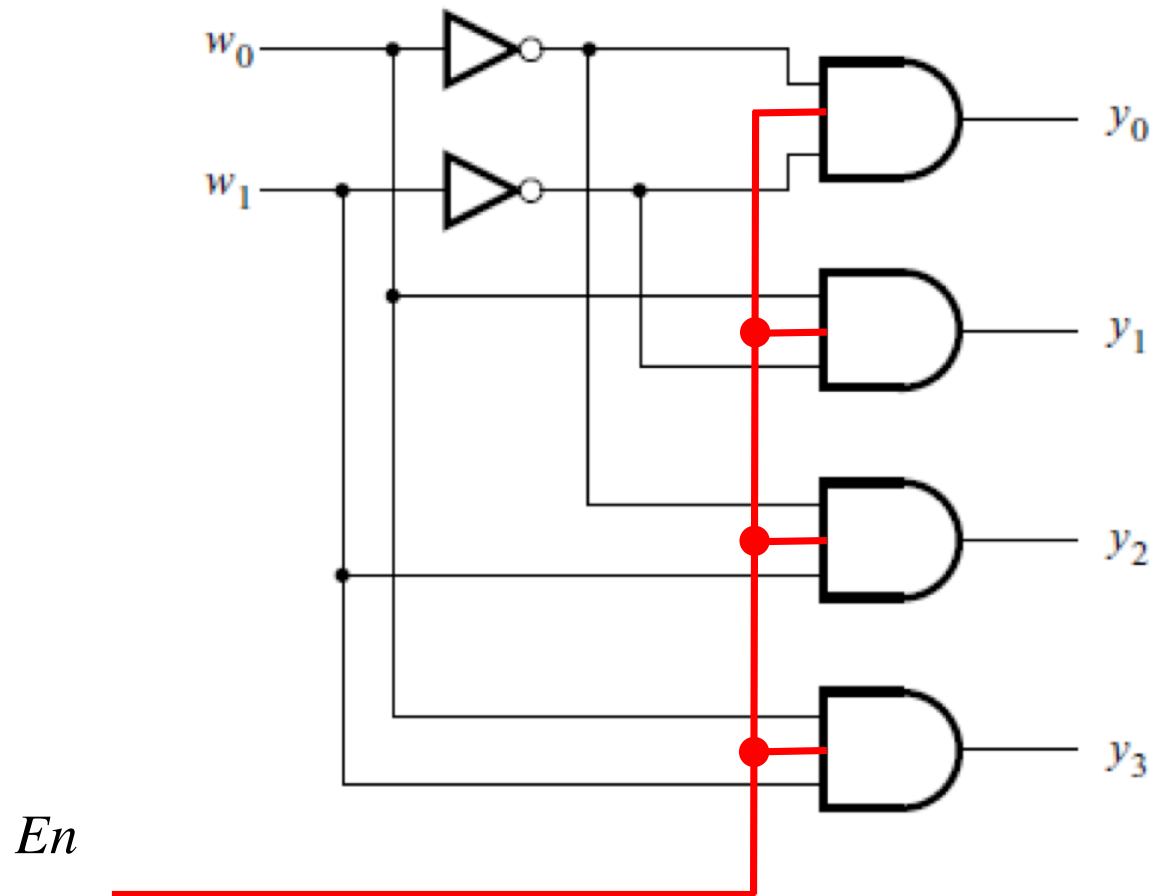
(b) Graphical symbol

Truth Logic Circuit for a 2-to-4 Decoder



[Figure 4.13c from the textbook]

Adding an Enable Input

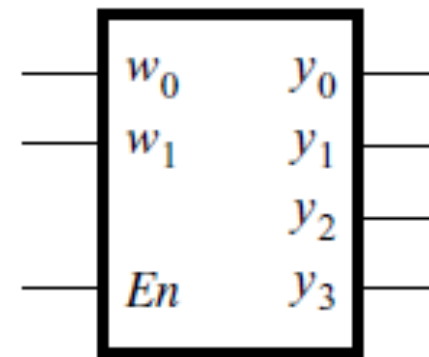


[Figure 4.13c from the textbook]

Truth Table and Graphical Symbol for a 2-to-4 Decoder with an Enable Input

En	w_1	w_0	y_0	y_1	y_2	y_3
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1
0	x	x	0	0	0	0

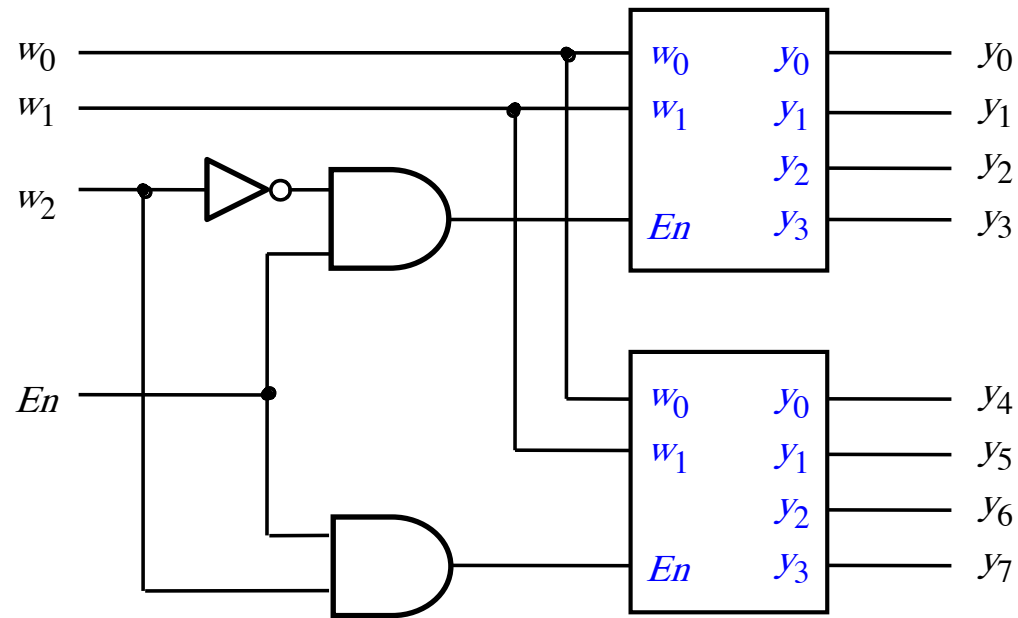
(a) Truth table



(b) Graphical symbol

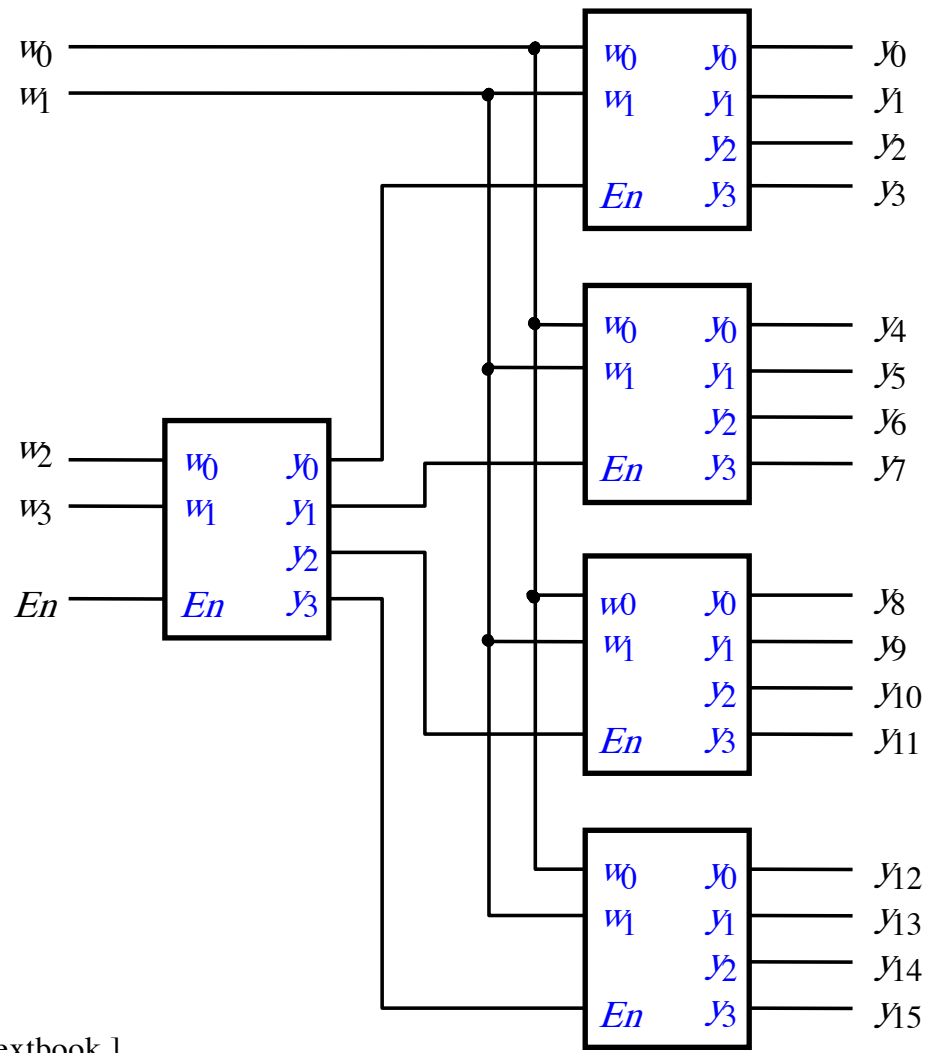
x indicates that it does not matter what the value of these variable is for this row of the truth table

A 3-to-8 decoder using two 2-to-4 decoders



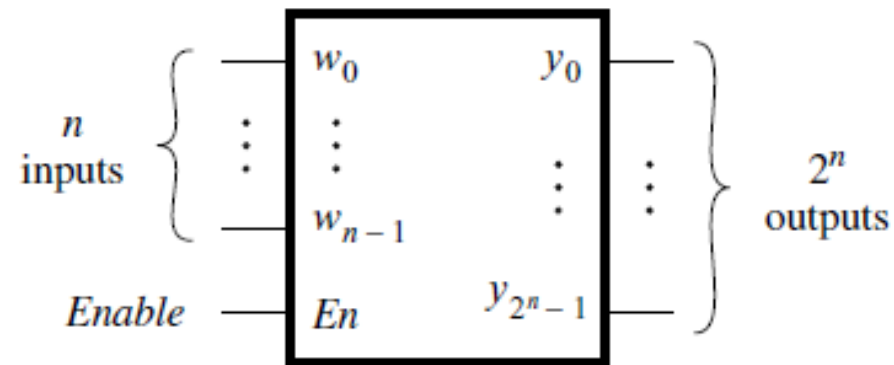
[Figure 4.15 from the textbook]

A 4-to-16 decoder built using a decoder tree



[Figure 4.16 from the textbook]

Graphical Symbol for a Binary n -to- 2^n Decoder with an Enable Input



(d) An n -to- 2^n decoder

A binary decoder with n inputs has 2^n outputs

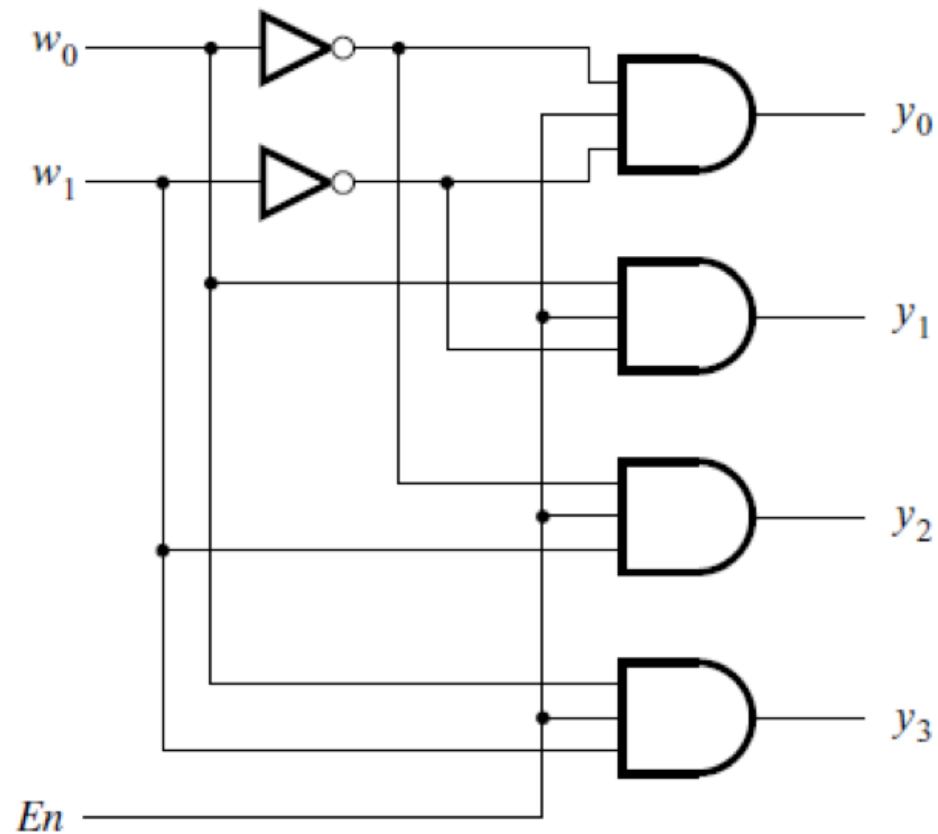
The outputs of an enabled binary decoder are “one-hot” encoded, meaning that only a single bit is set to 1, i.e., it is *hot*.

Demultiplexers

1-to-4 Demultiplexer (Definition)

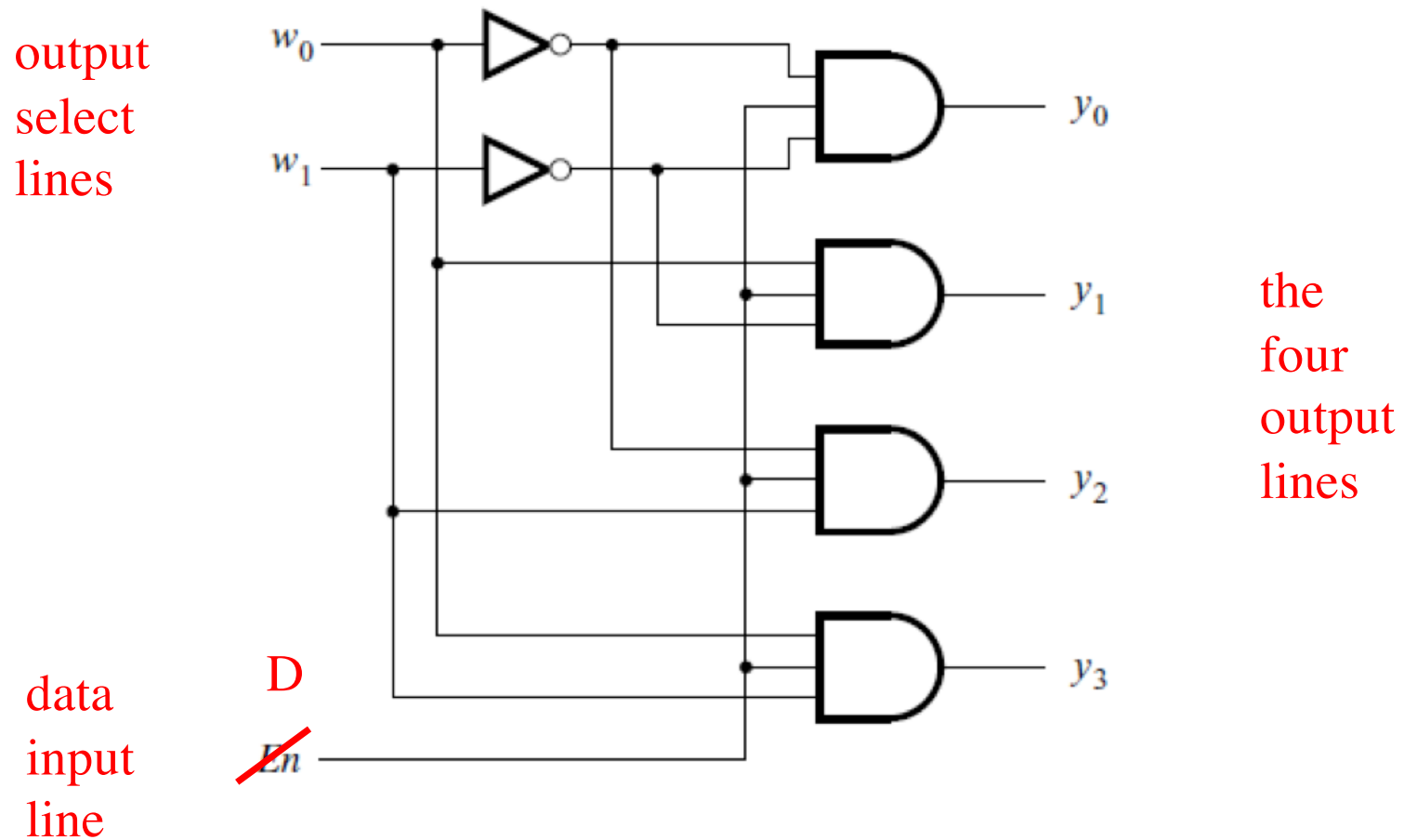
- **Has one data input line: D**
- **Has two output select lines: w_1 and w_0**
- **Has four outputs: y_0 , y_1 , y_2 , and y_3**
- **If $w_1=0$ and $w_0=0$, then the output y_0 is set to D**
- **If $w_1=0$ and $w_0=1$, then the output y_1 is set to D**
- **If $w_1=1$ and $w_0=0$, then the output y_2 is set to D**
- **If $w_1=1$ and $w_0=1$, then the output y_3 is set to D**
- **Only one output is set to D . All others are set to 0.**

A 1-to-4 demultiplexer built with a 2-to-4 decoder with enable



[Figure 4.14c from the textbook]

A 1-to-4 demultiplexer built with a 2-to-4 decoder with enable



[Figure 4.14c from the textbook]

Encoders

Binary Encoders

4-to-2 Binary Encoder (Definition)

- Has four inputs: w_3 , w_2 , w_1 , and w_0
- Has two outputs: y_1 and y_0
- Only one input is set to 1 (“one-hot” encoded). All others are set to 0.
- If $w_0=1$ then $y_1=0$ and $y_0=0$
- If $w_1=1$ then $y_1=0$ and $y_0=1$
- If $w_2=1$ then $y_1=1$ and $y_0=0$
- If $w_3=1$ then $y_1=1$ and $y_0=1$

Truth table for a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

The inputs are “one-hot” encoded

Expressions for 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	0		
0	0	0	1	0	0
0	0	1	0	0	1
0	0	1	1		
0	1	0	0	1	0
0	1	0	1		
0	1	1	0		
0	1	1	1		
1	0	0	0	1	1
1	0	0	1		
1	0	1	0		
1	0	1	1		
1	1	0	0		
1	1	0	1		
1	1	1	0		
1	1	1	1		

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

Expressions for 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	0	d	d
0	0	0	1	0	0
0	0	1	0	0	1
0	0	1	1	d	d
0	1	0	0	1	0
0	1	0	1	d	d
0	1	1	0	d	d
0	1	1	1	d	d
1	0	0	0	1	1
1	0	0	1	d	d
1	0	1	0	d	d
1	0	1	1	d	d
1	1	0	0	d	d
1	1	0	1	d	d
1	1	1	0	d	d
1	1	1	1	d	d

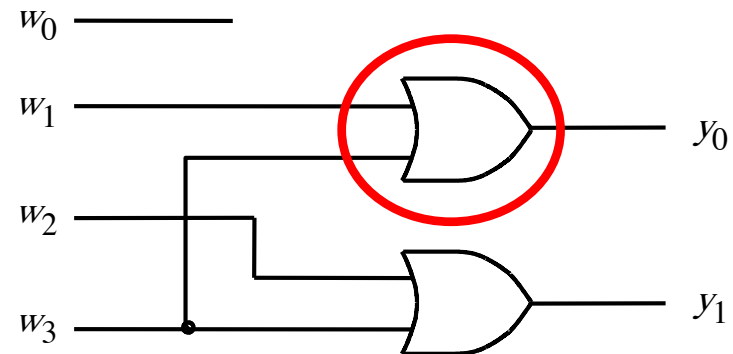
w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

Expressions for 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	0	d	d
0	0	0	1	0	0
0	0	1	0	0	1
0	0	1	1	d	d
0	1	0	0	1	0
0	1	0	1	d	d
0	1	1	0	d	d
0	1	1	1	d	d
1	0	0	0	1	1
1	0	0	1	d	d
1	0	1	0	d	d
1	0	1	1	d	d
1	1	0	0	d	d
1	1	0	1	d	d
1	1	1	0	d	d
1	1	1	1	d	d

$w_3 \ w_2 \ w_1 \ w_0$	00	01	11	10
00	d	0	d	1
01	0	d	d	d
11	d	d	d	d
10	1	d	d	d

$$y_0 = (w_1 + w_3)$$

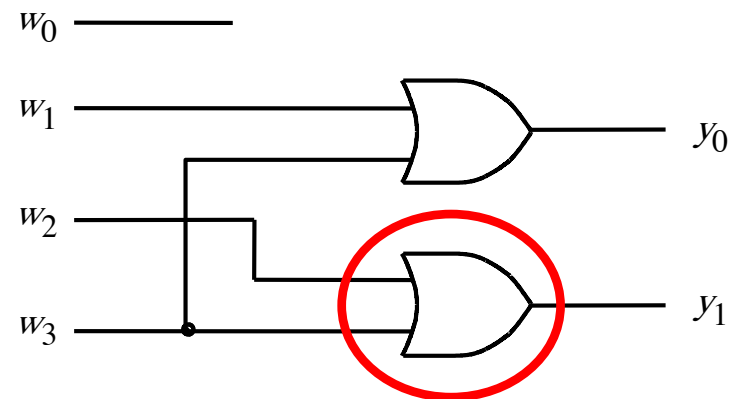


Expressions for 4-to-2 binary encoder

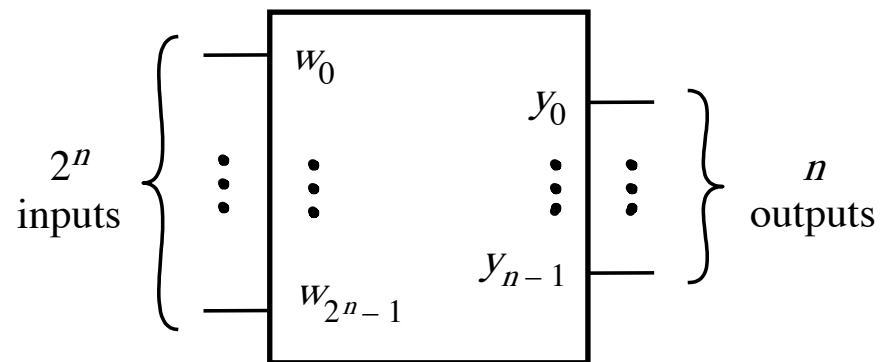
w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	0	d	d
0	0	0	1	0	0
0	0	1	0	0	1
0	0	1	1	d	d
0	1	0	0	1	0
0	1	0	1	d	d
0	1	1	0	d	d
0	1	1	1	d	d
1	0	0	0	1	1
1	0	0	1	d	d
1	0	1	0	d	d
1	0	1	1	d	d
1	1	0	0	d	d
1	1	0	1	d	d
1	1	1	0	d	d
1	1	1	1	d	d

		$w_3 w_2$			
		00	01	11	10
$w_1 w_0$	00	d	1	d	1
	01	0	d	d	d
	11	d	d	d	d
	10	0	d	d	d

$$y_1 = (w_3 + w_2)$$



A 2^n -to- n binary encoder



[Figure 4.18 from the textbook]

Priority Encoders

4-to-2 Priority Encoder (Definition)

- Has four inputs: w_3 , w_2 , w_1 , and w_0
- Has two primary outputs: y_1 and y_0
- Has one other output: z
- The inputs are NOT “one-hot” encoded.
- More than one input can be set to 1 but they have priorities associated with them: w_3 – highest priority and w_0 – lowest priority.
- $y_1=0$ and $y_0=0$ (if $w_0=1$ and $w_3=w_2=w_1=0$)
- $y_1=0$ and $y_0=1$ (if $w_1=1$ and $w_3=w_2=0$)
- $y_1=1$ and $y_0=0$ (if $w_2=1$ and $w_3=0$)
- $y_1=1$ and $y_0=1$ (if $w_3=1$)

4-to-2 Priority Encoder (Definition)

- Has four inputs: w_3 , w_2 , w_1 , and w_0
 - Has two primary outputs: y_1 and y_0
 - Has one other output: z
 - The inputs are NOT “one-hot” encoded.
 - More than one input can be set to 1 but they have priorities associated with them: w_3 – highest priority and w_0 – lowest priority.
 - $y_1=0$ and $y_0=0$ (if $w_0=1$ and $w_3=w_2=w_1=0$)
 - $y_1=0$ and $y_0=1$ (if $w_1=1$ and $w_3=w_2=0$) w_0
 - $y_1=1$ and $y_0=0$ (if $w_2=1$ and $w_3=0$) w_0 , w_1
 - $y_1=1$ and $y_0=1$ (if $w_3=1$) w_0 , w_1 , w_2
- these have lower priorities
and can be either 0 or 1.

4-to-2 Priority Encoder (Definition)

- Has four inputs: w_3 , w_2 , w_1 , and w_0
- Has two primary outputs: y_1 and y_0
- Has one other output: z
- The inputs are NOT “one-hot” encoded.
- More than one input can be set to 1 but they have priorities associated with them: w_3 – highest priority and w_0 – lowest priority.
- $y_1=0$ and $y_0=0$ (if $w_0=1$ and $w_3=w_2=w_1=0$)
- $y_1=0$ and $y_0=1$ (if $w_1=1$ and $w_3=w_2=0$)
- $y_1=1$ and $y_0=0$ (if $w_2=1$ and $w_3=0$)
- $y_1=1$ and $y_0=1$ (if $w_3=1$)
- $z = 0$ if $w_3=w_2=w_1=w_0=0$; otherwise $z=1$.

Truth table for a 4-to-2 priority encoder (abbreviated version)

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

[Figure 4.20 from the textbook]

Truth table for a 4-to-2 priority encoder (abbreviated version)

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

These two are don't cares, because the y values don't matter in this case.
In other words, when $z=0$ there are no priority inputs that are set to 1.

