

CprE 281: Digital Logic

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<http://www.ece.iastate.edu/~alexs/classes/>

NAND and NOR Logic Networks

*CprE 281: Digital Logic
Iowa State University, Ames, IA
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Administrative Stuff

- **HW2 is due today, Sep 9 @ 10pm**
- **Please write clearly on the first page the following three things:**
 - **Your First and Last Name**
 - **Your Student ID Number**
 - **Your Lab Section Letter**
- **Submit on Canvas as *one* PDF file.**
- **Please orient your pages such that the text can be read without the need to rotate the page.**

Administrative Stuff

- **This week we will start with Lab2**
- **Read the lab assignment and do the prelab at home.**
- **Complete the prelab on paper before you go to the lab. Otherwise, you'll lose 20% of your grade for that lab.**

Quick Review

Minterms
(a set of basis functions)

The Four Basis Functions

x	y	f₀₀
0	0	1
0	1	0
1	0	0
1	1	0

$f_{00}(x, y)$

x	y	f₀₁
0	0	0
0	1	1
1	0	0
1	1	0

$f_{01}(x, y)$

x	y	f₁₀
0	0	0
0	1	0
1	0	1
1	1	0

$f_{10}(x, y)$

x	y	f₁₁
0	0	0
0	1	0
1	0	0
1	1	1

$f_{11}(x, y)$

The Four Basis Functions

x	y	f_{00}
0	0	1
0	1	0
1	0	0
1	1	0

$f_{00}(x, y)$

x	y	f_{01}
0	0	0
0	1	1
1	0	0
1	1	0

$f_{01}(x, y)$

x	y	f_{10}
0	0	0
0	1	0
1	0	1
1	1	0

$f_{10}(x, y)$

x	y	f_{11}
0	0	0
0	1	0
1	0	0
1	1	1

$f_{11}(x, y)$

The Four Basis Functions

x	y	$f_{00}(x, y)$	$f_{01}(x, y)$	$f_{10}(x, y)$	$f_{11}(x, y)$
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

The Four Basis Functions

x	y	$\bar{x}\bar{y}$	$\bar{x}y$	$x\bar{y}$	xy
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

Expressions for the minterms

$$m_0 = \bar{x} \bar{y}$$

$$m_1 = \bar{x} y$$

$$m_2 = x \bar{y}$$

$$m_3 = x y$$

Expressions for the minterms

$$0 \ 0 \quad m_0 = \bar{x} \bar{y}$$

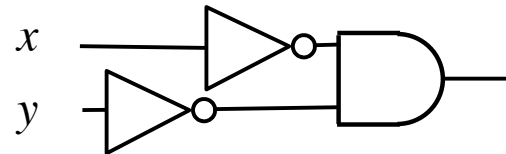
$$0 \ 1 \quad m_1 = \bar{x} y$$

$$1 \ 0 \quad m_2 = x \bar{y}$$

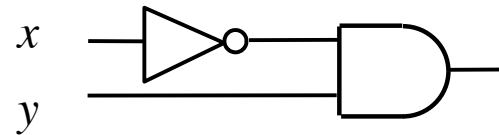
$$1 \ 1 \quad m_3 = x y$$

The bars coincide
with the 0's
in the binary expansion
of the minterm sub-index

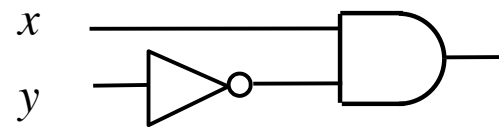
Circuits for the four basis functions



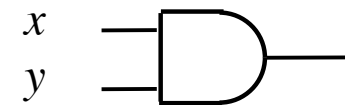
$$f_{00}(x, y) = \bar{x} \bar{y}$$



$$f_{01}(x, y) = \bar{x} y$$



$$f_{10}(x, y) = x \bar{y}$$



$$f_{11}(x, y) = x y$$

The Four Basis Functions

x	y	f ₀₀
0	0	1
0	1	0
1	0	0
1	1	0

$$f_{00}(x, y) = \bar{x} \bar{y}$$

x	y	f ₀₁
0	0	0
0	1	1
1	0	0
1	1	0

$$f_{01}(x, y) = \bar{x} y$$

x	y	f ₁₀
0	0	0
0	1	0
1	0	1
1	1	0

$$f_{10}(x, y) = x \bar{y}$$

x	y	f ₁₁
0	0	0
0	1	0
1	0	0
1	1	1

$$f_{11}(x, y) = x y$$

The Four Basis Functions (alternative names)

