

CprE 2810: Digital Logic

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http://www.ece.iastate.edu/~alexs/classes/

Synchronous Sequential Circuits Basic Design Steps

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Administrative Stuff

• **Homework 10 is due on Monday Nov. 11 @ 10pm.**

Administrative Stuff

• **We are starting with Chapter 6 from the textbook**

First Design Pattern: Moore Machines

Moore Machine: A Type of Finite State Machine (FSM)

- Finite number of states (nodes).
- Discrete state transitions (edges).
- Only "in" one state at a time.
- One reset state
- Every state has an outgoing state transition for each possible input.

The next state depends on both the current state and the current input.

The output depends only on the current state.

Let's do a simulation

Inferring the States

Inferring the States

What is a State?

It is not really a memory of every past input (We might run out of space to remember it all!)

Rather, it is a characterization or snapshot of the pattern of inputs that have come before.

Moore Machine Implementation

The state diagram is just an illustration to help us describe and reason about how the FSM will behave in each of its states.

So, how do we turn it into a circuit?

Moore Machine Implementation

Note: The *W* and *Z* lines need not be wires. They can be buses.

State Storage

Any usable "memory" of the preceding input sequence is encoded in the flip-flop array.

FSM States

The Flip-Flop array stores an encoding of the current state.

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State Encoding

Each of the states in our design is identified by a distinct code.

If we use *3* flip-flops, then the FSM can have up to $2^3 = 8$ distinct states.

So, when the flip-flop array contains the code *011*, we say that the machine is in state *011*.

Synchronous Design

Every *active clock edge* causes a state transition.

Synchronous Design

We expect the input signals to be stable before the *active clock edge* occurs.

Synchronous Design

There is a whole other class of sequential circuits that are asynchronous, but we will not study them in this course.

Sequential Circuits: Key Ideas

The current output depends on something about the preceding sequence of inputs (and maybe the current output).

Using *memory elements* (i.e., flip-flops), we design the circuit to remember some *relevant* information about the prior inputs.

Moore Machine Example

We need to find both the *next state logic* and the *output logic* implied by this machine.

How to represent the States?

One way is to encode each state with a 2-bit binary number

A ~ 00 B ~ 01 C ~ 10

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How many flip-flops do we need?

Let's use two flip-flops to hold the machine's state

Let's pick D Flip-Flops.

We will call y_1 and y_2 the *present state variables*.

We will call Y_1 and Y_2 the *next state variables*.

Two zeros on the output JOINTLY represent state A.

This flip-flop output pattern represents state B.

This flip-flop output pattern represents state C.

What does this flip-flop output pattern represent?

Clock

This would be state D, but we don't have one in this example. So, this is an impossible state.

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We need to find logic expressions for *Y₁(w, y₁, y₂), Y₂(w, y₁, y₂), and <i>z*(y₁, y₂).

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Suppose we encoded our states in the same order in which they were labeled:

> $A \sim 00$ $B \sim 01$ $C \sim 10$

 \Box
[Figure 6.6 from the textbook]

$$
Q(t) = y_2y_1 \text{ and } Q(t+1) = Y_2Y_1
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[Figure 6.6 from the textbook]

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Note that the textbook draws these K-Maps differently from all previous K-maps (the most significant bit indexes the rows).

 $Q(t) = y_2y_1$ and $Q(t+1) = Y_2Y_1$

Don't care conditions simplify the combinatorial logic

[Figure 6.7 from the textbook]

Moore Type

Don't Forget to Add the Reset Line

[[] Figure 6.8 from the textbook]

$$
State = y_2 y_1
$$

State A=00

[[] Figure 6.9 from the textbook]

Summary: Designing a Moore Machine

- Obtain the circuit specification.
- Derive a state diagram.
- Derive the state table.
- Decide on a state encoding.
- Encode the state table.
- Derive the output logic and next-state logic.
- Draw the circuit diagram
- Add a reset signal.

Questions?
THE END