Truth table for a 4-to-2 priority encoder (abbreviated version)

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

[Figure 4.20 from the textbook]

Truth table for a 4-to-2 priority encoder

	w ₃	w ₂	w ₁	w ₀	y ₁	y ₀	z
0 0 0 0	0	0	0	0	d	d	0
0 0 0 1	0	0	0	1	0	0	1
0 0 1 x	0	0	1	0	0	1	1
	0	0	1	1	0	1	1
0 1 x x	0	1	0	0	1	0	1
	0	1	0	1	1	0	1
	0	1	1	0	1	0	1
	0	1	1	1	1	0	1
1 x x x	1	0	0	0	1	1	1
	1	0	0	1	1	1	1
	1	0	1	0	1	1	1
	1	0	1	1	1	1	1
	1	1	0	0	1	1	1
	1	1	0	1	1	1	1
	1	1	1	0	1	1	1
	1	1	1	1	1	1	1

Expressions for 4-to-2 priority encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	0	0	1	1
0	0	1	1	0	1	1
0	1	0	0	1	0	1
0	1	0	1	1	0	1
0	1	1	0	1	0	1
0	1	1	1	1	0	1
1	0	0	0	1	1	1
1	0	0	1	1	1	1
1	0	1	0	1	1	1
1	0	1	1	1	1	1
1	1	0	0	1	1	1
1	1	0	1	1	1	1
1	1	1	0	1	1	1
1	1	1	1	1	1	1

		$w_3 w_2$			
$w_1 w_0$		00	01	11	10
	00	d	1	1	1
	01	0	1	1	1
	11	0	1	1	1
	10	0	1	1	1

$$y_1 = w_3 + w_2$$

Expressions for 4-to-2 priority encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	0	0	1	1
0	0	1	1	0	1	1
0	1	0	0	1	0	1
0	1	0	1	1	0	1
0	1	1	0	1	0	1
0	1	1	1	1	0	1
1	0	0	0	1	1	1
1	0	0	1	1	1	1
1	0	1	0	1	1	1
1	0	1	1	1	1	1
1	1	0	0	1	1	1
1	1	0	1	1	1	1
1	1	1	0	1	1	1
1	1	1	1	1	1	1

$w_3 w_2$					
$w_1 w_0$		00	01	11	10
	00	d	0	1	1
	01	0	0	1	1
	11	1	0	1	1
	10	1	0	1	1

$$y_0 = w_3 + w_1 \overline{w_2}$$

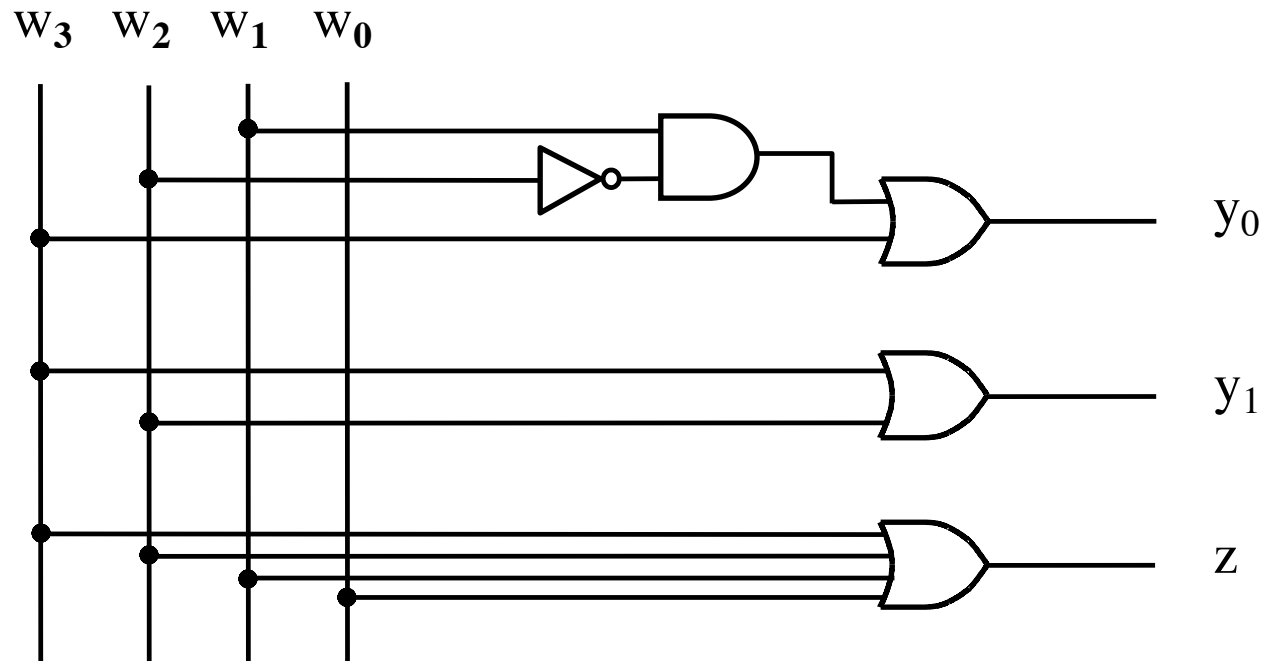
Expressions for 4-to-2 priority encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	0	0	1	1
0	0	1	1	0	1	1
0	1	0	0	1	0	1
0	1	0	1	1	0	1
0	1	1	0	1	0	1
0	1	1	1	1	0	1
1	0	0	0	1	1	1
1	0	0	1	1	1	1
1	0	1	0	1	1	1
1	0	1	1	1	1	1
1	1	0	0	1	1	1
1	1	0	1	1	1	1
1	1	1	0	1	1	1
1	1	1	1	1	1	1

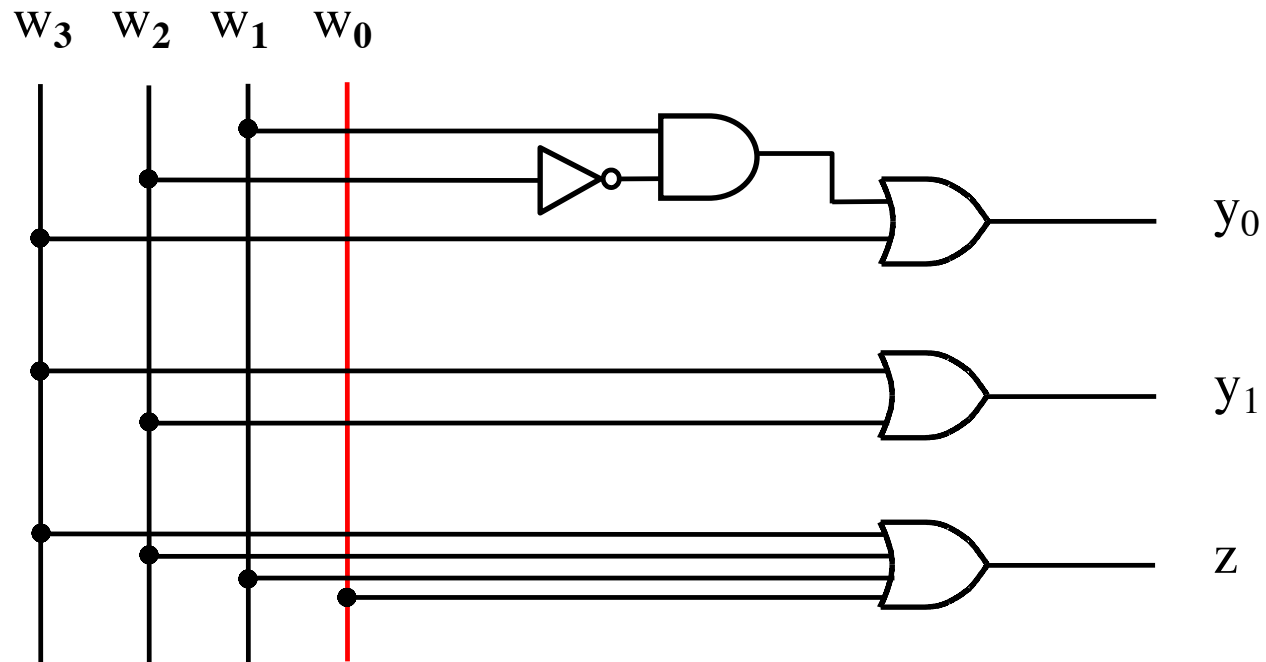
		$w_3 w_2$			
$w_1 w_0$		00	01	11	10
	00	0	1	1	1
	01	1	1	1	1
	11	1	1	1	1
	10	1	1	1	1

$$Z = w_3 + w_2 + w_1 + w_0$$

Circuit for the 4-to-2 priority encoder



Circuit for the 4-to-2 priority encoder



In this circuit w_0 is used.

The textbook derives a different circuit for the 4-to-2 priority encoder using a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

The textbook derives a different circuit for the 4-to-2 priority encoder using a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

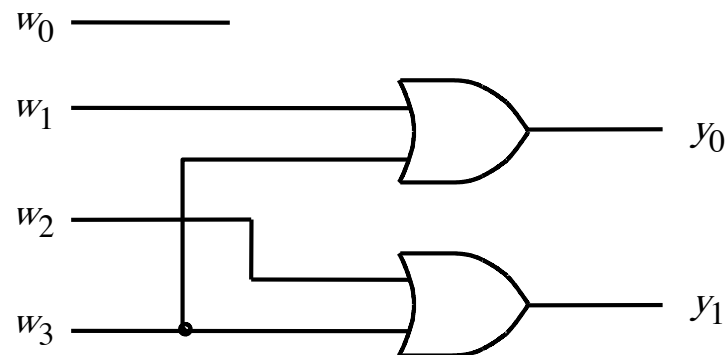
$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$



The textbook derives a different circuit for the 4-to-2 priority encoder using a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

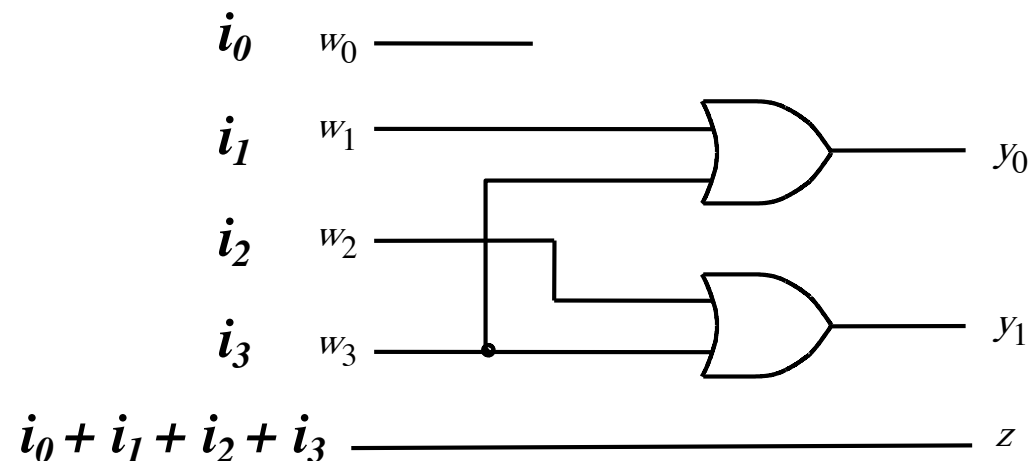
$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$



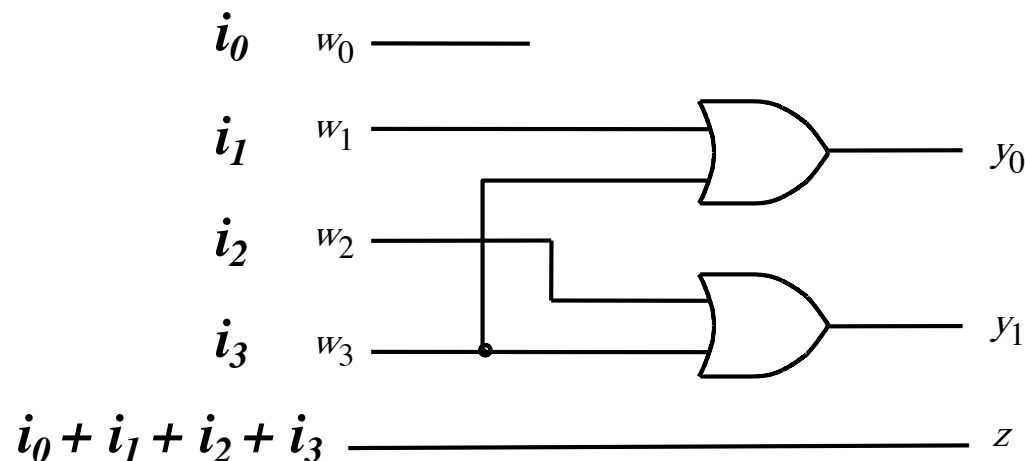
The textbook derives a different circuit for the 4-to-2 priority encoder using a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$\begin{aligned}
 i_0 &= \overline{w_3}\overline{w_2}\overline{w_1}w_0 \\
 i_1 &= \overline{w_3}\overline{w_2}w_1 \\
 i_2 &= \overline{w_3}w_2 \\
 i_3 &= w_3
 \end{aligned}$$

$$\begin{aligned}
 y_0 &= i_1 + i_3 \\
 y_1 &= i_2 + i_3
 \end{aligned}$$

$$z = i_0 + i_1 + i_2 + i_3$$



Try to prove that this is equivalent to the circuit that was derived above.

Let's prove this for z

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \bar{w}_3 \bar{w}_2 \bar{w}_1 w_0$$

$$i_1 = \bar{w}_3 \bar{w}_2 w_1$$

$$i_2 = \bar{w}_3 w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

		$w_3 w_2$			
$w_1 w_0$		00	01	11	10
	00	0	1	1	1
	01	1	1	1	1
	11	1	1	1	1
	10	1	1	1	1

$z = ?$

Let's prove this for z

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \bar{w}_3 \bar{w}_2 \bar{w}_1 w_0$$

$$i_1 = \bar{w}_3 \bar{w}_2 w_1$$

$$i_2 = \bar{w}_3 w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

		$w_3 w_2$			
		00	01	11	10
$w_1 w_0$	00	0	1	1	1
	01	1	1	1	1
	11	1	1	1	1
	10	1	1	1	1

$$z = (w_0 + w_1 + w_2 + w_3)$$

Let's prove this for y_0

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \bar{w}_3 \bar{w}_2 \bar{w}_1 w_0$$

$$i_1 = \bar{w}_3 \bar{w}_2 w_1$$

$$i_2 = \bar{w}_3 w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

$w_1 w_0$ \ $w_3 w_2$		$w_3 w_2$			
		00	01	11	10
00	00	0	0	1	1
01	00	0	0	1	1
11	00	1	0	1	1
10	00	1	0	1	1

$$y_0 = ?$$

Let's prove this for y_0

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

$w_3 w_2$ $w_1 w_0$		$w_3 w_2$			
		00	01	11	10
00	00	0	0	1	1
01	01	0	0	1	1
11	11	1	0	1	1
10	10	1	0	1	1

$$y_0 = w_3 + w_1 \overline{w_2}$$

Let's prove this for y_1

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \bar{w}_3 \bar{w}_2 \bar{w}_1 w_0$$

$$i_1 = \bar{w}_3 \bar{w}_2 w_1$$

$$i_2 = \bar{w}_3 w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

$w_3 w_2$ $w_1 w_0$		$w_3 w_2$			
		00	01	11	10
00	00	0	1	1	1
01	00	0	1	1	1
11	00	0	1	1	1
10	00	0	1	1	1

$y_1 = ?$

Let's prove this for y_1

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \bar{w}_3 \bar{w}_2 \bar{w}_1 w_0$$

$$i_1 = \bar{w}_3 \bar{w}_2 w_1$$

$$i_2 = \bar{w}_3 w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

$w_3 w_2$ $w_1 w_0$		$w_3 w_2$			
		00	01	11	10
00	00	0	1	1	1
01	01	0	1	1	1
11	11	0	1	1	1
10	10	0	1	1	1

$$y_1 = w_3 + w_2$$

Therefore, this circuit for the
4-to-2 priority encoder is equivalent to ...

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

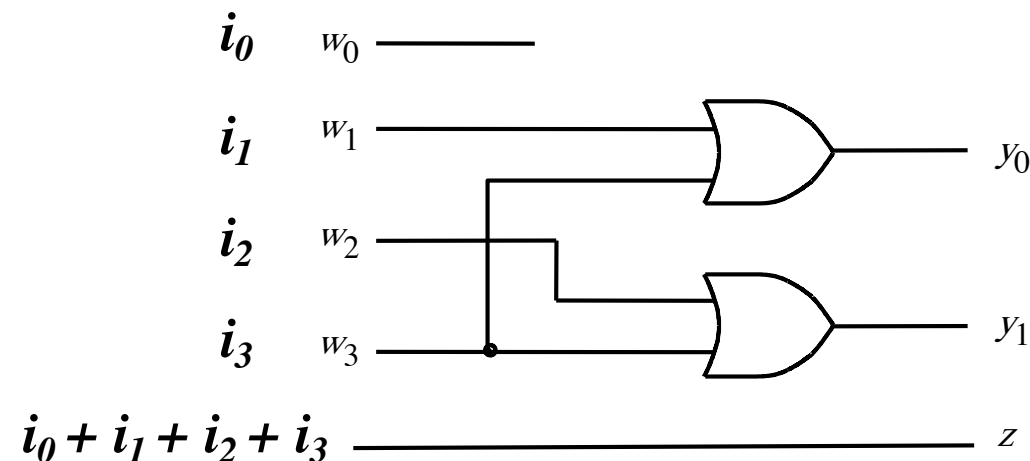
$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

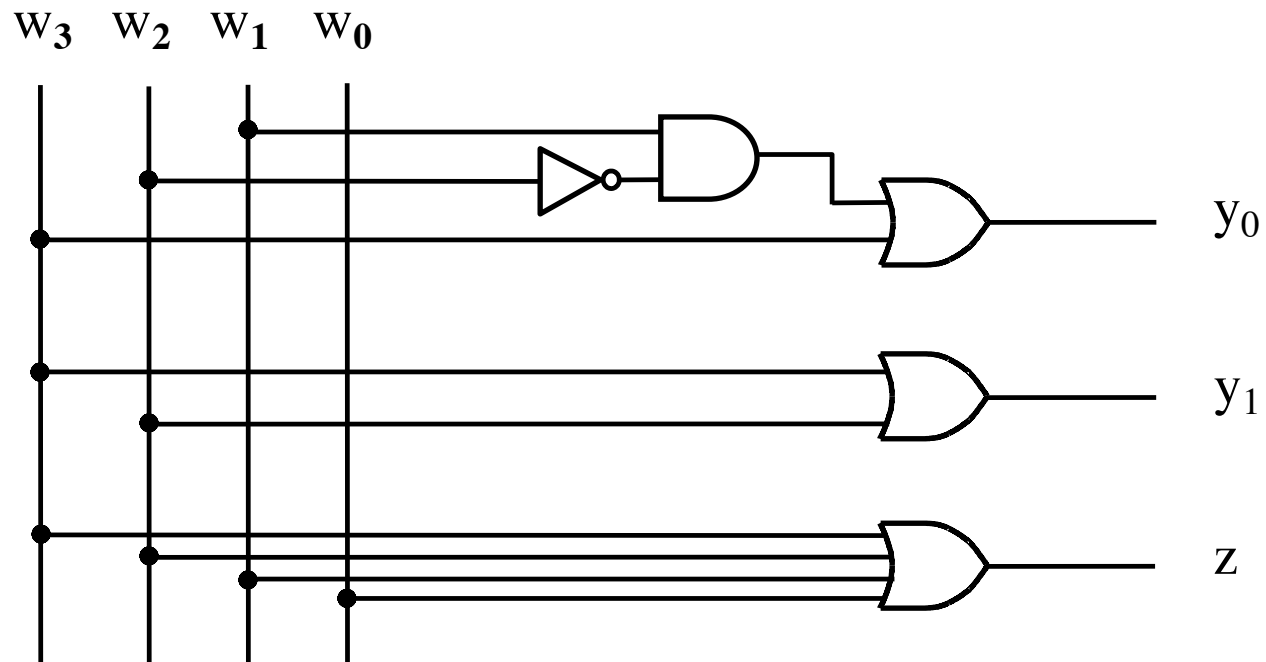
$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$



... this circuit for the 4-to-2 priority encoder



Code Converters

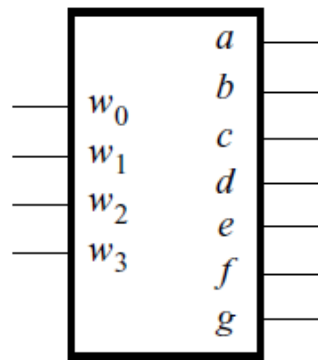
Code Converter (Definition)

- **Converts from one type of input encoding to a different type of output encoding.**

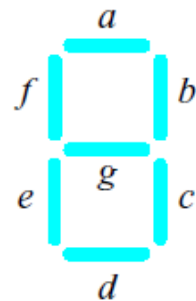
Code Converter (Definition)

- **Converts from one type of input encoding to a different type of output encoding.**
- **A decoder does that as well, but its outputs are always one-hot encoded so the output code is really only one type of output code.**
- **A binary encoder does that as well but its inputs are always one-hot encoded so the input code is really only one type of input code.**

A hex-to-7-segment display code converter



(a) Code converter

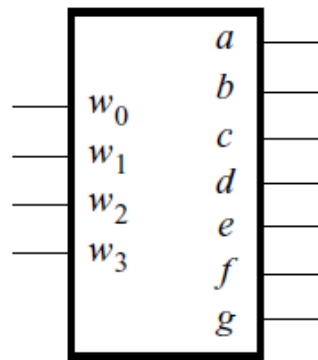


(b) 7-segment display

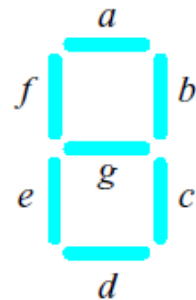
w_3	w_2	w_1	w_0	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	1	1	1	0	1	1	1
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	1	0	0	1	1	1	0
1	1	0	1	0	1	1	1	1	0	1
1	1	1	0	1	0	0	1	1	1	1
1	1	1	1	1	0	0	0	1	1	1

(c) Truth table

A hex-to-7-segment display code converter



(a) Code converter

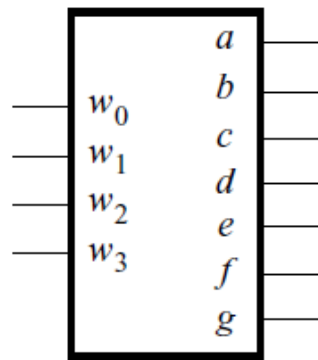


(b) 7-segment display

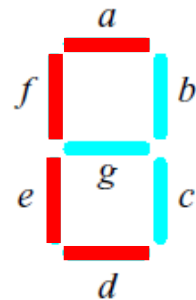
w_3	w_2	w_1	w_0	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	1	1	1	0	1	1	1
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	1	0	0	1	1	1	0
1	1	0	1	0	1	1	1	1	0	1
1	1	1	0	1	0	0	1	1	1	1
1	1	1	1	1	0	0	0	1	1	1

(c) Truth table

A hex-to-7-segment display code converter



(a) Code converter

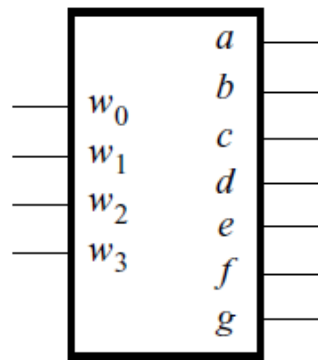


(b) 7-segment display

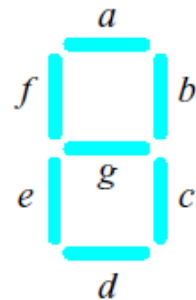
w_3	w_2	w_1	w_0	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	1	1	1	0	1	1	1
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	1	0	0	1	1	1	0
1	1	0	1	0	1	1	1	1	0	1
1	1	1	0	1	0	0	1	1	1	1
1	1	1	1	1	0	0	0	1	1	1

(c) Truth table

A hex-to-7-segment display code converter



(a) Code converter

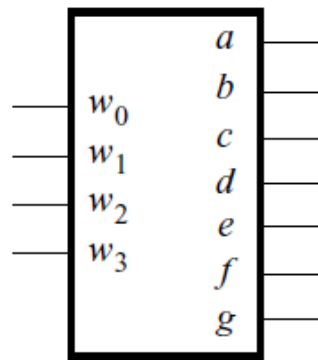


(b) 7-segment display

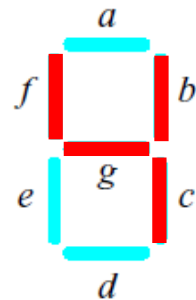
w_3	w_2	w_1	w_0	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	1	1	1	0	1	1	1
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	1	0	0	1	1	1	0
1	1	0	1	0	1	1	1	1	0	1
1	1	1	0	1	0	0	1	1	1	1
1	1	1	1	1	0	0	0	1	1	1

(c) Truth table

A hex-to-7-segment display code converter



(a) Code converter



(b) 7-segment display

w_3	w_2	w_1	w_0	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	1	1	1	0	1	1	1
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	1	0	0	1	1	1	0
1	1	0	1	0	1	1	1	1	0	1
1	1	1	0	1	0	0	1	1	1	1
1	1	1	1	1	0	0	0	1	1	1

(c) Truth table

What is the circuit for this code converter?

x_3	x_2	x_1	x_0	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	1	1	1	0	1	1	1
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	1	0	0	1	1	1	0
1	1	0	1	0	1	1	1	1	0	1
1	1	1	0	1	0	0	1	1	1	1
1	1	1	1	1	0	0	0	1	1	1

What is the circuit for this code converter?

x_3	x_2	x_1	x_0	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	1	1	1	0	1	1	1
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	1	0	0	1	1	1	0
1	1	0	1	0	1	1	1	1	0	1
1	1	1	0	1	0	0	1	1	1	1
1	1	1	1	1	0	0	0	1	1	1

$$f(x_1, x_2, x_3, x_4) = \sum m(0, 2, 3, 5, 6, 7, 8, 9, 10, 12, 14, 15)$$

What is the circuit for this code converter?

x_3	x_2	x_1	x_0	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	1	1	1	0	1	1	1
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	1	0	0	1	1	1	0
1	1	0	1	0	1	1	1	1	0	1
1	1	1	0	1	0	0	1	1	1	1
1	1	1	1	1	0	0	0	1	1	1

		$x_1 \ x_2$			
		00	01	11	10
$x_3 \ x_4$	00	1	0	1	1
	01	0	1	0	1
	11	1	1	1	0
	10	1	1	1	1

$$f(x_1, x_2, x_3, x_4) = \sum m(0, 2, 3, 5, 6, 7, 8, 9, 10, 12, 14, 15)$$

What is the circuit for this code converter?

x_3	x_2	x_1	x_0	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	1	1	1	0	1	1	1
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	1	0	0	1	1	1	0
1	1	0	1	0	1	1	1	1	0	1
1	1	1	0	1	0	0	1	1	1	1
1	1	1	1	1	0	0	0	1	1	1

What is the circuit for this code converter?

x_3	x_2	x_1	x_0	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	1	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	1	1	1	0	1	1	1
1	0	1	1	0	0	1	1	1	1	1
1	1	0	0	1	0	0	1	1	1	0
1	1	0	1	0	1	1	1	1	0	1
1	1	1	0	1	0	0	1	1	1	1
1	1	1	1	1	0	0	0	1	1	1

		$x_1 \ x_2$			
		00	01	11	10
$x_3 \ x_4$	00	1	0	1	1
	01	0	1	1	1
	11	1	0	0	1
	10	1	1	1	0

Diagram illustrating the Karnaugh map for the function $f(x_1, x_2, x_3, x_4)$. The map is a 4x4 grid with rows labeled $x_3 x_4$ (00, 01, 11, 10) and columns labeled $x_1 x_2$ (00, 01, 11, 10). The output values are 1 or 0. Red brackets indicate groupings for simplification: x_1 (top row), x_4 (right column), x_3 (left column), and x_2 (bottom row).

$$f(x_1, x_2, x_3, x_4) = \sum m(0, 2, 3, 5, 6, 8, 9, 11, 12, 13, 14)$$

Example Problems from Chapter 4

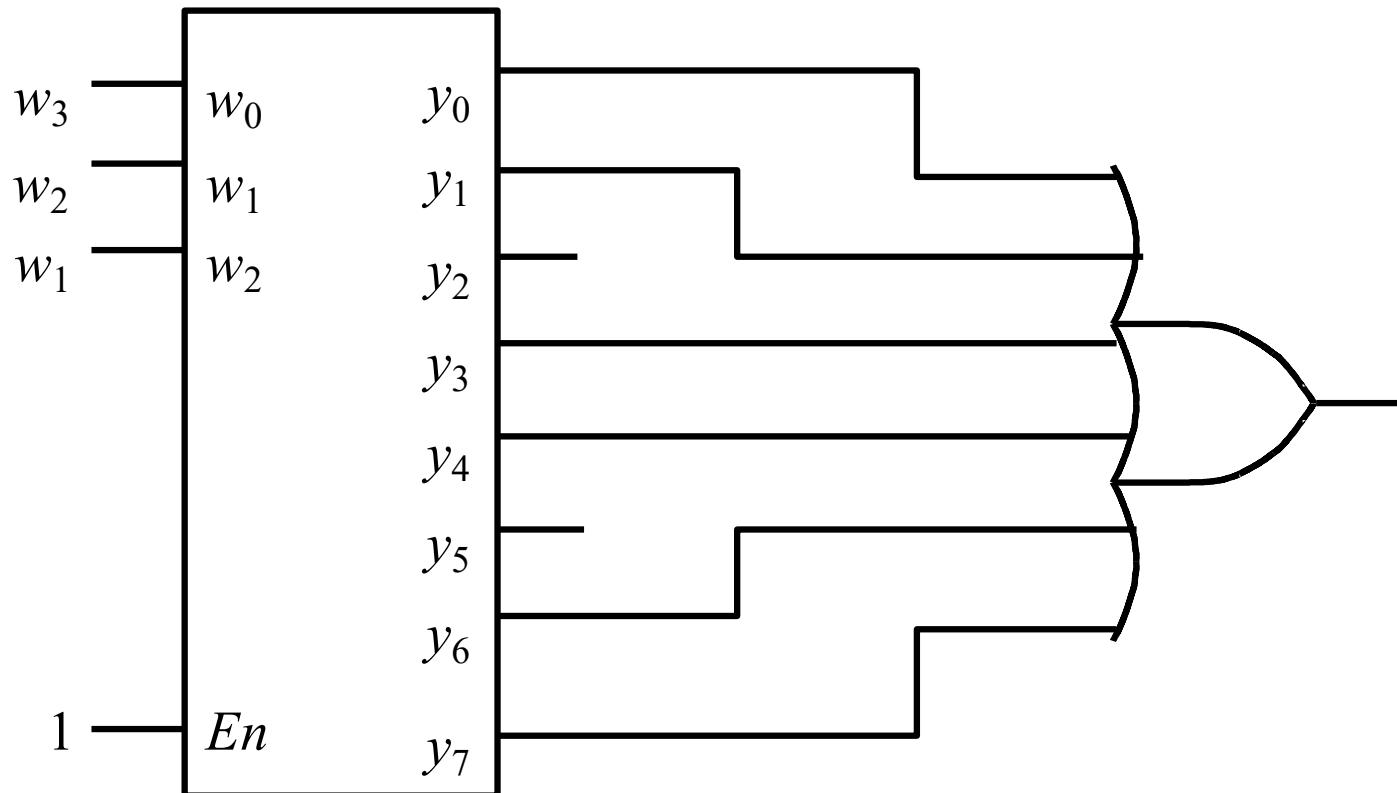
Example 1: SOP vs Decoders

Implement the function

$$f(w_1, w_2, w_3) = \sum m(0, 1, 3, 4, 6, 7)$$

by using a 3-to-8 binary decoder and one OR gate.

