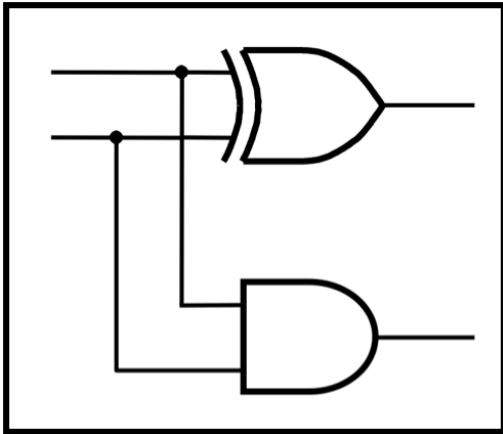




CprE 2810: Digital Logic

Instructor: Alexander Stoytchev

<http://www.ece.iastate.edu/~alexs/classes/>

Decoders and Encoders

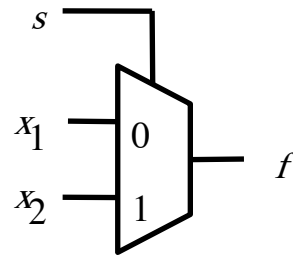
CprE 2810: Digital Logic
Iowa State University, Ames, IA
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Administrative Stuff

- **HW6 is due today**
- **HW7 is out. It is due on Monday Oct 20 @ 10pm.**

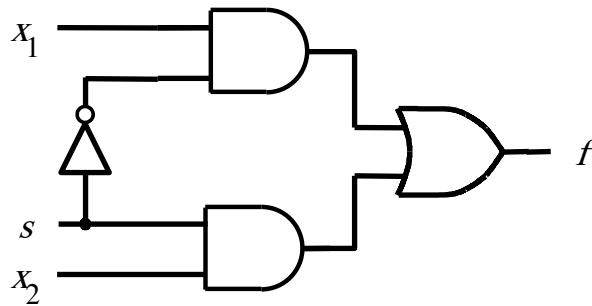
Quick Review

Graphical Symbol for a 2-1 Multiplexer

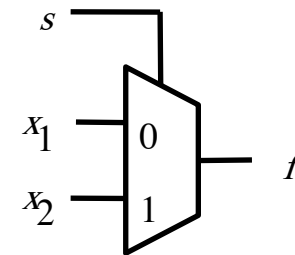


[Figure 2.33c from the textbook]

Circuit for 2-1 Multiplexer



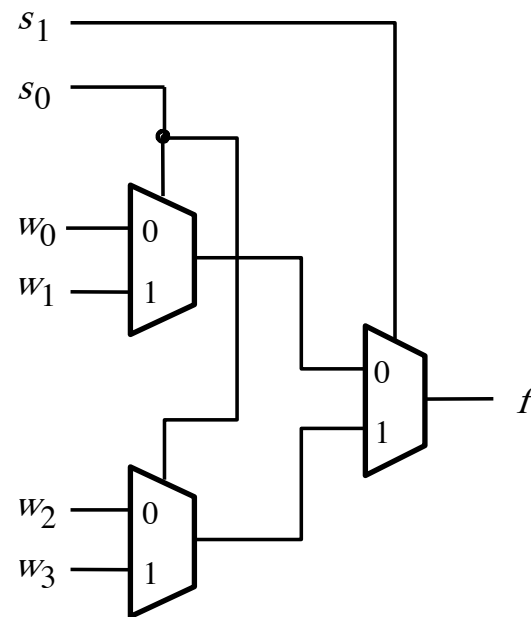
(b) Circuit



(c) Graphical symbol

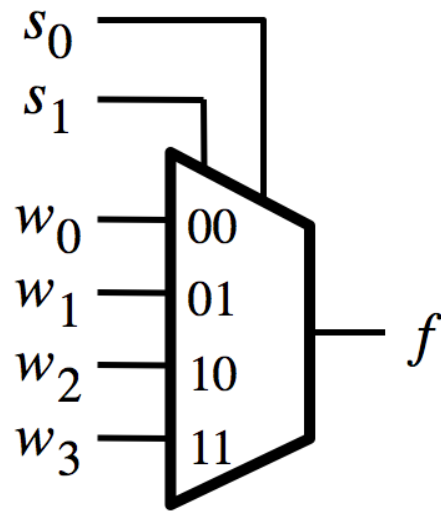
$$f(s, x_1, x_2) = \overline{s} x_1 + s x_2$$

Using three 2-to-1 multiplexers to build one 4-to-1 multiplexer



[Figure 4.3 from the textbook]

Graphical Symbol and Truth Table

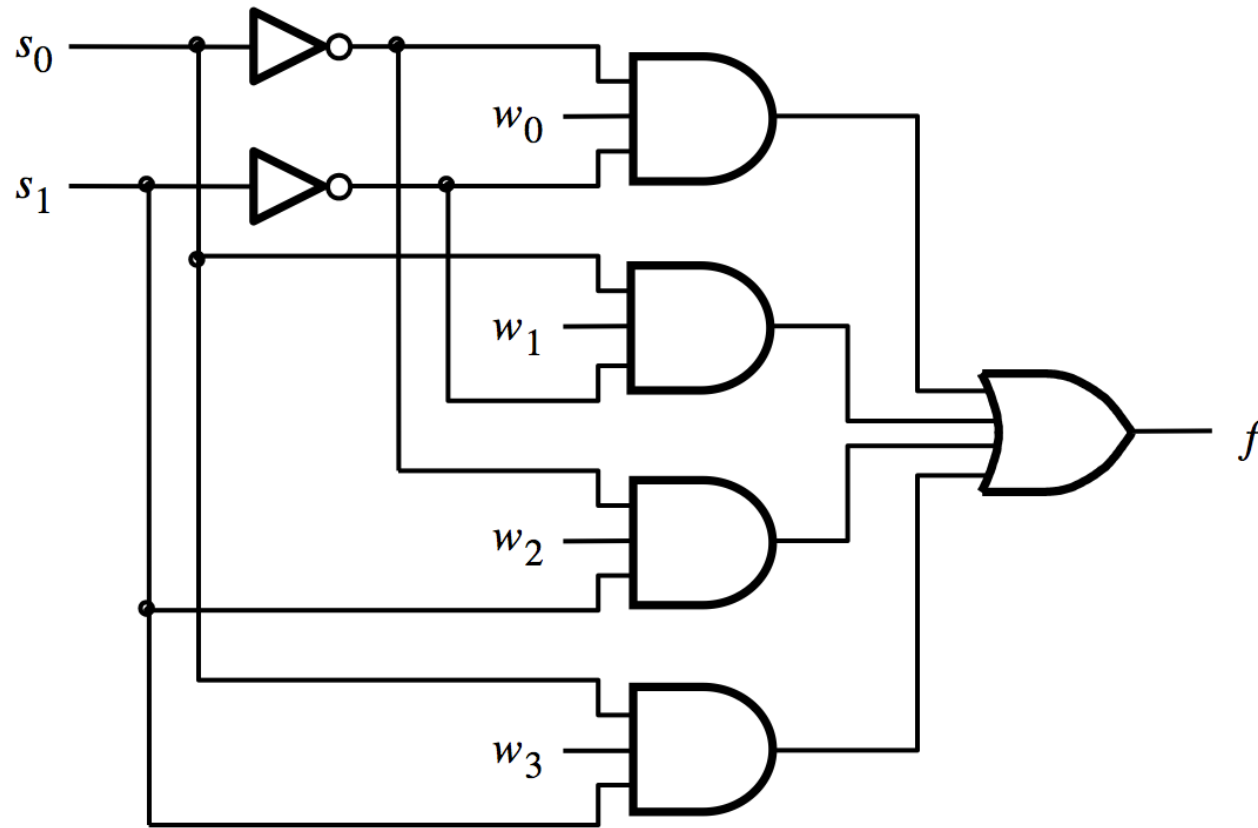


(a) Graphic symbol

s_1	s_0	f
0	0	w_0
0	1	w_1
1	0	w_2
1	1	w_3

(b) Truth table

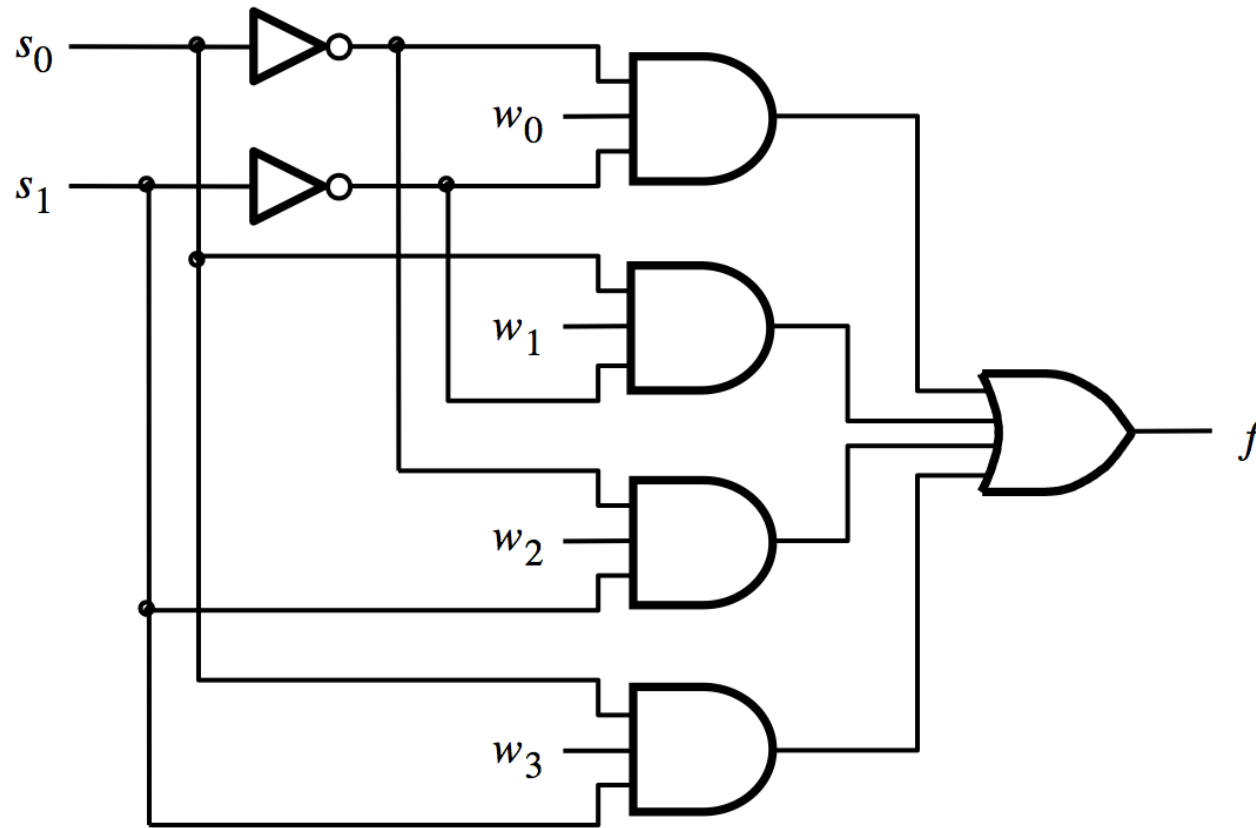
4-to-1 Multiplexer (SOP circuit)



$$f = \overline{s_1} \overline{s_0} w_0 + \overline{s_1} s_0 w_1 + s_1 \overline{s_0} w_2 + s_1 s_0 w_3$$

[Figure 4.2c from the textbook]

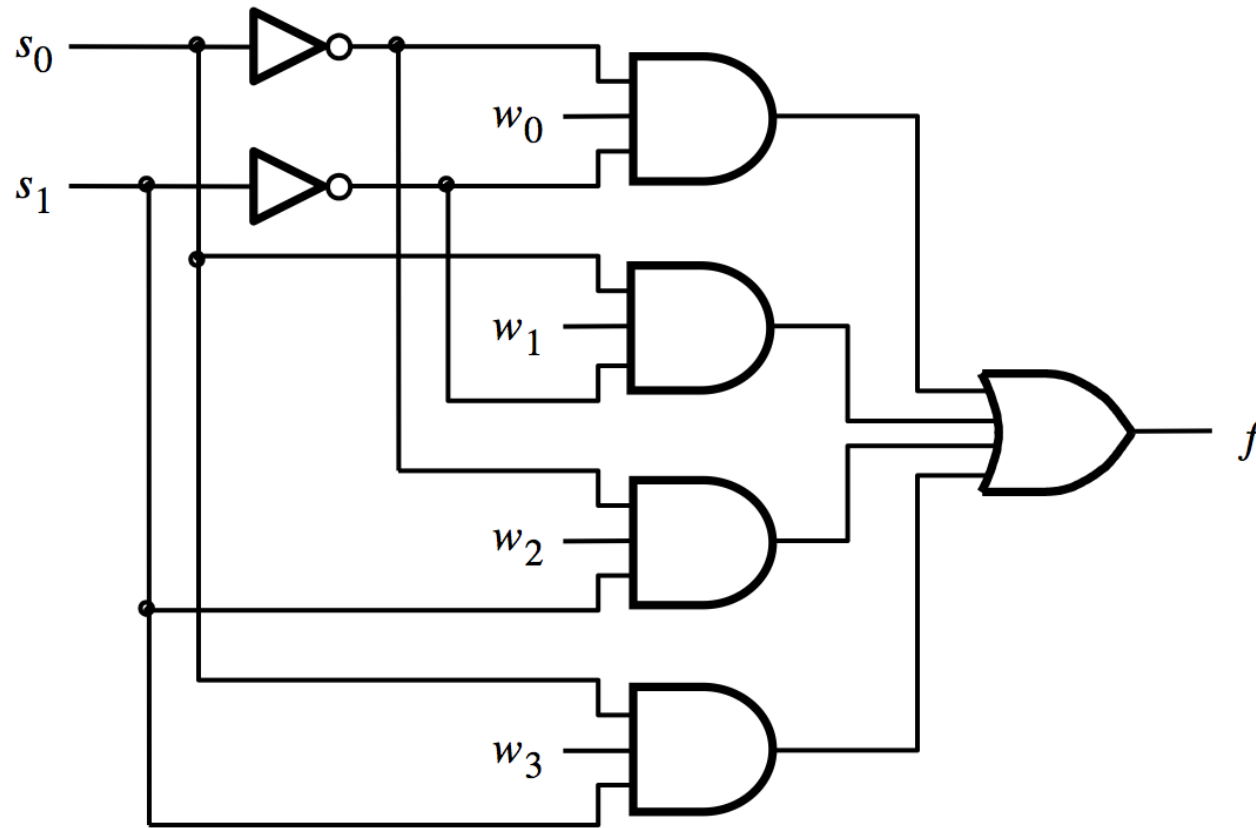
4-to-1 Multiplexer (SOP circuit)



$$f = \overline{s_1} \overline{s_0} w_0 + \overline{s_1} s_0 w_1 + s_1 \overline{s_0} w_2 + s_1 s_0 w_3$$

these are the four minterms in terms of s_1 and s_2

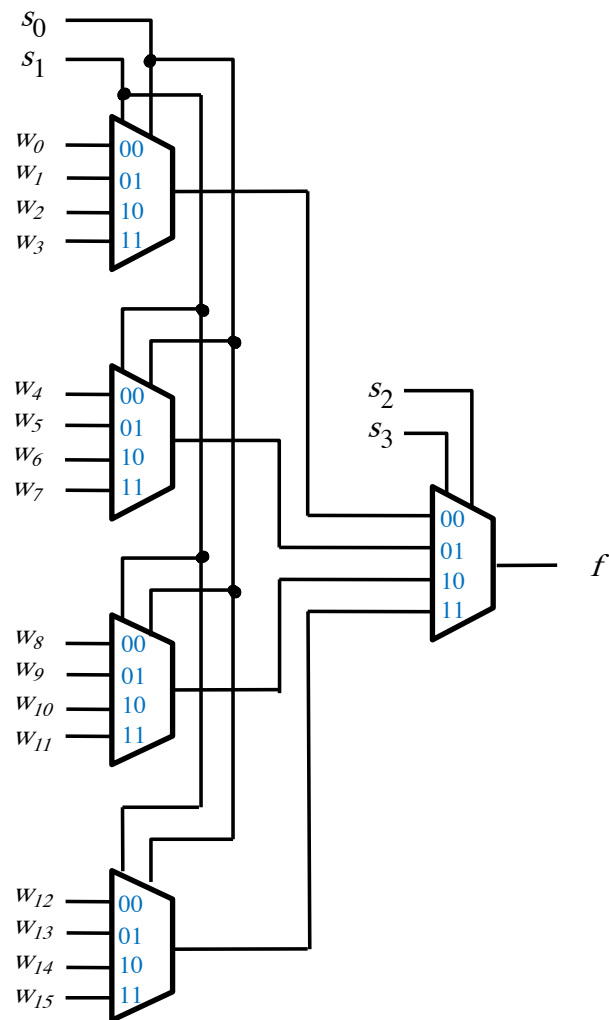
4-to-1 Multiplexer (SOP circuit)



$$f = \overline{s_1} \overline{s_0} w_0 + \overline{s_1} s_0 w_1 + s_1 \overline{s_0} w_2 + s_1 s_0 w_3$$

these are the four w inputs to the 4-to-1 MUX

16-1 Multiplexer with 4-to-1 Multiplexers



Shannon's Expansion Theorem

(general case with one select variable)

Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

$$f(w_1, w_2, \dots, w_n) = \overline{w}_1 \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$

This form is suitable for implementation with a 2-to-1 multiplexer.

Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

$$f(w_1, w_2, \dots, w_n) = \boxed{\overline{w}_1} \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$



w_1 set to 0

Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

$$f(w_1, w_2, \dots, w_n) = \overline{w}_1 \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$



w_1 set to 1

Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

$$f(w_1, w_2, \dots, w_n) = \overline{w}_1 \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$

Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

$$f(w_1, w_2, \dots, w_n) = \bar{w}_1 \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$

$$f = \bar{w}_1 f_{\bar{w}_1} + w_1 f_{w_1}$$

Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

$$f(w_1, w_2, \dots, w_n) = \boxed{\bar{w}_1} \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$

$$f = \boxed{\bar{w}_1} f_{\bar{w}_1} + w_1 f_{w_1}$$

select variable
(negated)

Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

$$f(w_1, w_2, \dots, w_n) = \bar{w}_1 \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$

$$f = \bar{w}_1 f_{\bar{w}_1} + w_1 f_{w_1}$$

select variable
(not negated)

Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

$$f(w_1, w_2, \dots, w_n) = \bar{w}_1 \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$

$$f = \bar{w}_1 f_{\bar{w}_1} + w_1 f_{w_1}$$

cofactor



Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

$$f(w_1, w_2, \dots, w_n) = \bar{w}_1 \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$

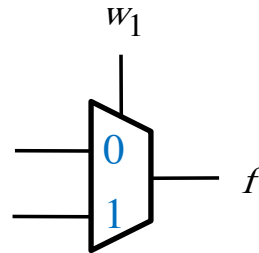
$$f = \bar{w}_1 f_{\bar{w}_1} + w_1 f_{w_1}$$

↑
cofactor

Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

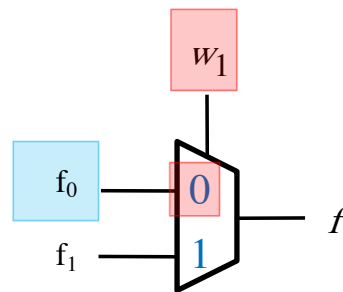
$$f(w_1, w_2, \dots, w_n) = \overline{w}_1 \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$



Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

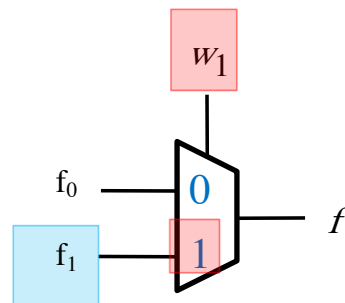
$$f(w_1, w_2, \dots, w_n) = \overline{w_1} \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$



Shannon's Expansion Theorem

Any Boolean function $f(w_1, \dots, w_n)$ can be rewritten in the form:

$$f(w_1, w_2, \dots, w_n) = \overline{w}_1 \cdot f(0, w_2, \dots, w_n) + w_1 \cdot f(1, w_2, \dots, w_n)$$



Shannon's Expansion Theorem
(example with only one select variable)
(used for 2-to-1 mux implementation)

Shannon's Expansion Theorem Example

$$f(w_1, w_2, w_3) = w_1 w_2 + w_1 w_3 + w_2 w_3$$

Shannon's Expansion Theorem Example

set $w_1 = 0$



$$f(w_1, w_2, w_3) = w_1 w_2 + w_1 w_3 + w_2 w_3$$

Shannon's Expansion Theorem Example

set $w_1 = 0$



$$f(w_1, w_2, w_3) = w_1 w_2 + w_1 w_3 + w_2 w_3$$

$$\underbrace{f(\underbrace{0}_{w_1}, w_2, w_3) = 0 \cdot w_2 + 0 \cdot w_3 + w_2 w_3}_{w_2 w_3}$$

$w_2 w_3$

Shannon's Expansion Theorem Example

set $w_1 = 1$



$$f(w_1, w_2, w_3) = w_1 w_2 + w_1 w_3 + w_2 w_3$$

Shannon's Expansion Theorem Example

set $w_1 = 1$



$$f(w_1, w_2, w_3) = w_1 w_2 + w_1 w_3 + w_2 w_3$$

$$f(\underset{w_1}{1}, w_2, w_3) = \underbrace{1 \cdot w_2 + 1 \cdot w_3 + w_2 w_3}$$

$$w_2 + w_3$$

Shannon's Expansion Theorem Example

started with this

$$f(w_1, w_2, w_3) = w_1 w_2 + w_1 w_3 + w_2 w_3$$

Shannon's Expansion Theorem Example

$$f(w_1, w_2, w_3) = w_1 w_2 + w_1 w_3 + w_2 w_3$$

$$\begin{aligned} f &= \overline{w}_1(0 \cdot w_2 + 0 \cdot w_3 + w_2 w_3) + w_1(1 \cdot w_2 + 1 \cdot w_3 + w_2 w_3) \\ &= \overline{w}_1(w_2 w_3) + w_1(w_2 + w_3) \end{aligned}$$

factored it in this form

Shannon's Expansion Theorem Example

$$f(w_1, w_2, w_3) = w_1 w_2 + w_1 w_3 + w_2 w_3$$

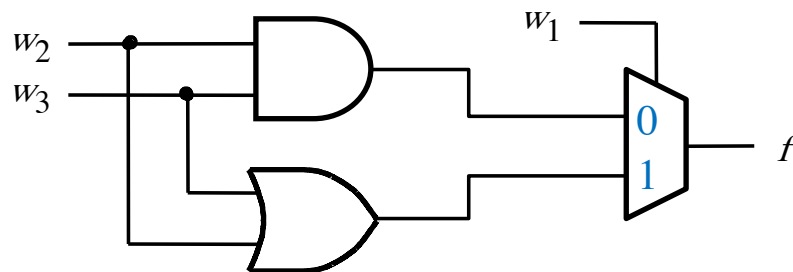
$$\begin{aligned} f &= \bar{w}_1(0 \cdot w_2 + 0 \cdot w_3 + w_2 w_3) + w_1(1 \cdot w_2 + 1 \cdot w_3 + w_2 w_3) \\ &= \bar{w}_1(w_2 w_3) + w_1(w_2 + w_3) \end{aligned}$$

$$f = \bar{w}_1 f_{\bar{w}_1} + w_1 f_{w_1}$$

which fits the general formula

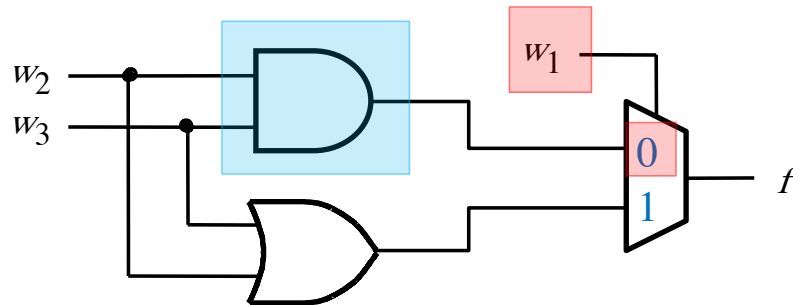
Shannon's Expansion Theorem Example

$$\begin{aligned} f &= \bar{w}_1(0 \cdot w_2 + 0 \cdot w_3 + w_2w_3) + w_1(1 \cdot w_2 + 1 \cdot w_3 + w_2w_3) \\ &= \bar{w}_1(w_2w_3) + w_1(w_2 + w_3) \end{aligned}$$