x	y	f ₀₀
0	0	1
0	1	0
1	0	0
1	1	0

$$f_{00}(x, y) = \bar{x} \bar{y}$$

m₀

x	y	f ₀₁
0	0	0
0	1	1
1	0	0
1	1	0

$$f_{01}(x, y) = \bar{x} y$$

m₁

x	y	f ₁₀
0	0	0
0	1	0
1	0	1
1	1	0

$$f_{10}(x, y) = x \bar{y}$$

m₂

x	y	f ₁₁
0	0	0
0	1	0
1	0	0
1	1	1

$$f_{11}(x, y) = x y$$

m₃

The Four Basis Functions (minterms)

x	y	m_0
0	0	1
0	1	0
1	0	0
1	1	0

$$f_{00}(x, y) = \bar{x} \bar{y}$$

m_0

x	y	m_1
0	0	0
0	1	1
1	0	0
1	1	0

$$f_{01}(x, y) = \bar{x} y$$

m_1

x	y	m_2
0	0	0
0	1	0
1	0	1
1	1	0

$$f_{10}(x, y) = x \bar{y}$$

m_2

x	y	m_3
0	0	0
0	1	0
1	0	0
1	1	1

$$f_{11}(x, y) = x y$$

m_3

Maxterms
(an alternative set of basis functions)

The Four Maxterms

x	y	M_0
0	0	0
0	1	1
1	0	1
1	1	1

$M_0(x, y)$

x	y	M_1
0	0	1
0	1	0
1	0	1
1	1	1

$M_1(x, y)$

x	y	M_2
0	0	1
0	1	1
1	0	0
1	1	1

$M_2(x, y)$

x	y	M_3
0	0	1
0	1	1
1	0	1
1	1	0

$M_3(x, y)$

The Four Maxterms

x	y	M_0
0	0	0
0	1	1
1	0	1
1	1	1

$M_0(x, y)$

x	y	M_1
0	0	1
0	1	0
1	0	1
1	1	1

$M_1(x, y)$

x	y	M_2
0	0	1
0	1	1
1	0	0
1	1	1

$M_2(x, y)$

x	y	M_3
0	0	1
0	1	1
1	0	1
1	1	0

$M_3(x, y)$

The Four Maxterms

x	y	$M_0(x, y)$	$M_1(x, y)$	$M_2(x, y)$	$M_3(x, y)$
0	0	0	1	1	1
0	1	1	0	1	1
1	0	1	1	0	1
1	1	1	1	1	0

The Four Maxterms

x	y		$x + y$	$x + \bar{y}$	$\bar{x} + y$	$\bar{x} + \bar{y}$
0	0		0	1	1	1
0	1		1	0	1	1
1	0		1	1	0	1
1	1		1	1	1	0

Expressions for the Maxterms

$$M_0 = x + y$$

$$M_1 = x + \bar{y}$$

$$M_2 = \bar{x} + y$$

$$M_3 = \bar{x} + \bar{y}$$

Expressions for the Maxterms

$$M_0 = x + y$$

$$M_1 = x + \bar{y}$$

$$M_2 = \bar{x} + y$$

$$M_3 = \bar{x} + \bar{y}$$

Note that these are now sums, not products.