Solution Circuit



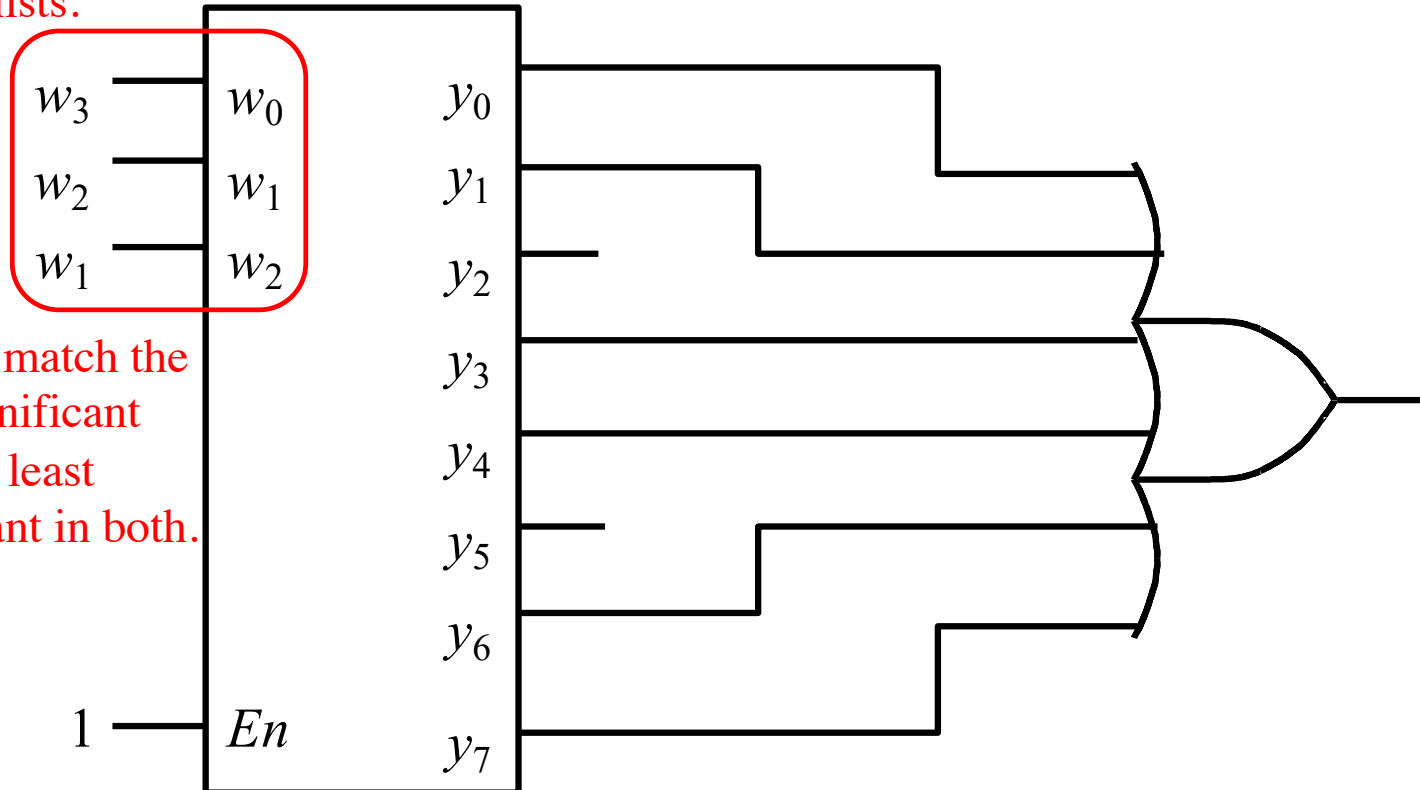
$$f(w_1, w_2, w_3) = \sum m(0, 1, 3, 4, 6, 7)$$

[Figure 4.44 from the textbook]

Solution Circuit

Notice this swap
of variables in
the two lists.

Need to match the
least significant
with the least
significant in both.



$$f(w_1, w_2, w_3) = \sum m(0, 1, 3, 4, 6, 7)$$

[Figure 4.44 from the textbook]

Example 2: Implement an 8-to-3 binary encoder

w_7	w_6	w_5	w_4	w_3	w_2	w_1	w_0	y_2	y_1	y_0
0	0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	1	0	0	0	1	0
0	0	0	0	1	0	0	0	0	1	1
0	0	0	1	0	0	0	0	1	0	0
0	0	1	0	0	0	0	0	1	0	1
0	1	0	0	0	0	0	0	1	1	0
1	0	0	0	0	0	0	0	1	1	1

[Figure 4.45 from the textbook]

Example 2: Implement an 8-to-3 binary encoder

w_7	w_6	w_5	w_4	w_3	w_2	w_1	w_0	y_2	y_1	y_0
0	0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	1	0	0	0	1	0
0	0	0	0	1	0	0	0	0	1	1
0	0	0	1	0	0	0	0	1	0	0
0	0	1	0	0	0	0	0	1	0	1
0	1	0	0	0	0	0	0	1	1	0
1	0	0	0	0	0	0	0	1	1	1

These formulas work due to the one-hot encoded property of the inputs.

$$y_0 = w_1 + w_3 + w_5 + w_7$$

$$y_1 = w_2 + w_3 + w_6 + w_7$$

$$y_2 = w_4 + w_5 + w_6 + w_7$$

[Figure 4.45 from the textbook]

Example 2: Implement an 8-to-3 binary encoder

w_7	w_6	w_5	w_4	w_3	w_2	w_1	w_0	y_2	y_1	y_0
0	0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	1	0	0	0	1	0
0	0	0	0	1	0	0	0	0	1	1
0	0	0	1	0	0	0	0	1	0	0
0	0	1	0	0	0	0	0	1	0	1
0	1	0	0	0	0	0	0	1	1	0
1	0	0	0	0	0	0	0	1	1	1

$$y_0 = w_1 + w_3 + w_5 + w_7$$

$$y_1 = w_2 + w_3 + w_6 + w_7$$

$$y_2 = w_4 + w_5 + w_6 + w_7$$

Example 2: Implement an 8-to-3 binary encoder

w_7	w_6	w_5	w_4	w_3	w_2	w_1	w_0	y_2	y_1	y_0
0	0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	1	0	0	0	1	0
0	0	0	0	1	0	0	0	0	1	1
0	0	0	1	0	0	0	0	1	0	0
0	0	1	0	0	0	0	0	1	0	1
0	1	0	0	0	0	0	0	1	1	0
1	0	0	0	0	0	0	0	1	1	1

$$y_0 = w_1 + w_3 + w_5 + w_7$$

$$y_1 = w_2 + w_3 + w_6 + w_7$$

$$y_2 = w_4 + w_5 + w_6 + w_7$$

Example 2: Implement an 8-to-3 binary encoder

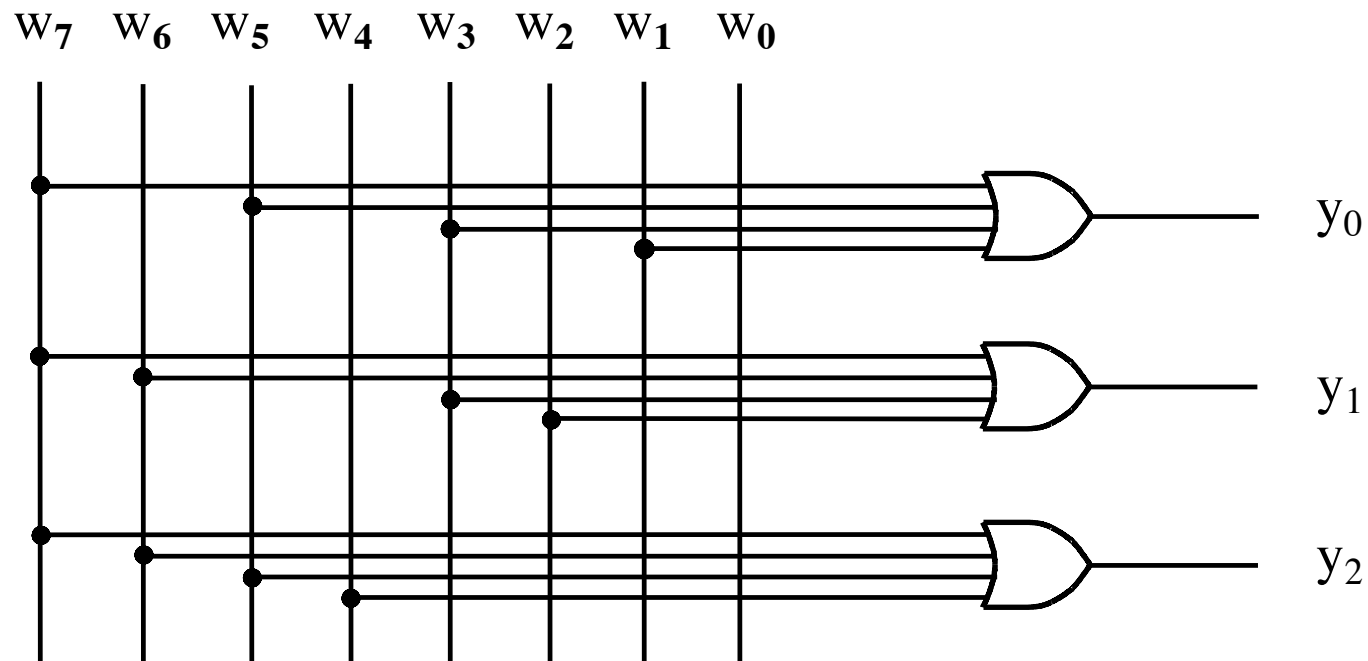
w_7	w_6	w_5	w_4	w_3	w_2	w_1	w_0	y_2	y_1	y_0
0	0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	1	0	0	0	1	0
0	0	0	0	1	0	0	0	0	1	1
0	0	0	1	0	0	0	0	1	0	0
0	0	1	0	0	0	0	0	1	0	1
0	1	0	0	0	0	0	0	1	1	0
1	0	0	0	0	0	0	0	1	1	1

$$y_0 = w_1 + w_3 + w_5 + w_7$$

$$y_1 = w_2 + w_3 + w_6 + w_7$$

$$y_2 = w_4 + w_5 + w_6 + w_7$$

Circuit for the 8-to-3 binary encoder



$$y_0 = w_1 + w_3 + w_5 + w_7$$

$$y_1 = w_2 + w_3 + w_6 + w_7$$

$$y_2 = w_4 + w_5 + w_6 + w_7$$

Example 3: Implement an 8-to-3 priority encoder

[illegible]

Example 3: Implement an 8-to-3 priority encoder

w_7	w_6	w_5	w_4	w_3	w_2	w_1	w_0		y_2	y_1	y_0
0	0	0	0	0	0	0	1		0	0	0
0	0	0	0	0	0	1	X		0	0	1
0	0	0	0	0	1	X	X		0	1	0
0	0	0	0	1	X	X	X		0	1	1
0	0	0	1	X	X	X	X		1	0	0
0	0	1	X	X	X	X	X		1	0	1
0	1	X	X	X	X	X	X		1	1	0
1	X	X	X	X	X	X	X		1	1	1

Example 3: Implement an 8-to-3 priority encoder

w_7	w_6	w_5	w_4	w_3	w_2	w_1	w_0	y_2	y_1	y_0	z
0	0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	x	0	0	1	1
0	0	0	0	0	1	x	x	0	1	0	1
0	0	0	0	1	x	x	x	0	1	1	1
0	0	0	1	x	x	x	x	1	0	0	1
0	0	1	x	x	x	x	x	1	0	1	1
0	1	x	x	x	x	x	x	1	1	0	1
1	x	x	x	x	x	x	x	1	1	1	1
0	0	0	0	0	0	0	0	d	d	d	0

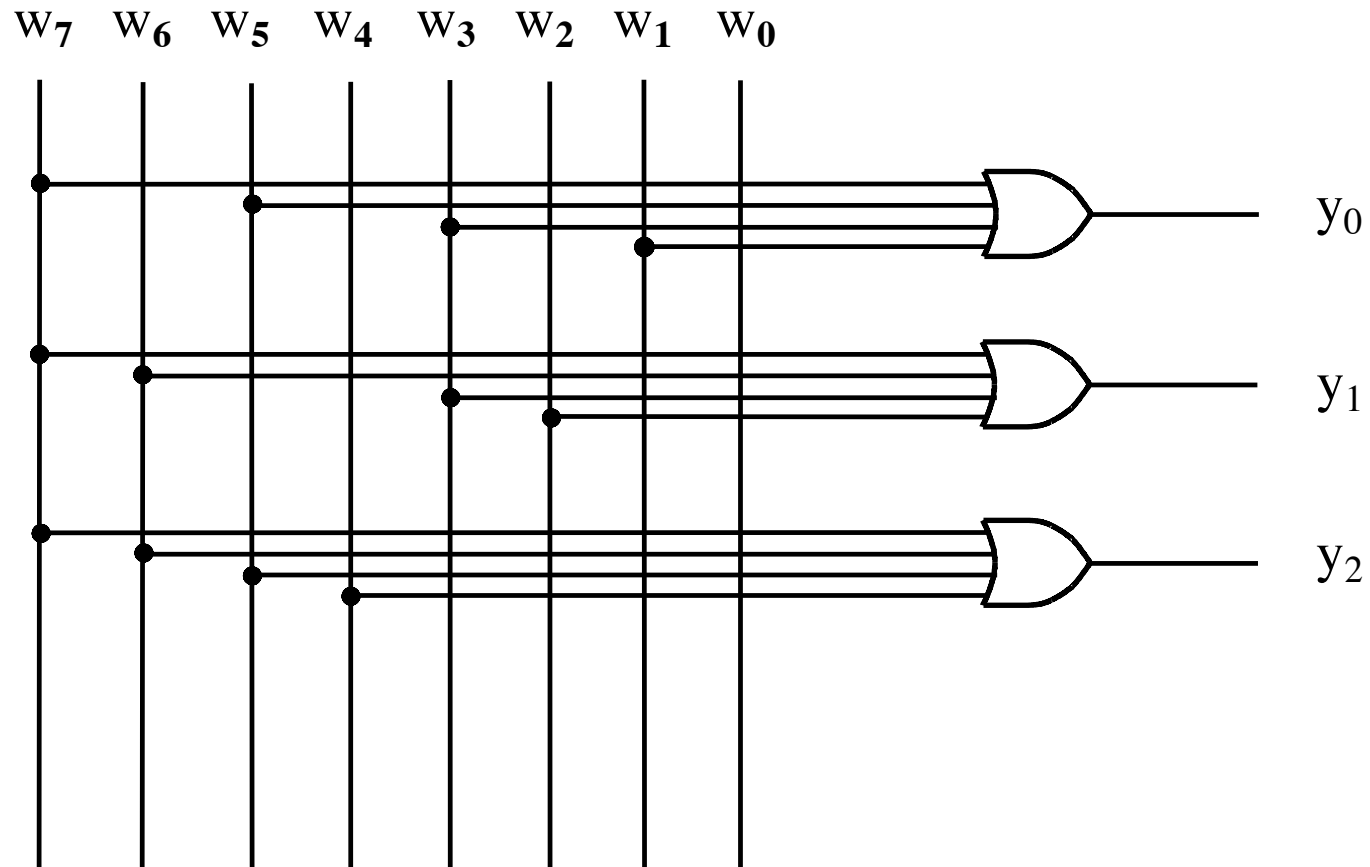
Example 3: Implement an 8-to-3 priority encoder

w_7	w_6	w_5	w_4	w_3	w_2	w_1	w_0	y_2	y_1	y_0	z
0	0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	x	0	0	1	1
0	0	0	0	0	1	x	x	0	1	0	1
0	0	0	0	1	x	x	x	0	1	1	1
0	0	0	1	x	x	x	x	1	0	0	1
0	0	1	x	x	x	x	x	1	0	1	1
0	1	x	x	x	x	x	x	1	1	0	1
1	x	x	x	x	x	x	x	1	1	1	1
0	0	0	0	0	0	0	0	d	d	d	0

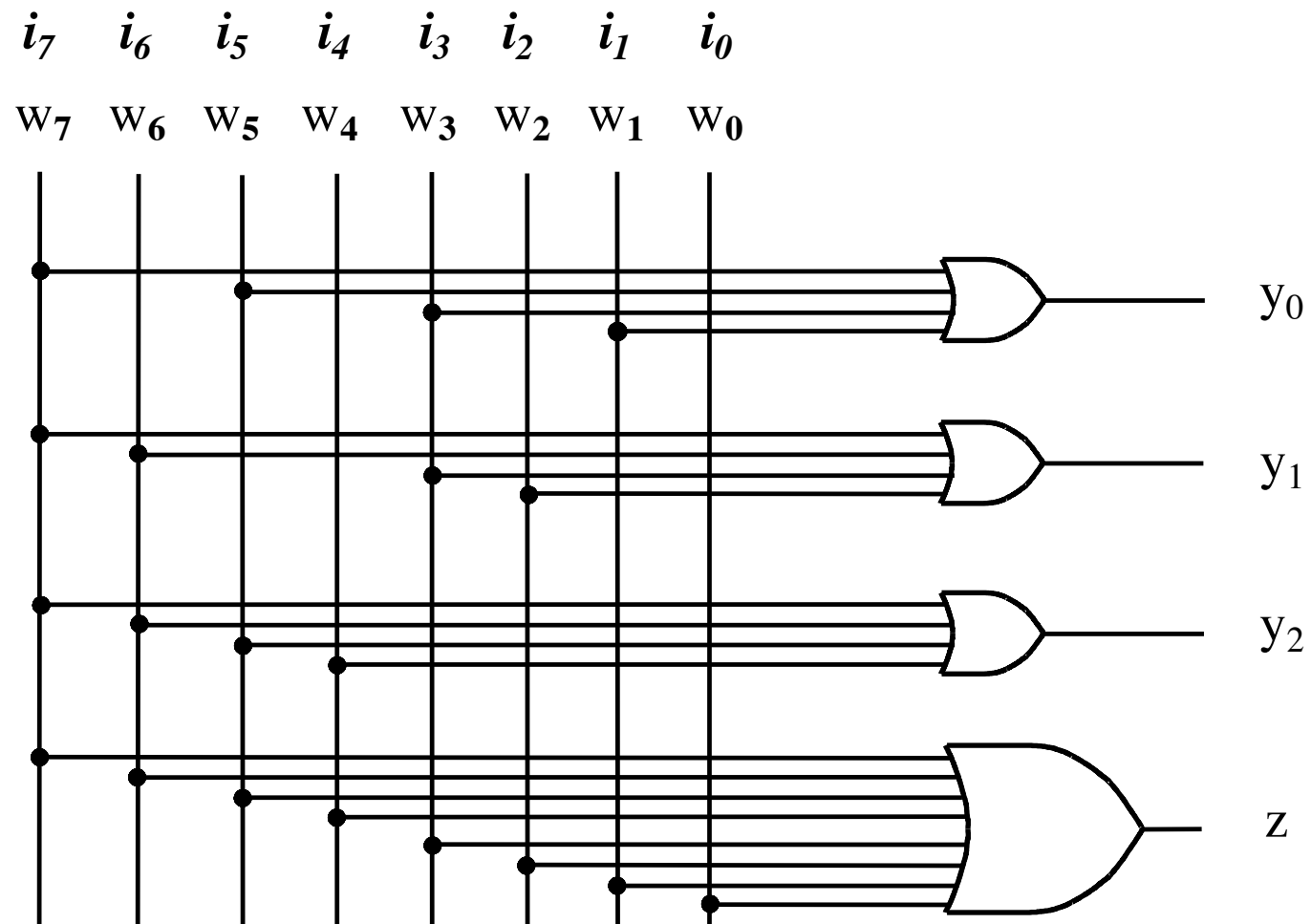
Example 3: Implement an 8-to-3 priority encoder

	w_7	w_6	w_5	w_4	w_3	w_2	w_1	w_0	y_2	y_1	y_0	z
$i_0 = \overline{w_7} \overline{w_6} \overline{w_5} \overline{w_4} \overline{w_3} \overline{w_2} \overline{w_1} w_0$	0	0	0	0	0	0	0	1	0	0	0	1
$i_1 = \overline{w_7} \overline{w_6} \overline{w_5} \overline{w_4} \overline{w_3} \overline{w_2} w_1$	0	0	0	0	0	0	1	x	0	0	1	1
$i_2 = \overline{w_7} \overline{w_6} \overline{w_5} \overline{w_4} \overline{w_3} w_2$	0	0	0	0	0	1	x	x	0	1	0	1
$i_3 = \overline{w_7} \overline{w_6} \overline{w_5} \overline{w_4} w_3$	0	0	0	0	1	x	x	x	0	1	1	1
$i_4 = \overline{w_7} \overline{w_6} \overline{w_5} w_4$	0	0	0	1	x	x	x	x	1	0	0	1
$i_5 = \overline{w_7} \overline{w_6} w_5$	0	0	1	x	x	x	x	x	1	0	1	1
$i_6 = \overline{w_7} w_6$	0	1	x	x	x	x	x	x	1	1	0	1
$i_7 = w_7$	1	x	x	x	x	x	x	x	1	1	1	1
$z = i_0 + i_1 + i_2 + i_3 + i_4 + i_5 + i_6 + i_7$	0	0	0	0	0	0	0	0	d	d	d	0

Circuit for the 8-to-3 binary encoder



Circuit for the 8-to-3 priority encoder



Example 4: Circuit implementation using a multiplexer

Implement the function

$$f(w_1, w_2, w_3, w_4, w_5) = \bar{w}_1 \bar{w}_2 \bar{w}_4 \bar{w}_5 + w_1 w_2 + w_1 w_3 + w_1 w_4 + w_3 w_4 w_5$$

using a 4-to-1 multiplexer

Some Boolean Algebra Leads To

$$\overline{w_1}\overline{w_2}\overline{w_4}\overline{w_5} + w_1w_2 + w_1w_3 + w_1w_4 + w_3w_4w_5$$

$$\overline{w_1}\overline{w_4}(\overline{w_5}\overline{w_2}) + w_4(w_3w_5) + w_1(w_2 + w_3) + w_1w_4(1)$$

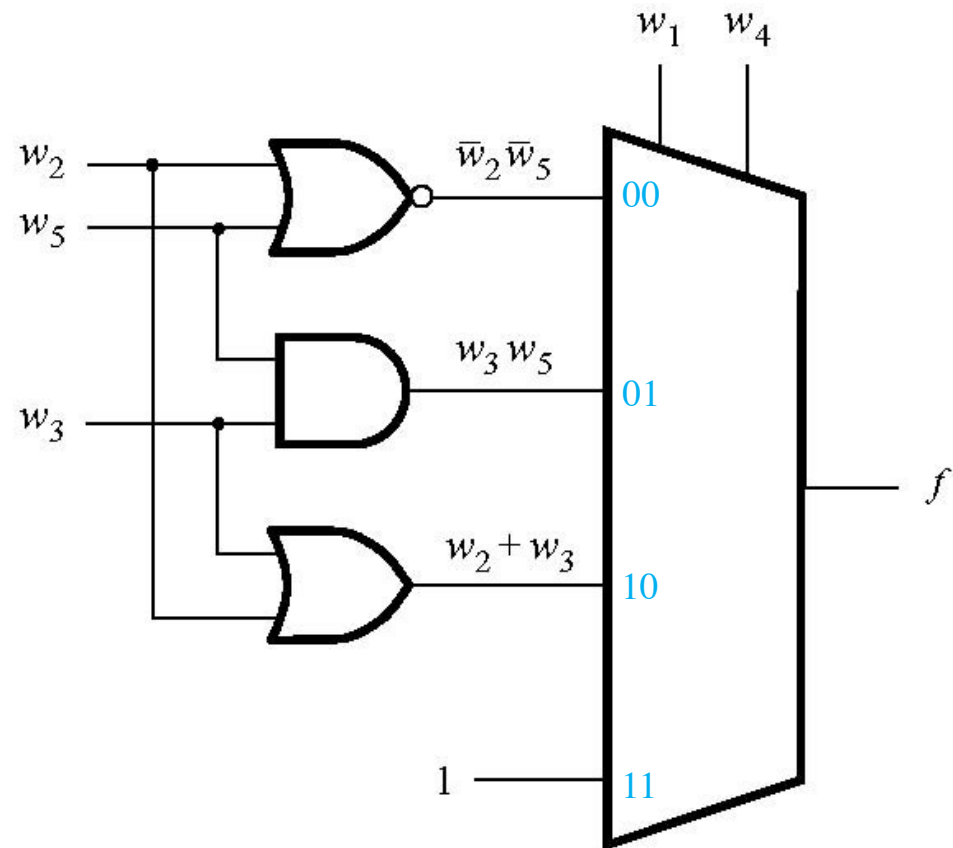
$$\overline{w_1}\overline{w_4}(\overline{w_5}\overline{w_2}) + (\overline{w_1}+w_1)w_4(w_3w_5) + w_1(\overline{w_4}+w_4)(w_2 + w_3) + w_1w_4(1)$$

$$\overline{w_1}\overline{w_4}(\overline{w_5}\overline{w_2}) + \overline{w_1}w_4(w_3w_5) + w_1\overline{w_4}(w_2 + w_3) + w_1w_4(w_3w_5 + (w_2+w_3) + 1)$$

$$\overline{w_1}\overline{w_4}(\overline{w_5}\overline{w_2}) + \overline{w_1}w_4(w_3w_5) + w_1\overline{w_4}(w_2 + w_3) + w_1w_4(1)$$

Note that the split is by w_1 and w_4 , not w_1 and w_2

Solution Circuit



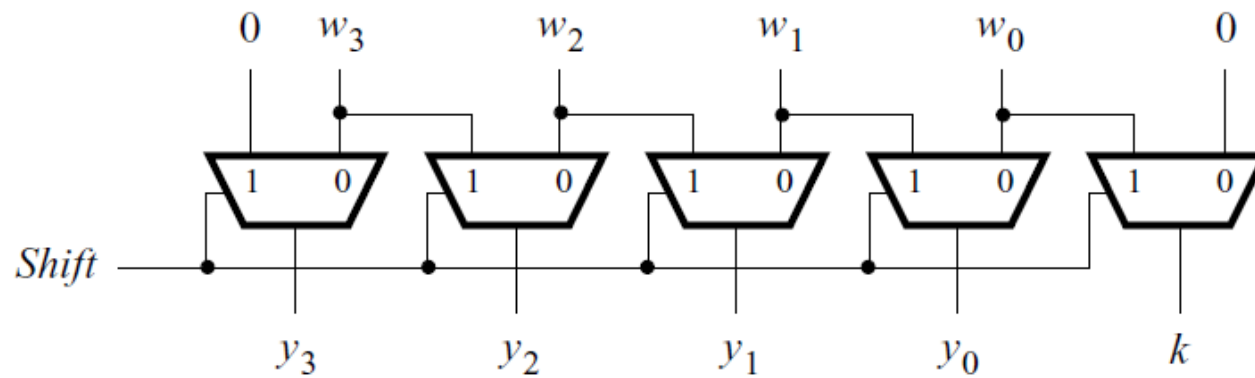
$$\bar{w}_1 \bar{w}_4 (\bar{w}_5 \bar{w}_2) + \bar{w}_1 w_4 (w_3 w_5) + w_1 \bar{w}_4 (w_2 + w_3) + w_1 w_4 (1)$$

[Figure 4.46 from the textbook]

Some Final Things from Chapter 4

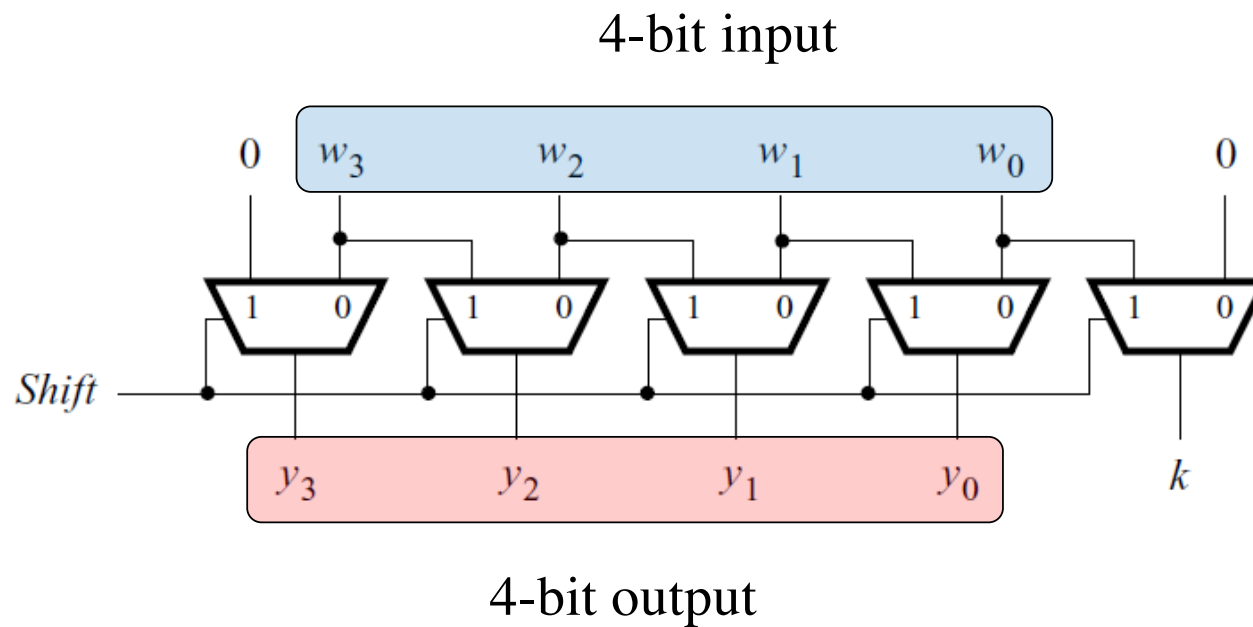
Shifter Circuit

Hold / Shift-Right Circuit



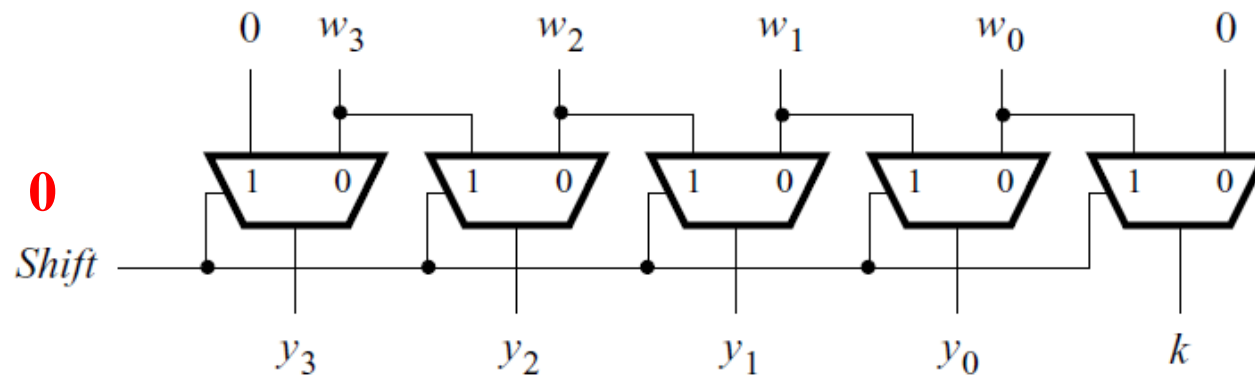
[Figure 4.50 from the textbook]