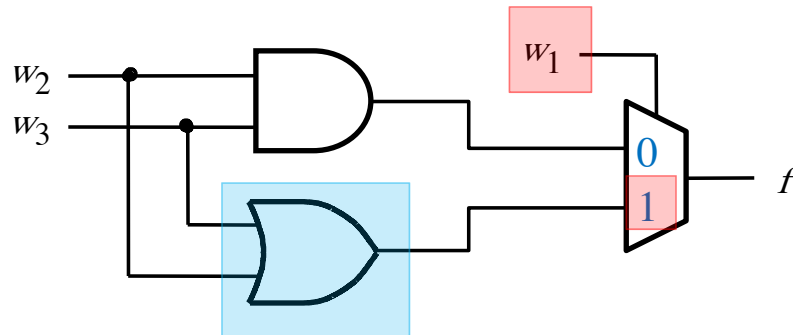
Shannon's Expansion Theorem Example

$$\begin{aligned} f &= \bar{w}_1(0 \cdot w_2 + 0 \cdot w_3 + w_2w_3) + w_1(1 \cdot w_2 + 1 \cdot w_3 + w_2w_3) \\ &= \bar{w}_1(w_2w_3) + w_1(w_2 + w_3) \end{aligned}$$



Shannon's Expansion Theorem Example

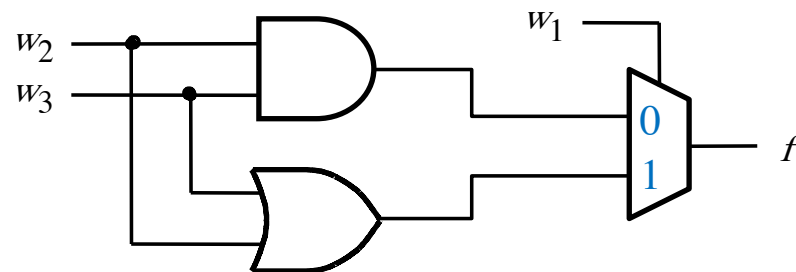
$$\begin{aligned} f &= \bar{w}_1(0 \cdot w_2 + 0 \cdot w_3 + w_2w_3) + w_1(1 \cdot w_2 + 1 \cdot w_3 + w_2w_3) \\ &= \bar{w}_1(w_2w_3) + \boxed{w_1}(w_2 + w_3) \end{aligned}$$



Earlier we derived the same result from the truth table using the divide-and-conquer method (a.k.a., Shannon's theorem)

w_1	w_2	w_3	f		w_1	f
0	0	0	0	}	0	$w_2 w_3$
0	0	1	0			
0	1	0	0			
0	1	1	1			
1	0	0	0	}	1	$w_2 + w_3$
1	0	1	1			
1	1	0	1			
1	1	1	1			

(b) Truth table



(b) Circuit

[Figure 4.10a from the textbook]

Shannon's Expansion Theorem

(general case with two select variables)

Shannon's Expansion Theorem (In terms of more than one variable)

$$\begin{aligned} f(w_1, \dots, w_n) = & \bar{w}_1 \bar{w}_2 \cdot f(0, 0, w_3, \dots, w_n) + \bar{w}_1 w_2 \cdot f(0, 1, w_3, \dots, w_n) \\ & + w_1 \bar{w}_2 \cdot f(1, 0, w_3, \dots, w_n) + w_1 w_2 \cdot f(1, 1, w_3, \dots, w_n) \end{aligned}$$

This form is suitable for implementation with a 4-to-1 multiplexer.

Shannon's Expansion Theorem

(In terms of more than one variable)

$$\begin{aligned} f(w_1, \dots, w_n) = & \overline{w_1} \overline{w_2} \cdot f(0, 0, w_3, \dots, w_n) + \overline{w_1} w_2 \cdot f(0, 1, w_3, \dots, w_n) \\ & + w_1 \overline{w_2} \cdot f(1, 0, w_3, \dots, w_n) + w_1 w_2 \cdot f(1, 1, w_3, \dots, w_n) \end{aligned}$$

These are the four possible minterms with two variables.

Shannon's Expansion Theorem

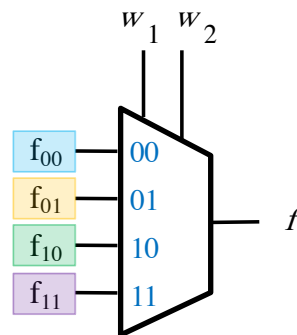
(In terms of more than one variable)

$$\begin{aligned} f(w_1, \dots, w_n) = & \bar{w}_1 \bar{w}_2 \cdot f(0, 0, w_3, \dots, w_n) + \bar{w}_1 w_2 \cdot f(0, 1, w_3, \dots, w_n) \\ & + w_1 \bar{w}_2 \cdot f(1, 0, w_3, \dots, w_n) + w_1 w_2 \cdot f(1, 1, w_3, \dots, w_n) \end{aligned}$$

These are the four cofactors, one for each of the minterms.

Shannon's Expansion Theorem (In terms of more than one variable)

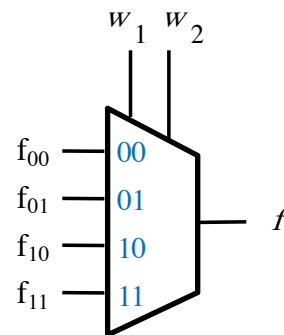
$$f(w_1, \dots, w_n) = \bar{w}_1 \bar{w}_2 \cdot f(0, 0, w_3, \dots, w_n) + \bar{w}_1 w_2 \cdot f(0, 1, w_3, \dots, w_n) \\ + w_1 \bar{w}_2 \cdot f(1, 0, w_3, \dots, w_n) + w_1 w_2 \cdot f(1, 1, w_3, \dots, w_n)$$



Shannon's Expansion Theorem

(In terms of more than one variable)

$$\begin{aligned} f(w_1, \dots, w_n) = & \bar{w}_1 \bar{w}_2 \cdot f(0, 0, w_3, \dots, w_n) + \bar{w}_1 w_2 \cdot f(0, 1, w_3, \dots, w_n) \\ & + w_1 \bar{w}_2 \cdot f(1, 0, w_3, \dots, w_n) + w_1 w_2 \cdot f(1, 1, w_3, \dots, w_n) \end{aligned}$$



Decoders

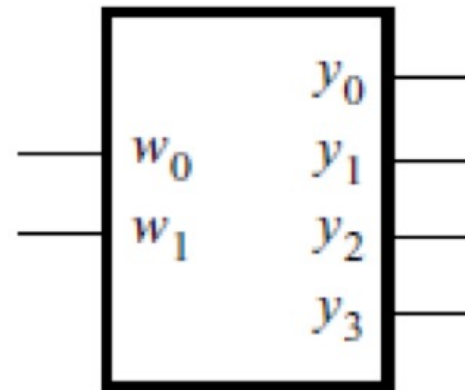
2-to-4 Decoder (Definition)

- Has two inputs: w_1 and w_0
- Has four outputs: y_0 , y_1 , y_2 , and y_3
- If $w_1=0$ and $w_0=0$, then the output y_0 is set to 1
- If $w_1=0$ and $w_0=1$, then the output y_1 is set to 1
- If $w_1=1$ and $w_0=0$, then the output y_2 is set to 1
- If $w_1=1$ and $w_0=1$, then the output y_3 is set to 1
- Only one output is set to 1. All others are set to 0.

Truth Table and Graphical Symbol for a 2-to-4 Decoder

w_1	w_0	y_0	y_1	y_2	y_3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

(a) Truth table



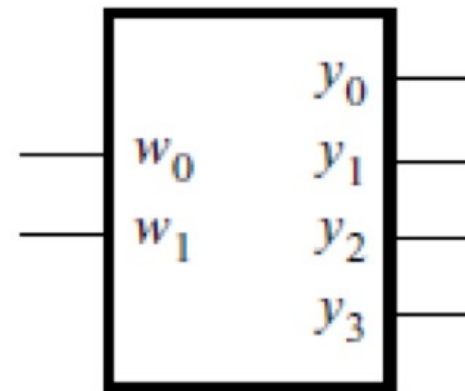
(b) Graphical symbol

Truth Table and Graphical Symbol for a 2-to-4 Decoder

w_1	w_0	y_0	y_1	y_2	y_3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

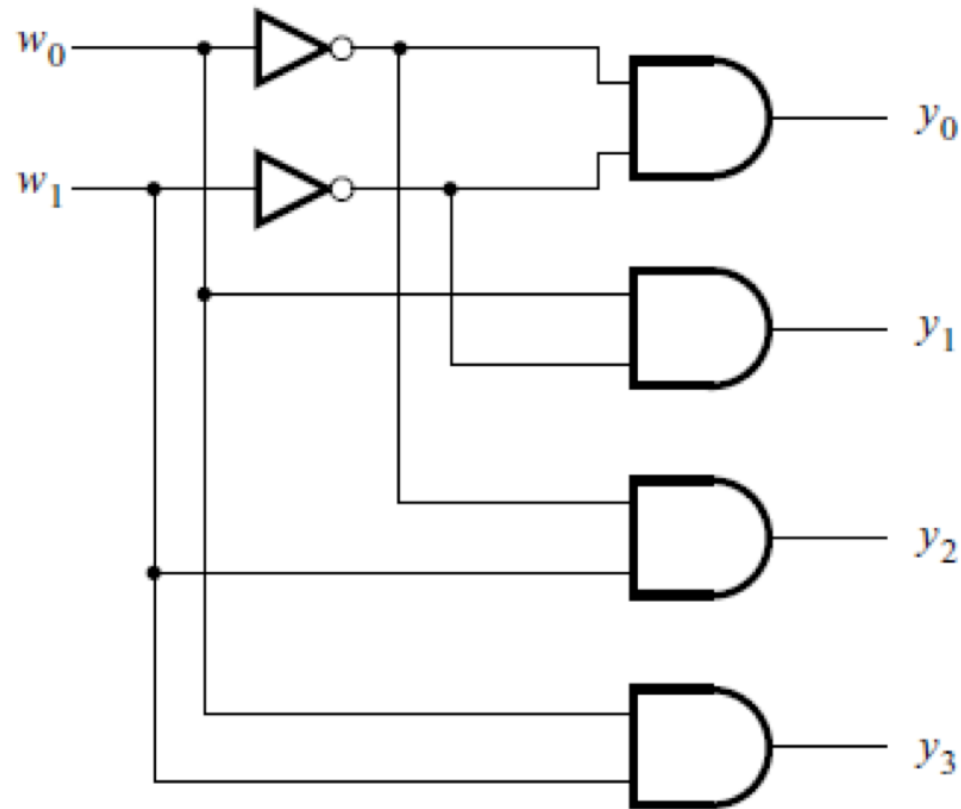
The outputs are “one-hot” encoded

(a) Truth table



(b) Graphical symbol

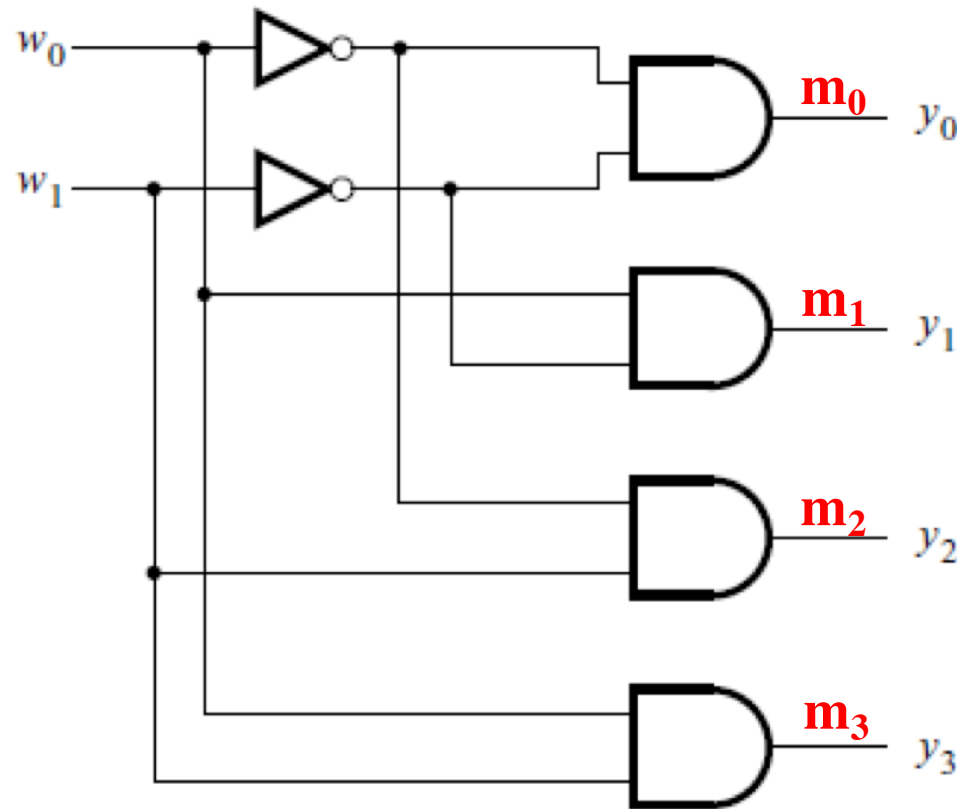
The Logic Circuit for a 2-to-4 Decoder



[Figure 4.13c from the textbook]

The Logic Circuit for a 2-to-4 Decoder

These AND gates
implement the four
minterm functions

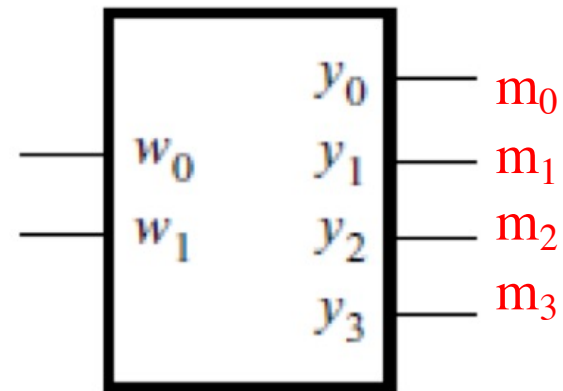


[Figure 4.13c from the textbook]

Truth Table and Graphical Symbol for a 2-to-4 Decoder

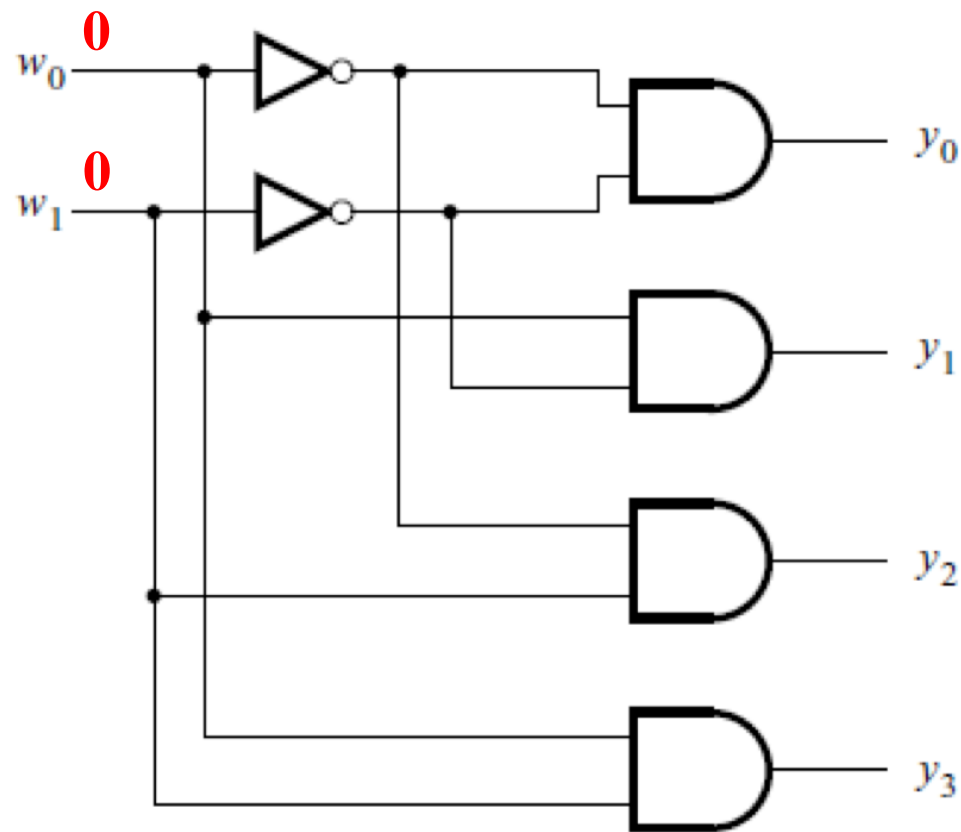
		m_0	m_1	m_2	m_3
w_1	w_0	y_0	y_1	y_2	y_3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

(a) Truth table



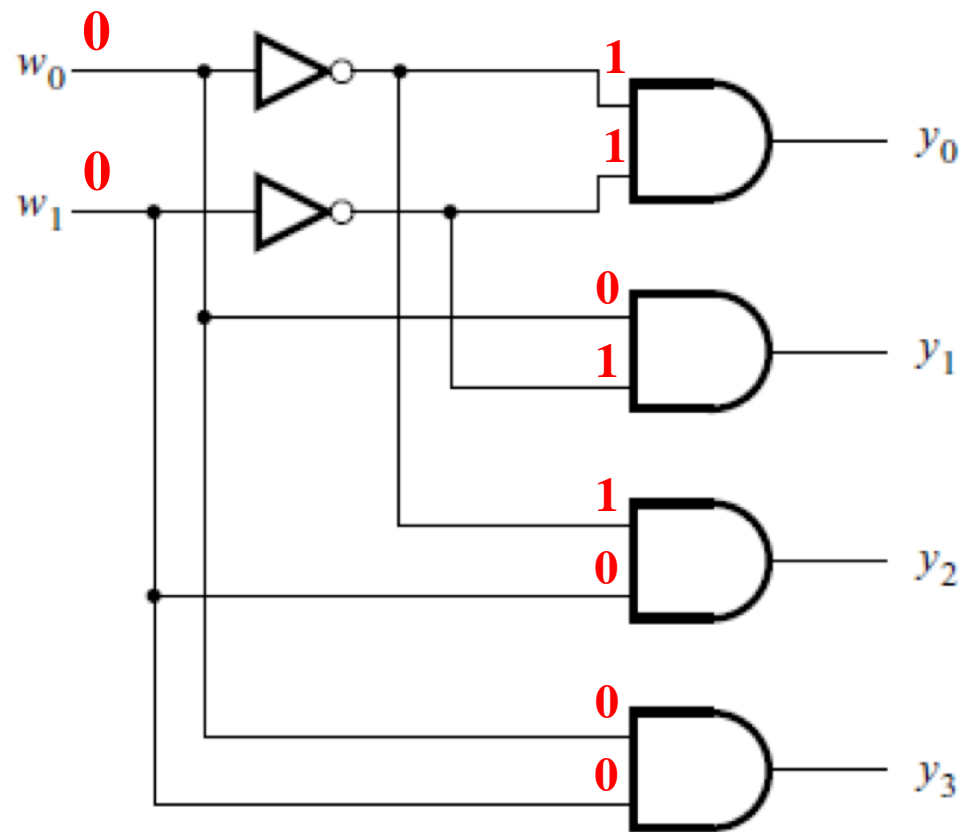
(b) Graphical symbol

The Logic Circuit for a 2-to-4 Decoder



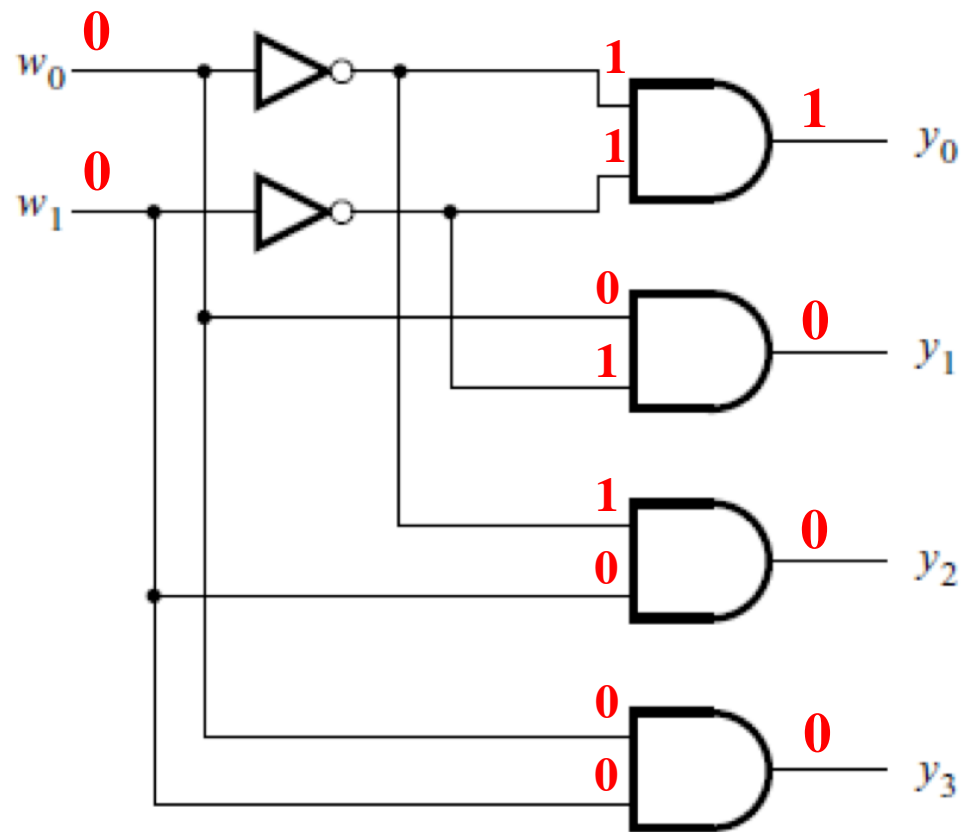
[Figure 4.13c from the textbook]