Expressions for the Maxterms

$$0 \ 0 \quad M_0 = x + y$$

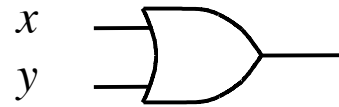
$$0 \ 1 \quad M_1 = x + \bar{y}$$

$$1 \ 0 \quad M_2 = \bar{x} + y$$

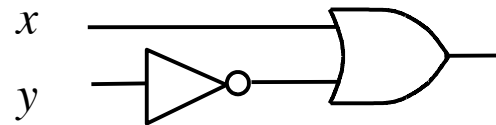
$$1 \ 1 \quad M_3 = \bar{x} + \bar{y}$$

The bars coincide
with the 1's
in the binary expansion
of the maxterm sub-index

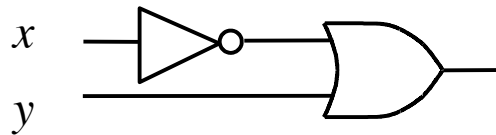
Circuits for the four Maxterms



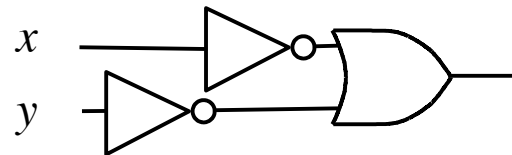
$$M_0(x, y) = x + y$$



$$M_1(x, y) = x + \bar{y}$$



$$M_2(x, y) = \bar{x} + y$$



$$M_3(x, y) = \bar{x} + \bar{y}$$

Minterms and Maxterms

Row number	x_1	x_2	Minterm	Maxterm
0	0	0	$m_0 = \bar{x}_1\bar{x}_2$	$M_0 = x_1 + x_2$
1	0	1	$m_1 = \bar{x}_1x_2$	$M_1 = x_1 + \bar{x}_2$
2	1	0	$m_2 = x_1\bar{x}_2$	$M_2 = \bar{x}_1 + x_2$
3	1	1	$m_3 = x_1x_2$	$M_3 = \bar{x}_1 + \bar{x}_2$

Minterms and Maxterms

Row number	x_1	x_2	Minterm	Maxterm
0	0	0	$m_0 = \bar{x}_1\bar{x}_2$	$M_0 = x_1 + x_2$
1	0	1	$m_1 = \bar{x}_1x_2$	$M_1 = x_1 + \bar{x}_2$
2	1	0	$m_2 = x_1\bar{x}_2$	$M_2 = \bar{x}_1 + x_2$
3	1	1	$m_3 = x_1x_2$	$M_3 = \bar{x}_1 + \bar{x}_2$

Use these for
Sum-of-Products
 Minimization
 (1's of the function)

Use these for
Product-of-Sums
 Minimization
 (0's of the function)

Sum-of-Products Form

(uses the **ones** of the function)

Sum-of-Products Form

Row number	x_1	x_2	Minterm	$f(x_1, x_2)$
0	0	0	$m_0 = \bar{x}_1\bar{x}_2$	1
1	0	1	$m_1 = \bar{x}_1x_2$	1
2	1	0	$m_2 = x_1\bar{x}_2$	0
3	1	1	$m_3 = x_1x_2$	1

Sum-of-Products Form

Row number	x_1	x_2	Minterm	$f(x_1, x_2)$
0	0	0	$m_0 = \bar{x}_1\bar{x}_2$	1
1	0	1	$m_1 = \bar{x}_1x_2$	1
2	1	0	$m_2 = x_1\bar{x}_2$	0
3	1	1	$m_3 = x_1x_2$	1

Sum-of-Products Form

Row number	x_1	x_2	Minterm	$f(x_1, x_2)$
0	0	0	$m_0 = \bar{x}_1\bar{x}_2$	1
1	0	1	$m_1 = \bar{x}_1x_2$	1
2	1	0	$m_2 = x_1\bar{x}_2$	0
3	1	1	$m_3 = x_1x_2$	1

$$\begin{aligned}f &= m_0 \cdot 1 + m_1 \cdot 1 + m_2 \cdot 0 + m_3 \cdot 1 \\&= m_0 + m_1 + m_3 \\&= \bar{x}_1\bar{x}_2 + \bar{x}_1x_2 + x_1x_2\end{aligned}$$

Product-of-Sums Form

(uses the **zeros** of the function)