Hold / Shift-Right Circuit

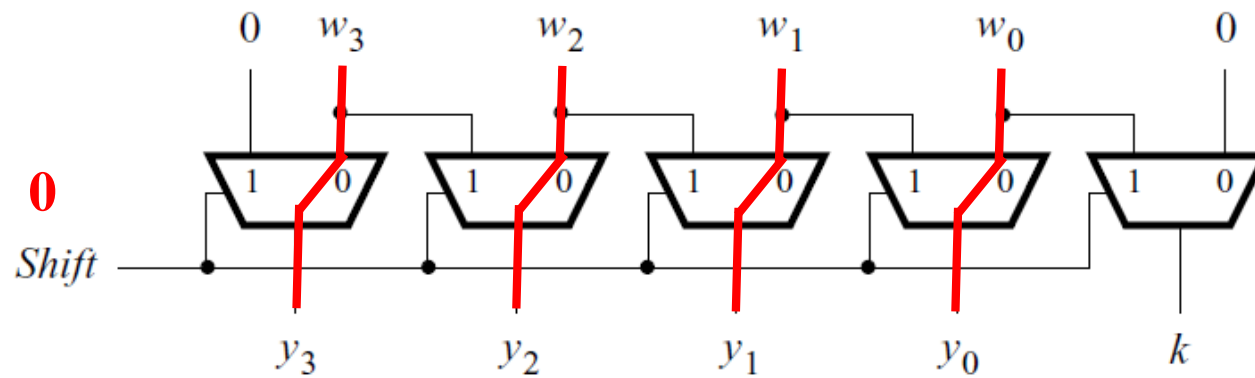


[Figure 4.50 from the textbook]

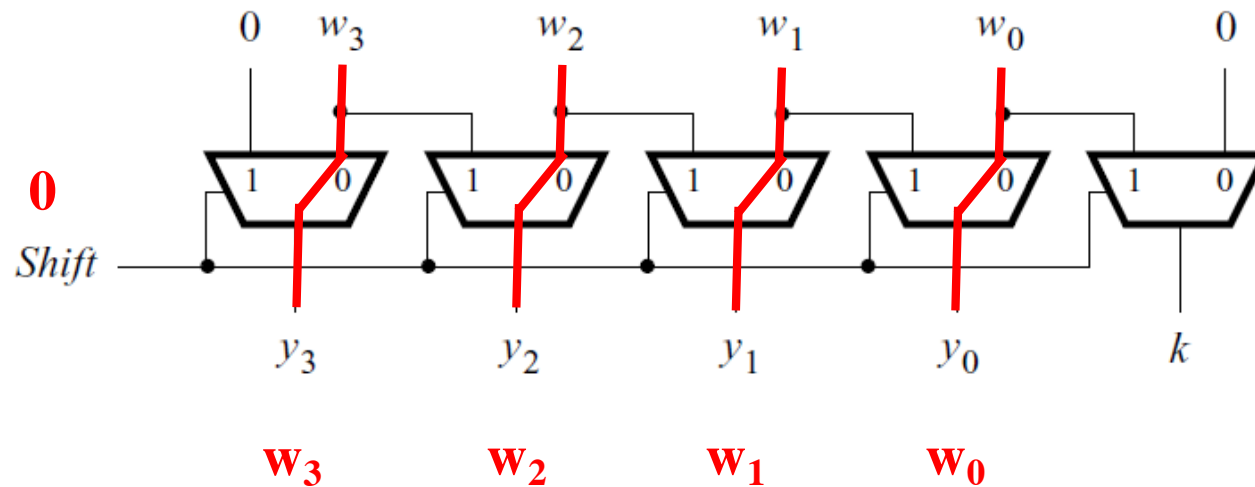
No shift if the control signal is 0



No shift if the control signal is 0

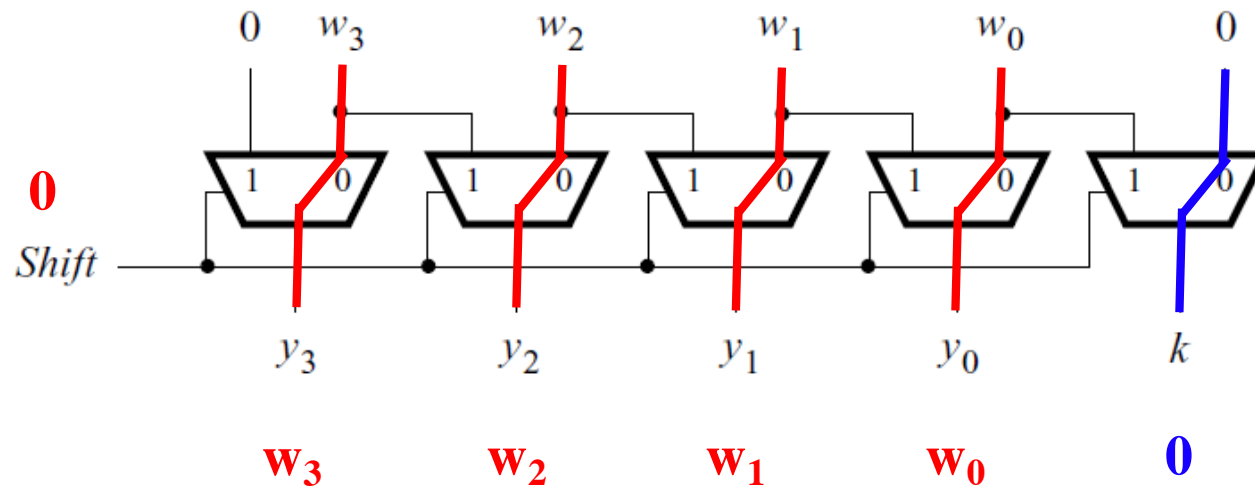


No shift if the control signal is 0



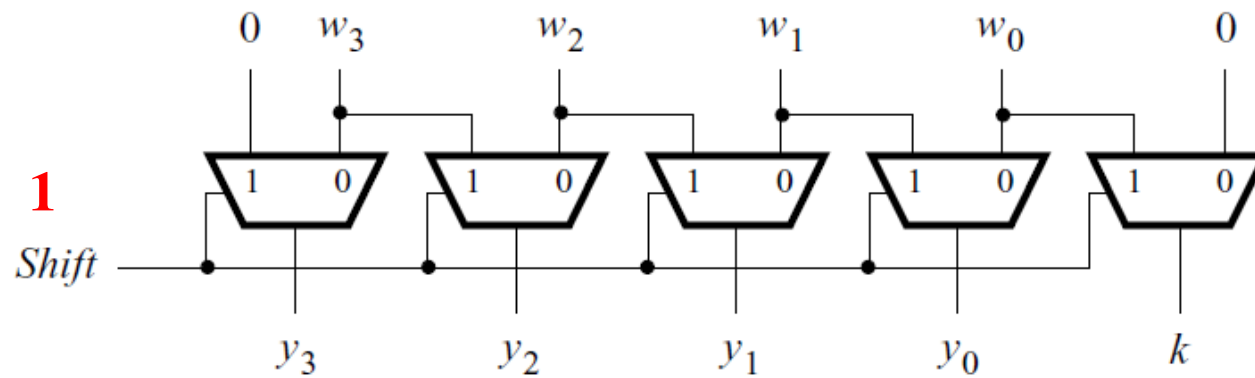
No shift in this case.

No shift if the control signal is 0

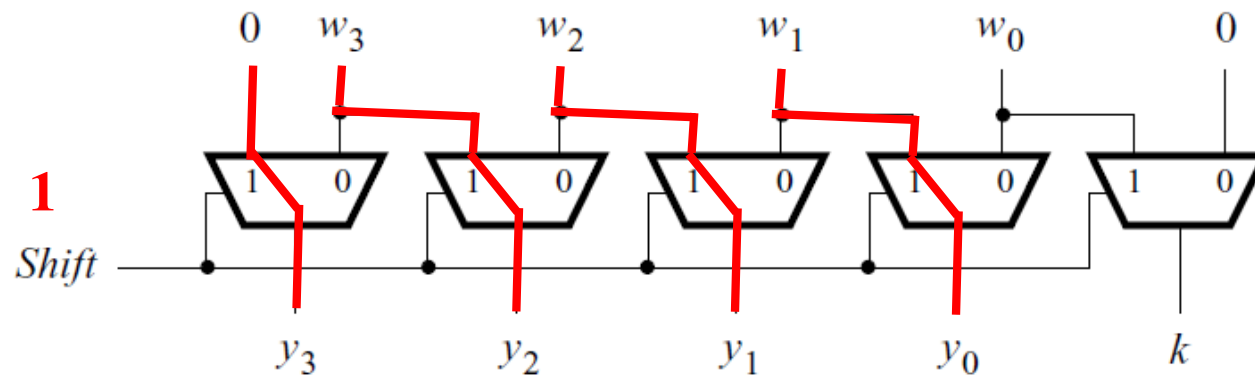


No shift in this case.

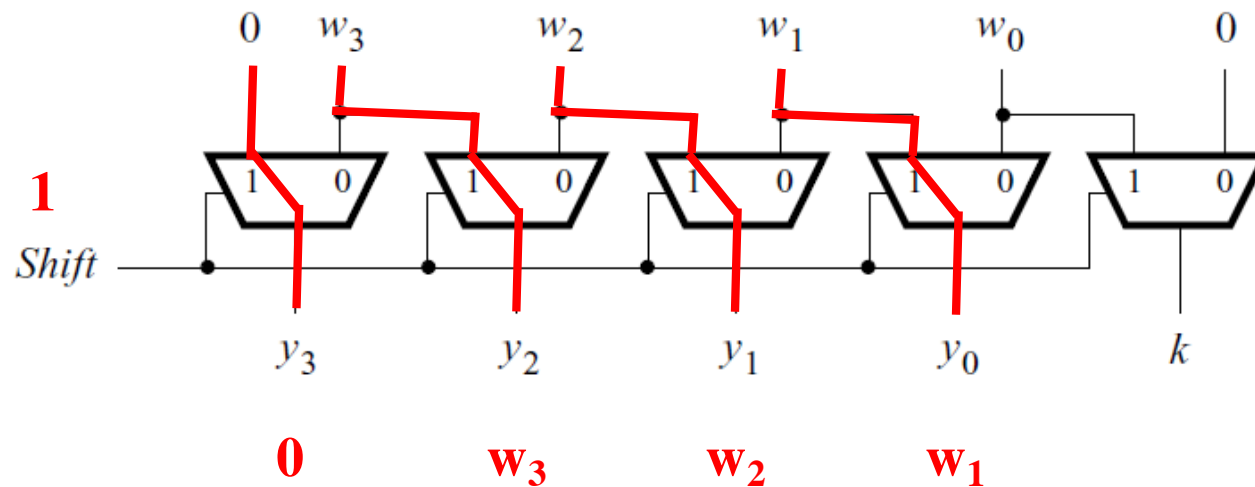
Shift right if the control signal is 1



Shift right if the control signal is 1

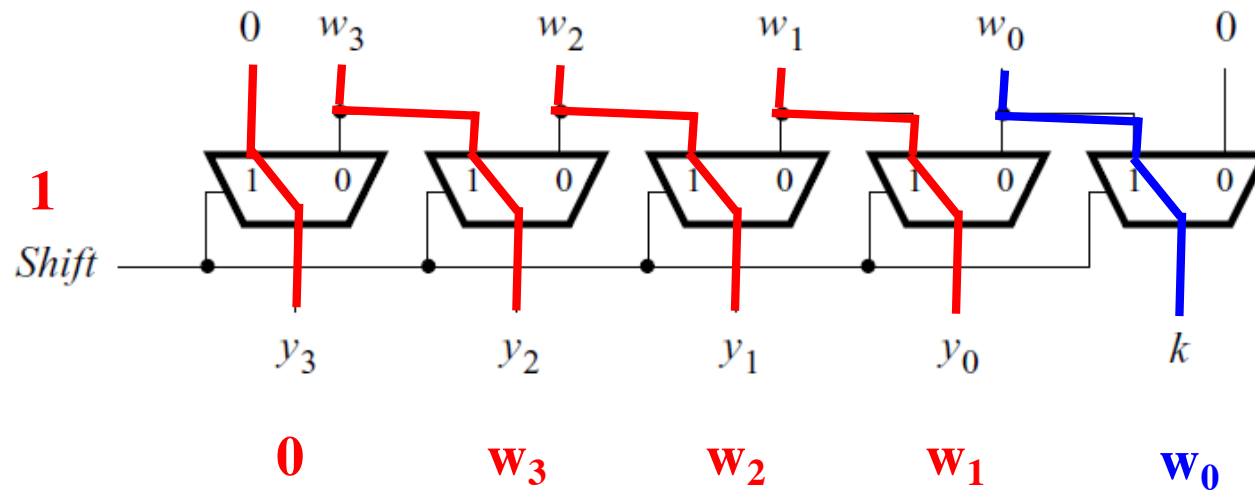


Shift right if the control signal is 1



Shift to the right by 1 bit

A shifter circuit



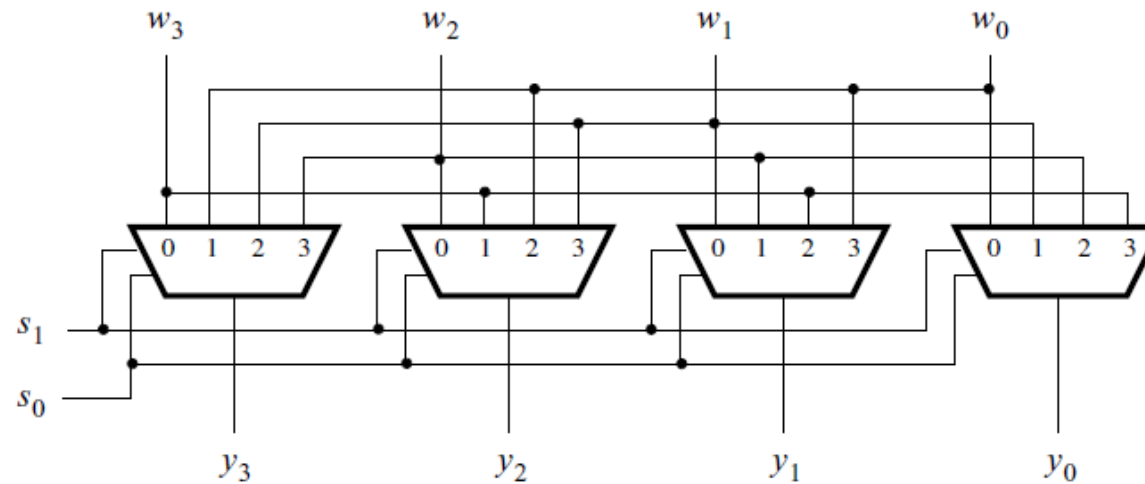
Shift to the right by 1 bit

Barrel Shifter

A barrel shifter circuit

s_1	s_0	y_3	y_2	y_1	y_0
0	0	w_3	w_2	w_1	w_0
0	1	w_0	w_3	w_2	w_1
1	0	w_1	w_0	w_3	w_2
1	1	w_2	w_1	w_0	w_3

(a) Truth table



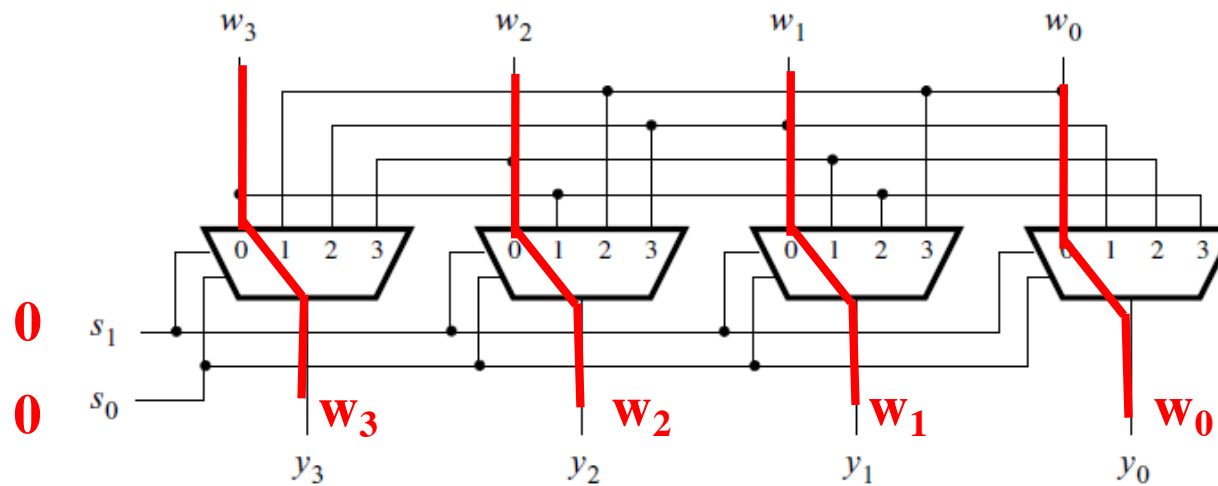
(b) Circuit

[Figure 4.51 from the textbook]

A barrel shifter circuit

s_1	s_0	y_3	y_2	y_1	y_0
0	0	w_3	w_2	w_1	w_0
0	1	w_0	w_3	w_2	w_1
1	0	w_1	w_0	w_3	w_2
1	1	w_2	w_1	w_0	w_3

(a) Truth table



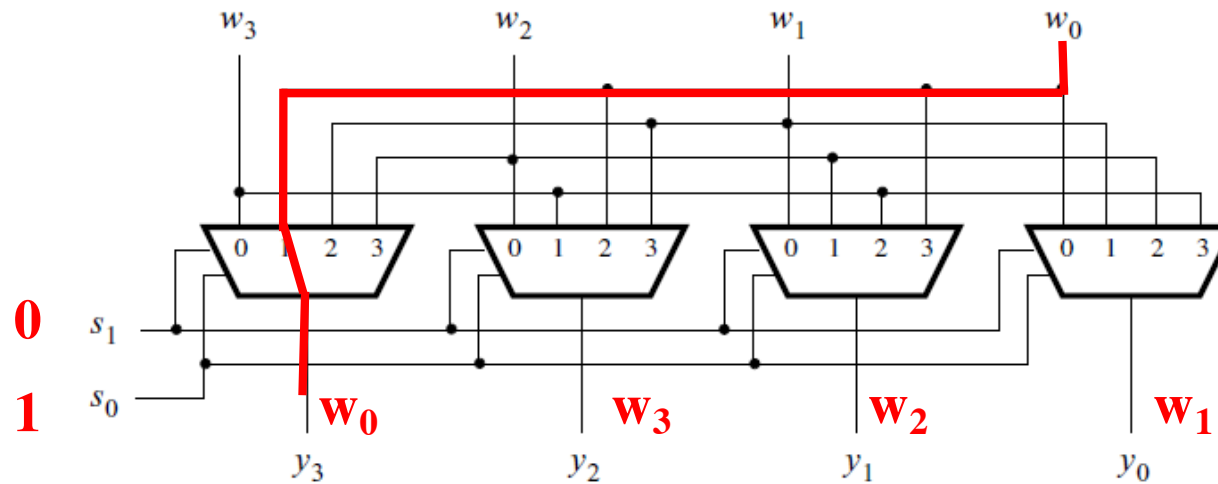
(b) Circuit

[Figure 4.51 from the textbook]

A barrel shifter circuit

s_1	s_0	y_3	y_2	y_1	y_0
0	0	w_3	w_2	w_1	w_0
0	1	w_0	w_3	w_2	w_1
1	0	w_1	w_0	w_3	w_2
1	1	w_2	w_1	w_0	w_3

(a) Truth table



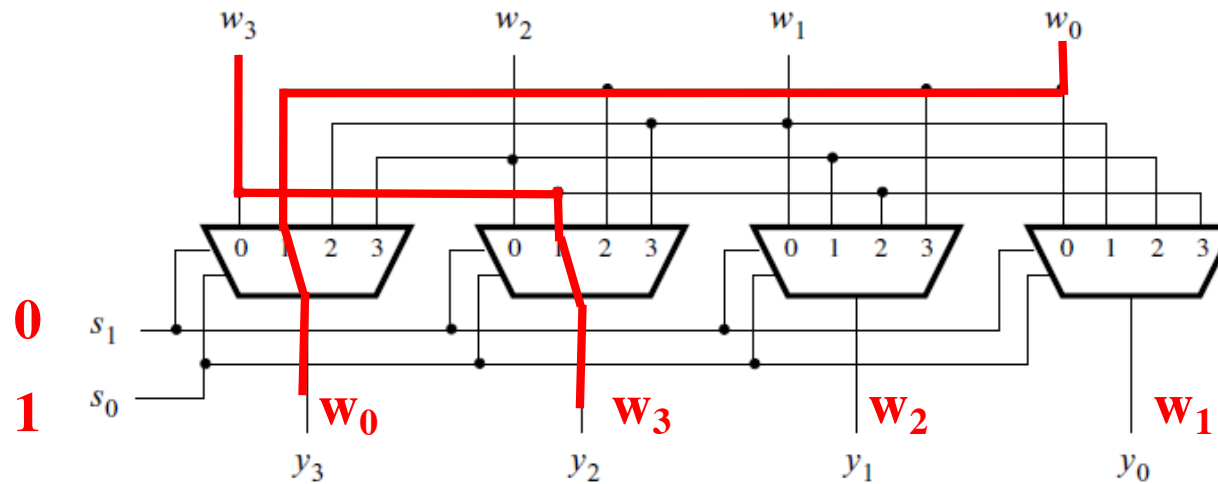
(b) Circuit

[Figure 4.51 from the textbook]

A barrel shifter circuit

s_1	s_0	y_3	y_2	y_1	y_0
0	0	w_3	w_2	w_1	w_0
0	1	w_0	w_3	w_2	w_1
1	0	w_1	w_0	w_3	w_2
1	1	w_2	w_1	w_0	w_3

(a) Truth table



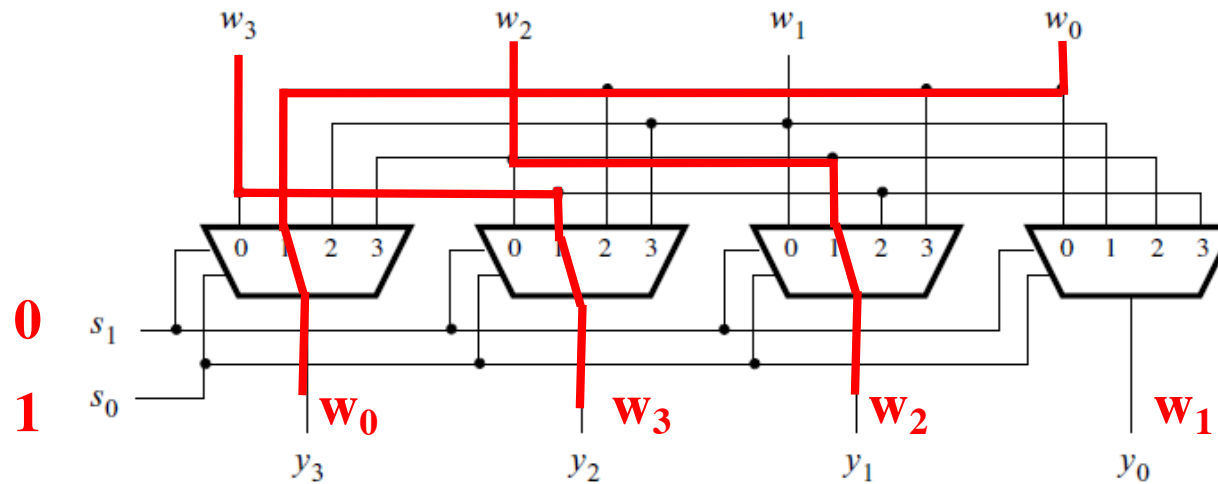
(b) Circuit

[Figure 4.51 from the textbook]

A barrel shifter circuit

s_1	s_0	y_3	y_2	y_1	y_0
0	0	w_3	w_2	w_1	w_0
0	1	w_0	w_3	w_2	w_1
1	0	w_1	w_0	w_3	w_2
1	1	w_2	w_1	w_0	w_3

(a) Truth table



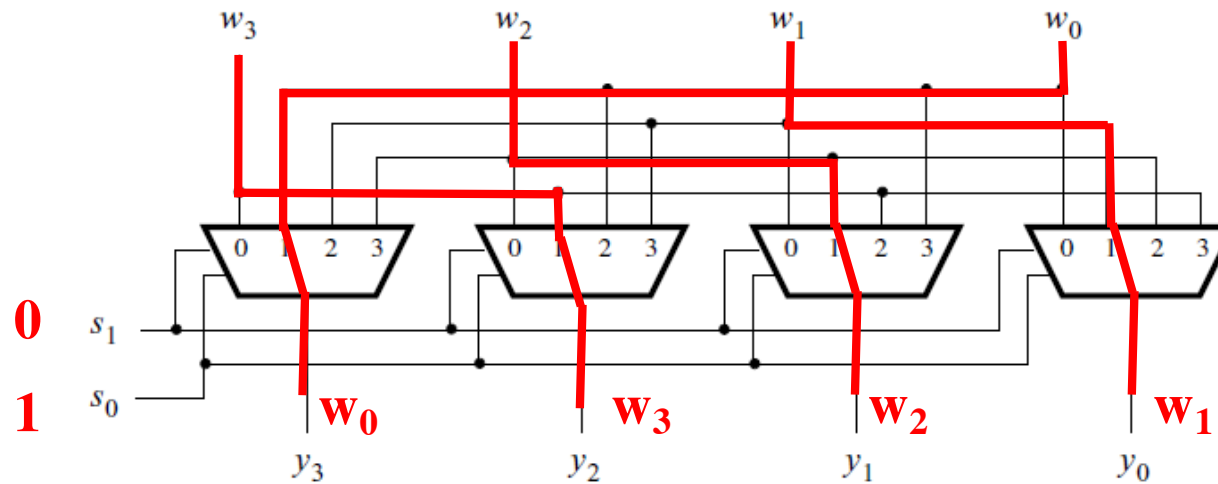
(b) Circuit

[Figure 4.51 from the textbook]

A barrel shifter circuit

s_1	s_0	y_3	y_2	y_1	y_0
0	0	w_3	w_2	w_1	w_0
0	1	w_0	w_3	w_2	w_1
1	0	w_1	w_0	w_3	w_2
1	1	w_2	w_1	w_0	w_3

(a) Truth table



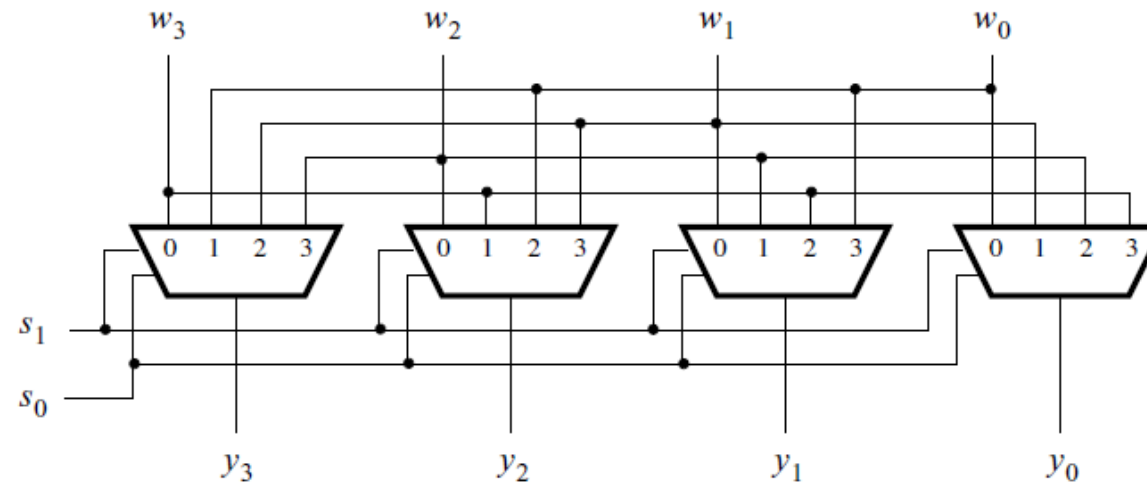
(b) Circuit

[Figure 4.51 from the textbook]