The Logic Circuit for a 2-to-4 Decoder



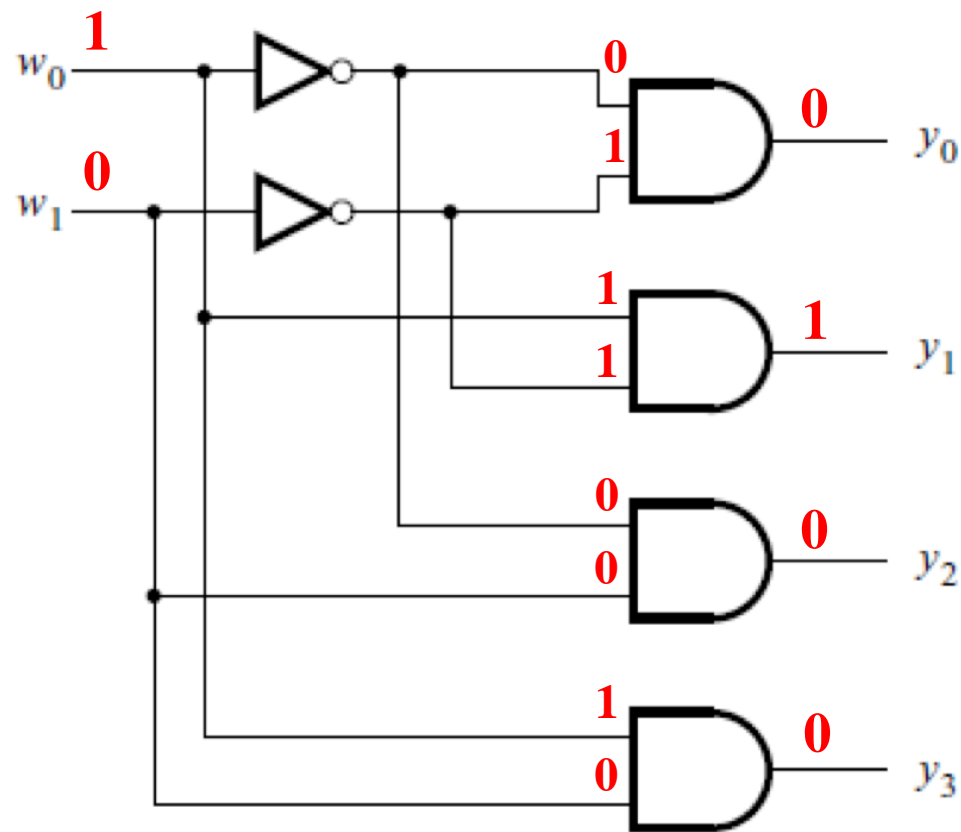
[Figure 4.13c from the textbook]

The Logic Circuit for a 2-to-4 Decoder



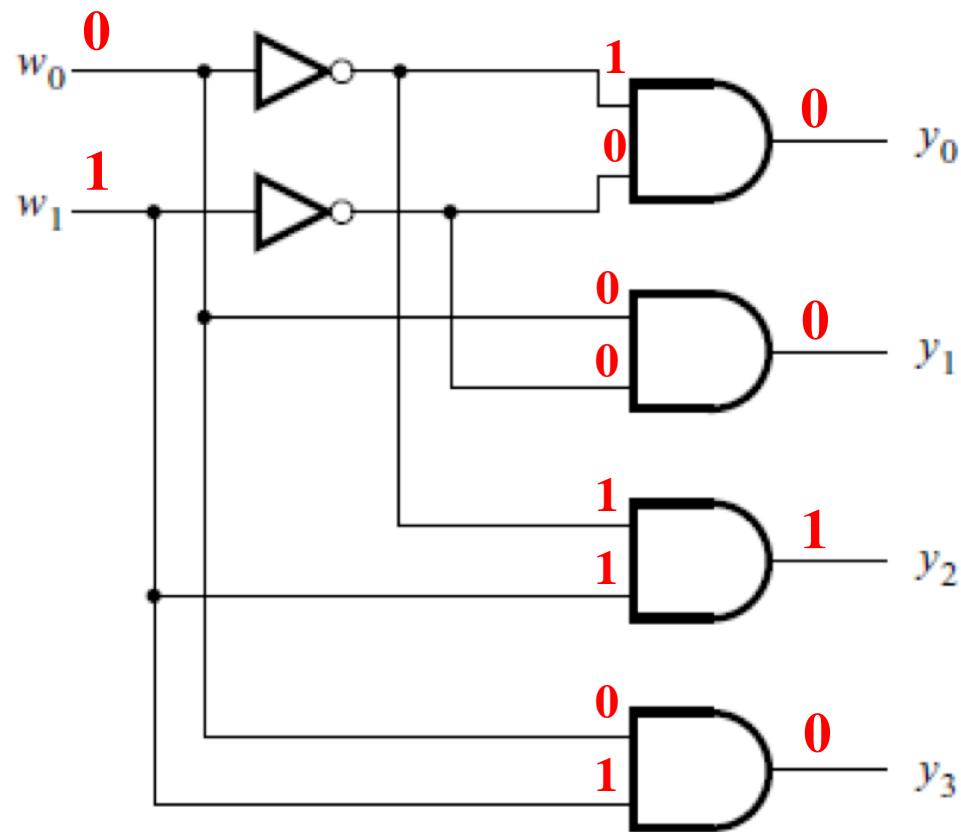
[Figure 4.13c from the textbook]

The Logic Circuit for a 2-to-4 Decoder



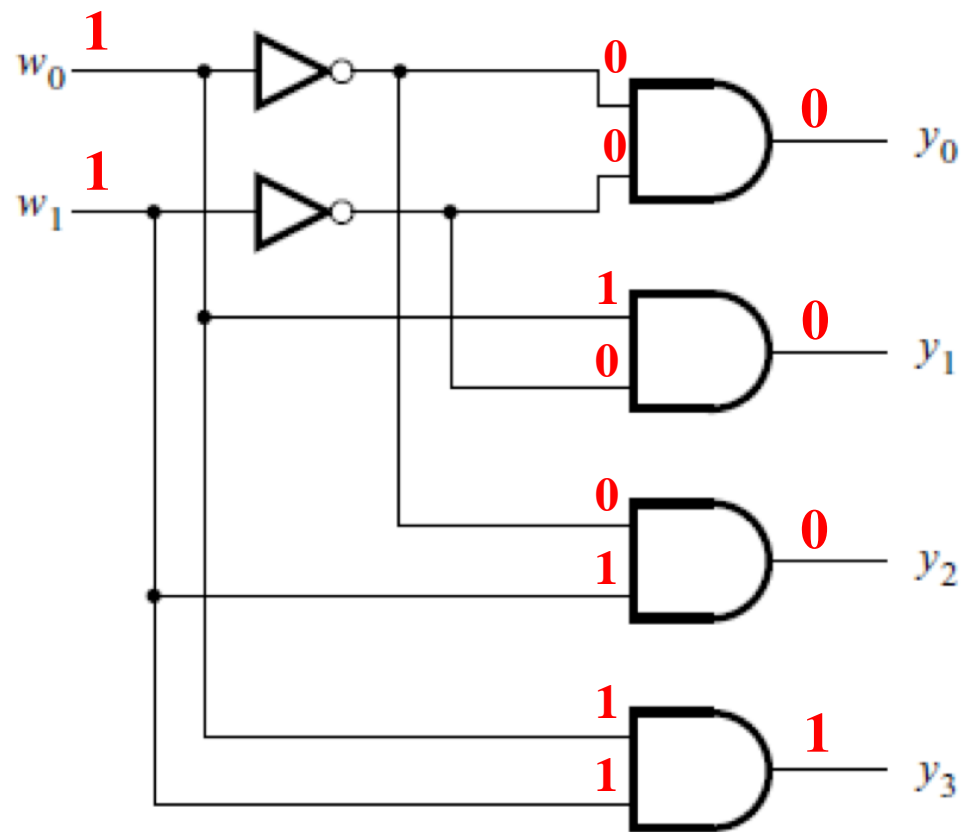
[Figure 4.13c from the textbook]

The Logic Circuit for a 2-to-4 Decoder



[Figure 4.13c from the textbook]

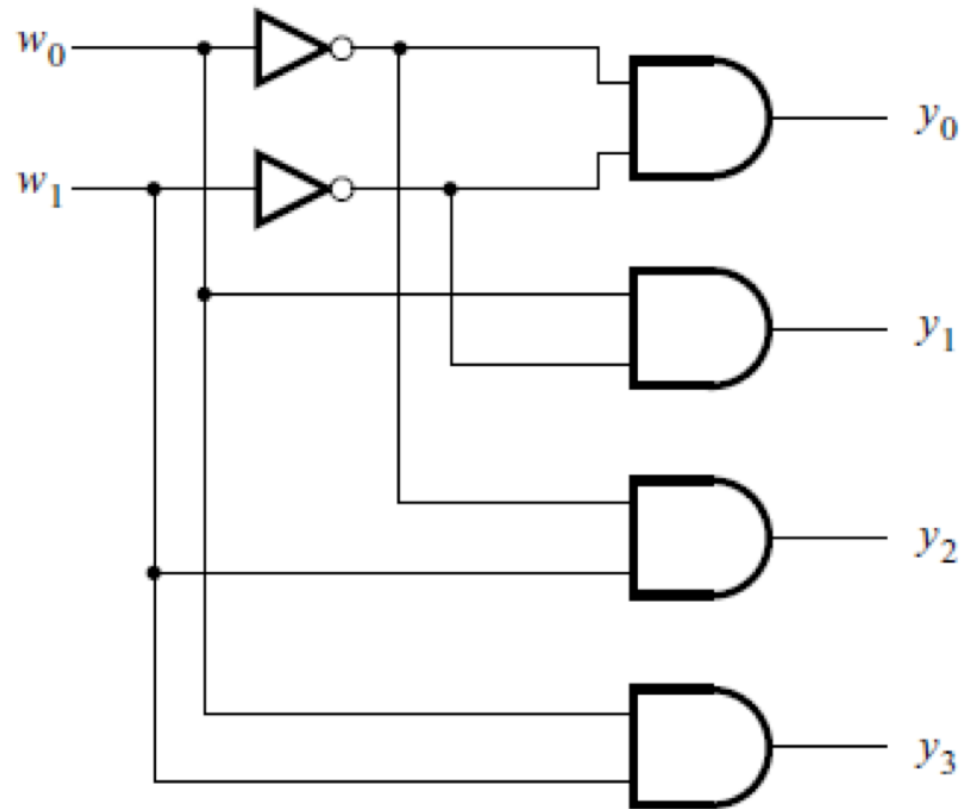
The Logic Circuit for a 2-to-4 Decoder



[Figure 4.13c from the textbook]

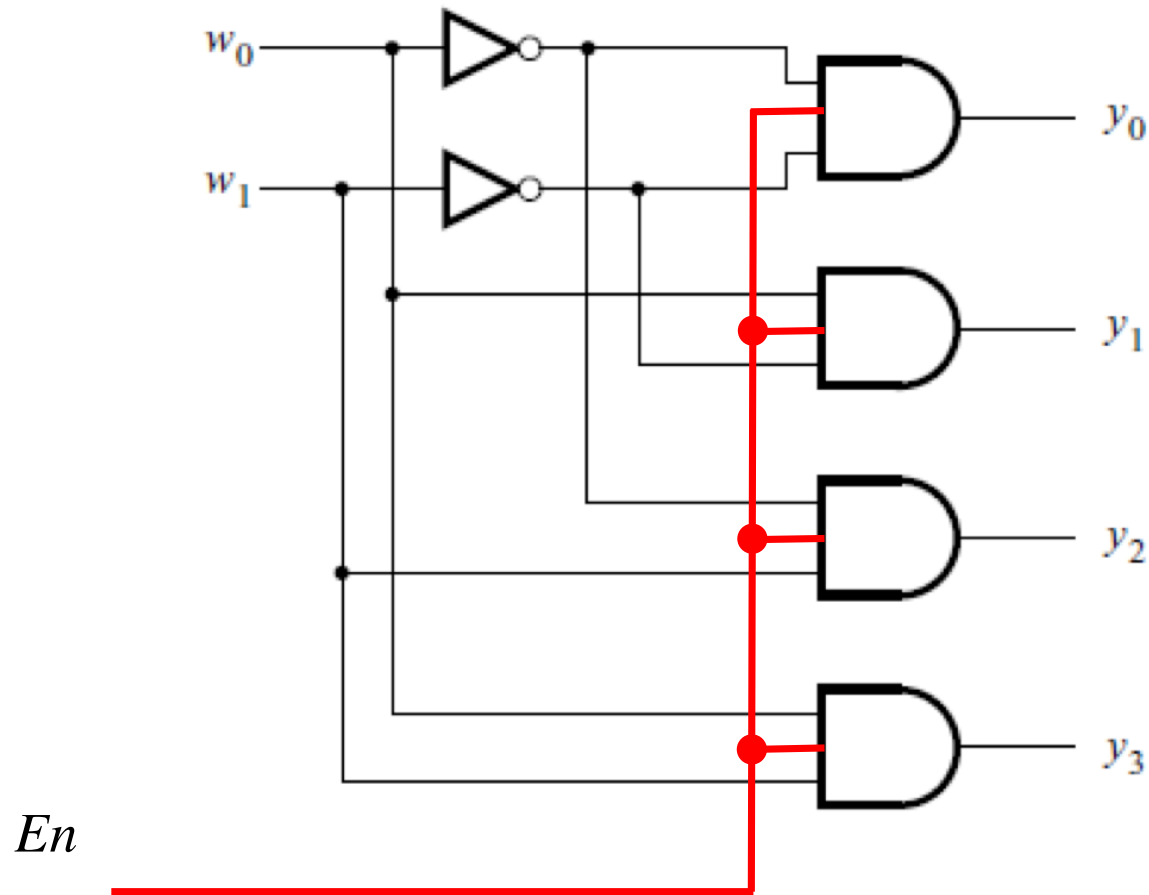
Decoder with an Enable Input

Adding an Enable Input



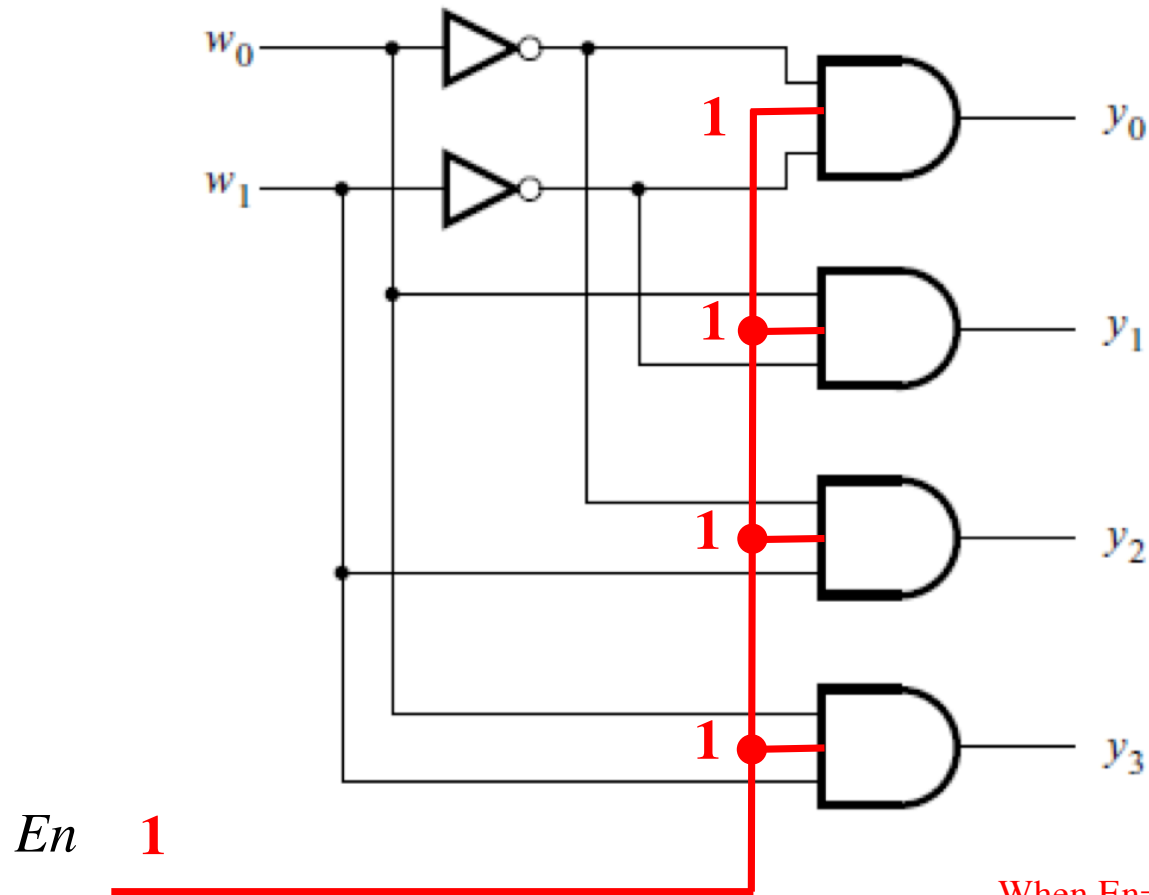
[Figure 4.13c from the textbook]

Adding an Enable Input



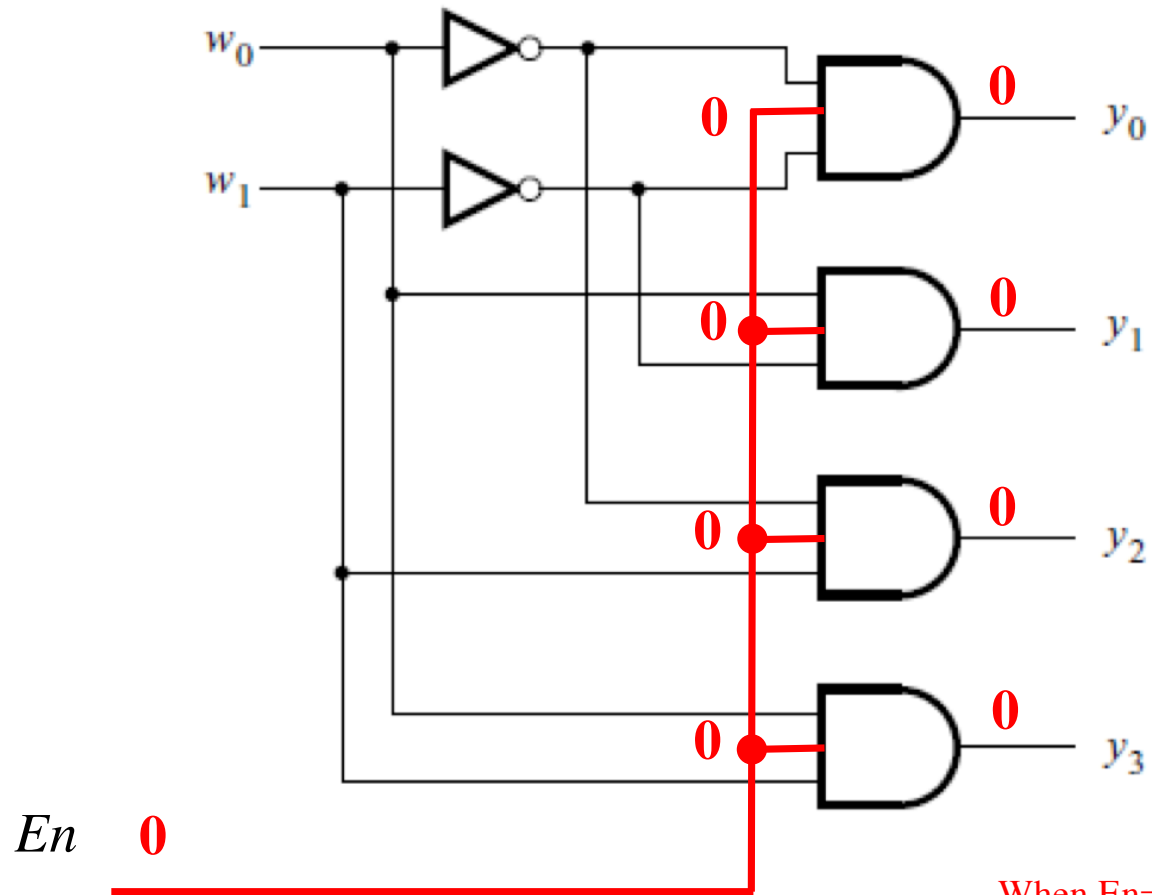
[Figure 4.14c from the textbook]

2-to-4 Decoder with an Enable Input



When $En=1$ this circuit behaves like the one without enable.

2-to-4 Decoder with an Enable Input

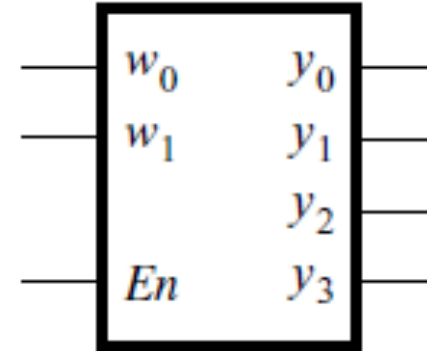


When $En=0$ all outputs are set to 0 and the decoder is disabled.

Truth Table and Graphical Symbol for a 2-to-4 Decoder with an Enable Input

En	w_1	w_0	y_0	y_1	y_2	y_3
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1
0	x	x	0	0	0	0

(a) Truth table



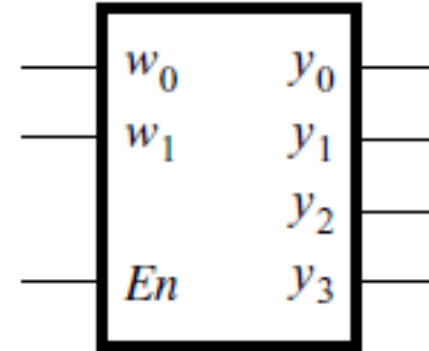
(b) Graphical symbol

Truth Table and Graphical Symbol for a 2-to-4 Decoder with an Enable Input

En	w_1	w_0	y_0	y_1	y_2	y_3
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1
0	x	x	0	0	0	0

(a) Truth table

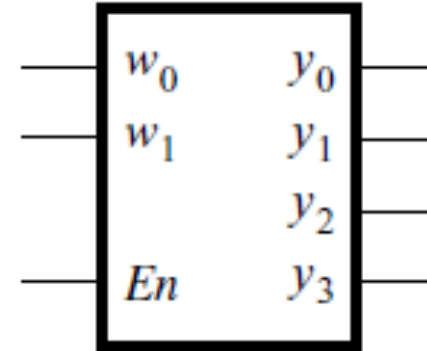
x indicates that it does not matter what the value of this variable is for this row of the truth table



(b) Graphical symbol

Truth Table and Graphical Symbol for a 2-to-4 Decoder with an Enable Input

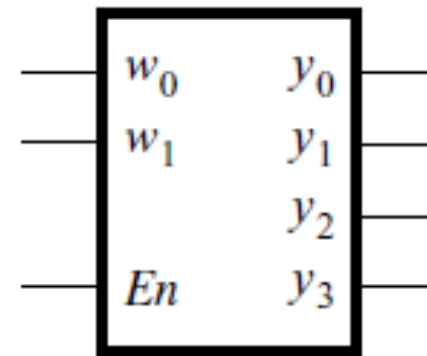
En	w_1	w_0	y_0	y_1	y_2	y_3
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1
0	0	0	0	0	0	0
0	0	1	0	0	0	0
0	1	0	0	0	0	0
0	1	1	0	0	0	0



(b) Graphical symbol

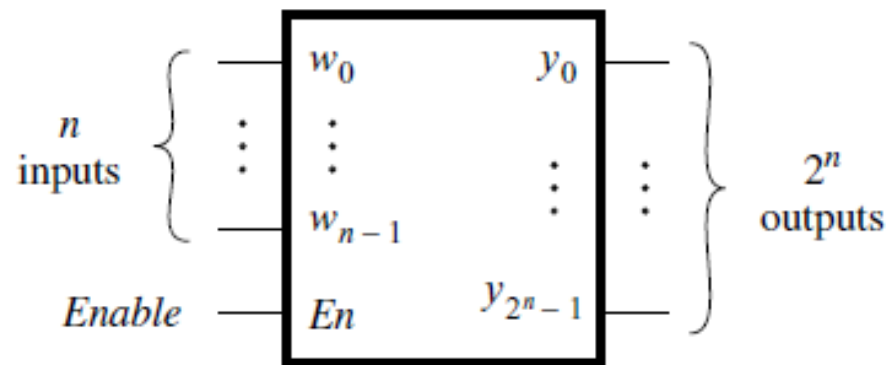
Truth Table and Graphical Symbol for a 2-to-4 Decoder with an Enable Input

En	w_1	w_0	y_0	y_1	y_2	y_3
0	0	0	0	0	0	0
0	0	1	0	0	0	0
0	1	0	0	0	0	0
0	1	1	0	0	0	0
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1



(b) Graphical symbol

Graphical Symbol for a Binary n -to- 2^n Decoder with an Enable Input



(d) An n -to- 2^n decoder

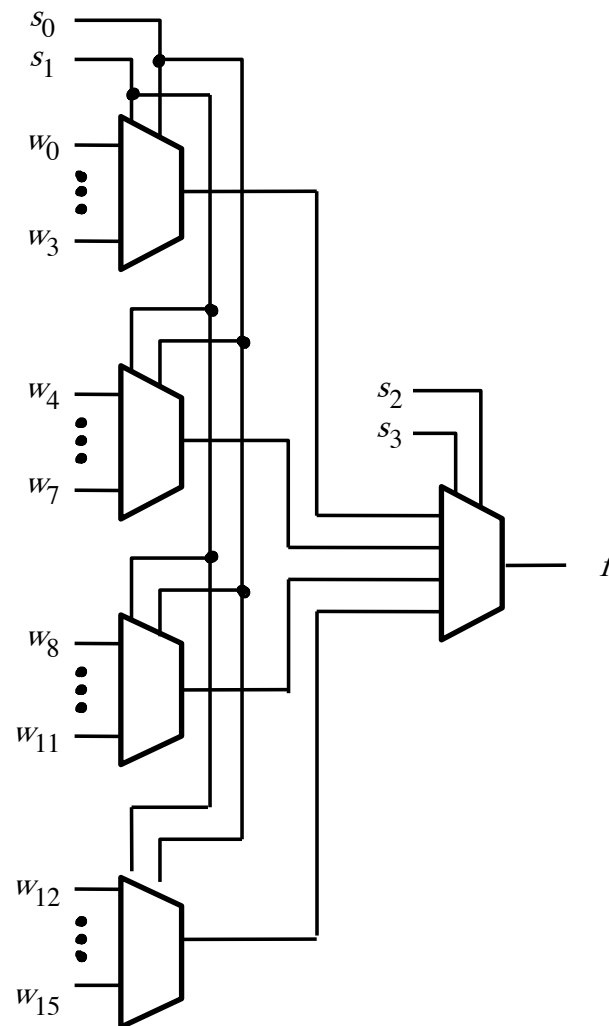
A binary decoder with n inputs has 2^n outputs

The outputs of an enabled binary decoder are “one-hot” encoded, meaning that only a single bit is set to 1, i.e., it is *hot*.