Product-of-Sums Form

(for this logic function)

Row number	x_1	x_2	Maxterm	$f(x_1, x_2)$
0	0	0	$M_0 = x_1 + x_2$	0
1	0	1	$M_1 = x_1 + \bar{x}_2$	1
2	1	0	$M_2 = \bar{x}_1 + x_2$	0
3	1	1	$M_3 = \bar{x}_1 + \bar{x}_2$	1

Product-of-Sums Form (for this logic function)

Row number	x_1	x_2	Maxterm	$f(x_1, x_2)$
0	0	0	$M_0 = x_1 + x_2$	0
1	0	1	$M_1 = x_1 + \bar{x}_2$	1
2	1	0	$M_2 = \bar{x}_1 + x_2$	0
3	1	1	$M_3 = \bar{x}_1 + \bar{x}_2$	1

Product-of-Sums Form (for this logic function)

Row number	x_1	x_2	Maxterm	$f(x_1, x_2)$
0	0	0	$M_0 = x_1 + x_2$	0
1	0	1	$M_1 = x_1 + \bar{x}_2$	1
2	1	0	$M_2 = \bar{x}_1 + x_2$	0
3	1	1	$M_3 = \bar{x}_1 + \bar{x}_2$	1

$$f(x_1, x_2) = M_0 \cdot M_2 = (x_1 + x_2) \cdot (\bar{x}_1 + x_2)$$

Shorthand Notation

- **Sum-of-Products (SOP)**

$$f(x_1, x_2, x_3) = \sum (m_1, m_4, m_5, m_6)$$

or

$$f(x_1, x_2, x_3) = \sum m(1, 4, 5, 6)$$

- **Product-of-Sums (POS)**

$$f(x_1, x_2, x_3) = \Pi(M_0, M_2, M_3, M_7)$$

or

$$f(x_1, x_2, x_3) = \Pi M(0, 2, 3, 7)$$

Shorthand Notation for SOP

Row number	x_1	x_2	x_3	$f(x_1, x_2, x_3)$
0	0	0	0	0
1	0	0	1	1
2	0	1	0	0
3	0	1	1	0
4	1	0	0	1
5	1	0	1	1
6	1	1	0	1
7	1	1	1	0

$$f(x_1, x_2, x_3) = \sum (m_1, m_4, m_5, m_6)$$

or

$$f(x_1, x_2, x_3) = \sum m(1, 4, 5, 6)$$

Shorthand Notation

Row number	x_1	x_2	x_3	Minterm	Maxterm
0	0	0	0	$m_0 = \bar{x}_1\bar{x}_2\bar{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \bar{x}_1\bar{x}_2x_3$	$M_1 = x_1 + x_2 + \bar{x}_3$
2	0	1	0	$m_2 = \bar{x}_1x_2\bar{x}_3$	$M_2 = x_1 + \bar{x}_2 + x_3$
3	0	1	1	$m_3 = \bar{x}_1x_2x_3$	$M_3 = x_1 + \bar{x}_2 + \bar{x}_3$
4	1	0	0	$m_4 = x_1\bar{x}_2\bar{x}_3$	$M_4 = \bar{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1\bar{x}_2x_3$	$M_5 = \bar{x}_1 + x_2 + \bar{x}_3$
6	1	1	0	$m_6 = x_1x_2\bar{x}_3$	$M_6 = \bar{x}_1 + \bar{x}_2 + x_3$
7	1	1	1	$m_7 = x_1x_2x_3$	$M_7 = \bar{x}_1 + \bar{x}_2 + \bar{x}_3$

$$f(x_1, x_2, x_3) = \sum (m_1, m_4, m_5, m_6)$$

$$f(x_1, x_2, x_3) = \sum m(1, 4, 5, 6)$$

Shorthand Notation for POS

Row number	x_1	x_2	x_3	$f(x_1, x_2, x_3)$
0	0	0	0	0
1	0	0	1	1
2	0	1	0	0
3	0	1	1	0
4	1	0	0	1
5	1	0	1	1
6	1	1	0	1
7	1	1	1	0