A barrel shifter circuit

s_1	s_0	y_3	y_2	y_1	y_0
0	0	w_3	w_2	w_1	w_0
0	1	w_0	w_3	w_2	w_1
1	0	w_1	w_0	w_3	w_2
1	1	w_2	w_1	w_0	w_3

(a) Truth table

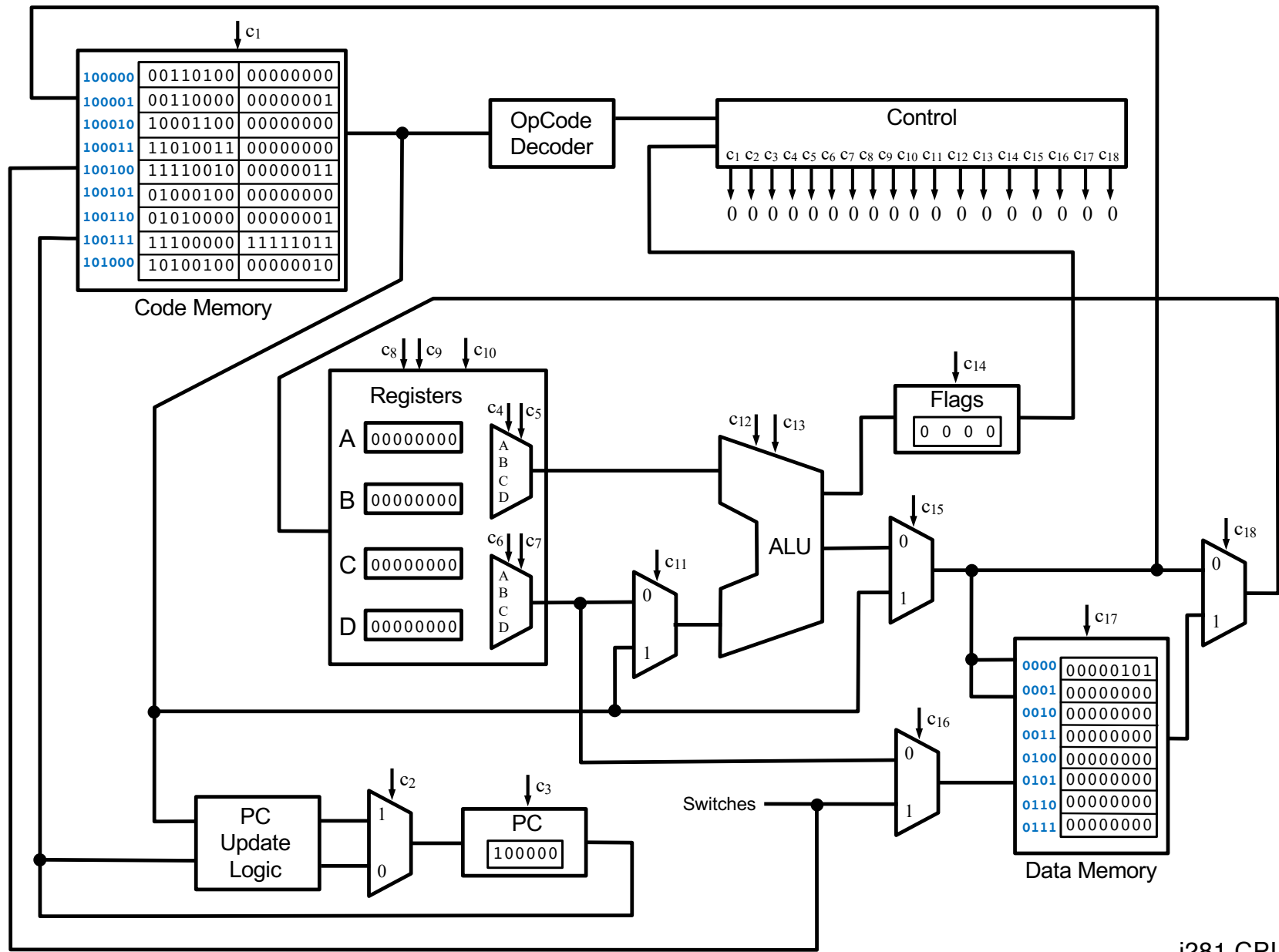


(b) Circuit

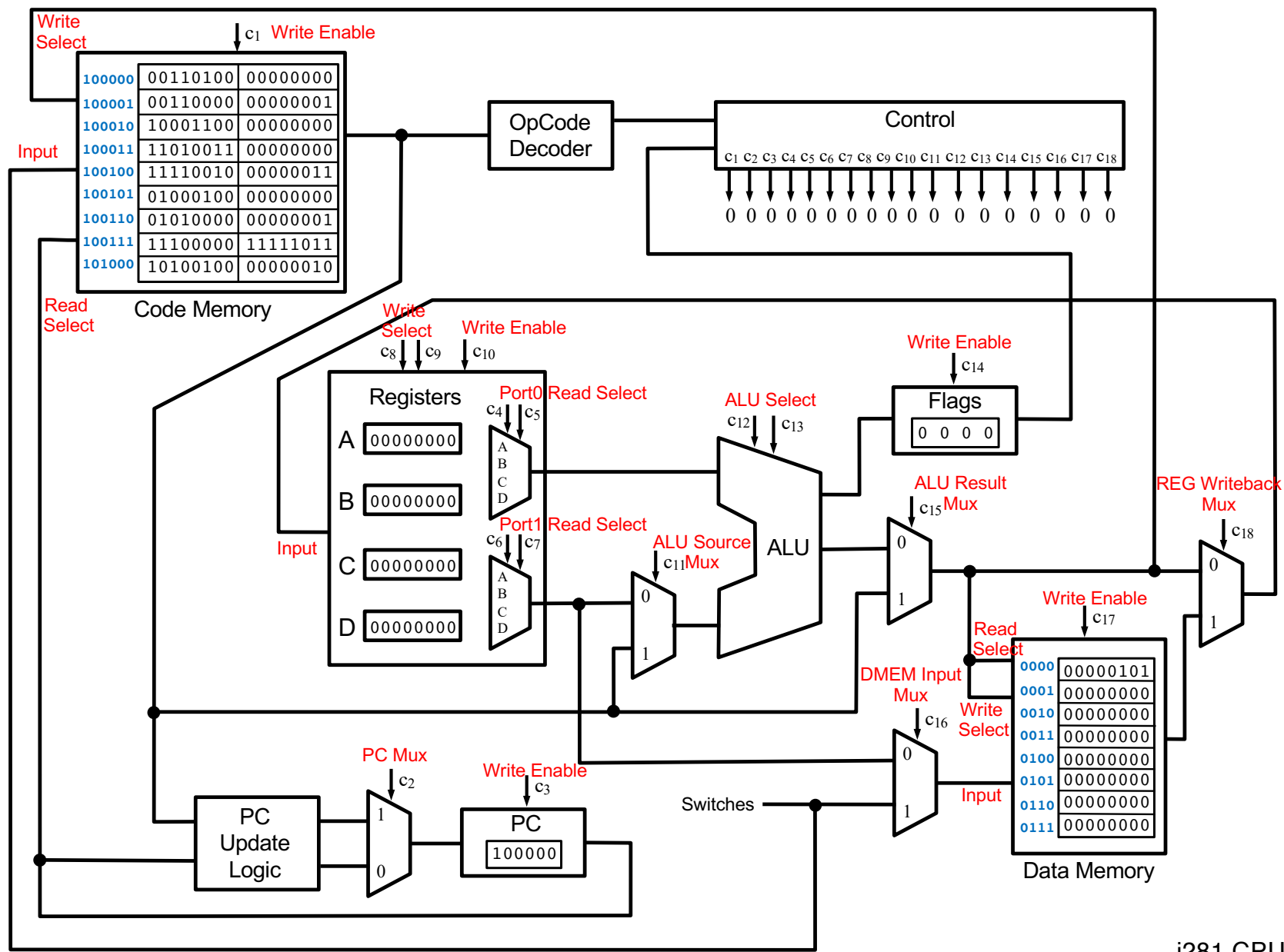
[Figure 4.51 from the textbook]

Arithmetic Logic Unit (ALU)