How can we build larger decoders?

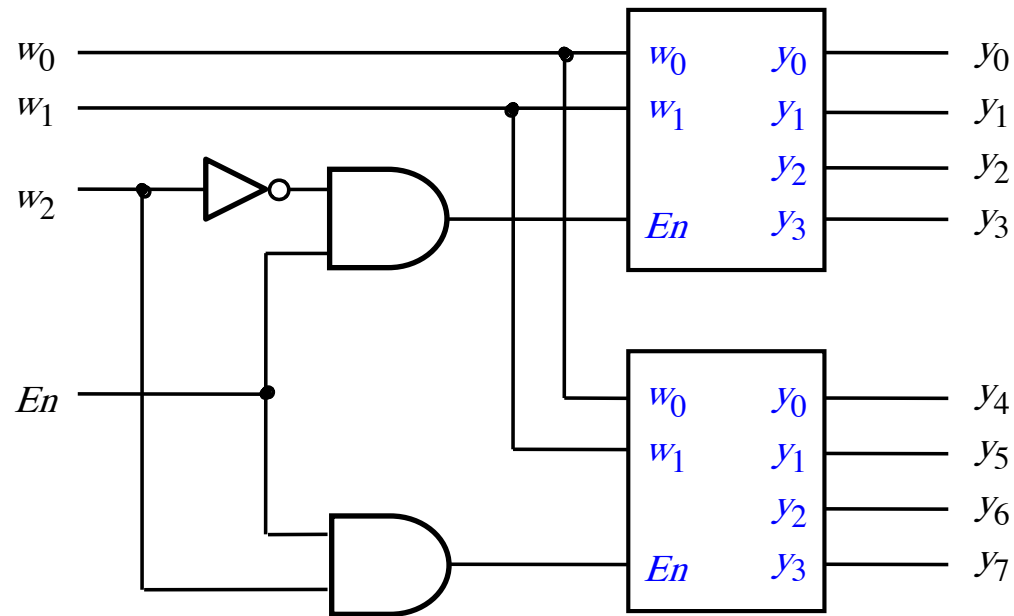
- 3-to-8 ?
- 4-to-16?
- 5-to-??

Hint: How did we build a 16-1 Multiplexer



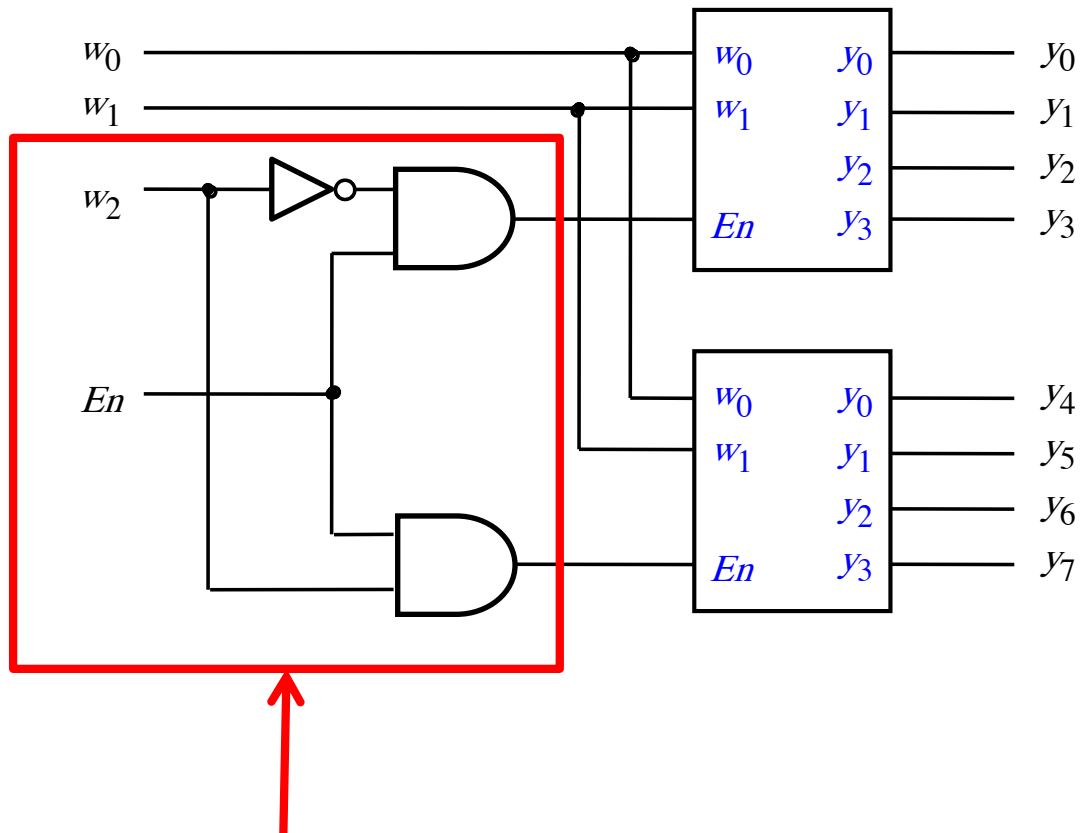
[Figure 4.4 from the textbook]

A 3-to-8 decoder using two 2-to-4 decoders



[Figure 4.15 from the textbook]

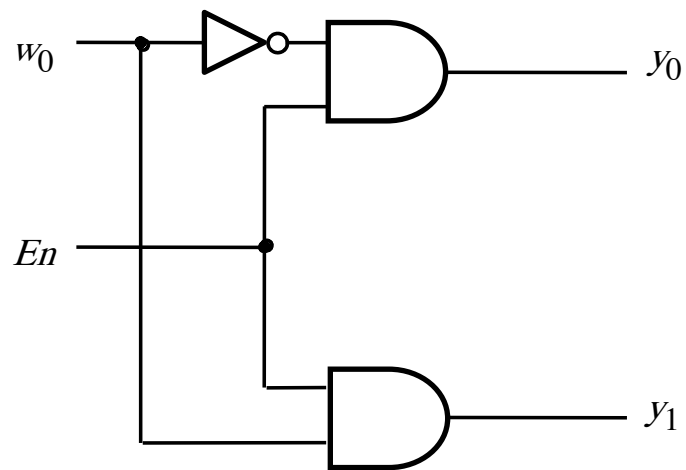
A 3-to-8 decoder using two 2-to-4 decoders



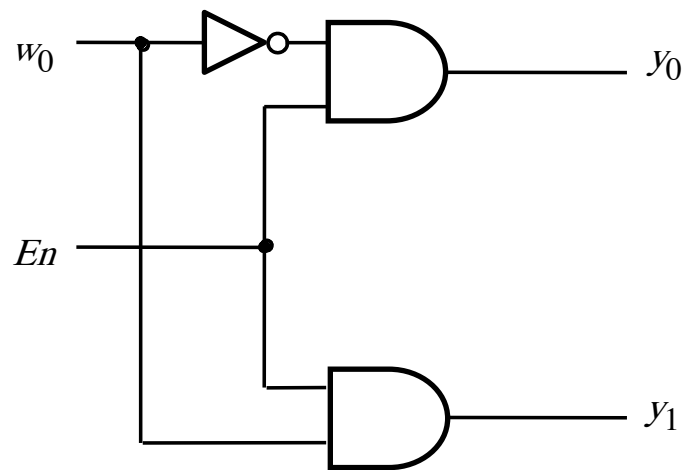
What is this?

[Figure 4.15 from the textbook]

What is this?

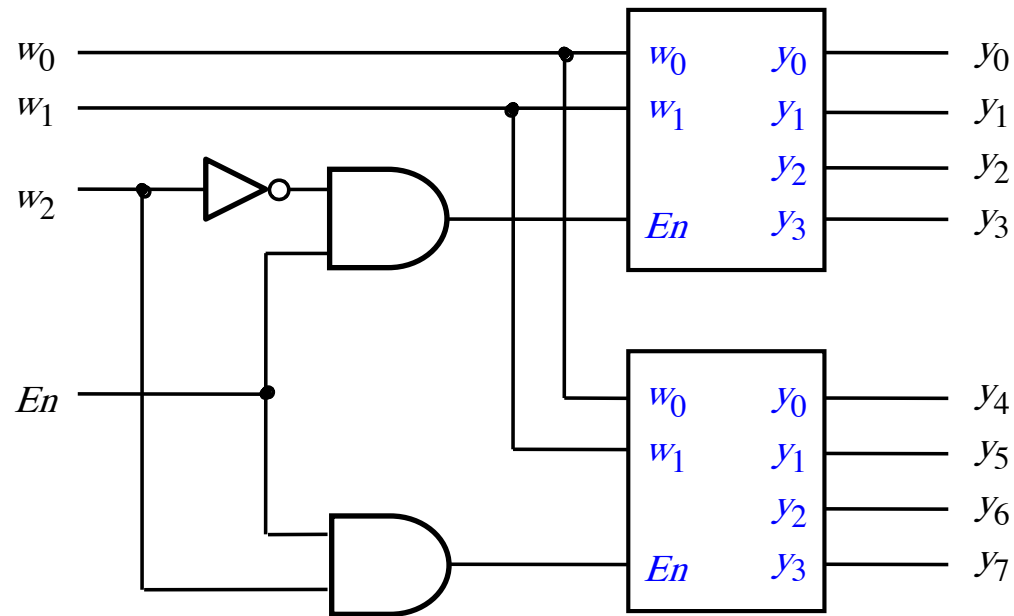


What is this?



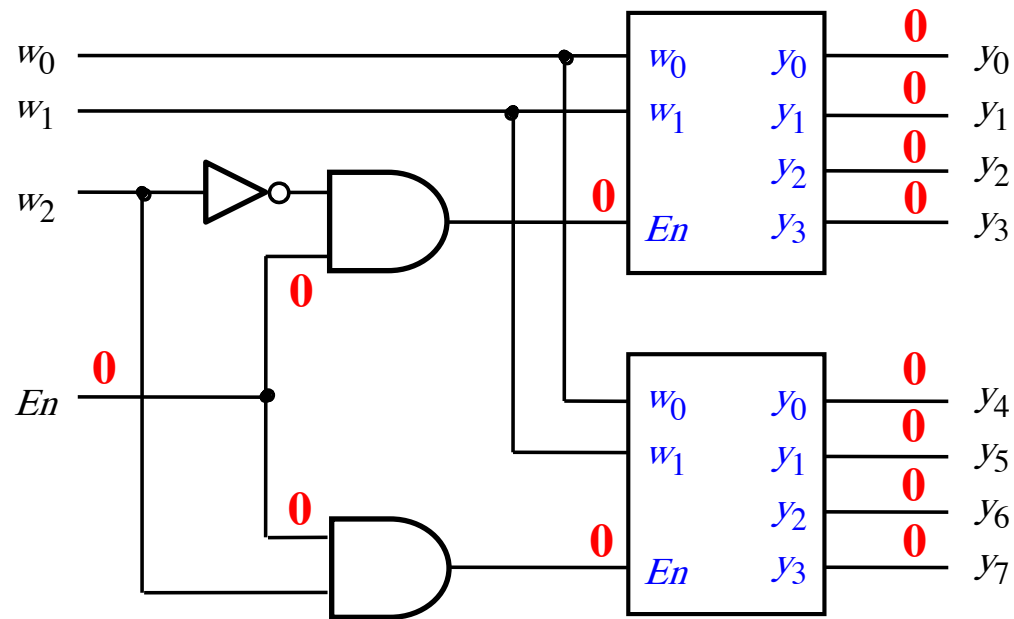
This is a 1-to-2 decoder with an enable input.

A 3-to-8 decoder using two 2-to-4 decoders



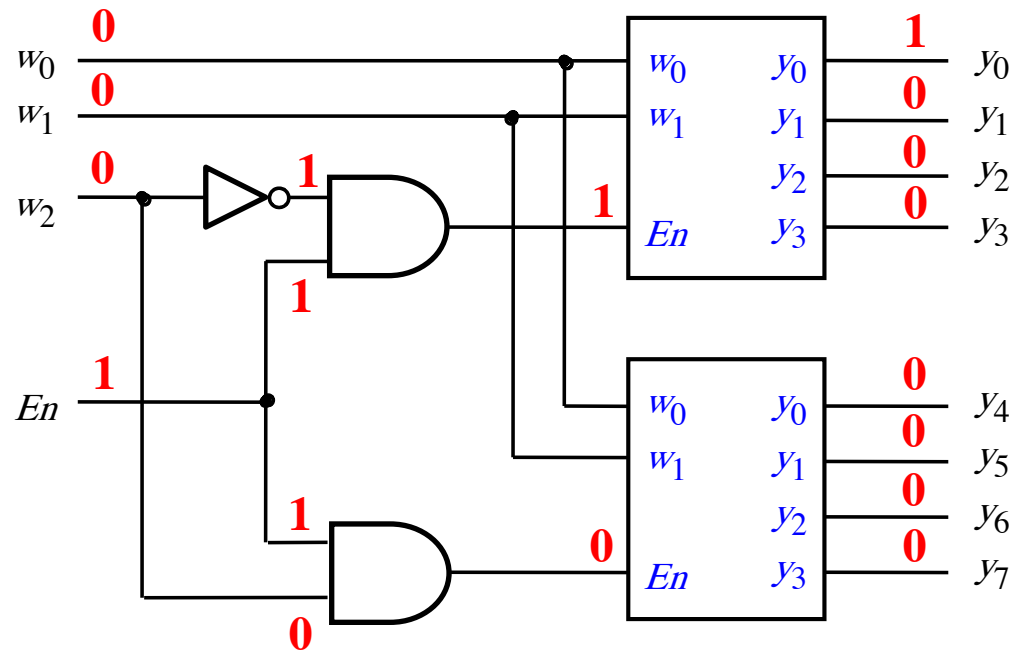
[Figure 4.15 from the textbook]

A 3-to-8 decoder using two 2-to-4 decoders



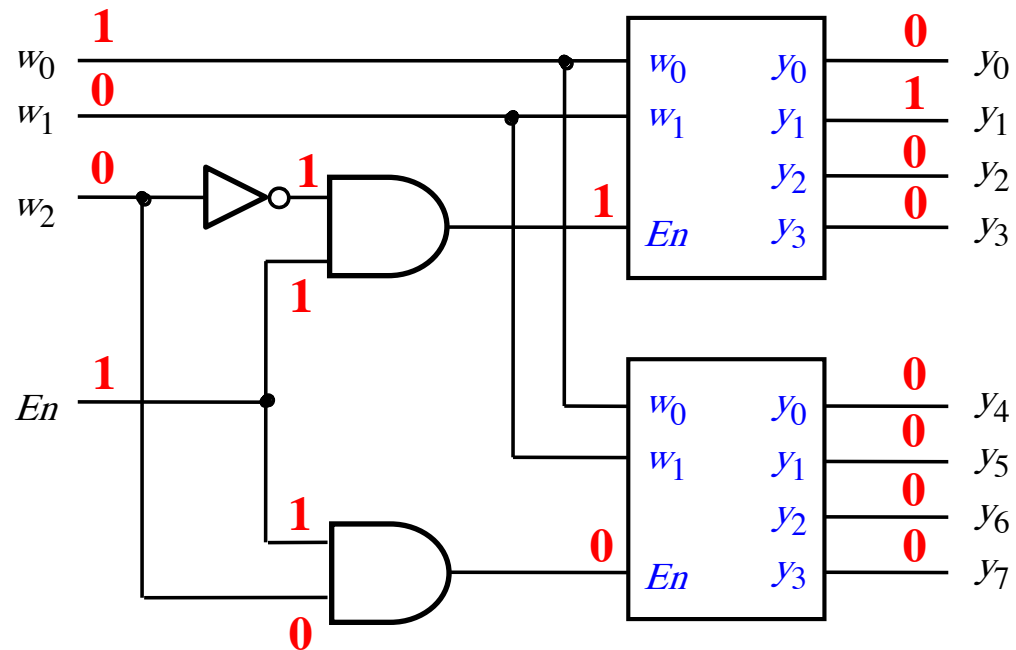
[Figure 4.15 from the textbook]

A 3-to-8 decoder using two 2-to-4 decoders



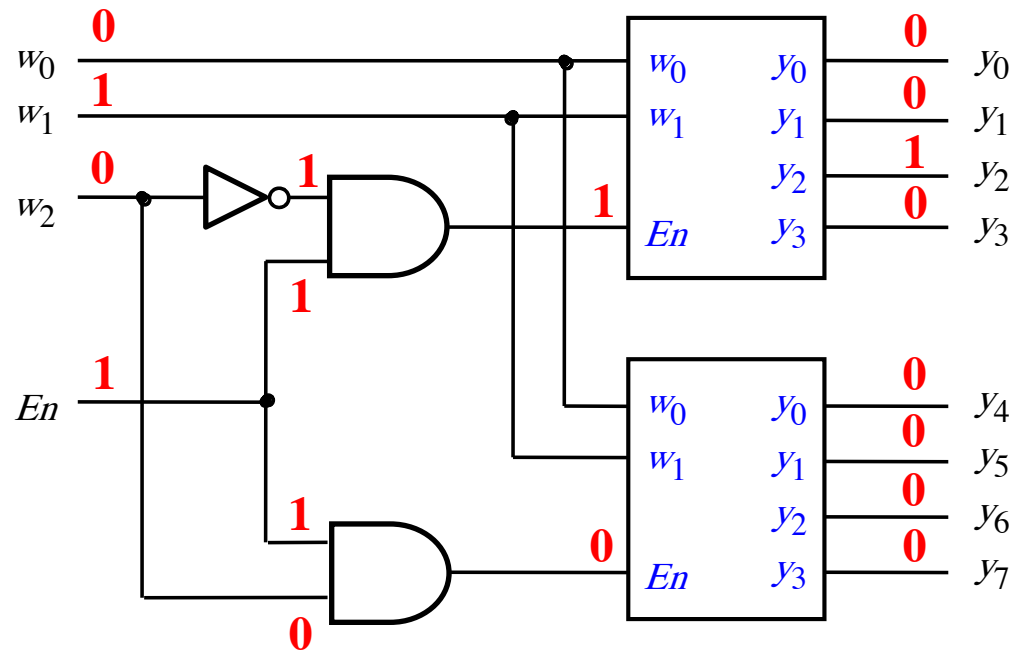
[Figure 4.15 from the textbook]

A 3-to-8 decoder using two 2-to-4 decoders



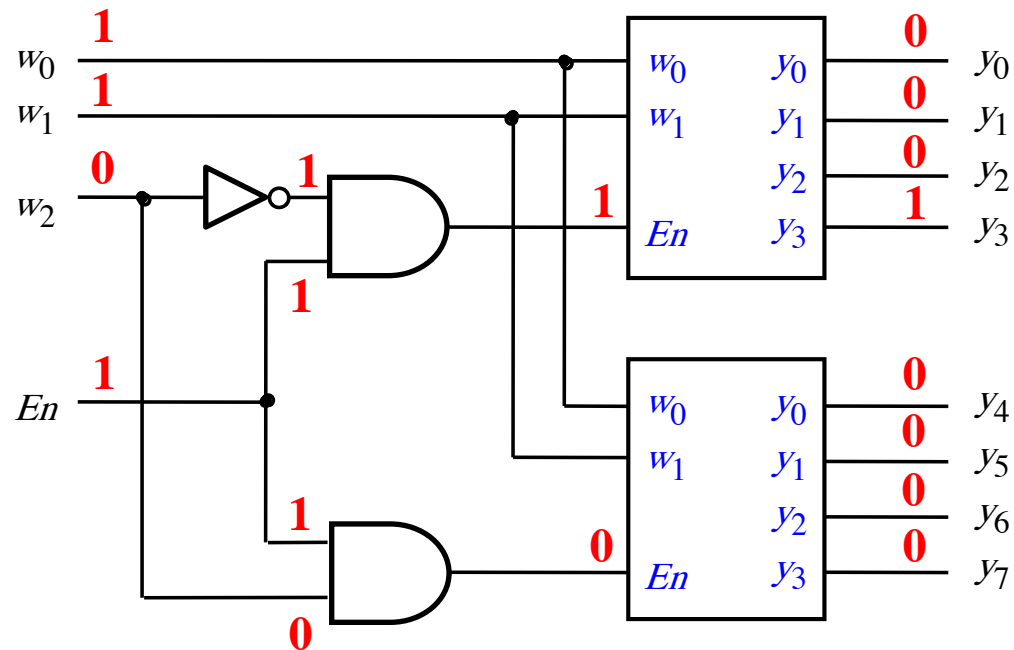
[Figure 4.15 from the textbook]