$$f(x_1, x_2, x_3) = \Pi(M_0, M_2, M_3, M_7)$$

or

$$f(x_1, x_2, x_3) = \Pi M(0, 2, 3, 7)$$

Shorthand Notation

Row number	x_1	x_2	x_3	Minterm	Maxterm
0	0	0	0	$m_0 = \bar{x}_1\bar{x}_2\bar{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \bar{x}_1\bar{x}_2x_3$	$M_1 = x_1 + x_2 + \bar{x}_3$
2	0	1	0	$m_2 = \bar{x}_1x_2\bar{x}_3$	$M_2 = x_1 + \bar{x}_2 + x_3$
3	0	1	1	$m_3 = \bar{x}_1x_2x_3$	$M_3 = x_1 + \bar{x}_2 + \bar{x}_3$
4	1	0	0	$m_4 = x_1\bar{x}_2\bar{x}_3$	$M_4 = \bar{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1\bar{x}_2x_3$	$M_5 = \bar{x}_1 + x_2 + \bar{x}_3$
6	1	1	0	$m_6 = x_1x_2\bar{x}_3$	$M_6 = \bar{x}_1 + \bar{x}_2 + x_3$
7	1	1	1	$m_7 = x_1x_2x_3$	$M_7 = \bar{x}_1 + \bar{x}_2 + \bar{x}_3$

$$f(x_1, x_2, x_3) = \Pi(M_0, M_2, M_3, M_7)$$

$$f(x_1, x_2, x_3) = \Pi M(0, 2, 3, 7)$$

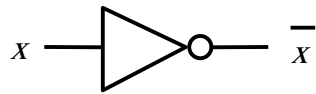
Shorthand Notation

Row number	x_1	x_2	x_3	Minterm	Maxterm
0	0	0	0	$m_0 = \bar{x}_1\bar{x}_2\bar{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \bar{x}_1\bar{x}_2x_3$	$M_1 = x_1 + x_2 + \bar{x}_3$
2	0	1	0	$m_2 = \bar{x}_1x_2\bar{x}_3$	$M_2 = x_1 + \bar{x}_2 + x_3$
3	0	1	1	$m_3 = \bar{x}_1x_2x_3$	$M_3 = x_1 + \bar{x}_2 + \bar{x}_3$
4	1	0	0	$m_4 = x_1\bar{x}_2\bar{x}_3$	$M_4 = \bar{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1\bar{x}_2x_3$	$M_5 = \bar{x}_1 + x_2 + \bar{x}_3$
6	1	1	0	$m_6 = x_1x_2\bar{x}_3$	$M_6 = \bar{x}_1 + \bar{x}_2 + x_3$
7	1	1	1	$m_7 = x_1x_2x_3$	$M_7 = \bar{x}_1 + \bar{x}_2 + \bar{x}_3$

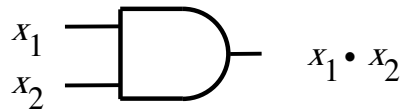
Notice that the red and the green are nicely separated and that they cover all possible rows (no gaps).