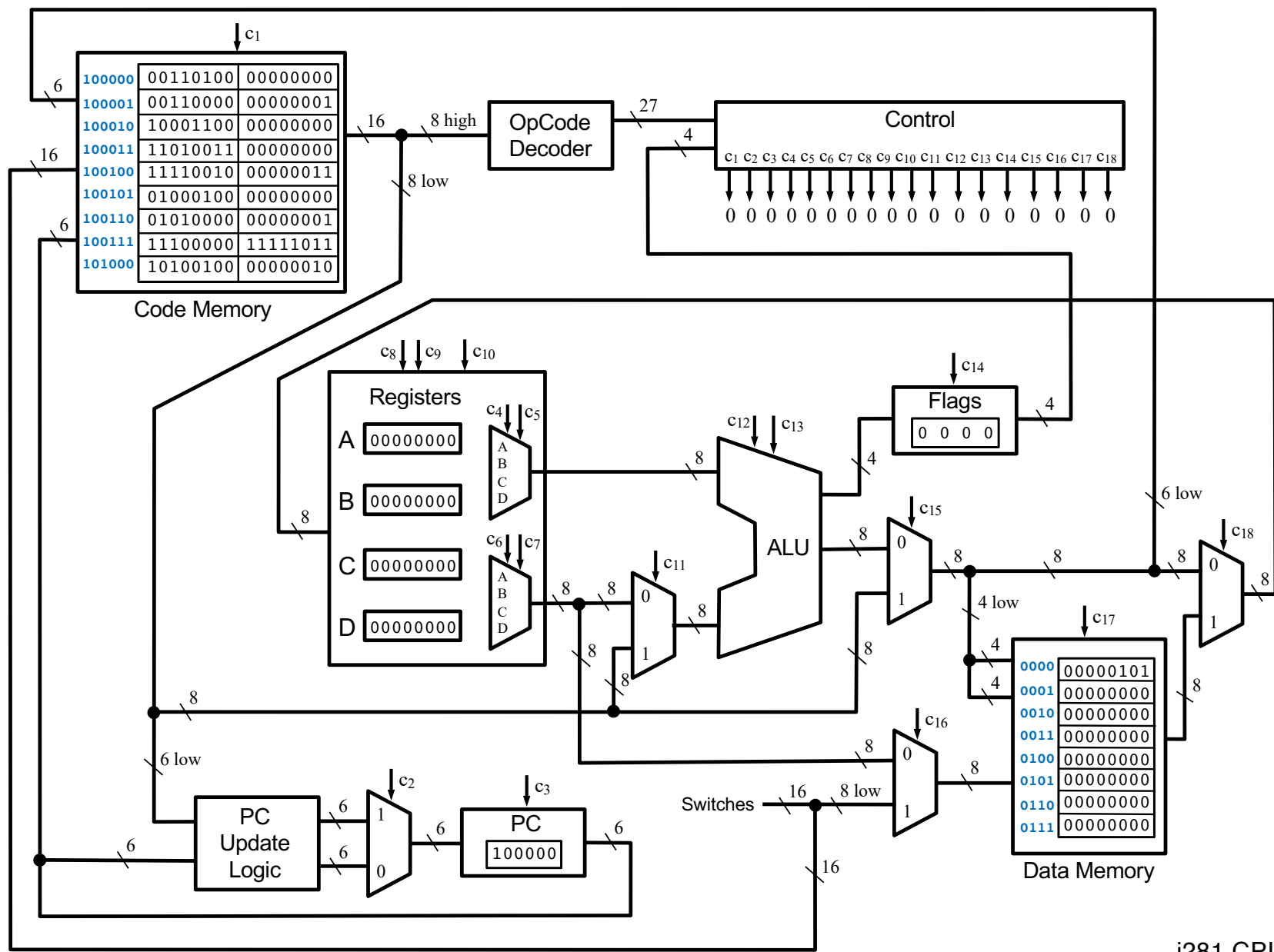
Shifter Circuit of the i281 CPU



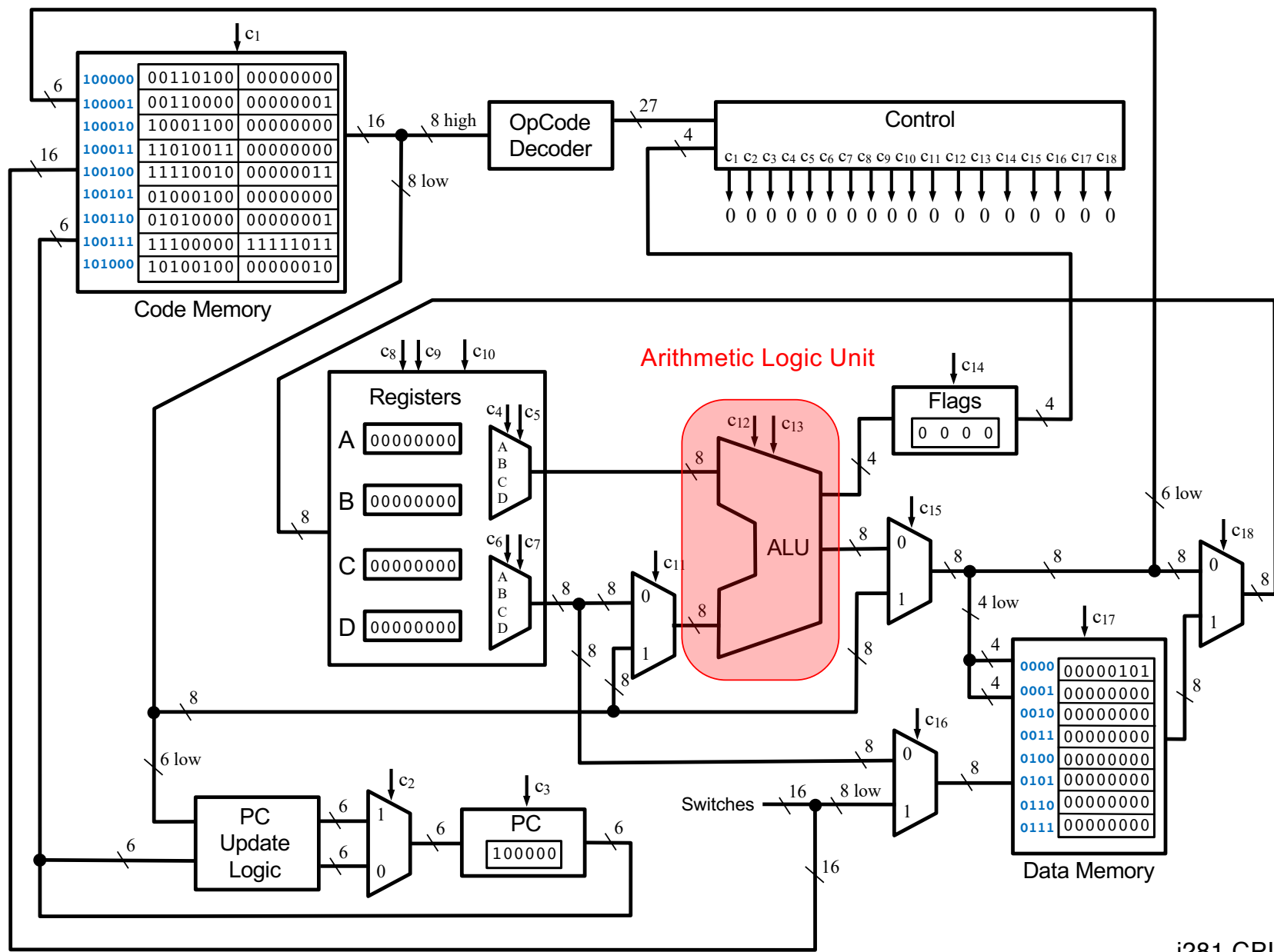
i281 CPU



i281 CPU

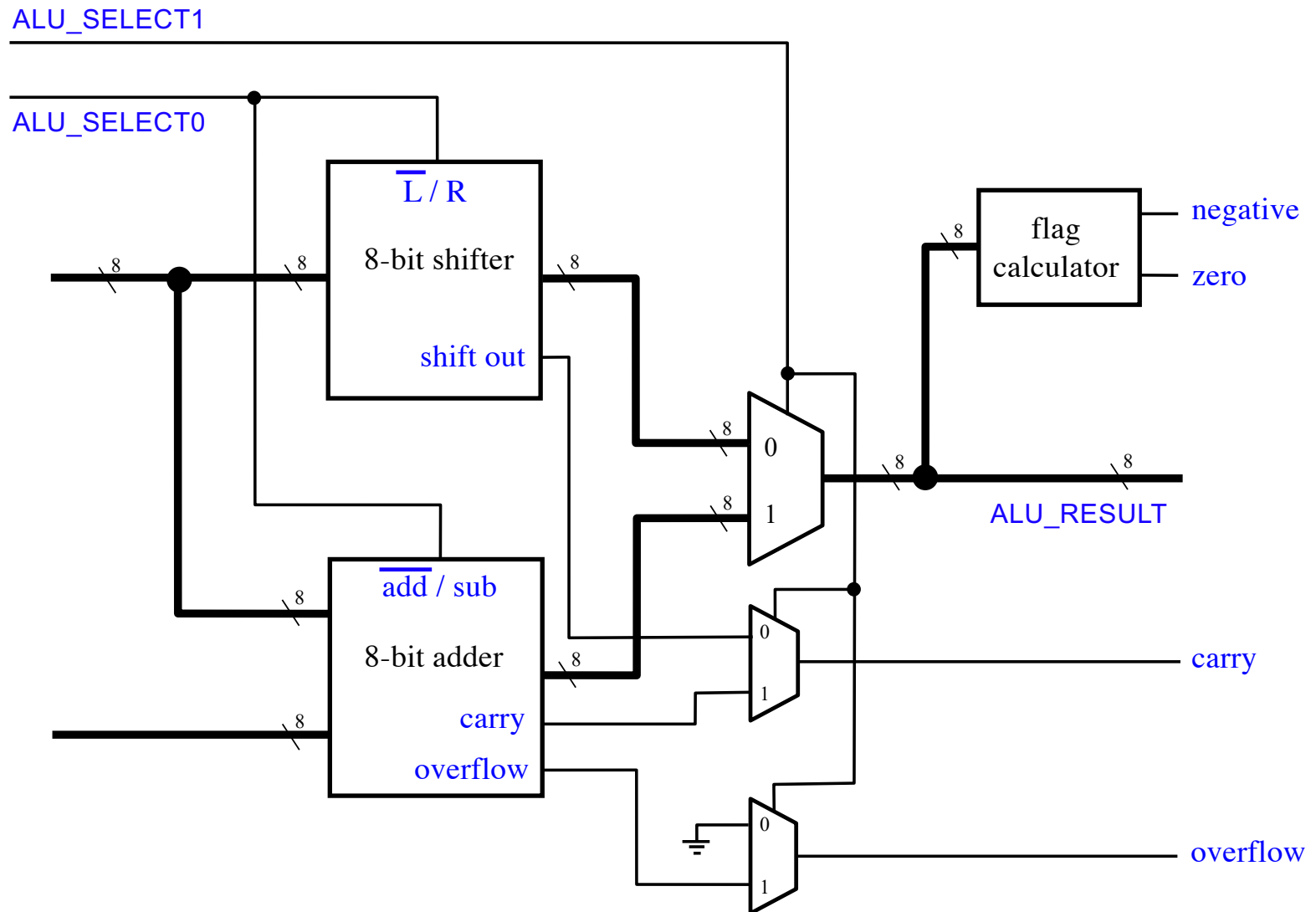


i281 CPU

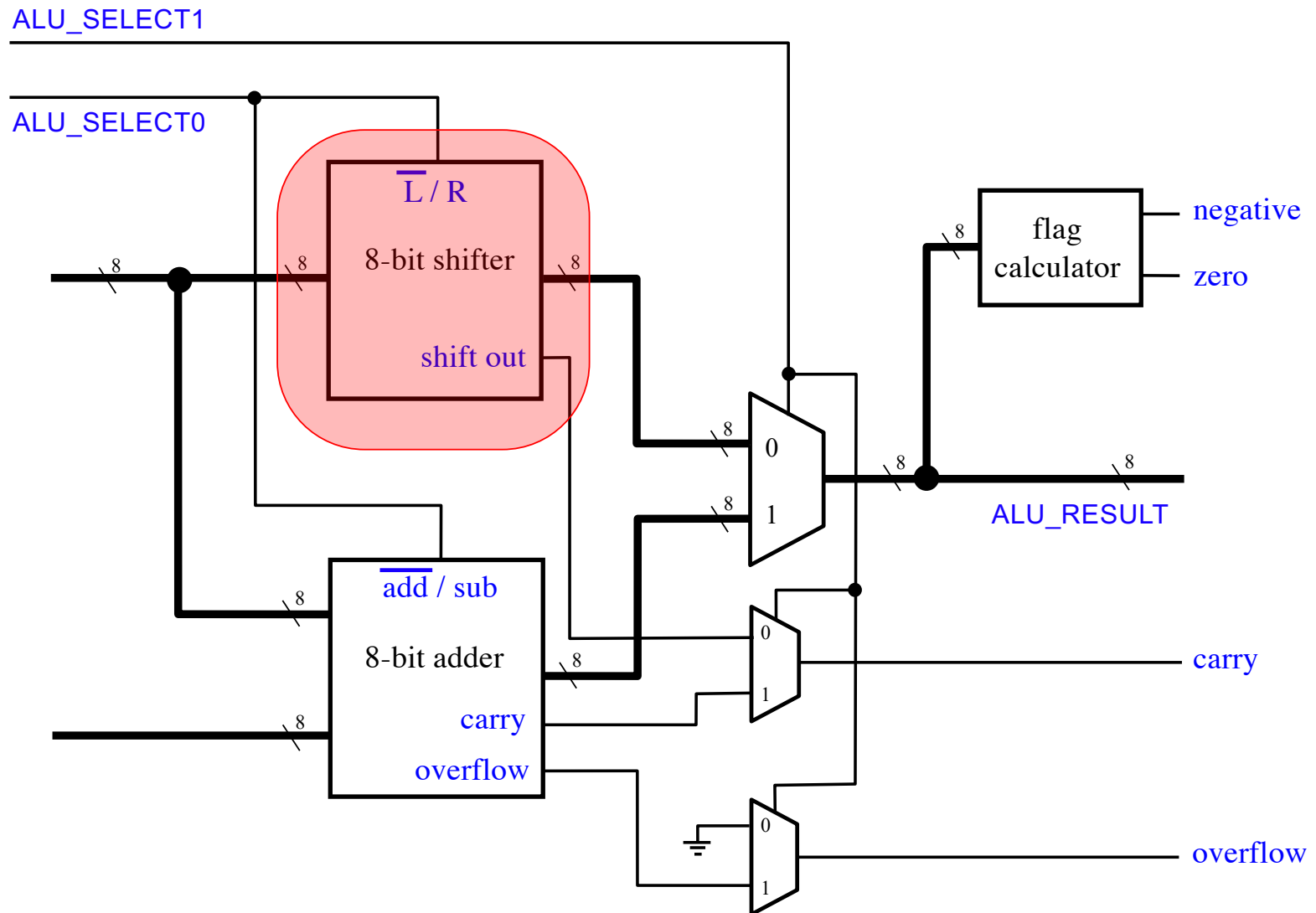


i281 CPU

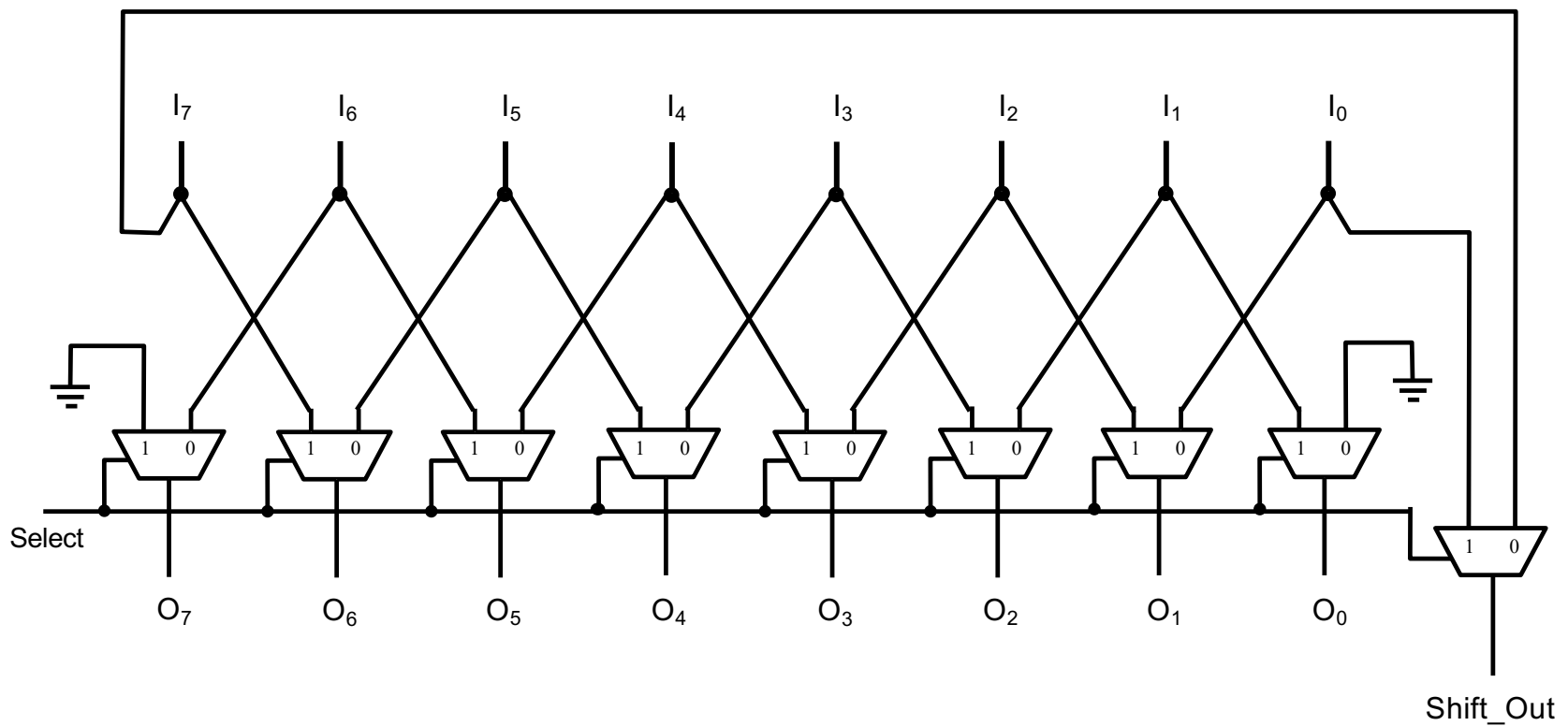
The ALU



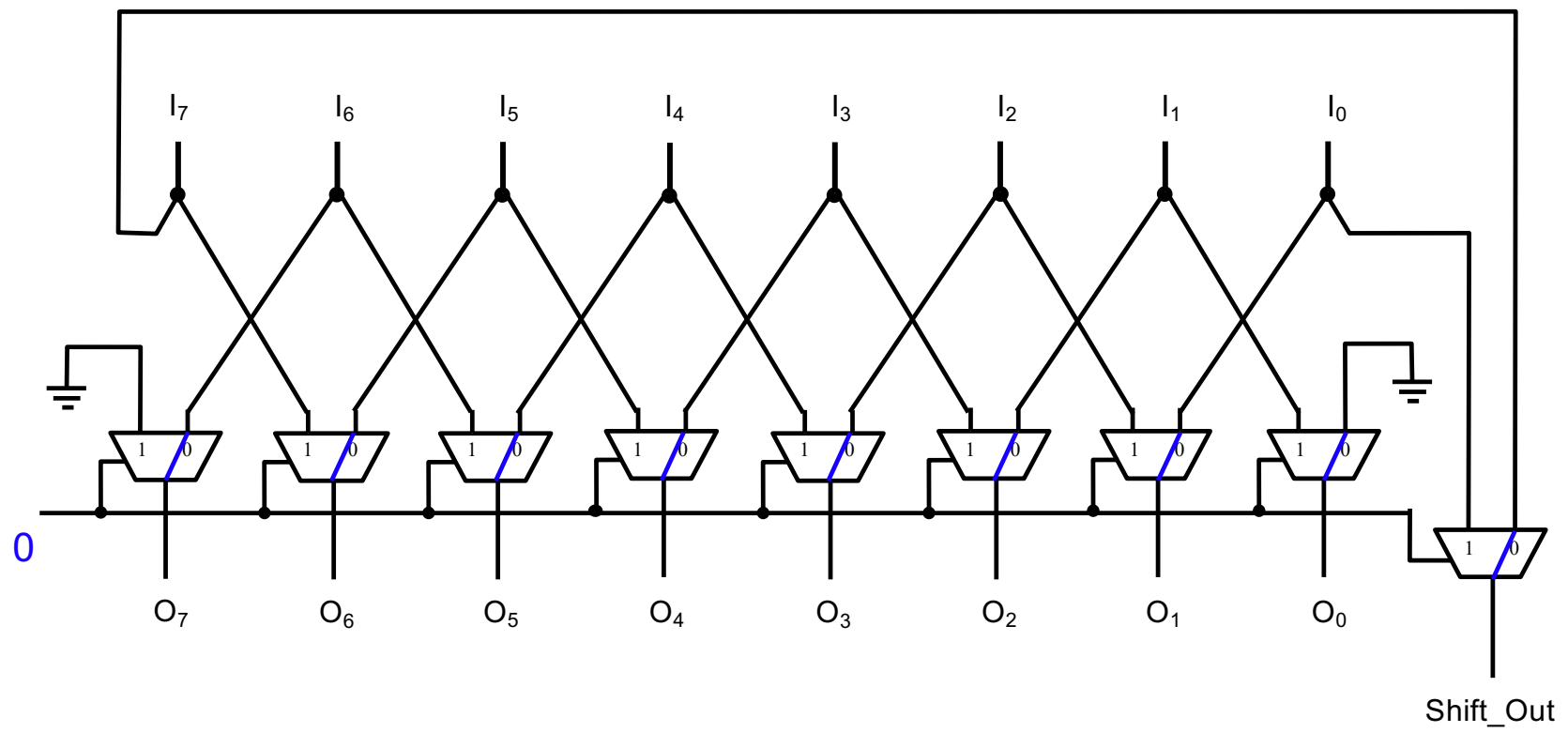
The Shifter Circuit



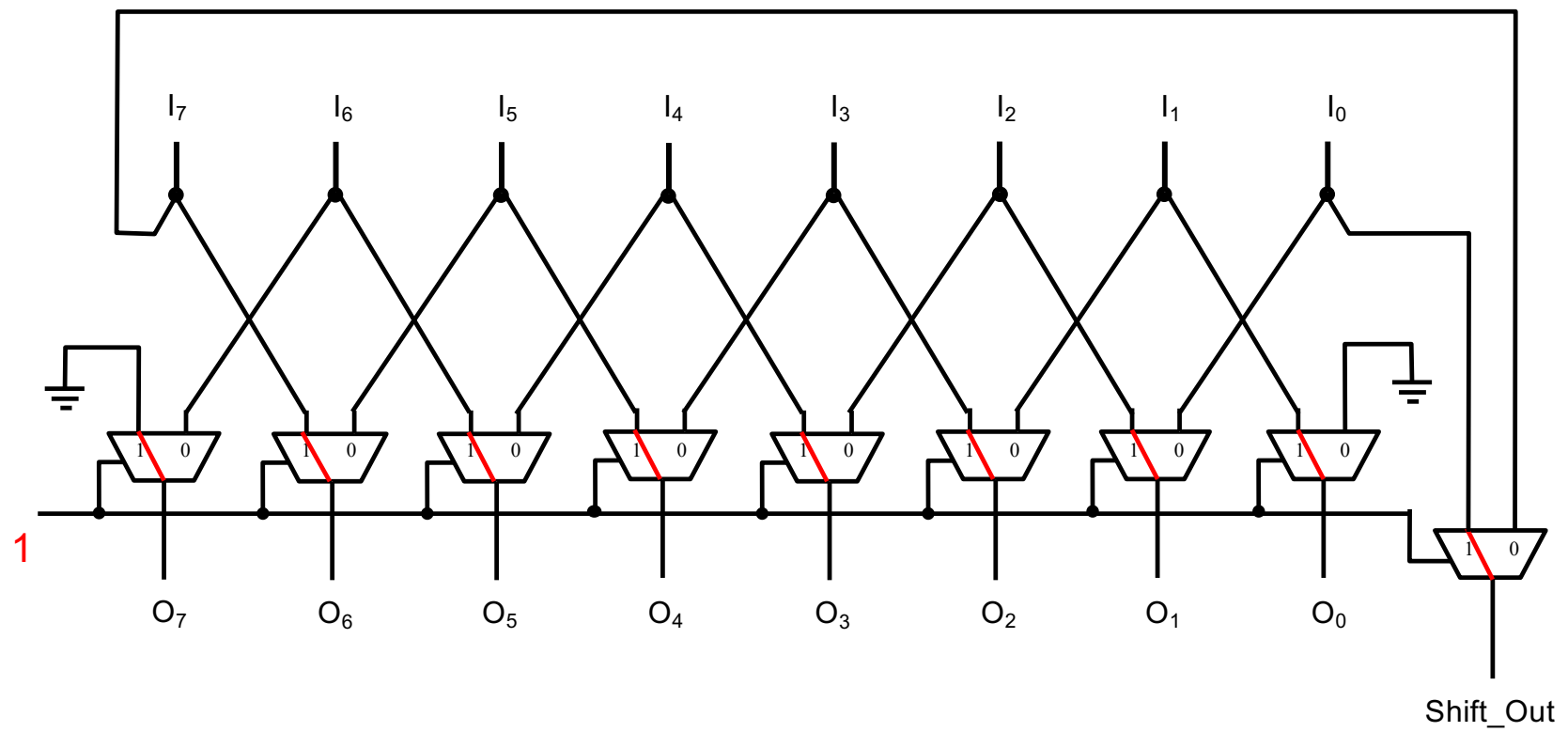
The i281 CPU Shifter Circuit



i281 CPU

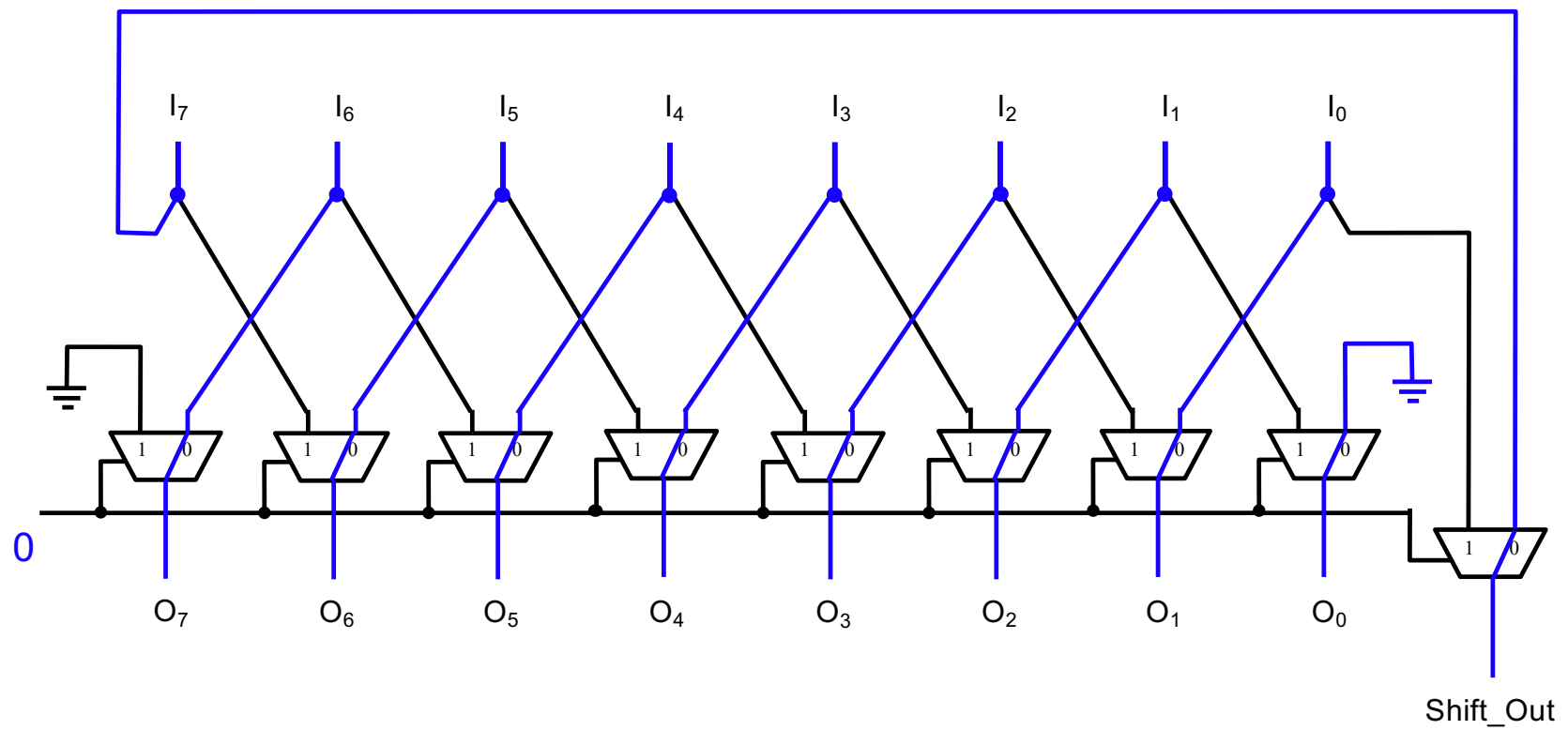


i281 CPU



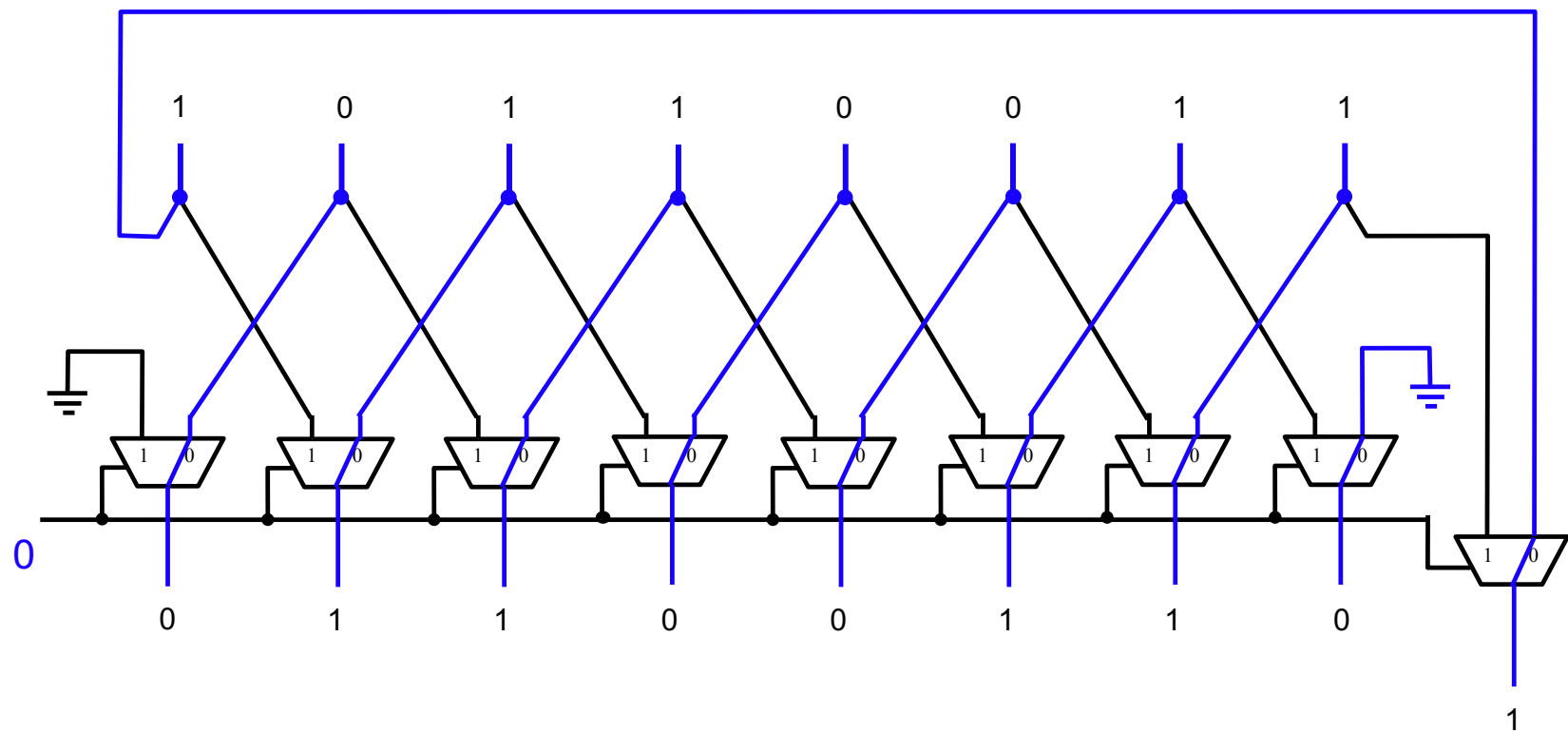
i281 CPU

Shift Left



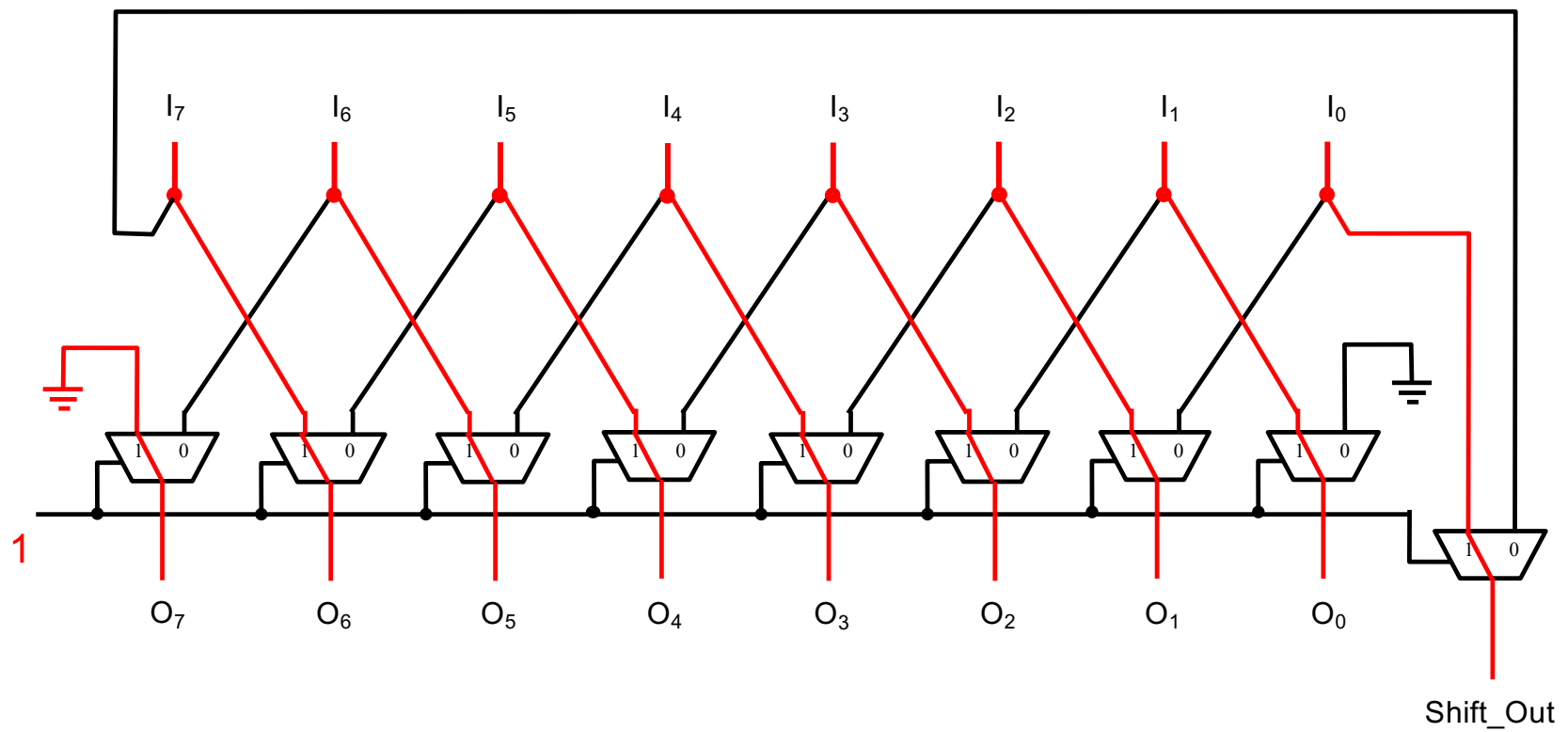
i281 CPU

Shift Left



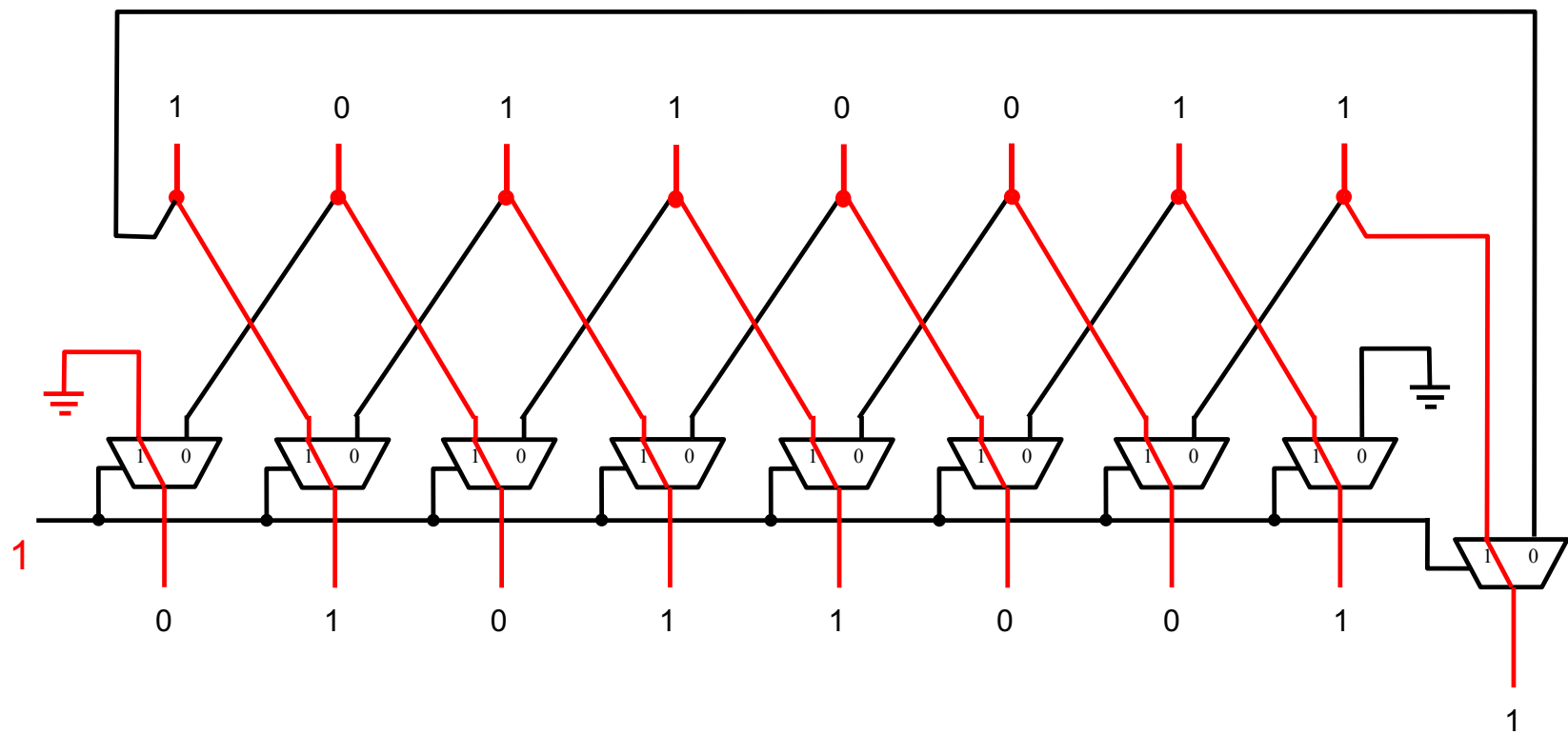
i281 CPU

Shift Right



i281 CPU

Shift Right



i281 CPU

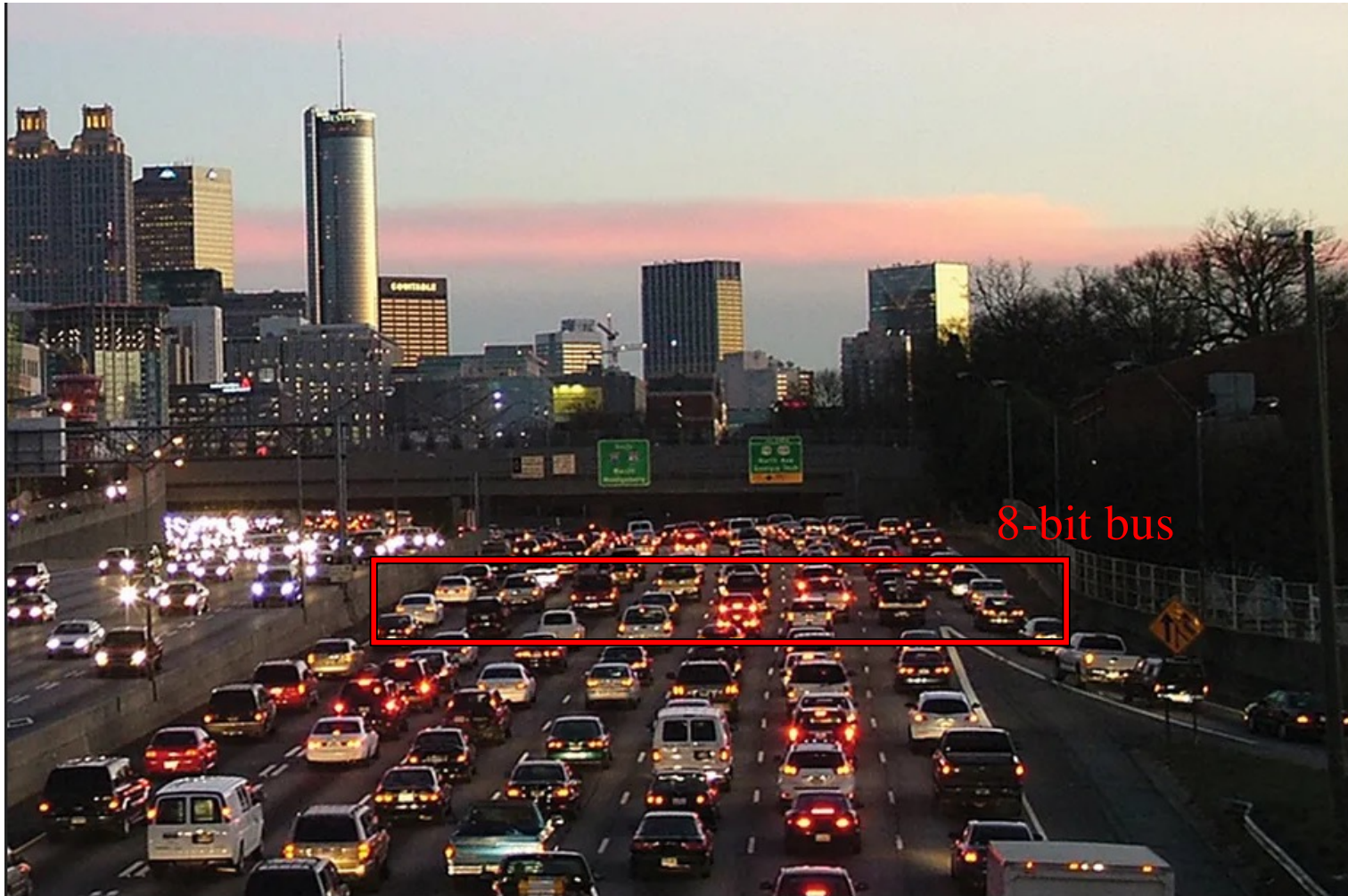
Bus Multiplexer

Bus Analogy



[<https://atlanta.curbed.com/2018/10/8/17949734/study-evening-rush-hour-slows-traffic-23-percent>]

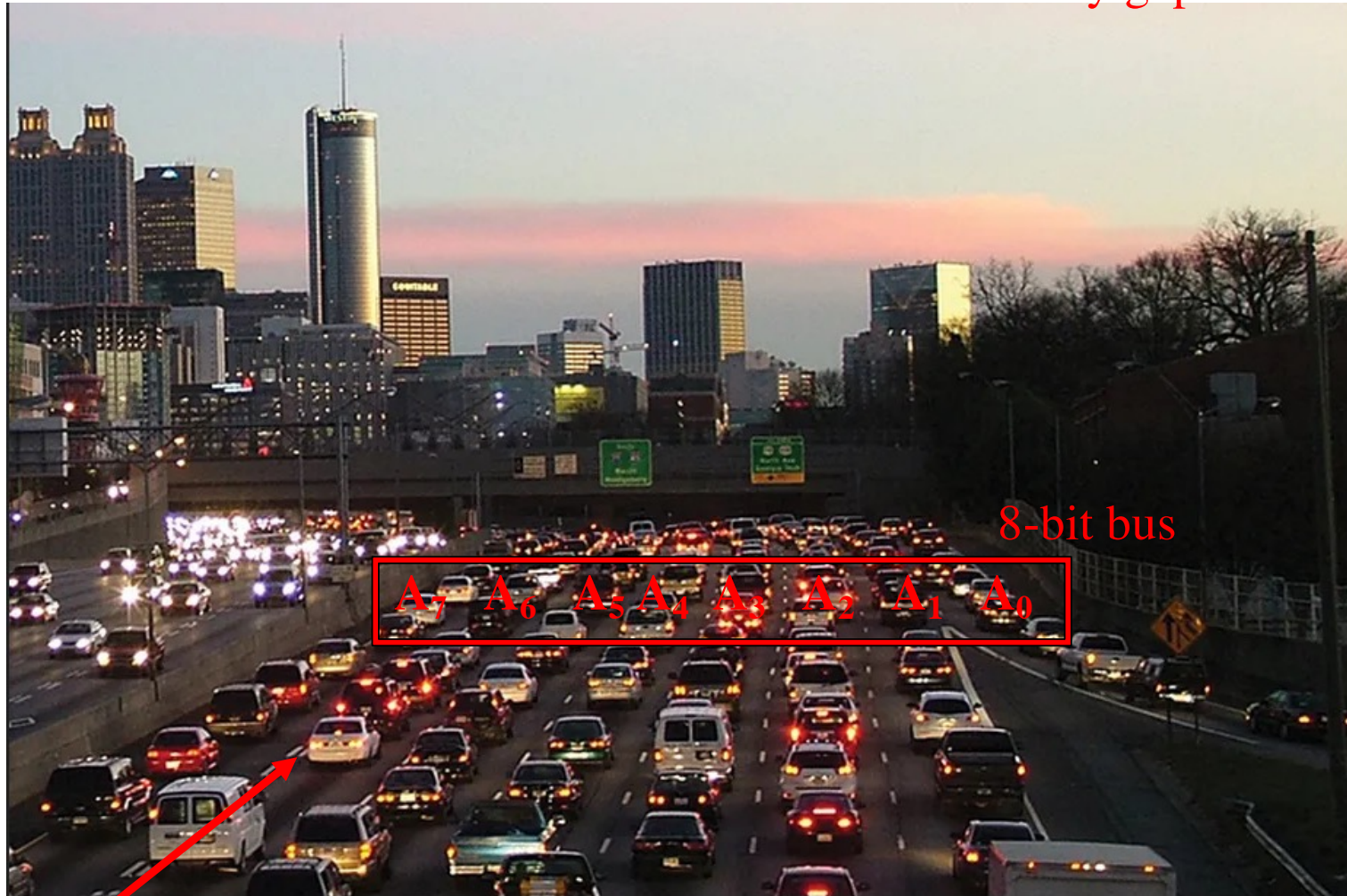
Bus Analogy



[<https://atlanta.curbed.com/2018/10/8/17949734/study-evening-rush-hour-slows-traffic-23-percent>]

Bus Analogy

For each lane:
every car is a 1,
every gap is a 0.