A 3-to-8 decoder using two 2-to-4 decoders



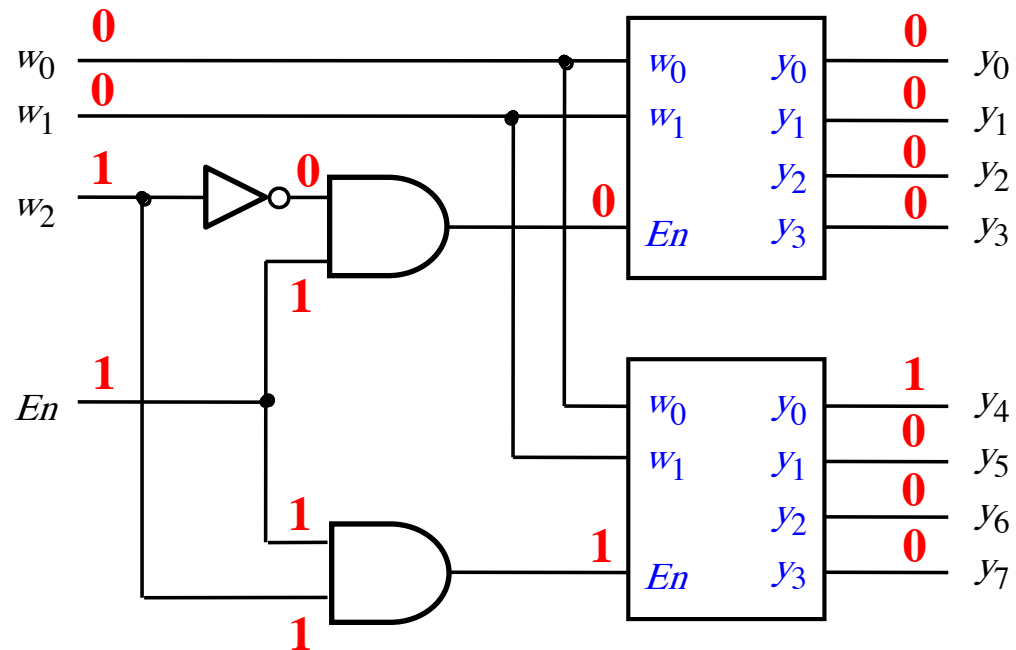
[Figure 4.15 from the textbook]

A 3-to-8 decoder using two 2-to-4 decoders



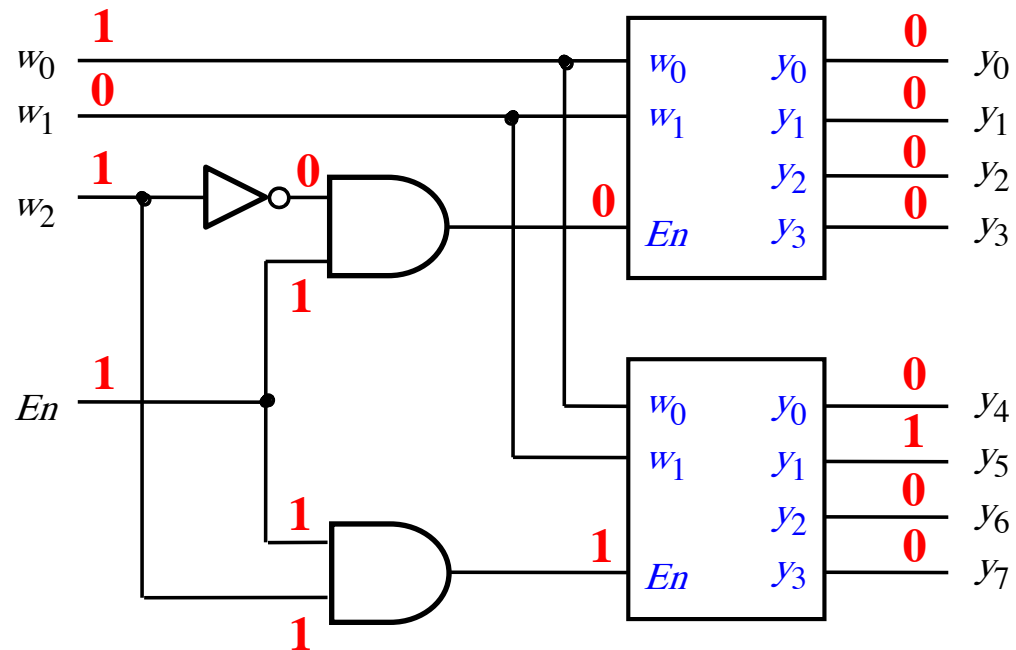
[Figure 4.15 from the textbook]

A 3-to-8 decoder using two 2-to-4 decoders



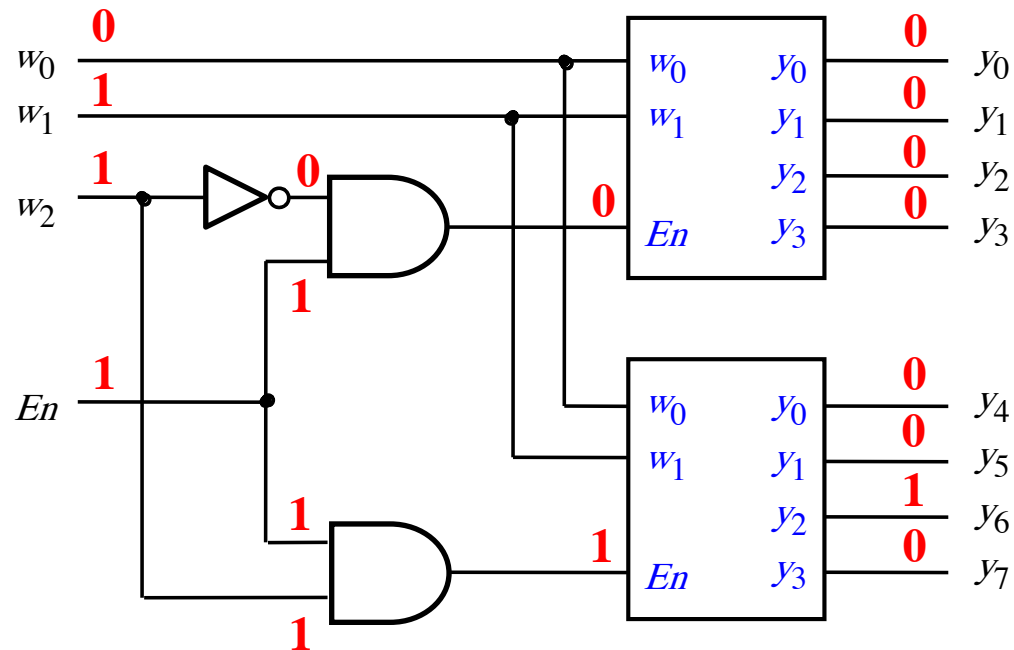
[Figure 4.15 from the textbook]

A 3-to-8 decoder using two 2-to-4 decoders



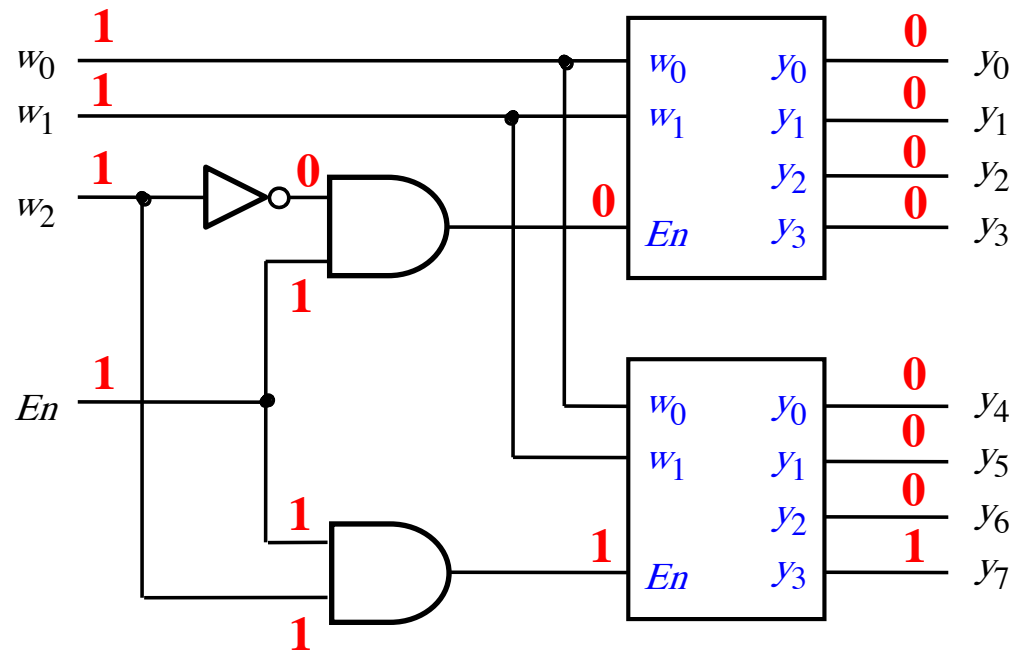
[Figure 4.15 from the textbook]

A 3-to-8 decoder using two 2-to-4 decoders



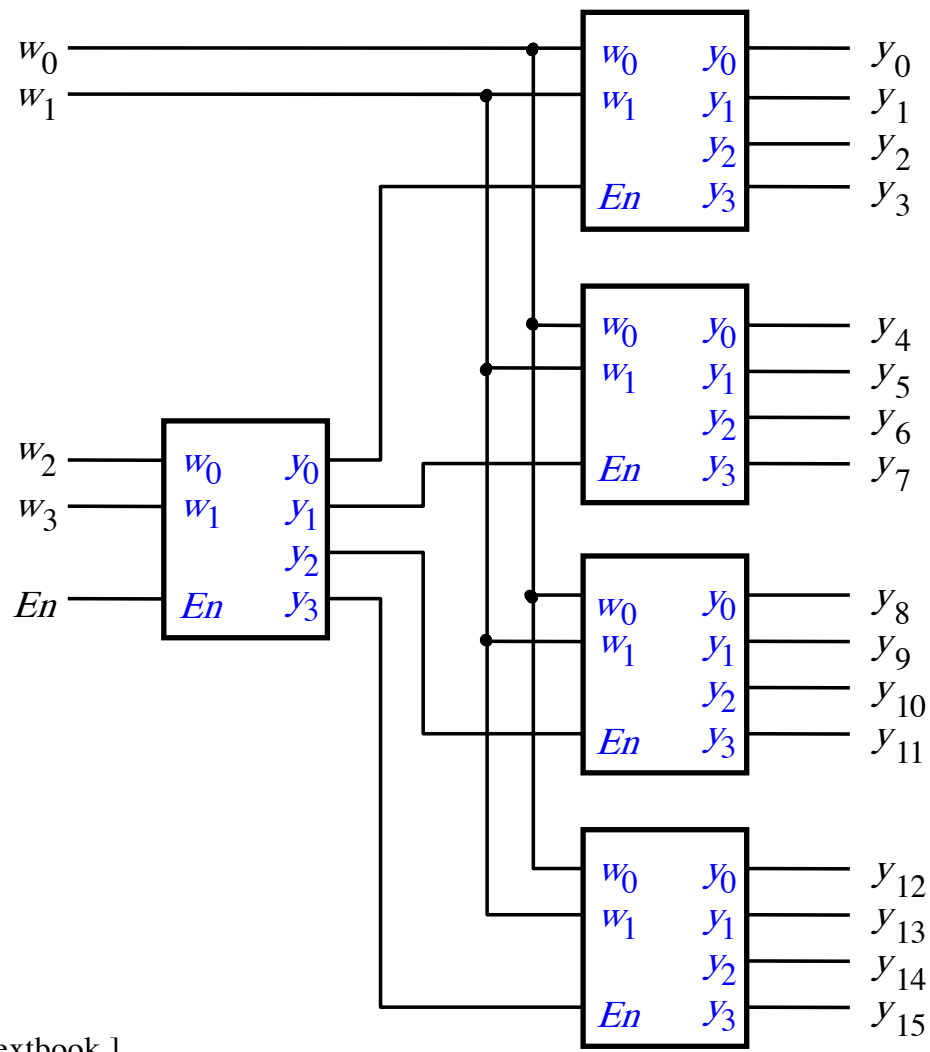
[Figure 4.15 from the textbook]

A 3-to-8 decoder using two 2-to-4 decoders



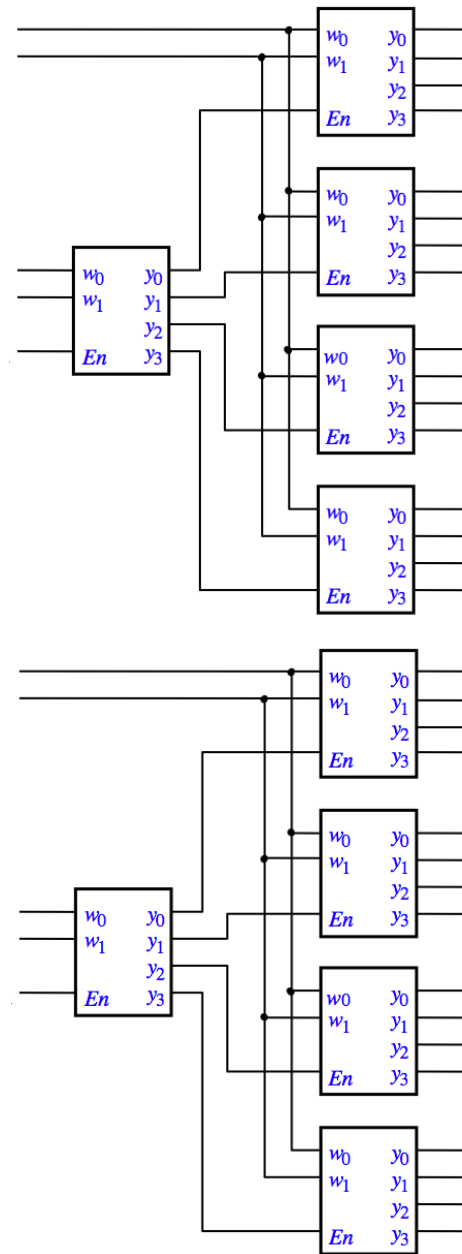
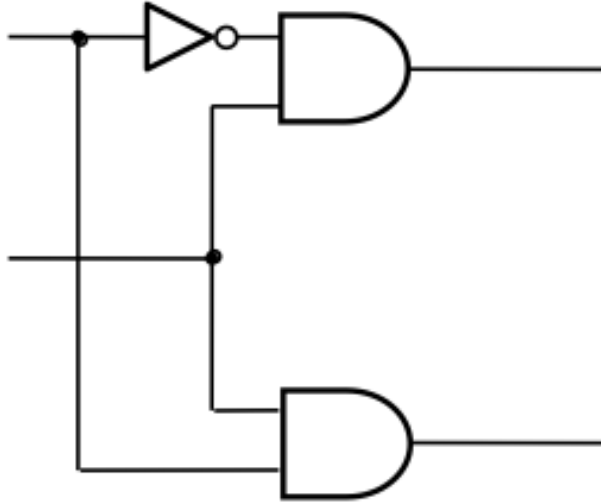
[Figure 4.15 from the textbook]

4-to-16 decoder built using a decoder tree

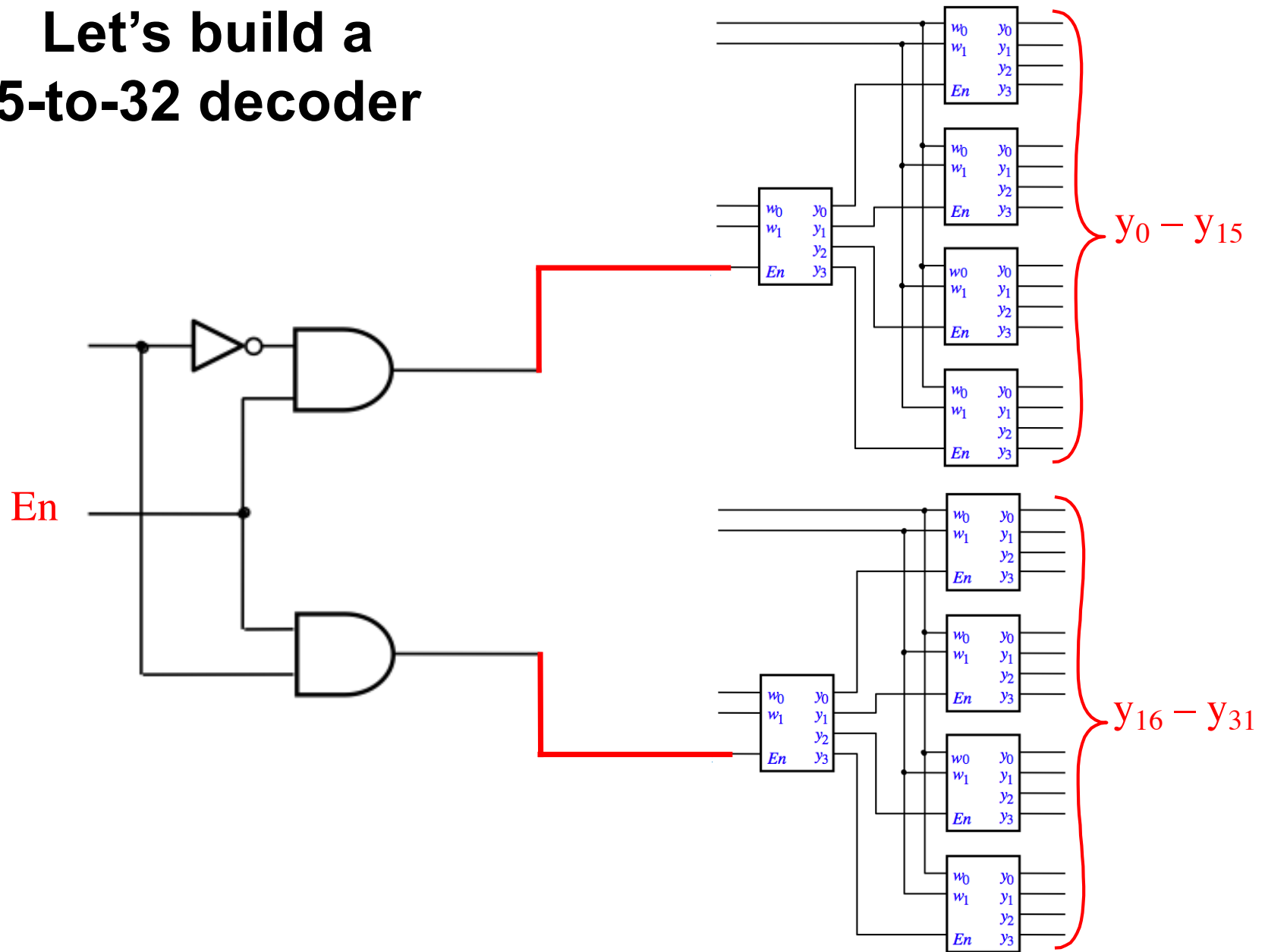


[Figure 4.16 from the textbook]

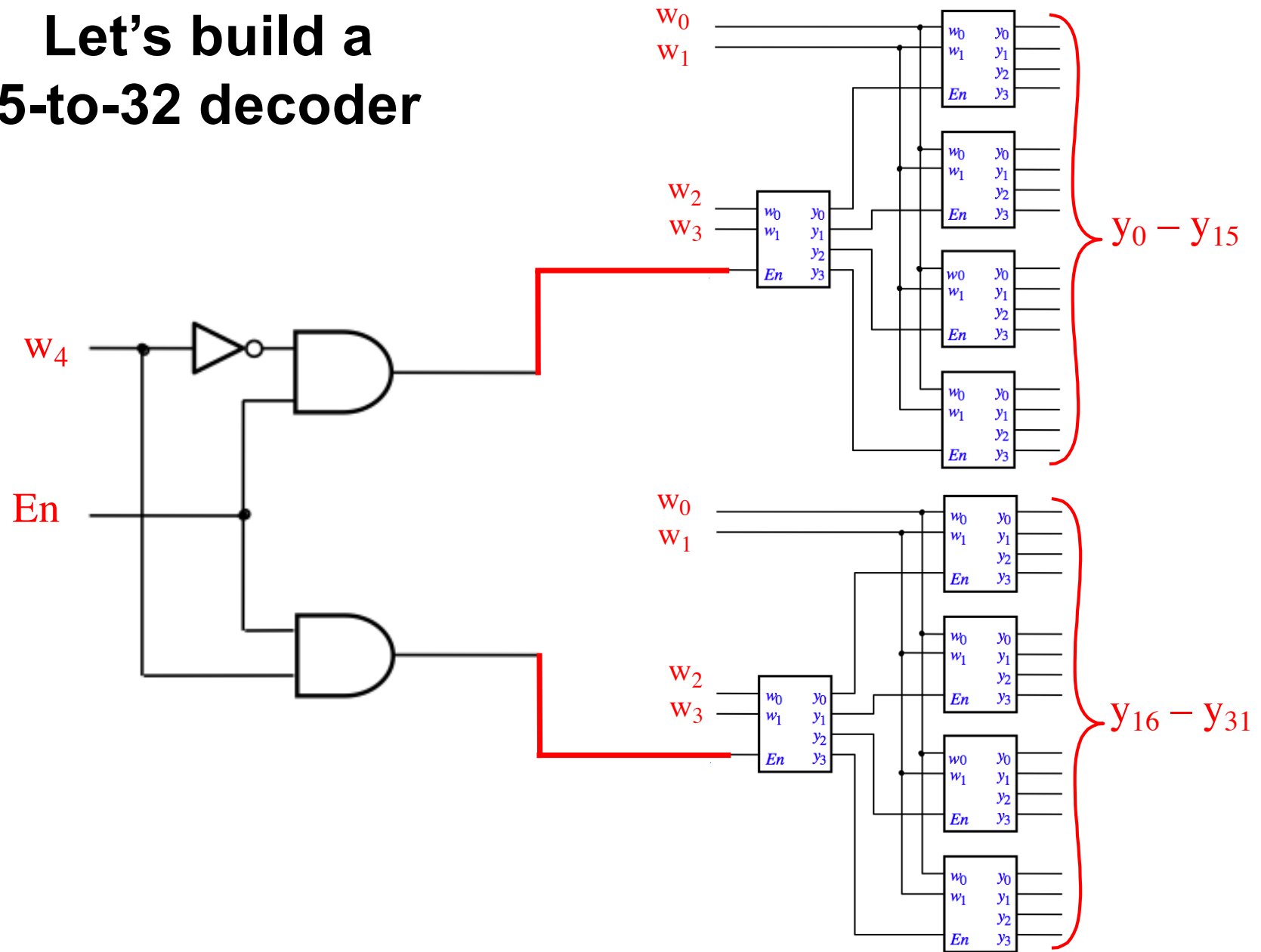
Let's build a 5-to-32 decoder



Let's build a
5-to-32 decoder



Let's build a 5-to-32 decoder

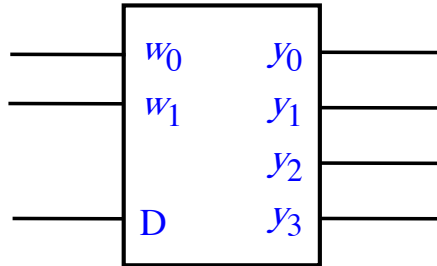


Demultiplexers

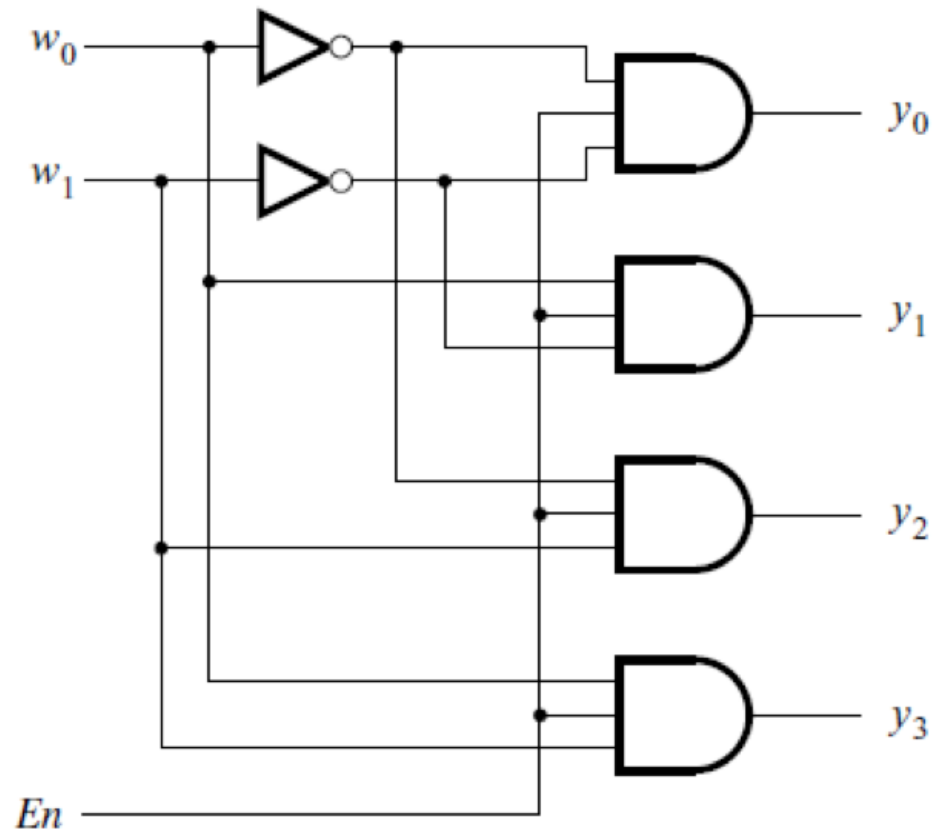
1-to-4 Demultiplexer (Definition)

- Has one data input line: D
- Has two output select lines: w_1 and w_0
- Has four outputs: y_0 , y_1 , y_2 , and y_3
- If $w_1=0$ and $w_0=0$, then the output y_0 is set to D
- If $w_1=0$ and $w_0=1$, then the output y_1 is set to D
- If $w_1=1$ and $w_0=0$, then the output y_2 is set to D
- If $w_1=1$ and $w_0=1$, then the output y_3 is set to D
- Only one output is set to D . All others are set to 0.

Graphical Symbol for a 1-to-4 Demultiplexer

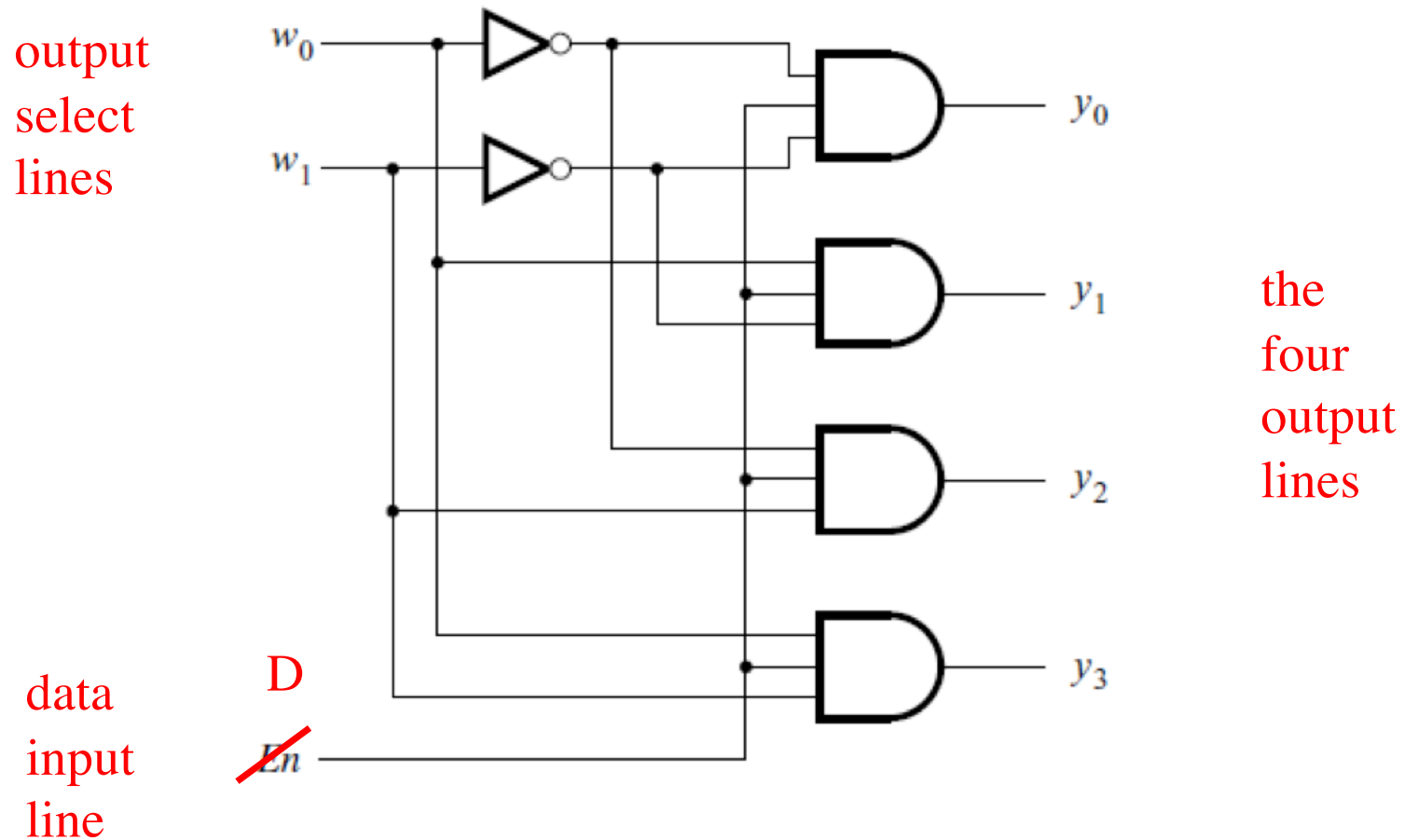


A 1-to-4 demultiplexer built with a 2-to-4 decoder with enable



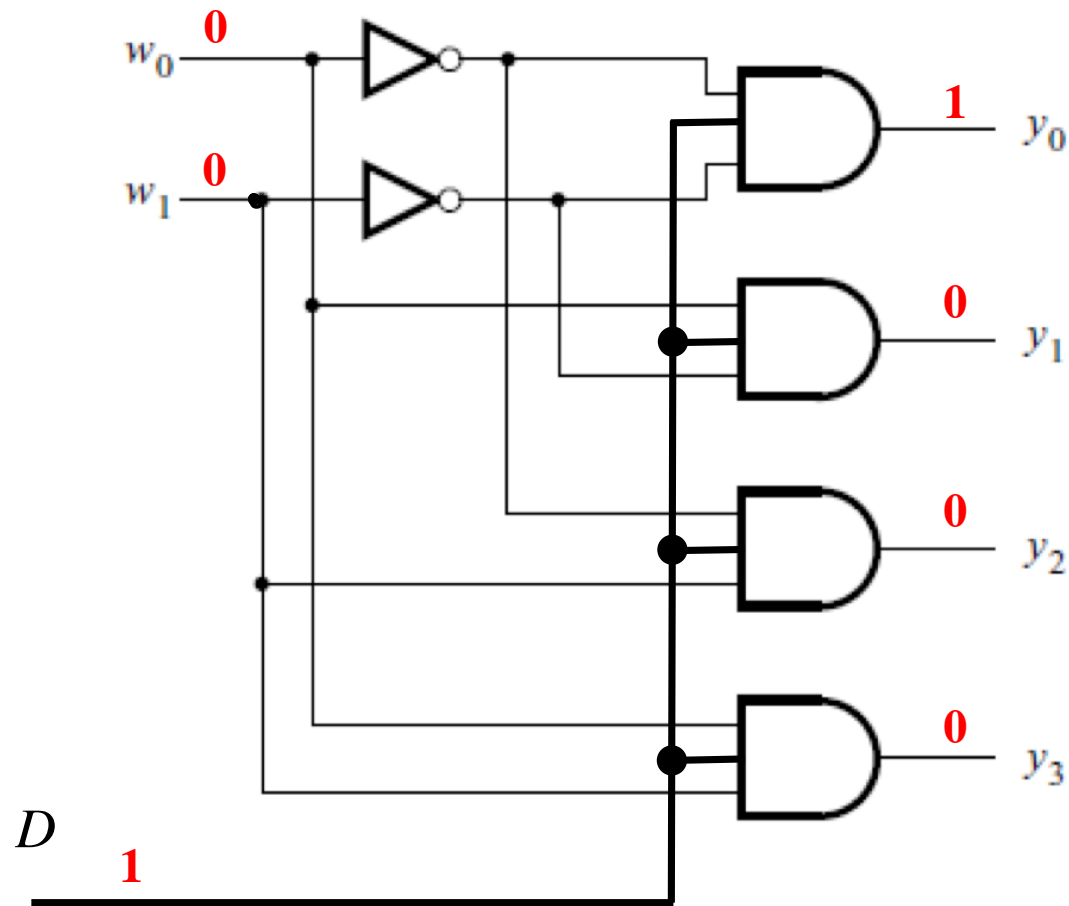
[Figure 4.14c from the textbook]

A 1-to-4 demultiplexer built with a 2-to-4 decoder with enable

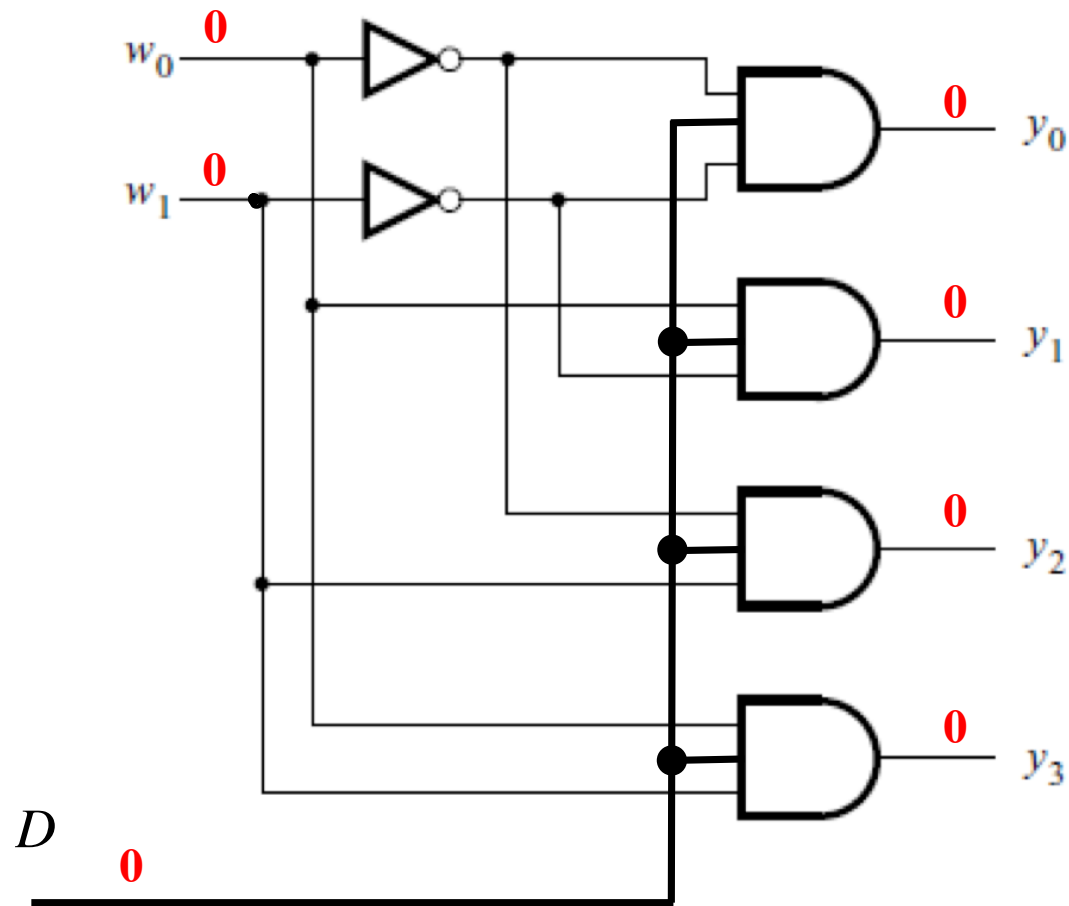


[Figure 4.14c from the textbook]

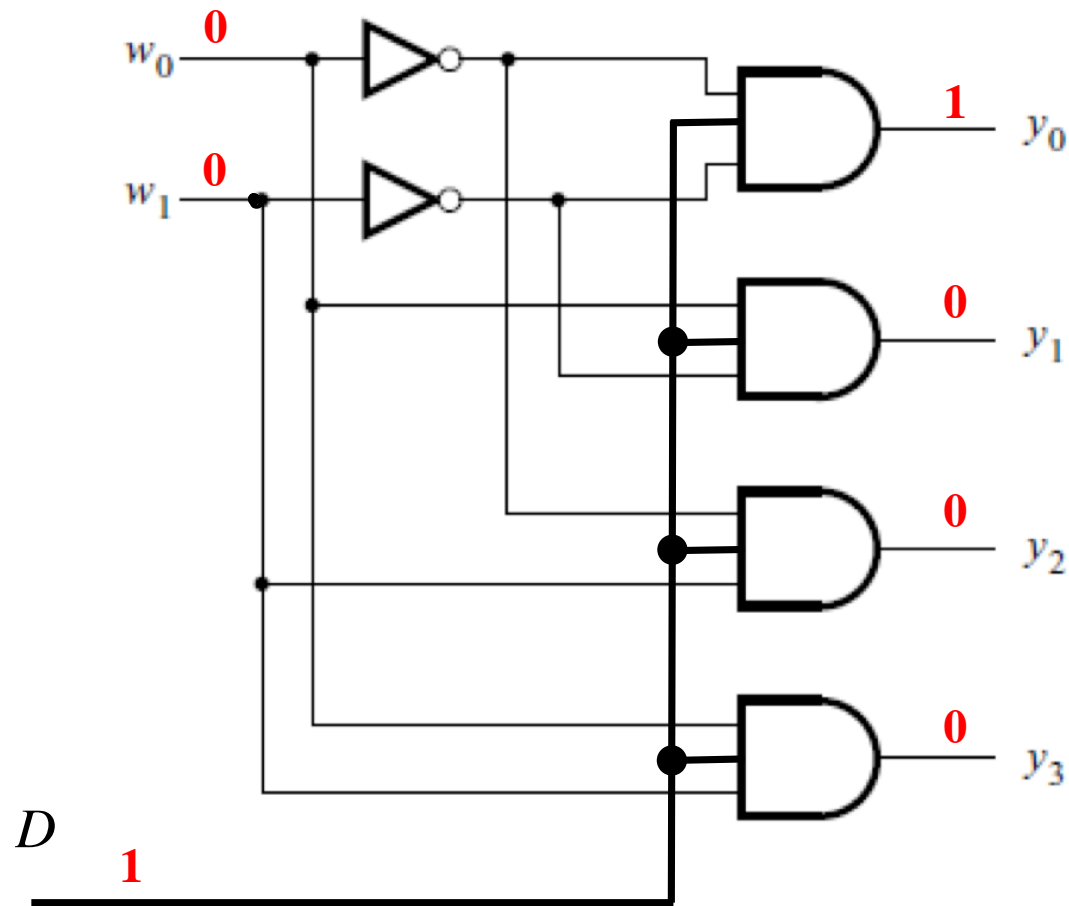
A 1-to-4 demultiplexer built with a 2-to-4 decoder with enable



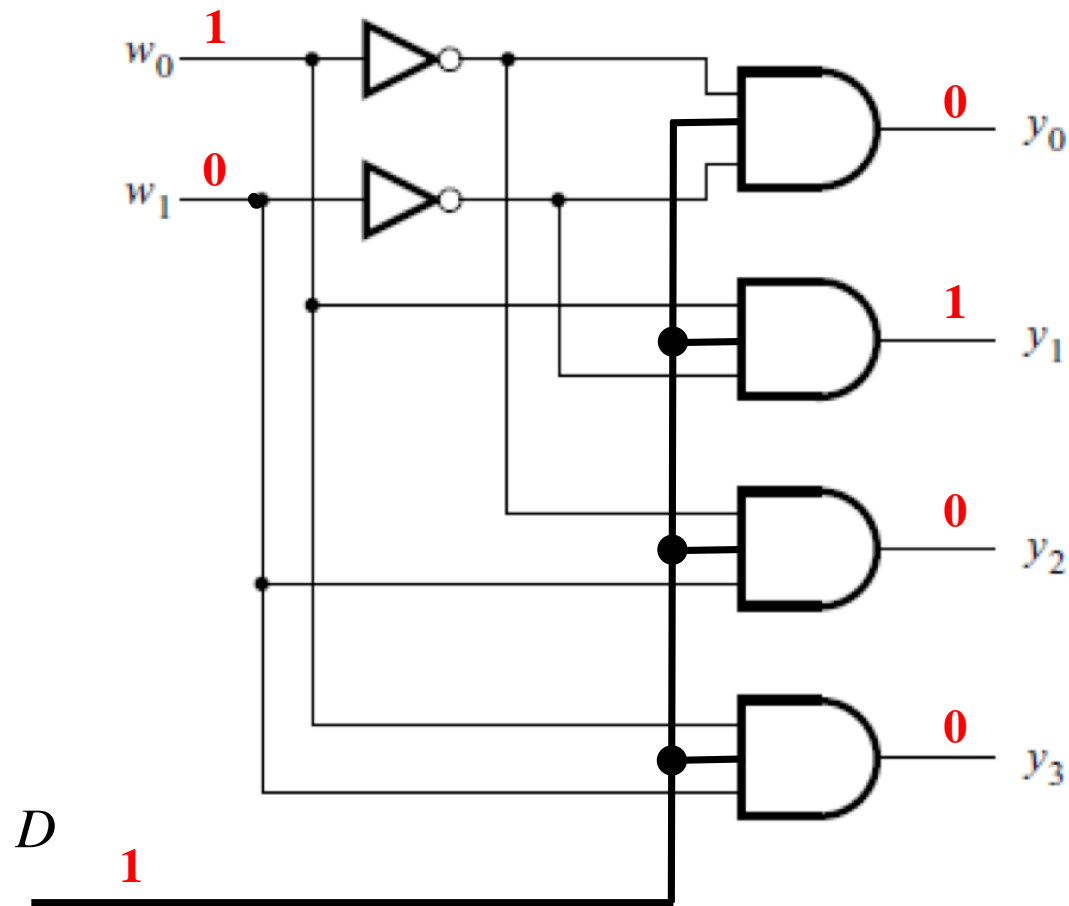
A 1-to-4 demultiplexer built with a 2-to-4 decoder with enable



A 1-to-4 demultiplexer built with a 2-to-4 decoder with enable

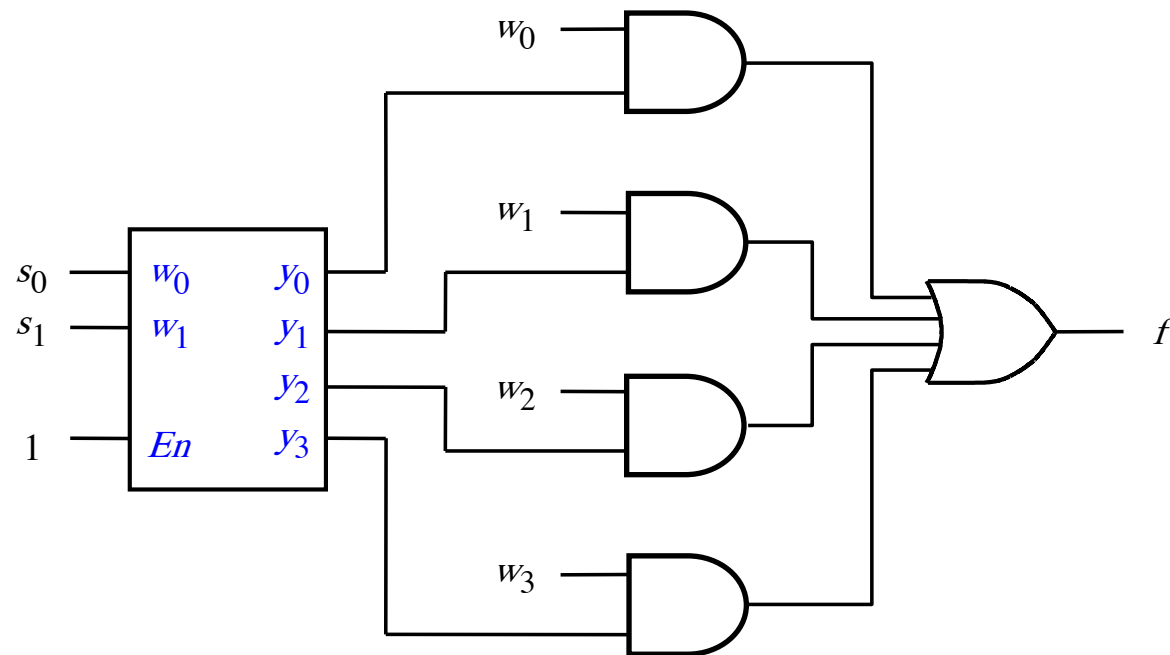


A 1-to-4 demultiplexer built with a 2-to-4 decoder with enable



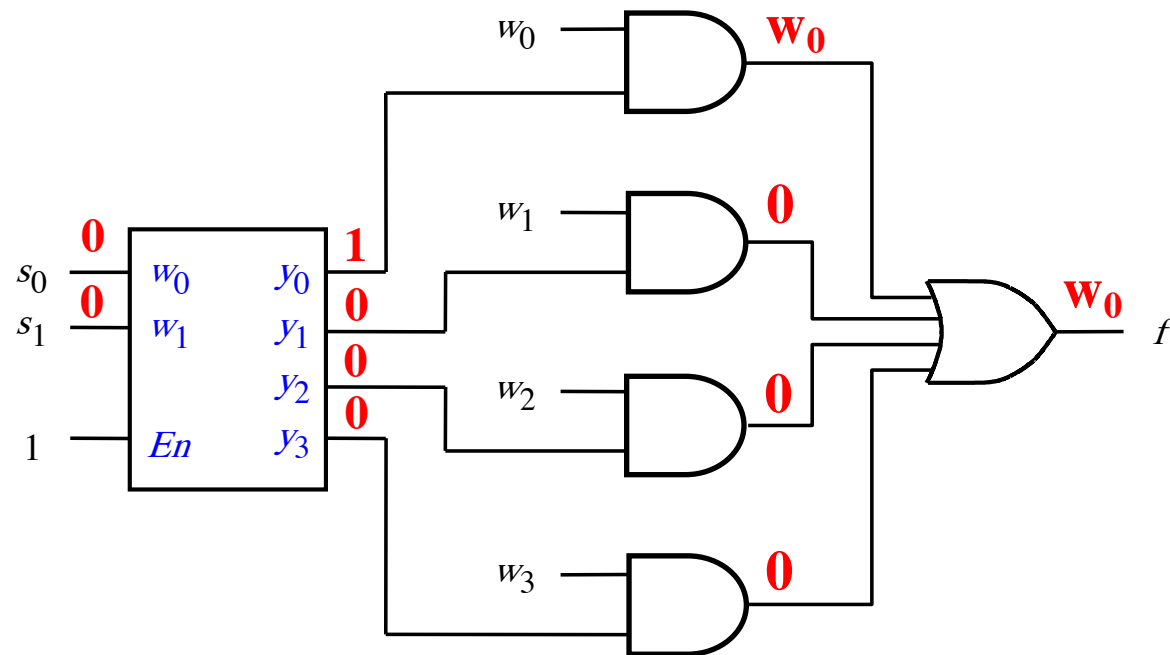
Multiplexers (Implemented with Decoders)