A_7 A_6 A_5 A_4 A_3 A_2 A_1 A_0

8-bit bus

Time

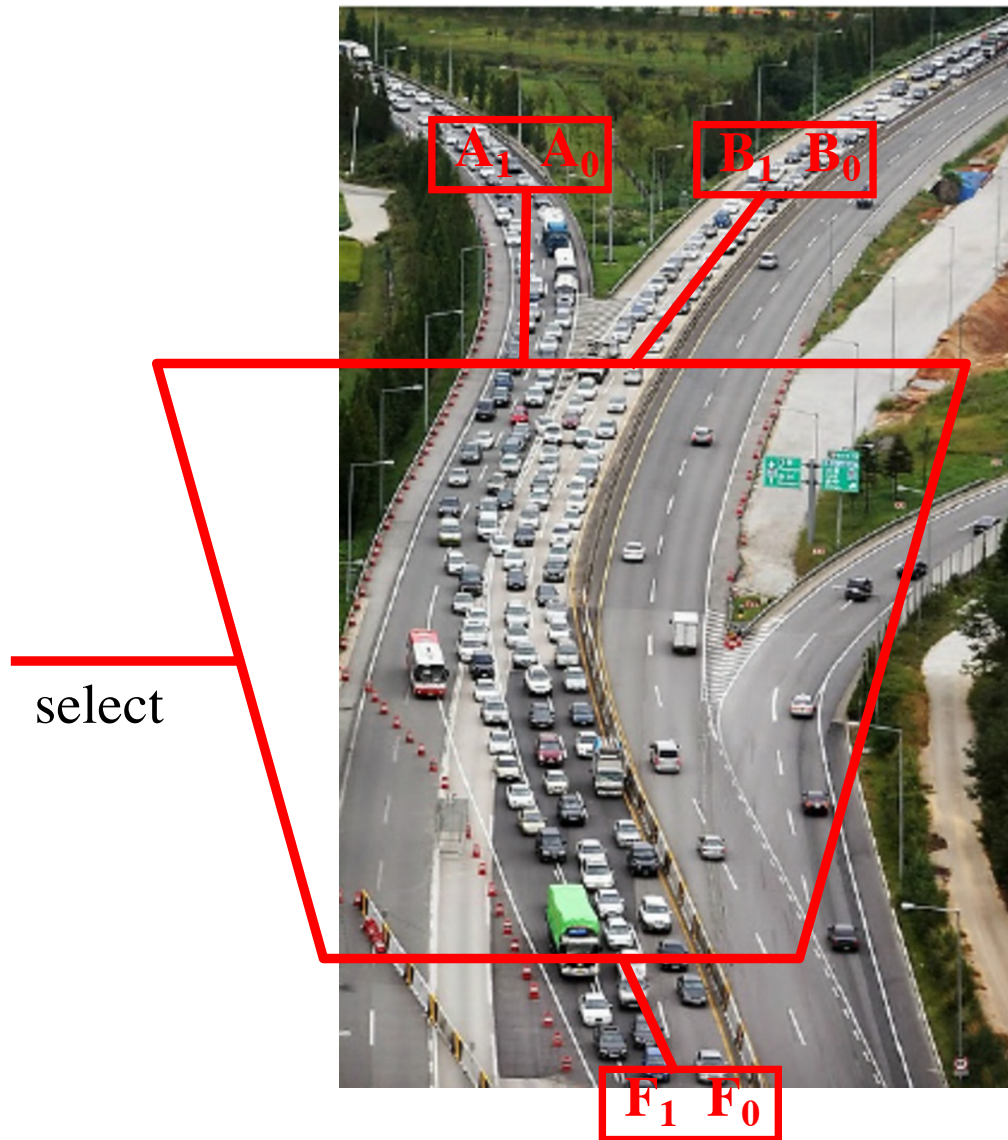
[<https://atlanta.curbed.com/2018/10/8/17949734/study-evening-rush-hour-slows-traffic-23-percent>]

Bus Multiplexer Analogy

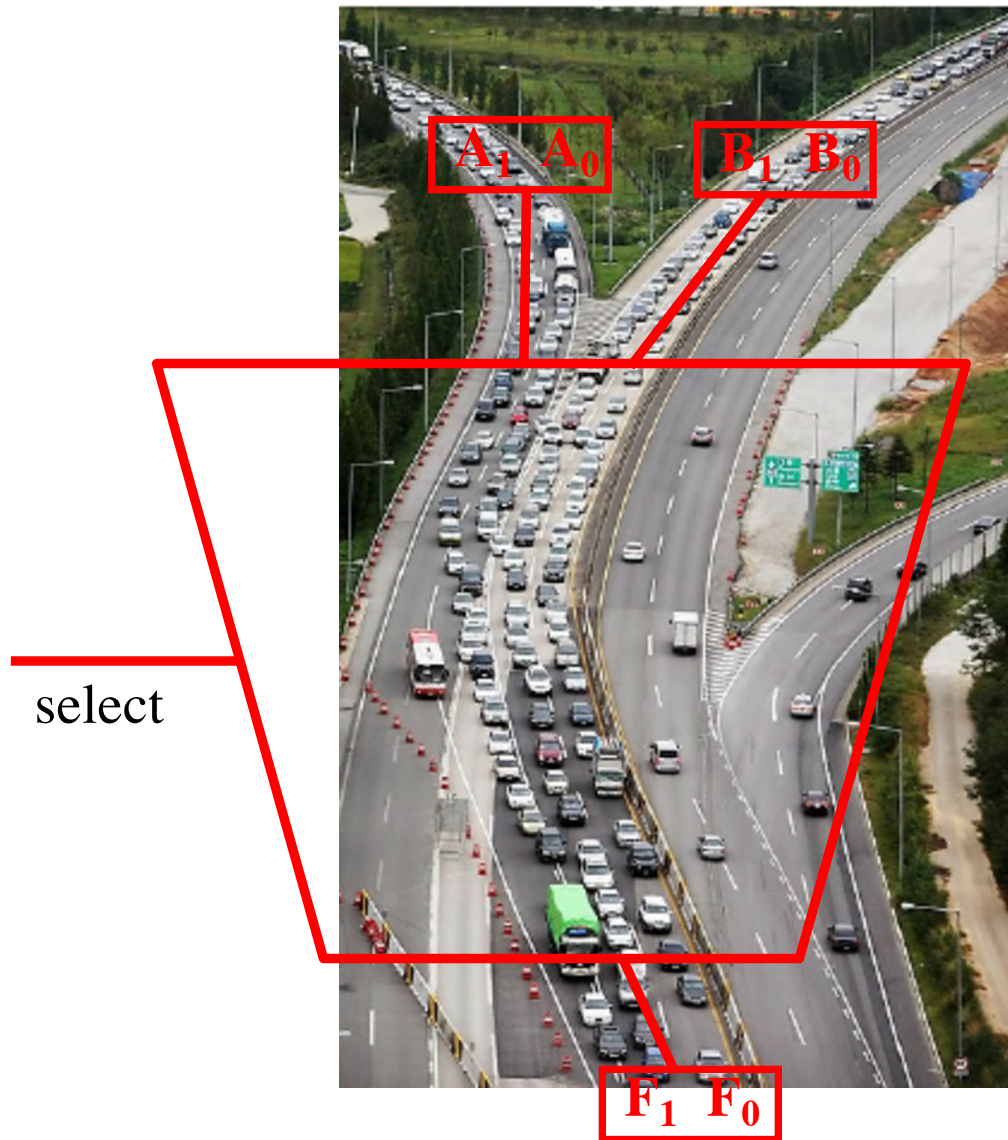


[https://www.researchgate.net/figure/An-example-of-traffic-jam-around-a-junction-where-several-lanes-join_fig1_221567960]]

Bus Multiplexer Analogy

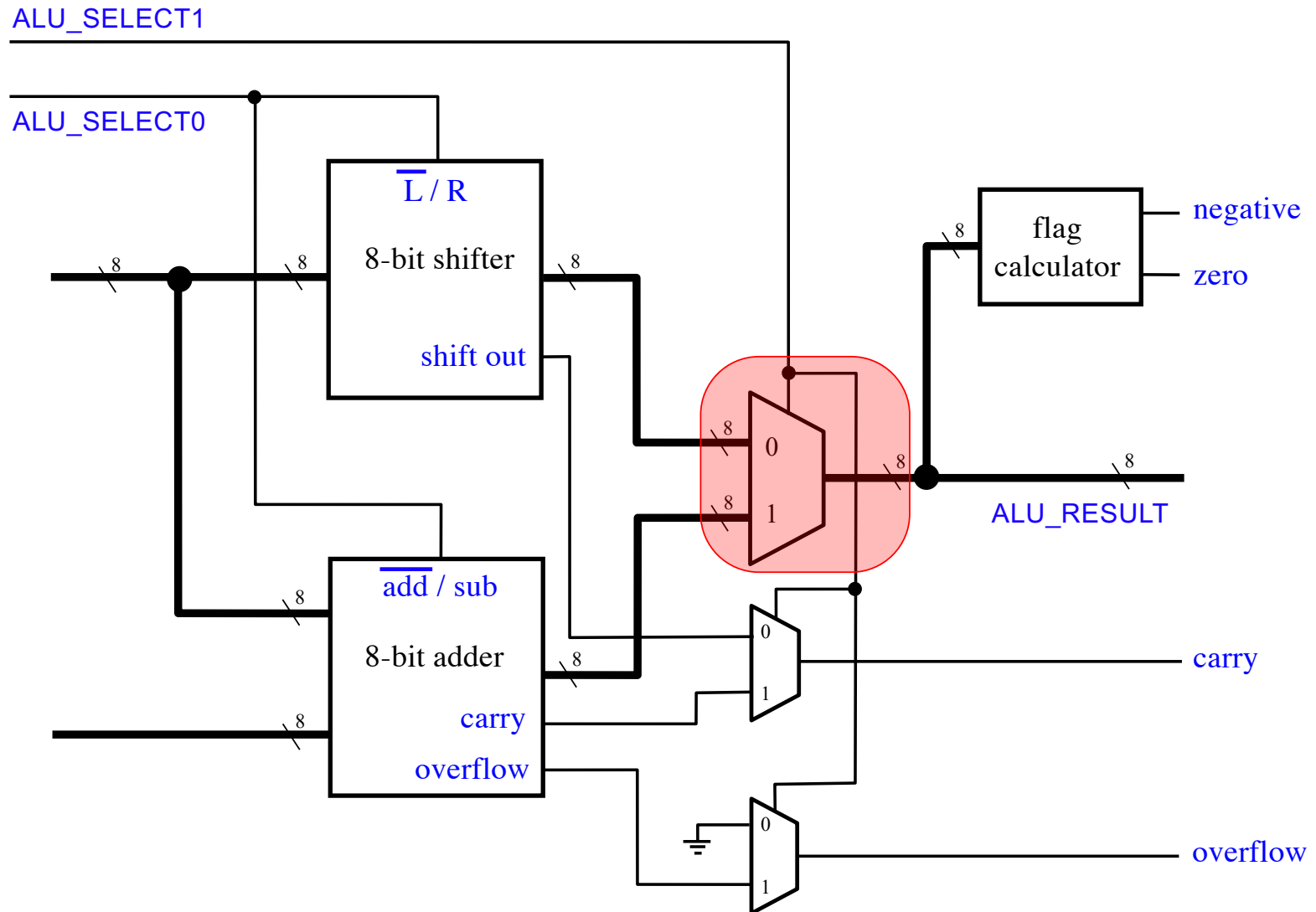


Bus Multiplexer Analogy

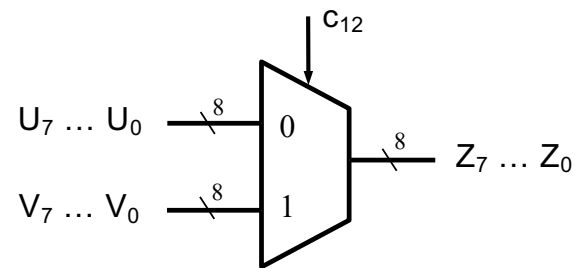


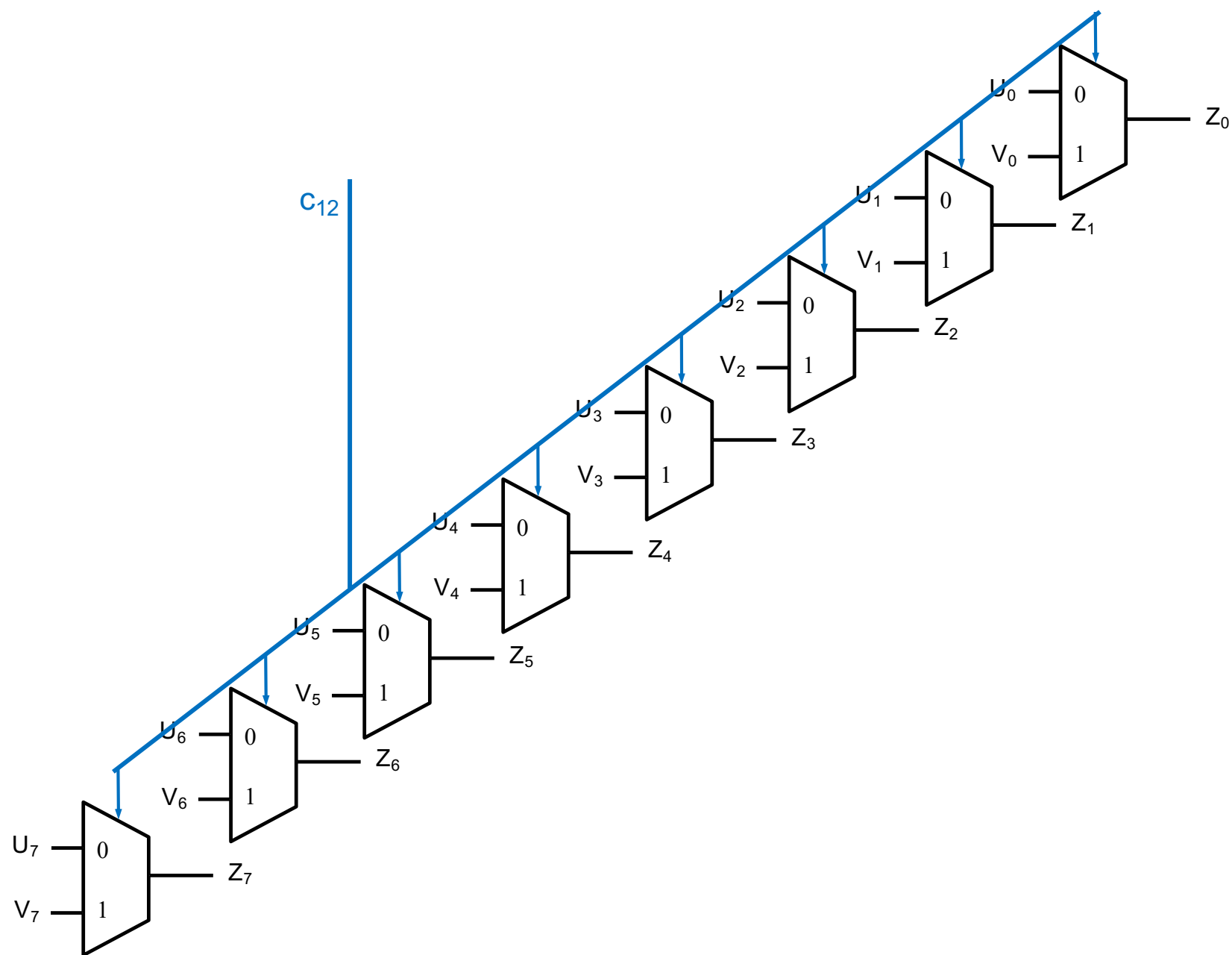
One restriction of digital multiplexers is that they cannot mix bits from A and B when outputting F. At each time step it is only A or only B.

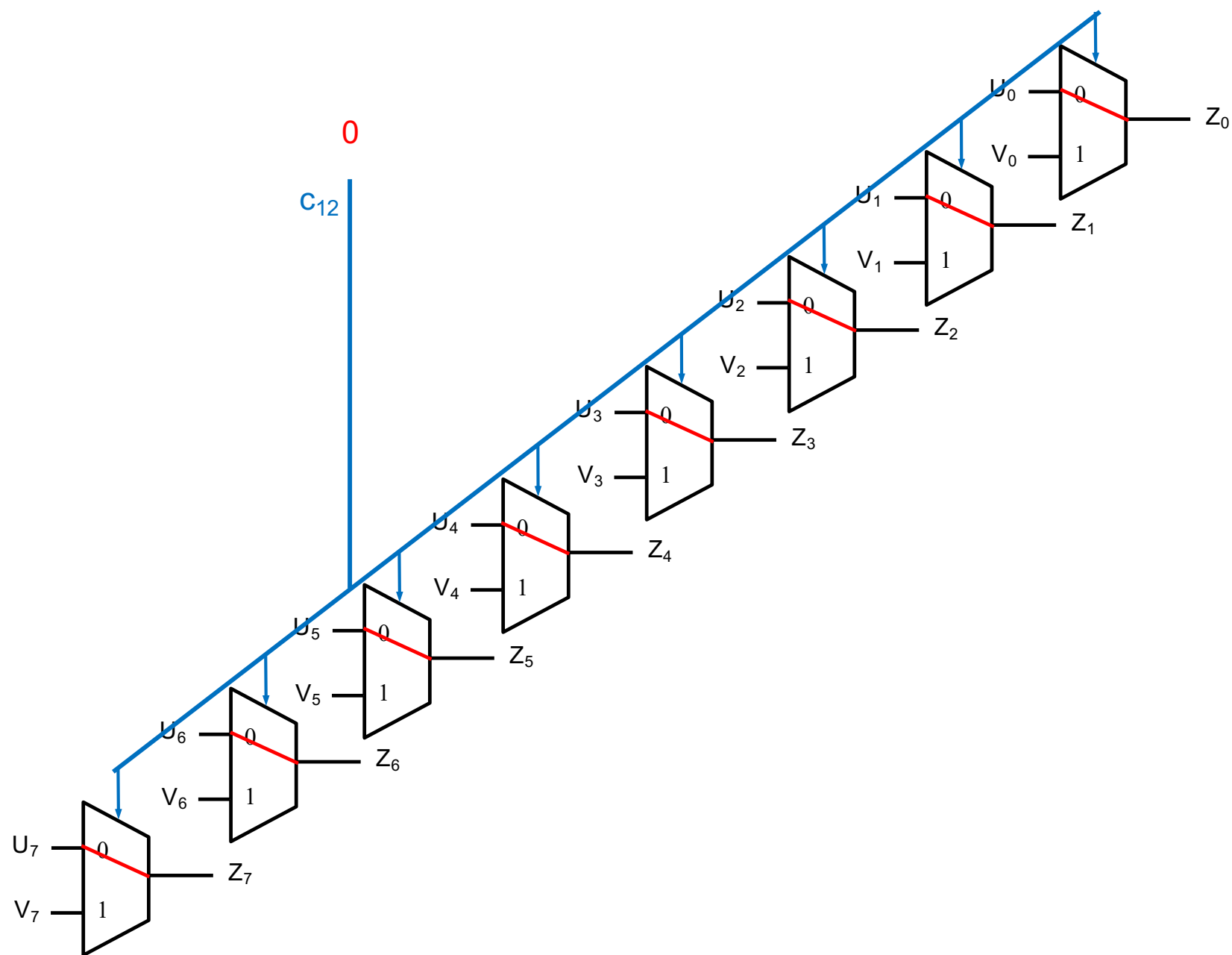
The Internal ALU Bus Multiplexer

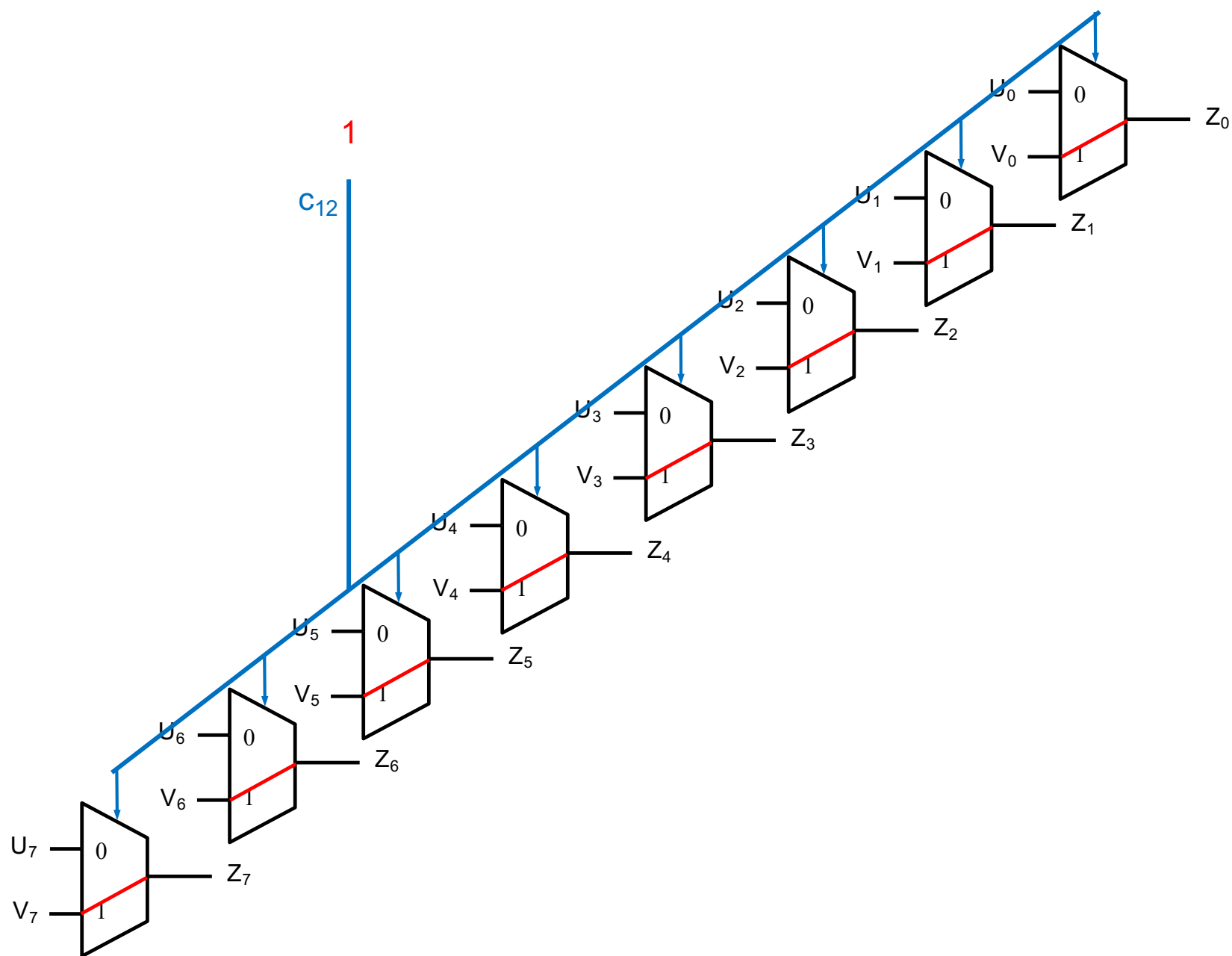


2-to-1 Bus Multiplexer (with 8-bit lines)