A 4-to-1 multiplexer built using a 2-to-4 decoder



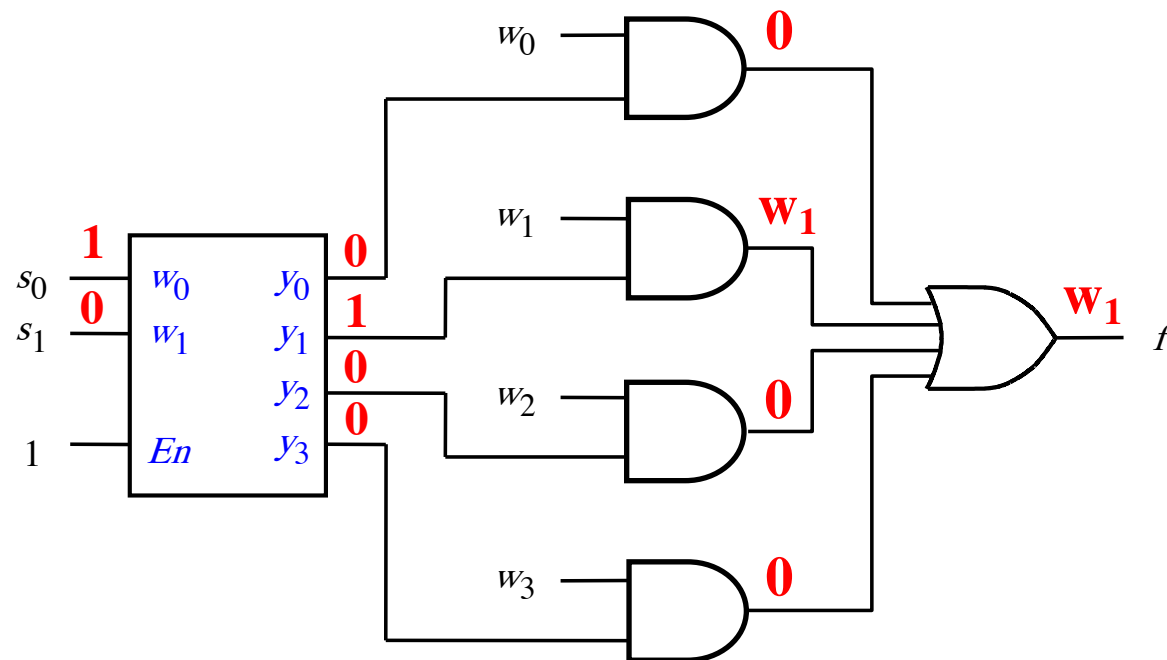
[Figure 4.17 from the textbook]

A 4-to-1 multiplexer built using a 2-to-4 decoder



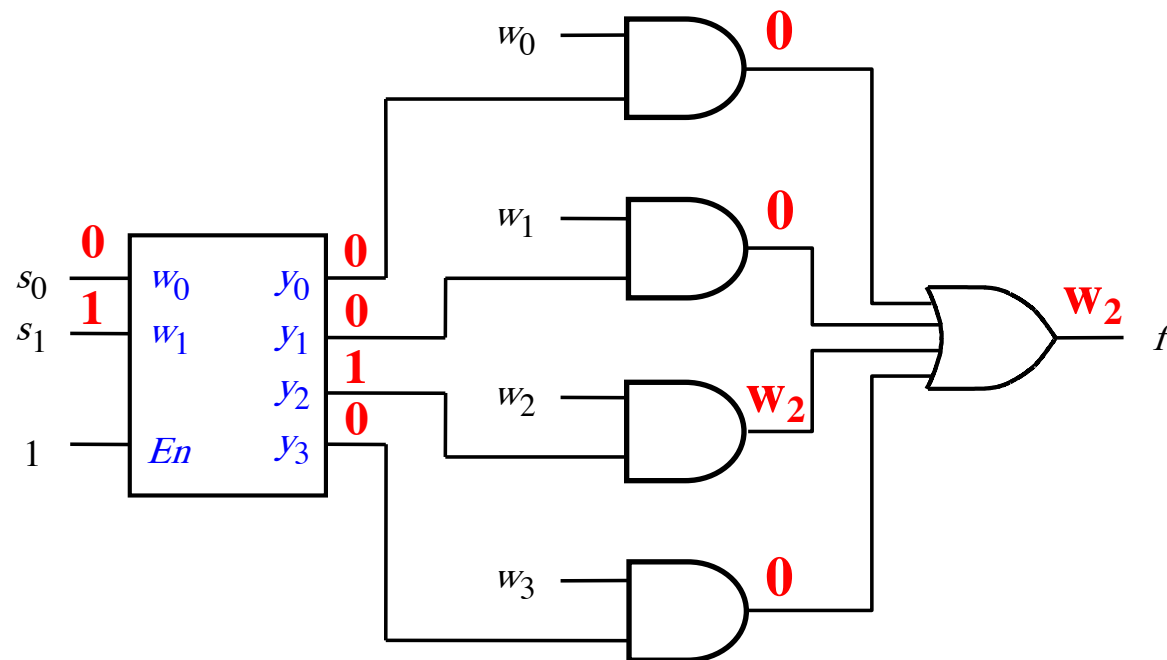
[Figure 4.17 from the textbook]

A 4-to-1 multiplexer built using a 2-to-4 decoder



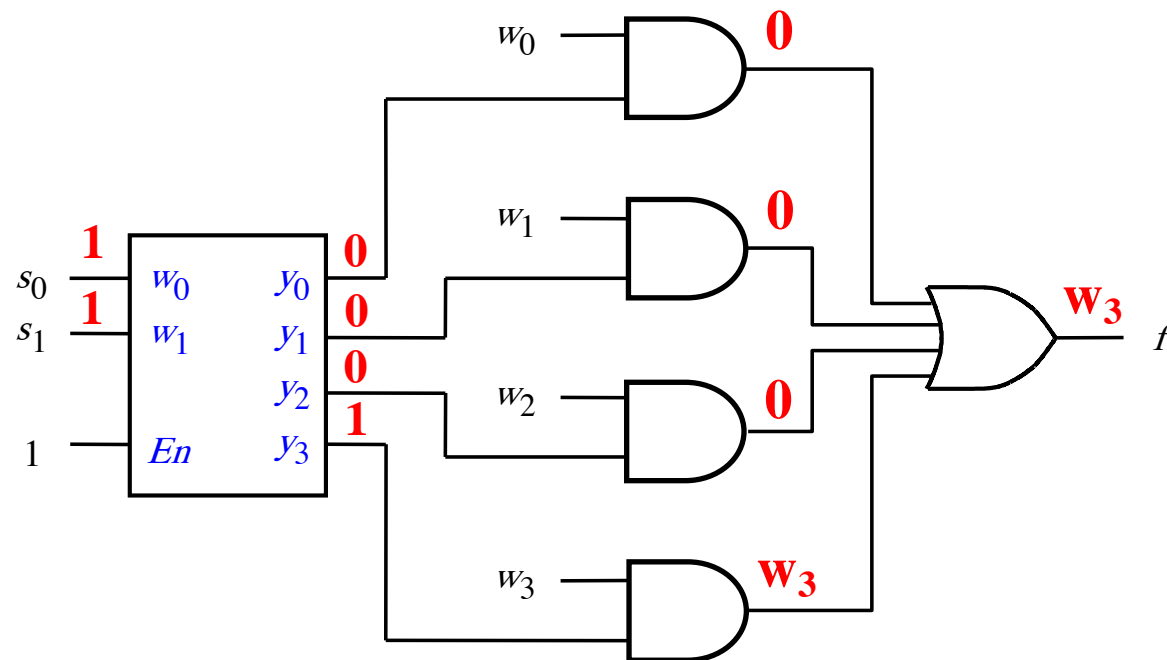
[Figure 4.17 from the textbook]

A 4-to-1 multiplexer built using a 2-to-4 decoder



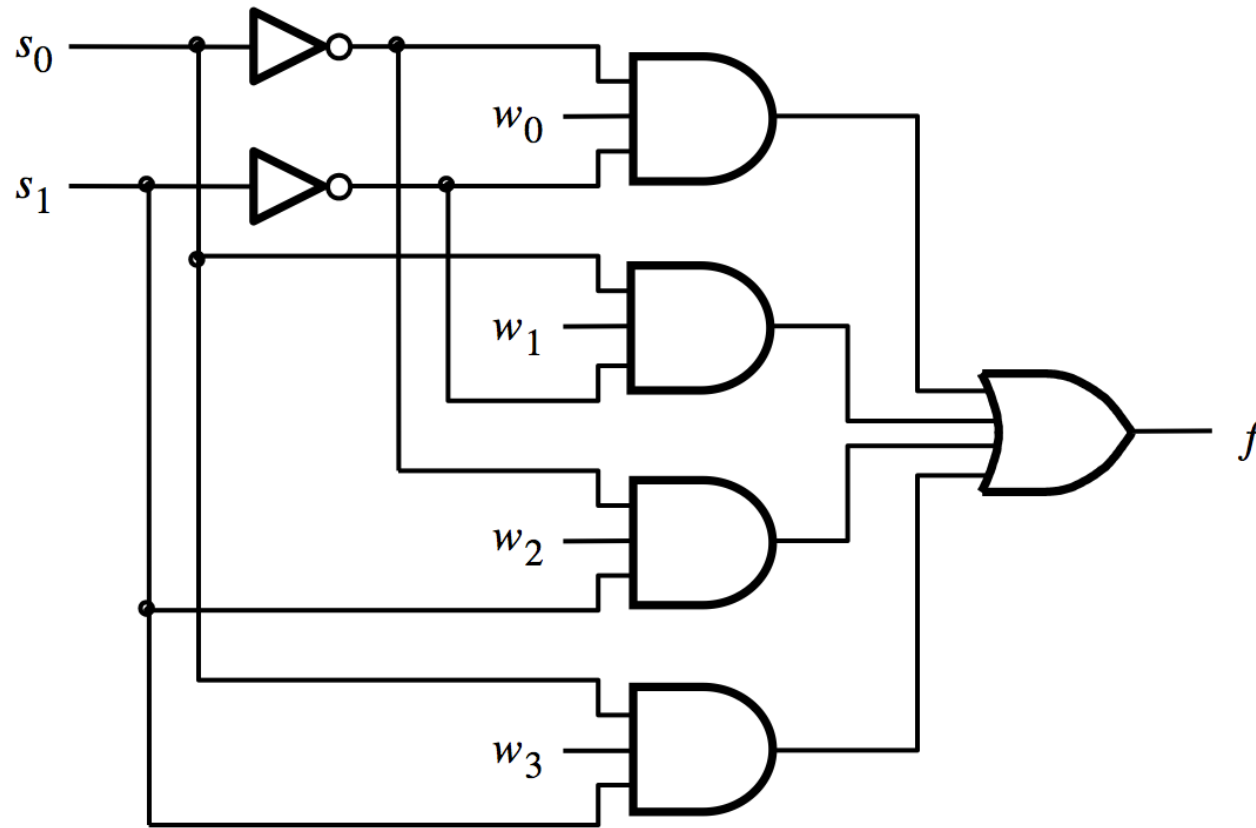
[Figure 4.17 from the textbook]

A 4-to-1 multiplexer built using a 2-to-4 decoder



[Figure 4.17 from the textbook]

4-to-1 Multiplexer (SOP circuit)



$$f = \overline{s_1} \overline{s_0} w_0 + \overline{s_1} s_0 w_1 + s_1 \overline{s_0} w_2 + s_1 s_0 w_3$$

[Figure 4.2c from the textbook]

Encoders

Encoders
(there are several types)

Binary Encoders

4-to-2 Binary Encoder (Definition)

- Has four inputs: w_3 , w_2 , w_1 , and w_0
- Has two outputs: y_1 and y_0
- Only one input is set to 1 (“one-hot” encoded). All others are set to 0.
- If $w_0=1$ then $y_1=0$ and $y_0=0$
- If $w_1=1$ then $y_1=0$ and $y_0=1$
- If $w_2=1$ then $y_1=1$ and $y_0=0$
- If $w_3=1$ then $y_1=1$ and $y_0=1$

Truth table for a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

[Figure 4.19 from the textbook]

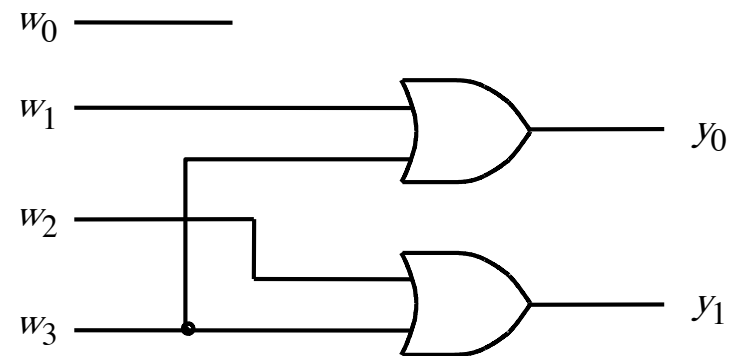
Truth table for a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

The inputs are “one-hot” encoded

Circuit for a 4-to-2 binary encoder

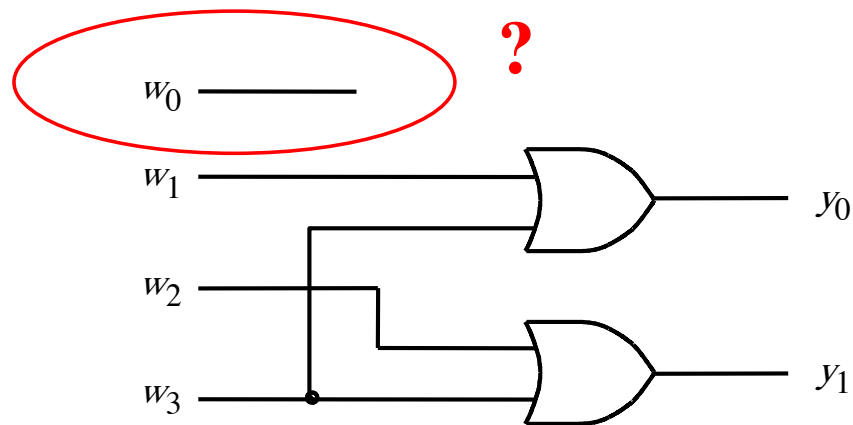
w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1



[Figure 4.19 from the textbook]

Circuit for a 4-to-2 binary encoder

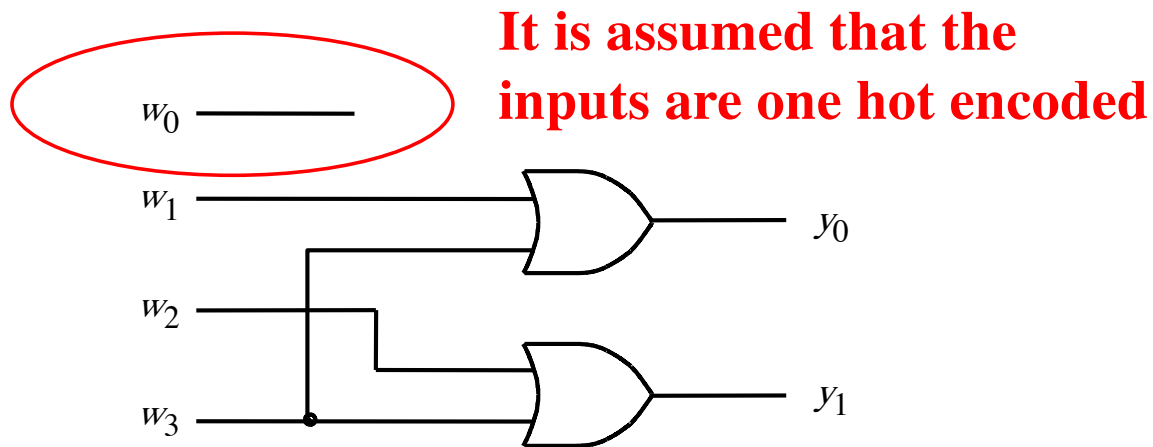
w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1



[Figure 4.19 from the textbook]

Circuit for a 4-to-2 binary encoder

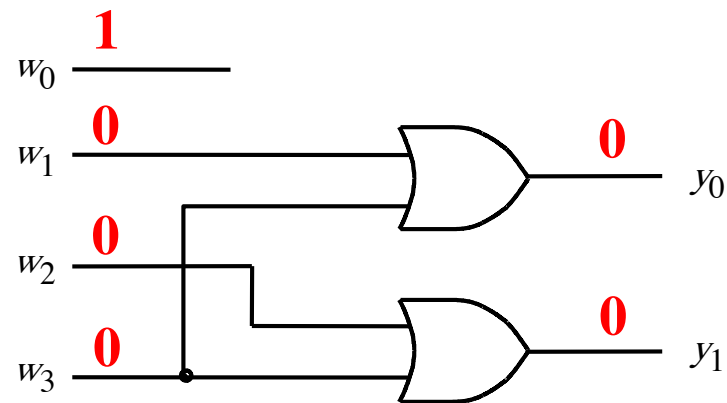
w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1



[Figure 4.19 from the textbook]

Circuit for a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

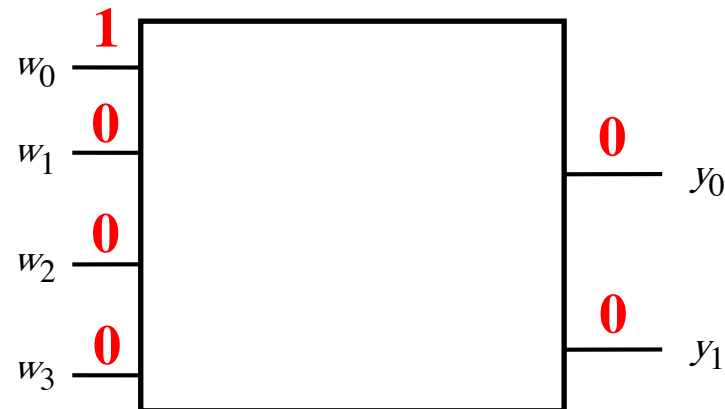


[Figure 4.19 from the textbook]

Circuit for a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

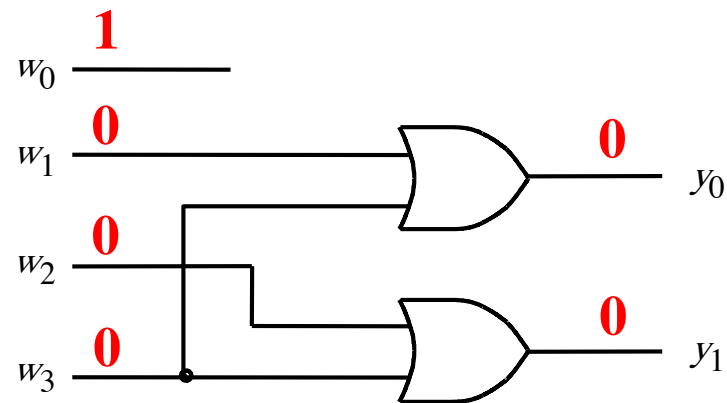
As this level of abstraction, we need that w_0 input for this to be a proper 4-to-2 binary encoder.



[Figure 4.19 from the textbook]

Circuit for a 4-to-2 binary encoder

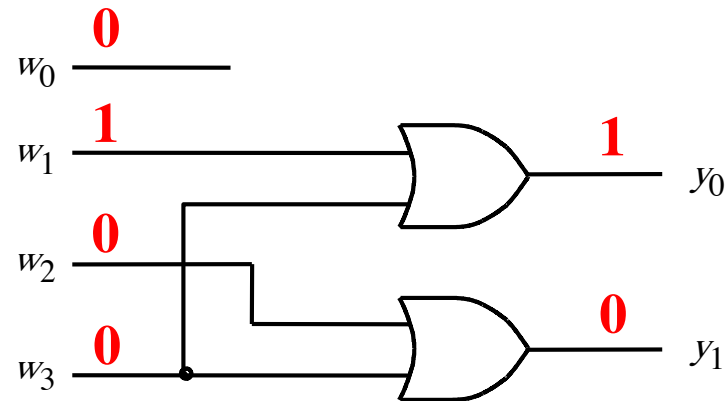
w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1



[Figure 4.19 from the textbook]

Circuit for a 4-to-2 binary encoder

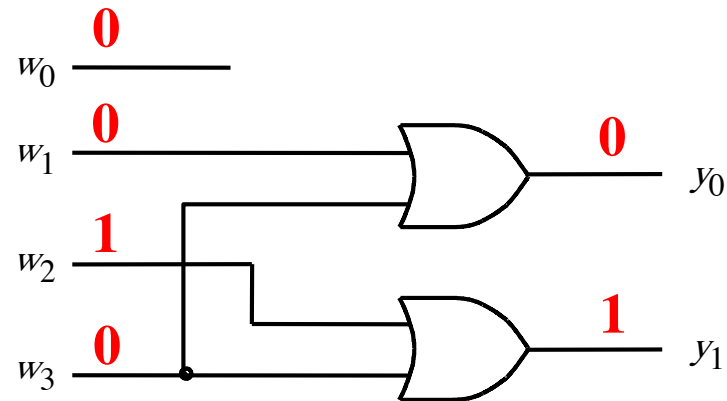
w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1



[Figure 4.19 from the textbook]

Circuit for a 4-to-2 binary encoder

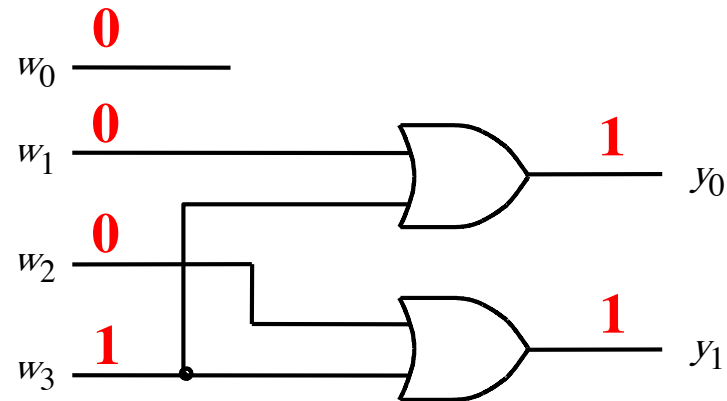
w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1



[Figure 4.19 from the textbook]

Circuit for a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1



[Figure 4.19 from the textbook]

Expressions for 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	0		
0	0	0	1	0	0
0	0	1	0	0	1
0	0	1	1		
0	1	0	0	1	0
0	1	0	1		
0	1	1	0		
0	1	1	1		
1	0	0	0	1	1
1	0	0	1		
1	0	1	0		
1	0	1	1		
1	1	0	0		
1	1	0	1		
1	1	1	0		
1	1	1	1		

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

Expressions for 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	0	d	d
0	0	0	1	0	0
0	0	1	0	0	1
0	0	1	1	d	d
0	1	0	0	1	0
0	1	0	1	d	d
0	1	1	0	d	d
0	1	1	1	d	d
1	0	0	0	1	1
1	0	0	1	d	d
1	0	1	0	d	d
1	0	1	1	d	d
1	1	0	0	d	d
1	1	0	1	d	d
1	1	1	0	d	d
1	1	1	1	d	d

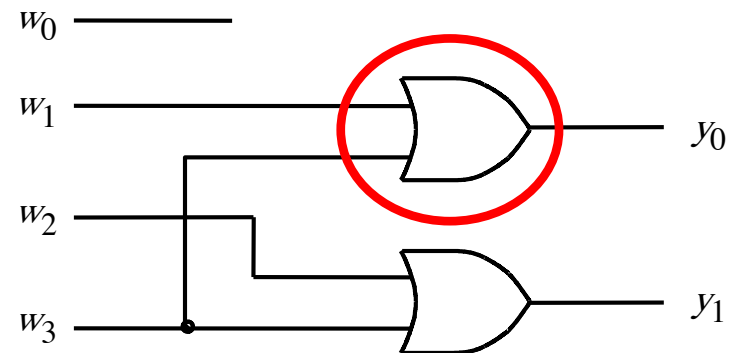
w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

Expressions for 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	0	d	d
0	0	0	1	0	0
0	0	1	0	0	1
0	0	1	1	d	d
0	1	0	0	1	0
0	1	0	1	d	d
0	1	1	0	d	d
0	1	1	1	d	d
1	0	0	0	1	1
1	0	0	1	d	d
1	0	1	0	d	d
1	0	1	1	d	d
1	1	0	0	d	d
1	1	0	1	d	d
1	1	1	0	d	d
1	1	1	1	d	d

$w_3 \ w_2 \ w_1 \ w_0$	00	01	11	10
00	d	0	d	1
01	0	d	d	d
11	d	d	d	d
10	1	d	d	d

$$y_0 = (w_1 + w_3)$$

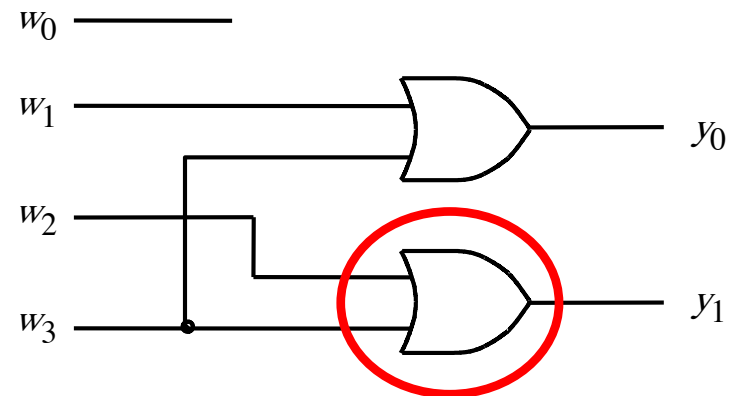


Expressions for 4-to-2 binary encoder

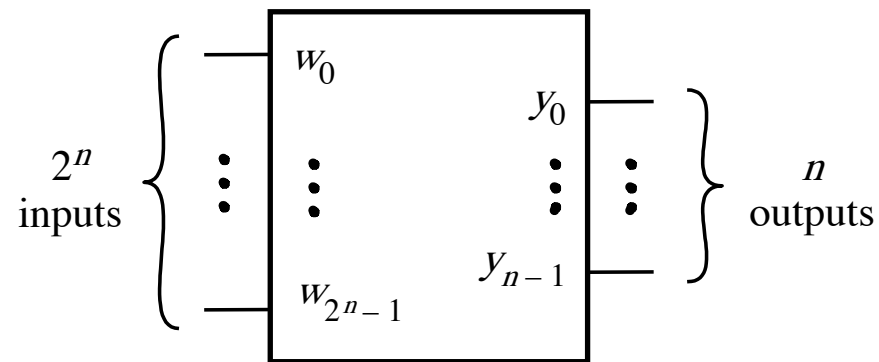
w_3	w_2	w_1	w_0	y_1	y_0
0	0	0	0	d	d
0	0	0	1	0	0
0	0	1	0	0	1
0	0	1	1	d	d
0	1	0	0	1	0
0	1	0	1	d	d
0	1	1	0	d	d
0	1	1	1	d	d
1	0	0	0	1	1
1	0	0	1	d	d
1	0	1	0	d	d
1	0	1	1	d	d
1	1	0	0	d	d
1	1	0	1	d	d
1	1	1	0	d	d
1	1	1	1	d	d

		$w_3 w_2$					
		$w_1 w_0$	00	01	11	10	
$w_1 w_0$	00	d	1	d	1		
	01	0	d	d	d		
	11	d	d	d	d		
	10	0	d	d	d		

$$y_1 = (w_3 + w_2)$$



The Most General Case: 2^n -to- n binary encoder



[Figure 4.18 from the textbook]

Priority Encoders

4-to-2 Priority Encoder (Definition)

- Has four inputs: w_3 , w_2 , w_1 , and w_0
- Has two primary outputs: y_1 and y_0
- Has one other output: z
- The inputs are NOT “one-hot” encoded.
- More than one input can be set to 1 but they have priorities associated with them: w_3 – highest priority and w_0 – lowest priority.
- $y_1=0$ and $y_0=0$ (if $w_0=1$ and $w_3=w_2=w_1=0$)
- $y_1=0$ and $y_0=1$ (if $w_1=1$ and $w_3=w_2=0$)
- $y_1=1$ and $y_0=0$ (if $w_2=1$ and $w_3=0$)
- $y_1=1$ and $y_0=1$ (if $w_3=1$)

4-to-2 Priority Encoder (Definition)

- Has four inputs: w_3 , w_2 , w_1 , and w_0
 - Has two primary outputs: y_1 and y_0
 - Has one other output: z
 - The inputs are NOT “one-hot” encoded.
 - More than one input can be set to 1 but they have priorities associated with them: w_3 – highest priority and w_0 – lowest priority.
 - $y_1=0$ and $y_0=0$ (if $w_0=1$ and $w_3=w_2=w_1=0$)
 - $y_1=0$ and $y_0=1$ (if $w_1=1$ and $w_3=w_2=0$) w_0
 - $y_1=1$ and $y_0=0$ (if $w_2=1$ and $w_3=0$) w_0 , w_1
 - $y_1=1$ and $y_0=1$ (if $w_3=1$) w_0 , w_1 , w_2
- these have lower priorities
and can be either 0 or 1.

4-to-2 Priority Encoder (Definition)

- Has four inputs: w_3 , w_2 , w_1 , and w_0
- Has two primary outputs: y_1 and y_0
- Has one other output: z
- The inputs are NOT “one-hot” encoded.
- More than one input can be set to 1 but they have priorities associated with them: w_3 – highest priority and w_0 – lowest priority.
- $y_1=0$ and $y_0=0$ (if $w_0=1$ and $w_3=w_2=w_1=0$)
- $y_1=0$ and $y_0=1$ (if $w_1=1$ and $w_3=w_2=0$)
- $y_1=1$ and $y_0=0$ (if $w_2=1$ and $w_3=0$)
- $y_1=1$ and $y_0=1$ (if $w_3=1$)
- $z = 0$ if $w_3=w_2=w_1=w_0=0$; otherwise, $z=1$.

Truth table for a 4-to-2 priority encoder (abbreviated version)

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

[Figure 4.20 from the textbook]

Truth table for a 4-to-2 priority encoder (abbreviated version)

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

[Figure 4.20 from the textbook]

Truth table for a 4-to-2 priority encoder

	w ₃	w ₂	w ₁	w ₀	y ₁	y ₀	z
0 0 0 0	0	0	0	0	d	d	0
0 0 0 1	0	0	0	1	0	0	1
0 0 1 x	0	0	1	0	0	1	1
	0	0	1	1	0	1	1
0 1 x x	0	1	0	0	1	0	1
	0	1	0	1	1	0	1
	0	1	1	0	1	0	1
	0	1	1	1	1	0	1
1 x x x	1	0	0	0	1	1	1
	1	0	0	1	1	1	1
	1	0	1	0	1	1	1
	1	0	1	1	1	1	1
	1	1	0	0	1	1	1
	1	1	0	1	1	1	1
	1	1	1	0	1	1	1
	1	1	1	1	1	1	1

Expressions for 4-to-2 priority encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	0	0	1	1
0	0	1	1	0	1	1
0	1	0	0	1	0	1
0	1	0	1	1	0	1
0	1	1	0	1	0	1
0	1	1	1	1	0	1
1	0	0	0	1	1	1
1	0	0	1	1	1	1
1	0	1	0	1	1	1
1	0	1	1	1	1	1
1	1	0	0	1	1	1
1	1	0	1	1	1	1
1	1	1	0	1	1	1
1	1	1	1	1	1	1

		$w_3 \ w_2$			
$w_1 \ w_0$		00	01	11	10
	00	d	1	1	1
	01	0	1	1	1
	11	0	1	1	1
	10	0	1	1	1

$$y_1 = w_3 + w_2$$

Expressions for 4-to-2 priority encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	0	0	1	1
0	0	1	1	0	1	1
0	1	0	0	1	0	1
0	1	0	1	1	0	1
0	1	1	0	1	0	1
0	1	1	1	1	0	1
1	0	0	0	1	1	1
1	0	0	1	1	1	1
1	0	1	0	1	1	1
1	0	1	1	1	1	1
1	1	0	0	1	1	1
1	1	0	1	1	1	1
1	1	1	0	1	1	1
1	1	1	1	1	1	1

$w_3 \backslash w_2 \backslash w_1 w_0$	00	01	11	10
00	d	0	1	1
01	0	0	1	1
11	1	0	1	1
10	1	0	1	1

$$y_0 = w_3 + w_1 \overline{w_2}$$

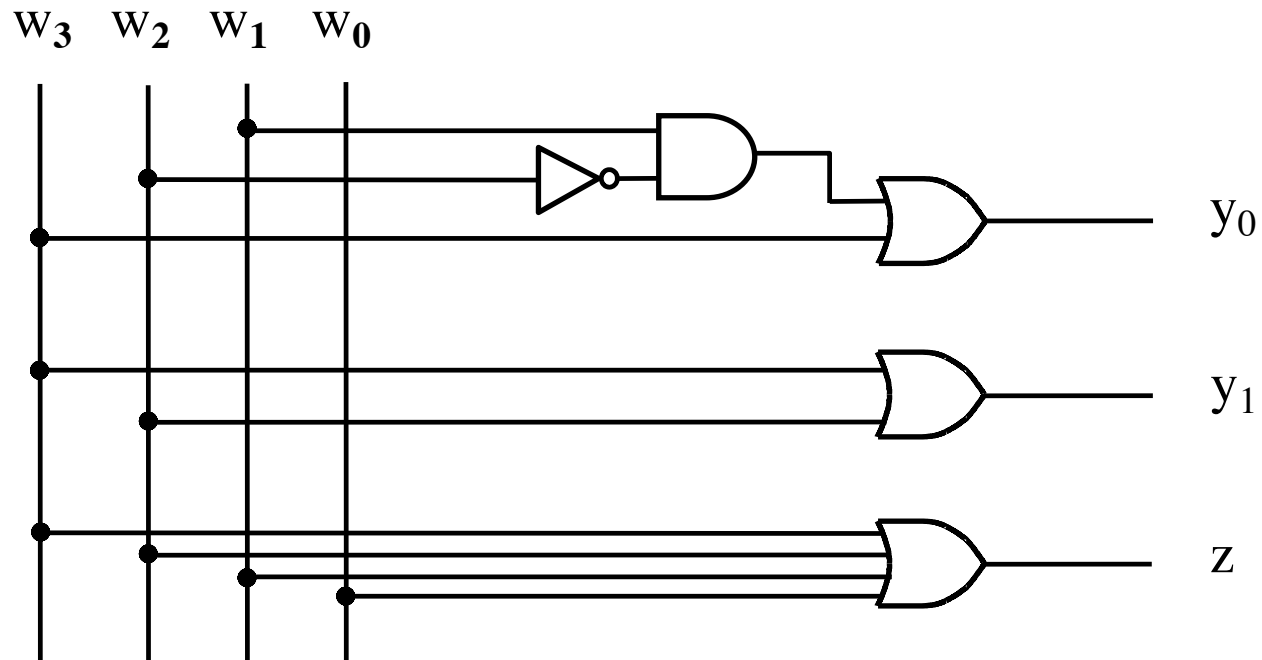
Expressions for 4-to-2 priority encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	0	0	1	1
0	0	1	1	0	1	1
0	1	0	0	1	0	1
0	1	0	1	1	0	1
0	1	1	0	1	0	1
0	1	1	1	1	0	1
1	0	0	0	1	1	1
1	0	0	1	1	1	1
1	0	1	0	1	1	1
1	0	1	1	1	1	1
1	1	0	0	1	1	1
1	1	0	1	1	1	1
1	1	1	0	1	1	1
1	1	1	1	1	1	1

		$w_3 w_2$			
$w_1 w_0$		00	01	11	10
	00	0	1	1	1
	01	1	1	1	1
	11	1	1	1	1
	10	1	1	1	1

$$Z = w_3 + w_2 + w_1 + w_0$$

Circuit for the 4-to-2 priority encoder



The textbook derives a different circuit for the 4-to-2 priority encoder using a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

The textbook derives a different circuit for the 4-to-2 priority encoder using a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

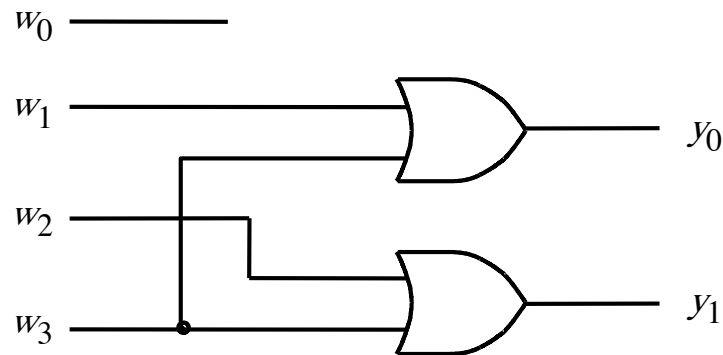
$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$



The textbook derives a different circuit for the 4-to-2 priority encoder using a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

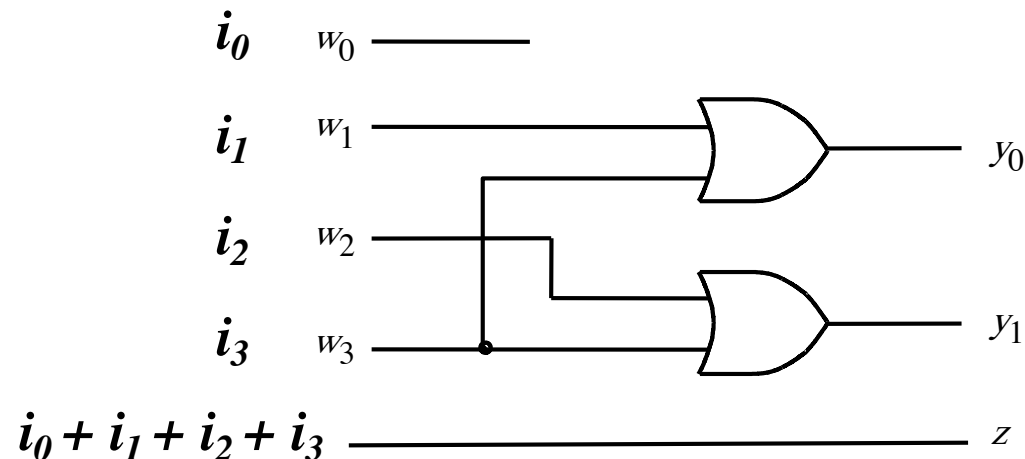
$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$



The textbook derives a different circuit for the 4-to-2 priority encoder using a 4-to-2 binary encoder

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

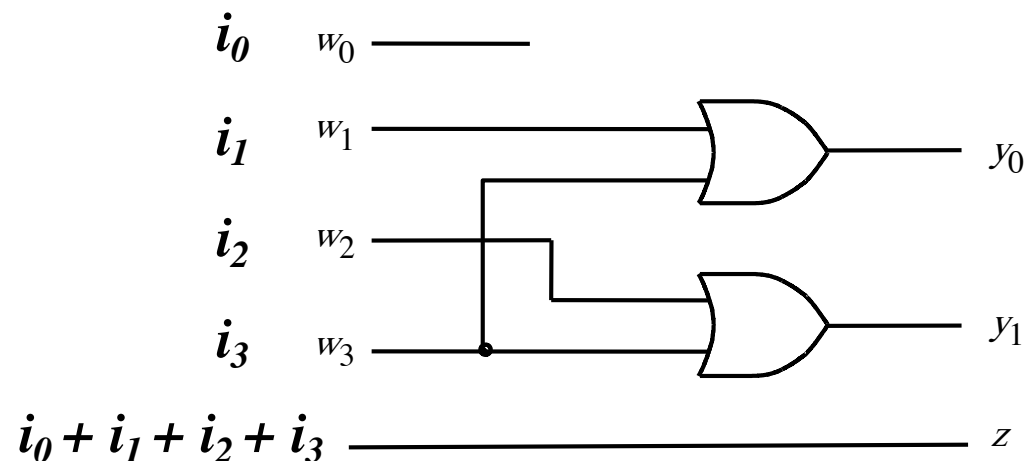
$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$



Try to prove that this is equivalent to the circuit that was derived above.

Let's prove this for z

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

		$w_3 w_2$			
		00	01	11	10
$w_1 w_0$	00	0	1	1	1
	01	1	1	1	1
	11	1	1	1	1
	10	1	1	1	1

$z = ?$

Let's prove this for z

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

		$w_3 w_2$			
w_1	w_0	00	01	11	10
		0	1	1	1
0	1	1	1	1	1
1	1	1	1	1	1
1	0	1	1	1	1

$$z = (w_0 + w_1 + w_2 + w_3)$$

Let's prove this for y_0

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

		$w_3 w_2$			
		00	01	11	10
$w_1 w_0$	00	0	0	1	1
	01	0	0	1	1
	11	1	0	1	1
	10	1	0	1	1

$y_0 = ?$

Let's prove this for y_0

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

		$w_3 w_2$			
w_1	w_0	00	01	11	10
		0	0	1	1
00		0	0	1	1
01		0	0	1	1
11		1	0	1	1
10		1	0	1	1

$$y_0 = w_3 + w_1 \overline{w_2}$$

Let's prove this for y_1

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

		$w_3 w_2$			
		00	01	11	10
$w_1 w_0$	00	0	1	1	1
	01	0	1	1	1
	11	0	1	1	1
	10	0	1	1	1

$y_1 = ?$

Let's prove this for y_1

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$

		$w_3 w_2$			
$w_1 w_0$		00	01	11	10
00	0	1	1	1	
01	0	1	1	1	
11	0	1	1	1	
10	0	1	1	1	

$$y_1 = w_3 + w_2$$

Therefore, this circuit for the
4-to-2 priority encoder is equivalent to ...

w_3	w_2	w_1	w_0	y_1	y_0	z
0	0	0	0	d	d	0
0	0	0	1	0	0	1
0	0	1	x	0	1	1
0	1	x	x	1	0	1
1	x	x	x	1	1	1

$$i_0 = \overline{w_3}\overline{w_2}\overline{w_1}w_0$$

$$i_1 = \overline{w_3}\overline{w_2}w_1$$

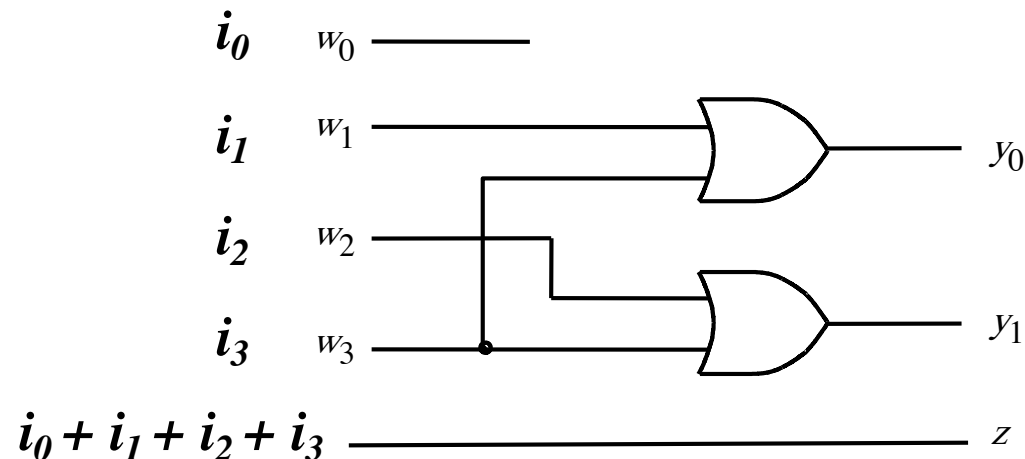
$$i_2 = \overline{w_3}w_2$$

$$i_3 = w_3$$

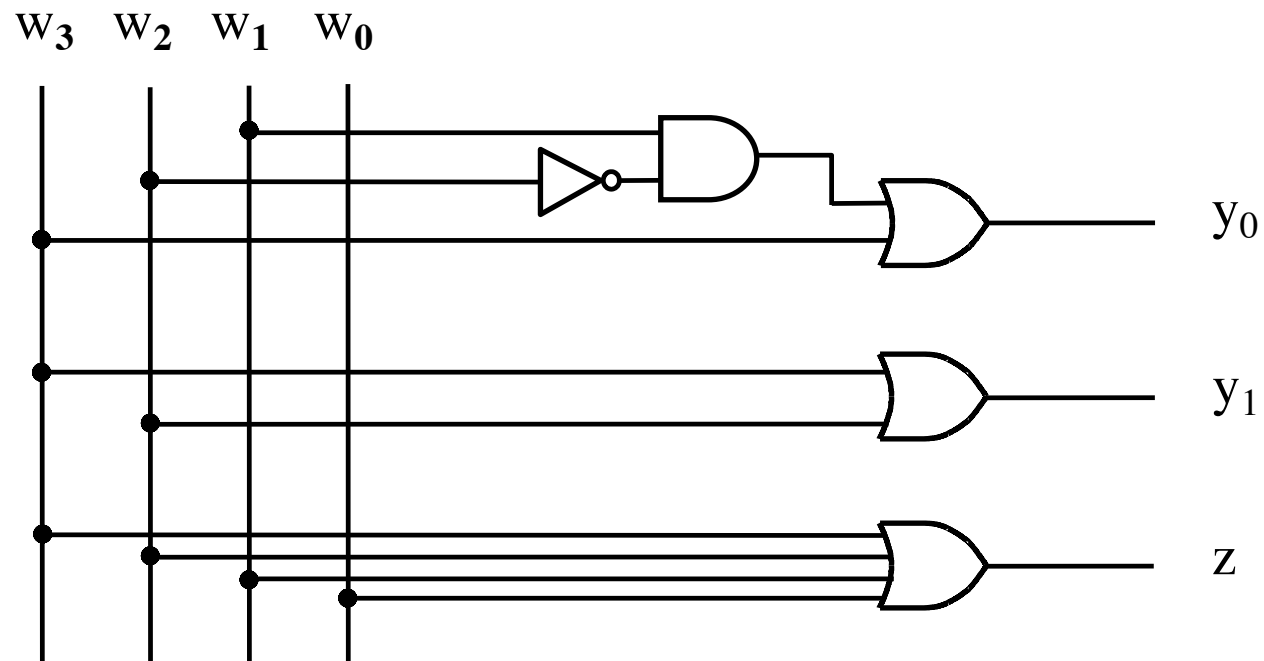
$$y_0 = i_1 + i_3$$

$$y_1 = i_2 + i_3$$

$$z = i_0 + i_1 + i_2 + i_3$$



... this circuit for the 4-to-2 priority encoder



Questions?

THE END