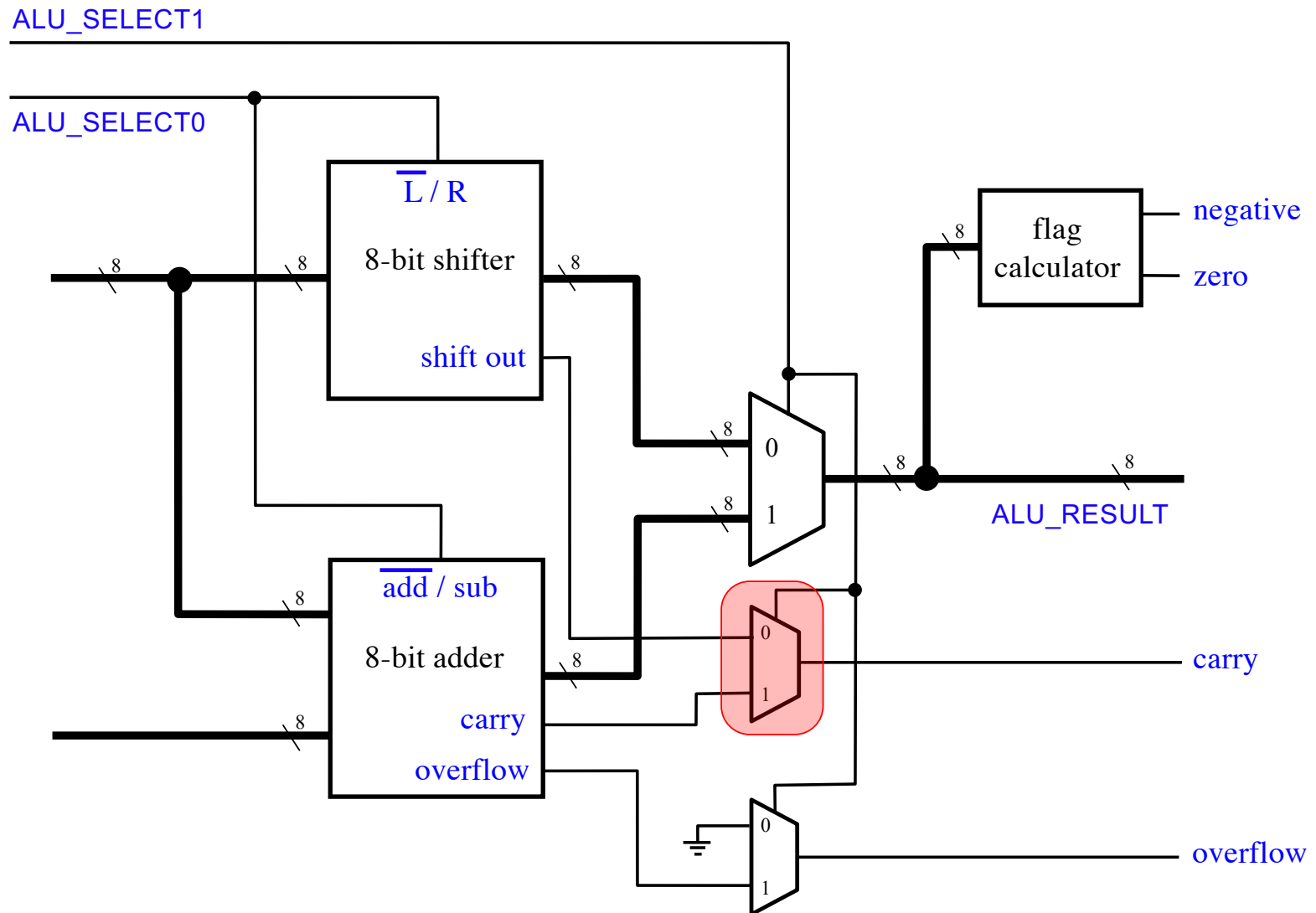




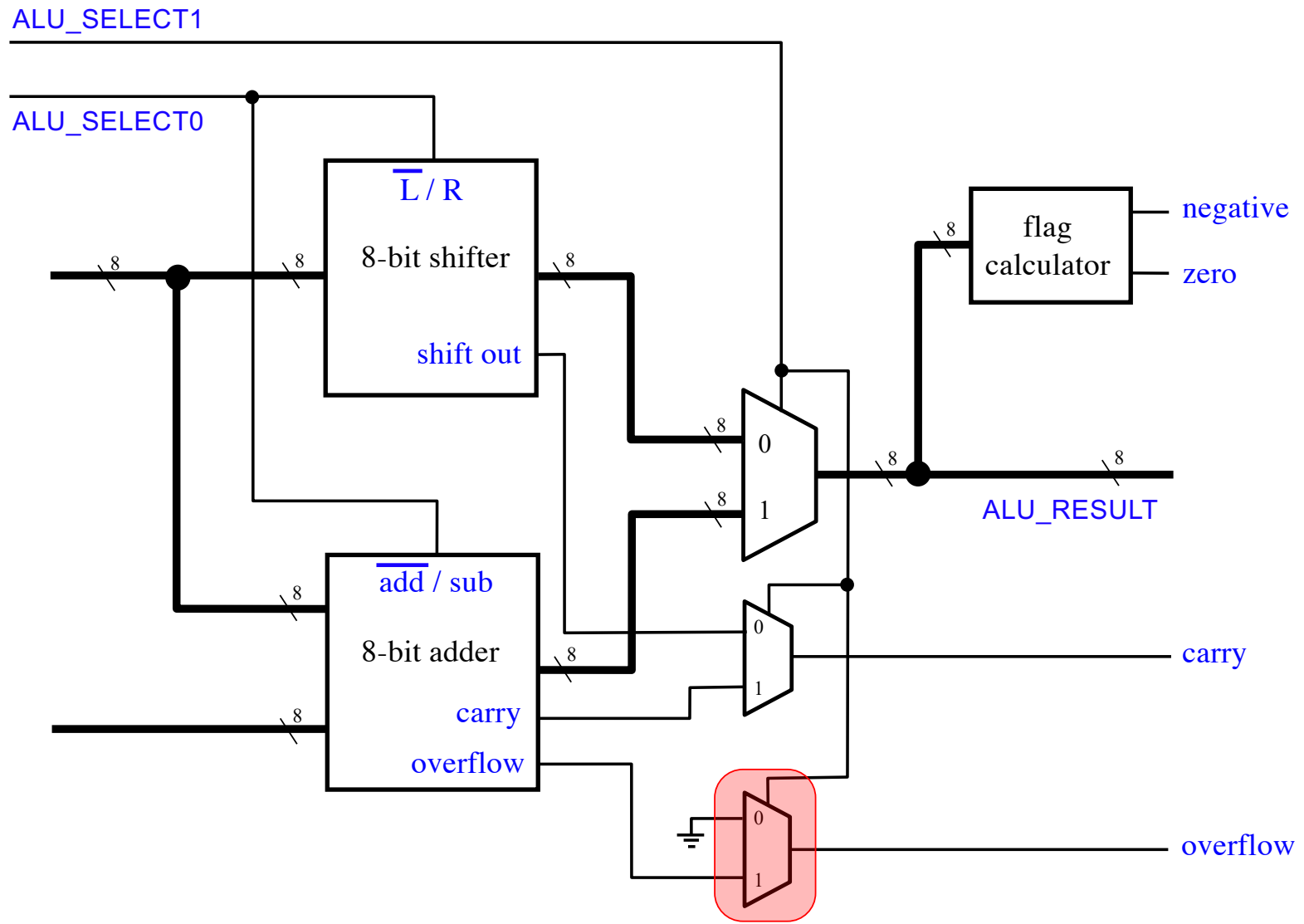




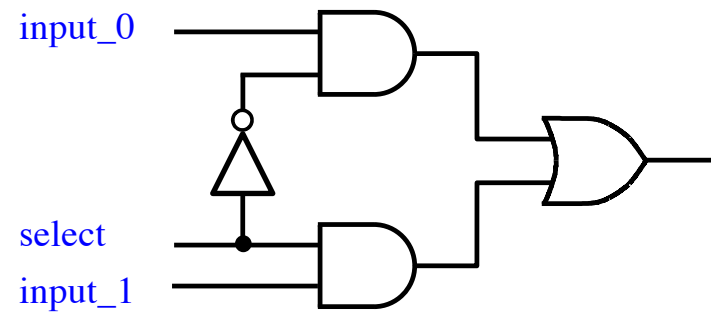
2-to-1 Multiplexer



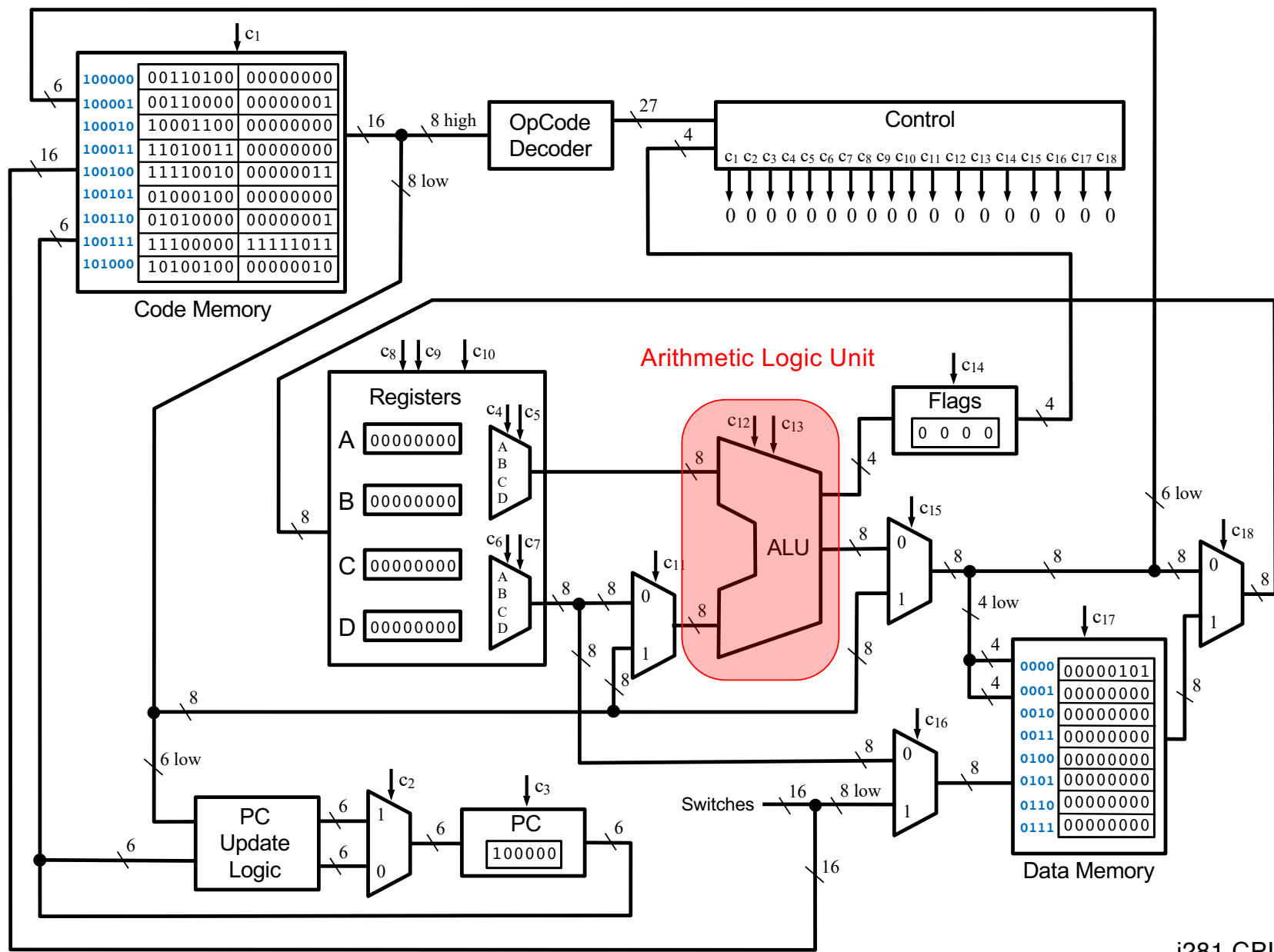
2-to-1 Multiplexer



2-to-1 Multiplexer

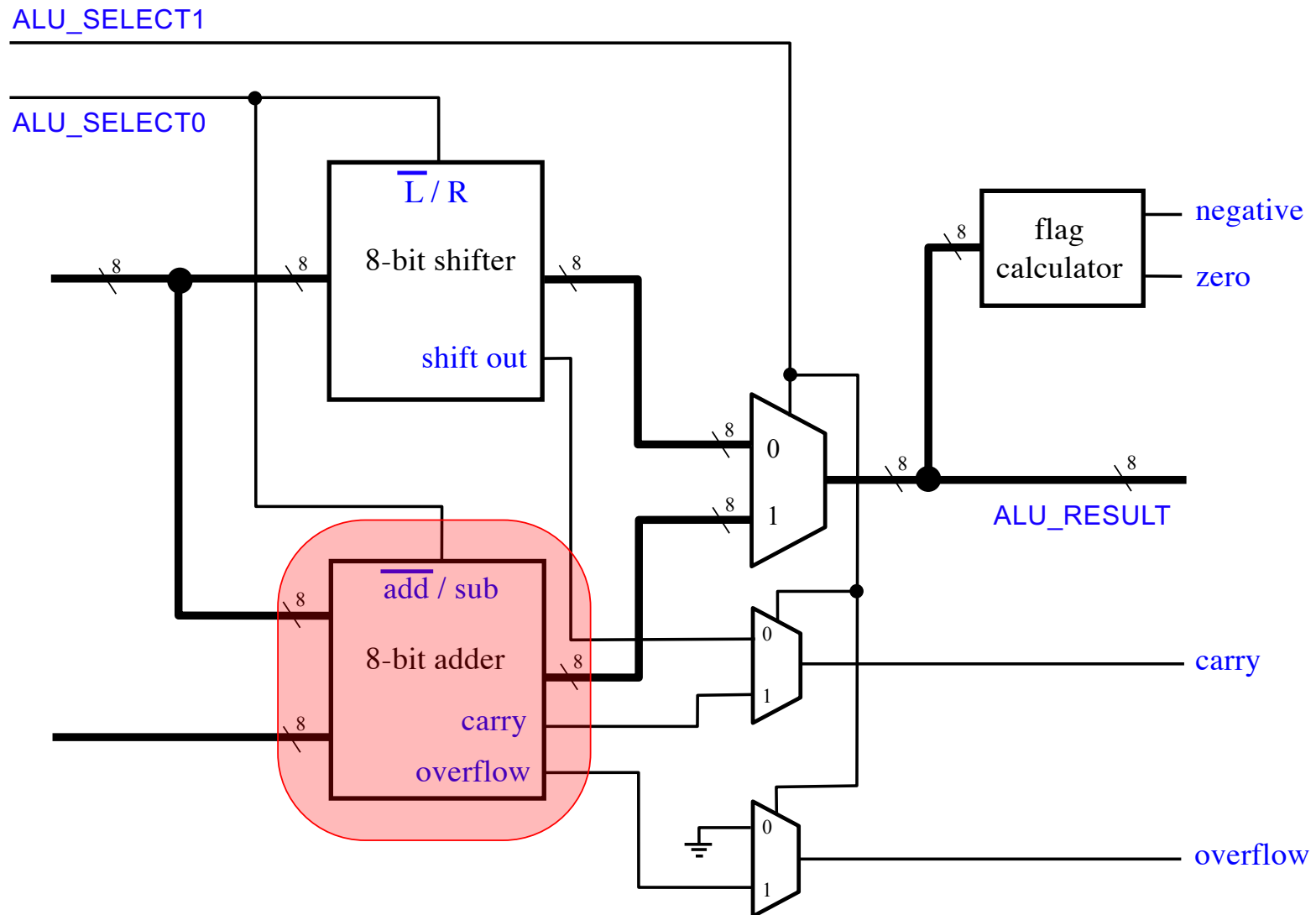


The adder/subtractor of the i281 CPU



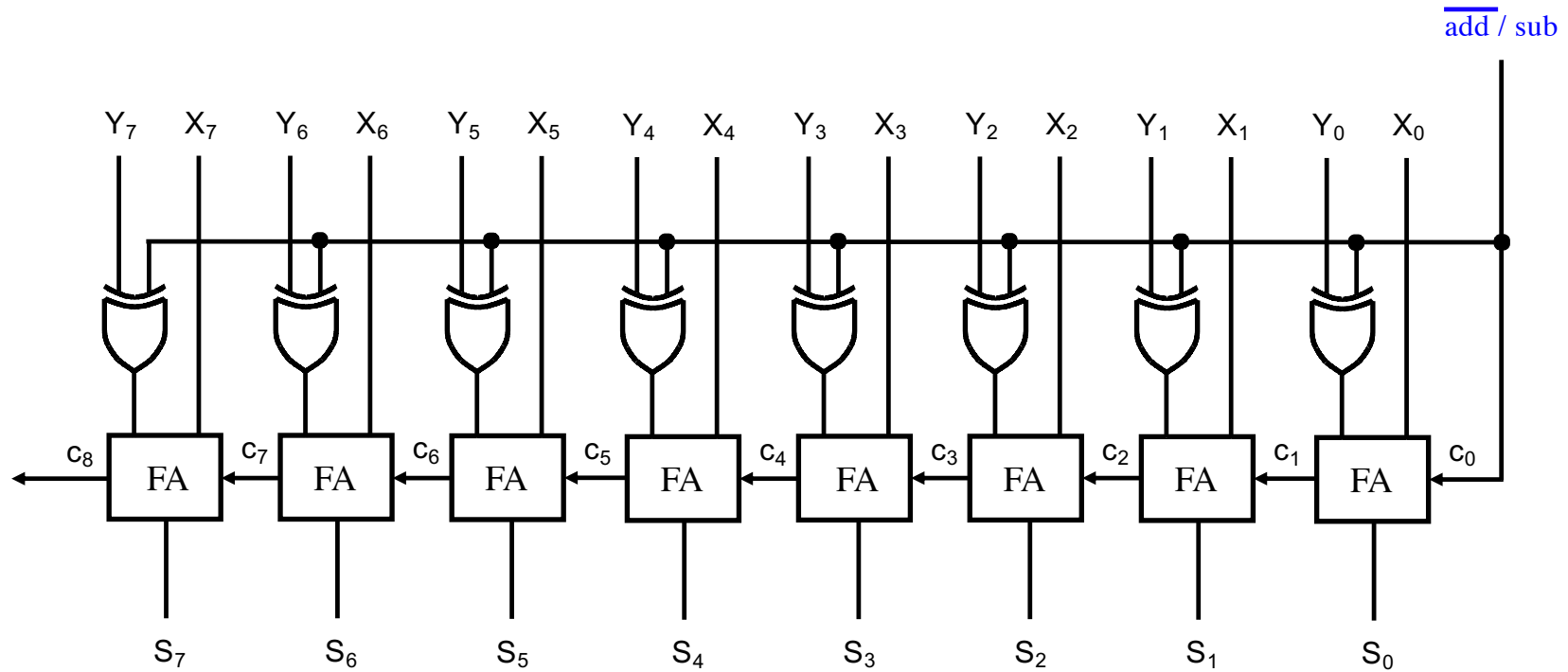
i281 CPU

The Adder / Subtractor



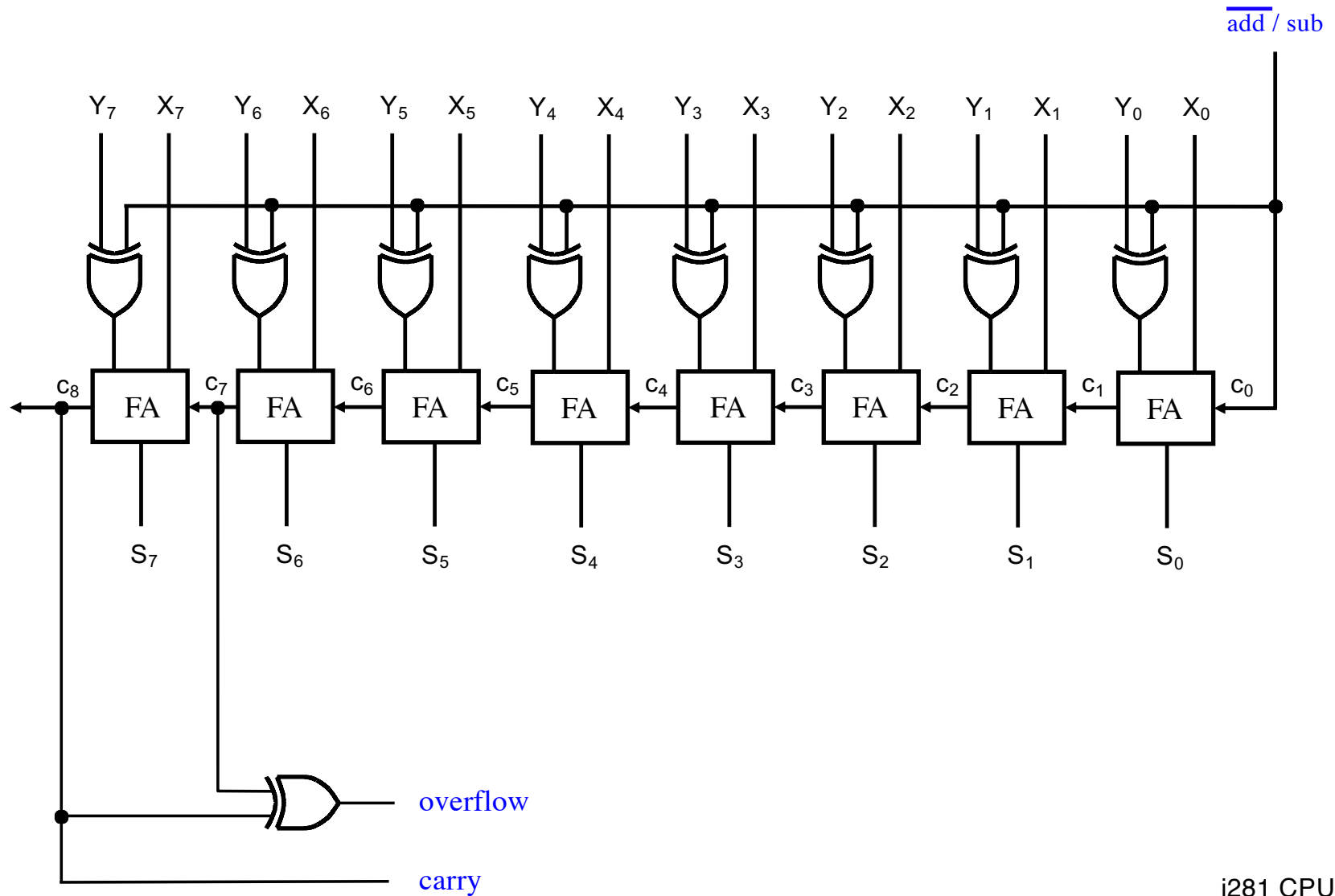
i281 CPU

The Adder / Subtractor

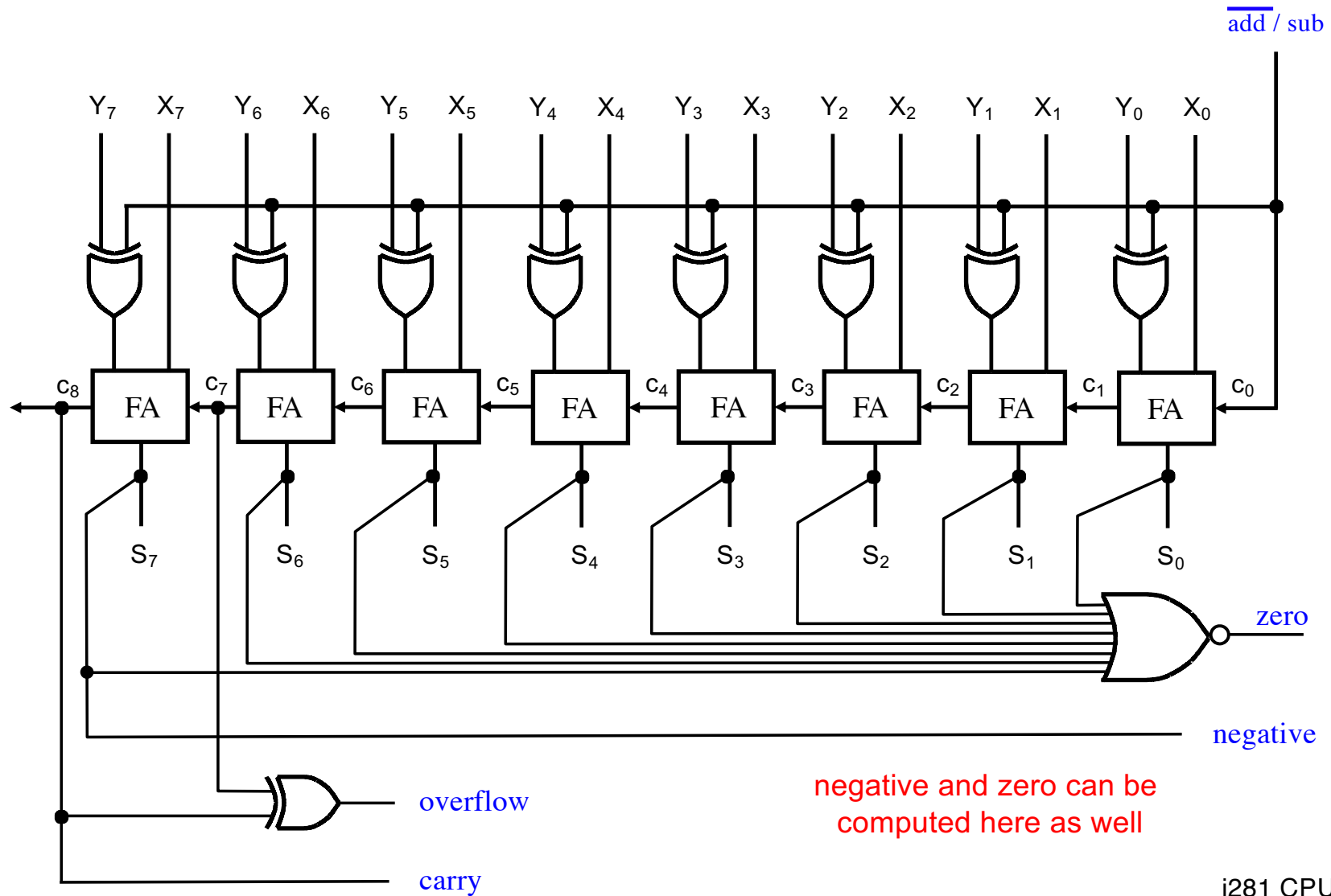


This is an 8-bit ripple-carry adder. Note that the X and Y lines are swapped.

The Adder / Subtractor



The Adder / Subtractor



Abbreviations for the Flags

- **Carry Flag (CF)**
- **Overflow Flag (OF)**
- **Negative Flag (NF)**
- **Zero Flag (ZF)**

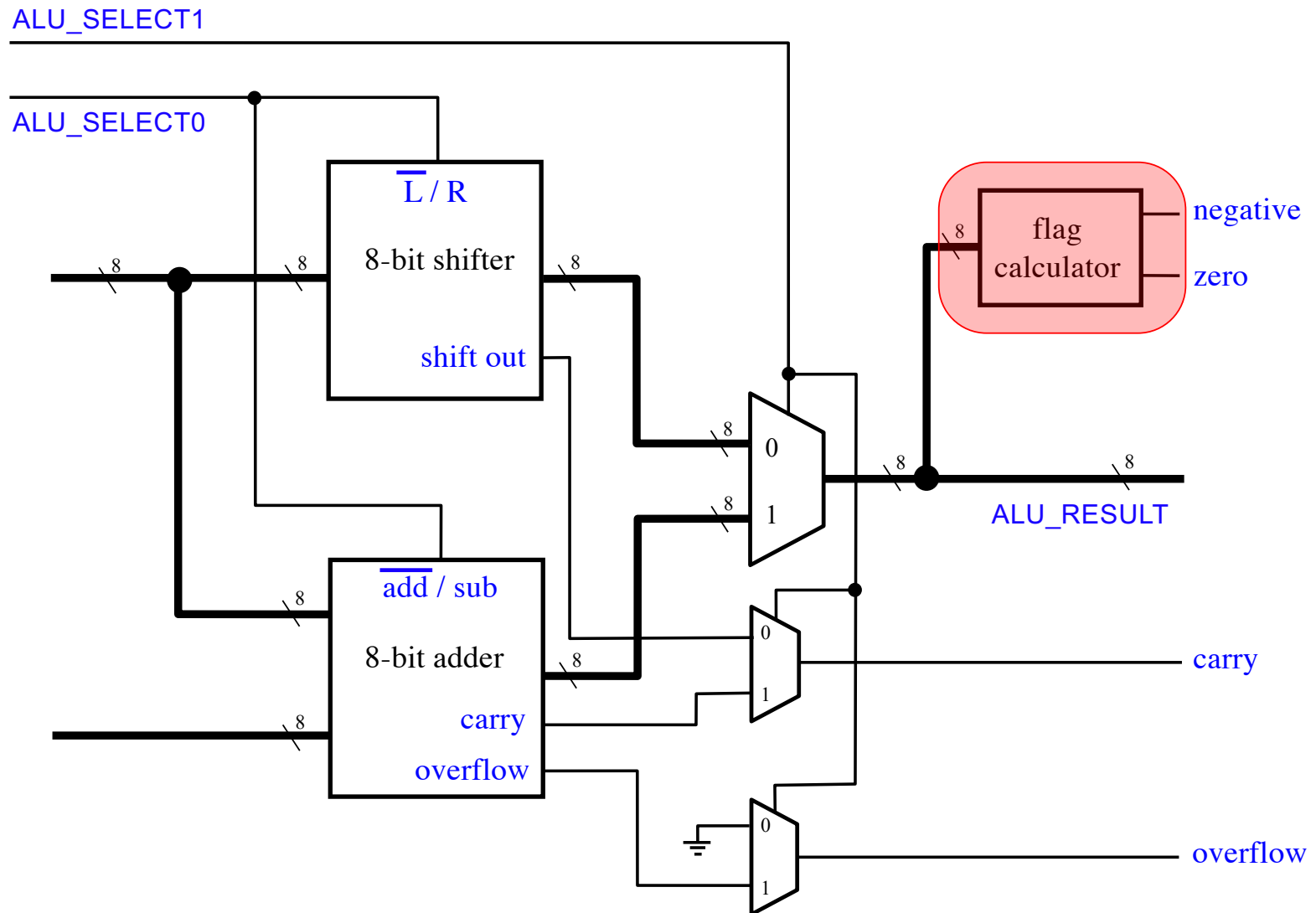
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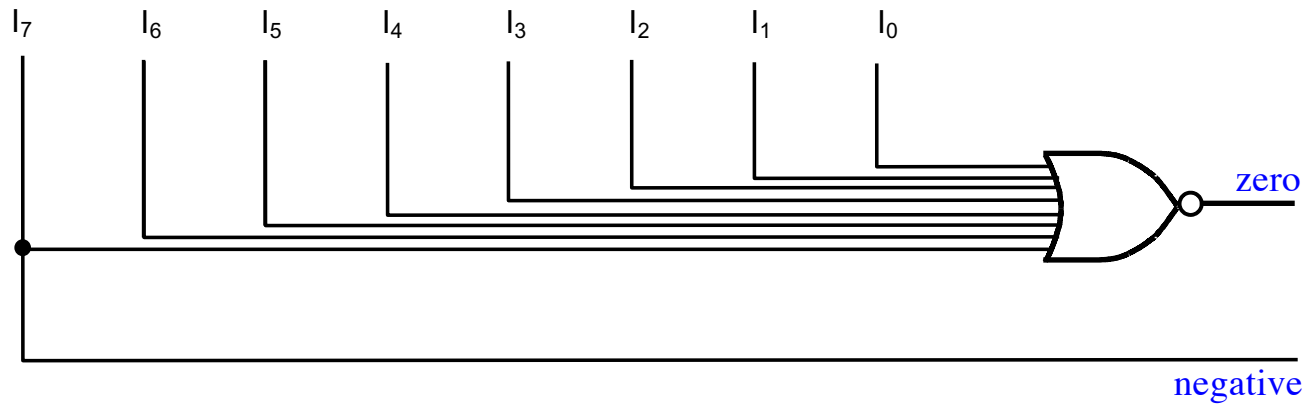
In some CPU architectures the carry flag means borrow. And it could be inverted relative to the previous diagram.

The flag calculator of the i281 CPU

The ALU Flag Calculator



The ALU Flag Calculator

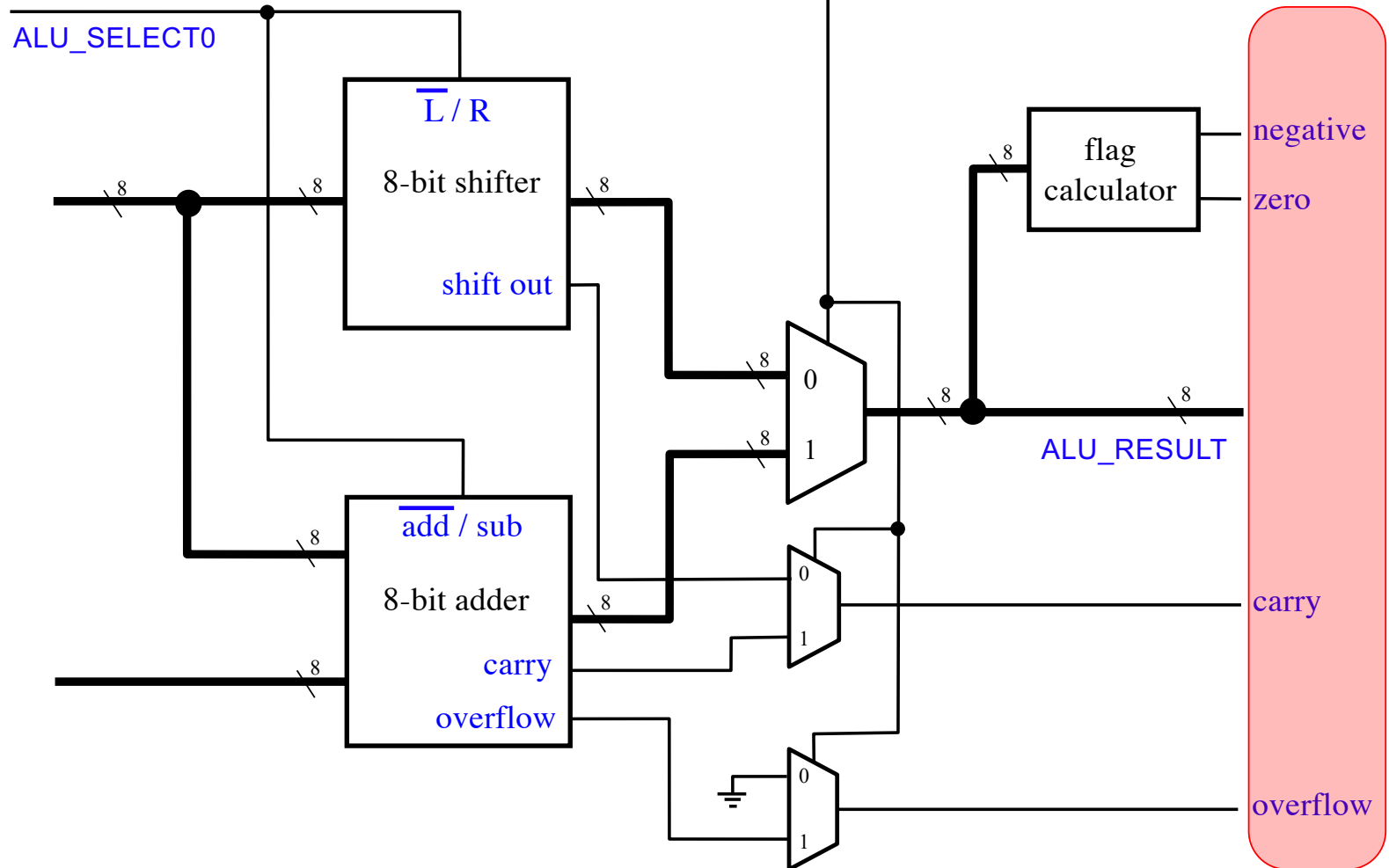


ALU Outputs to the Flags Register

ALU_SELECT1

ALU_SELECT0

4 Outputs to the
Flags Register



Questions?

THE END