Two New Logic Gates

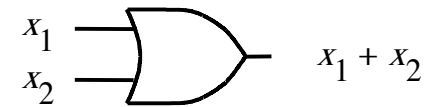
The Three Basic Logic Gates



NOT gate

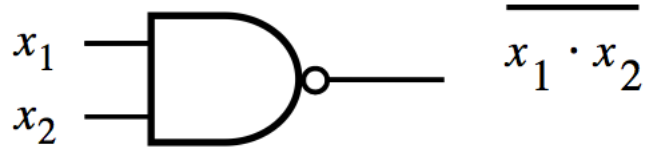


AND gate



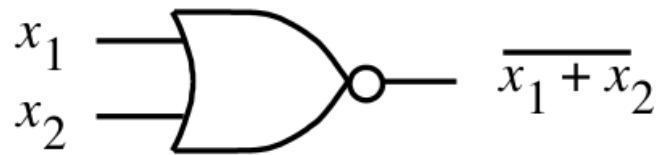
OR gate

NAND Gate



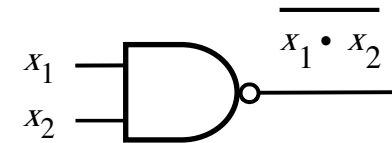
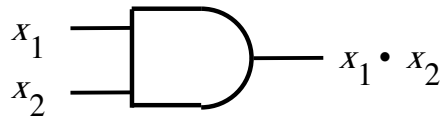
x_1	x_2	f
0	0	1
0	1	1
1	0	1
1	1	0

NOR Gate



x_1	x_2	f
0	0	1
0	1	0
1	0	0
1	1	0

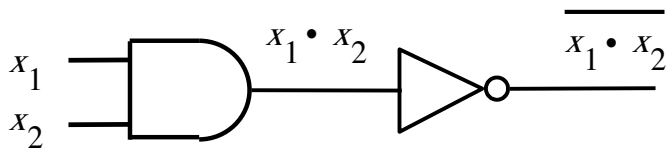
AND vs NAND



x_1	x_2	f
0	0	0
0	1	0
1	0	0
1	1	1

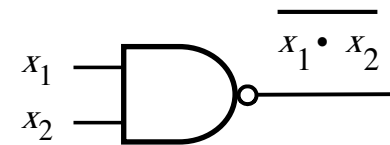
x_1	x_2	f
0	0	1
0	1	1
1	0	1
1	1	0

AND followed by NOT = NAND



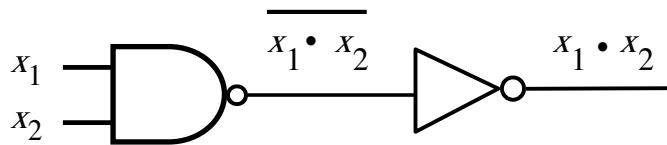
x_1	x_2	f
0	0	0
0	1	0
1	0	0
1	1	1

f
1
1
1
0



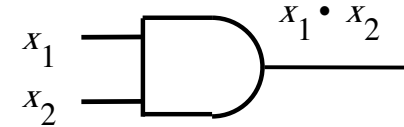
x_1	x_2	f
0	0	1
0	1	1
1	0	1
1	1	0

NAND followed by NOT = AND



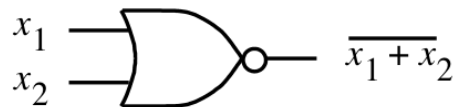
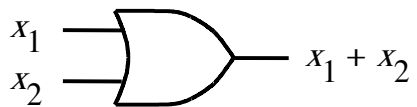
x_1	x_2	f
0	0	1
0	1	1
1	0	1
1	1	0

f
0
0
0
1



x_1	x_2	f
0	0	0
0	1	0
1	0	0
1	1	1

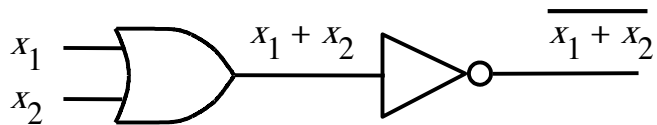
OR vs NOR



x_1	x_2	f
0	0	0
0	1	1
1	0	1
1	1	1

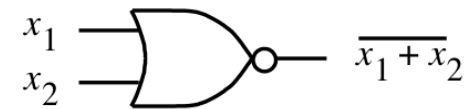
x_1	x_2	f
0	0	1
0	1	0
1	0	0
1	1	0

OR followed by NOT = NOR



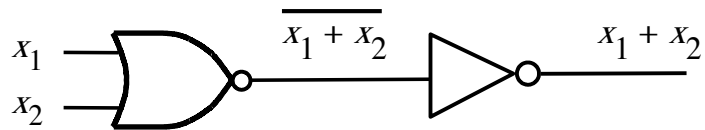
x_1	x_2	f
0	0	0
0	1	1
1	0	1
1	1	1

f
1
0
0
0



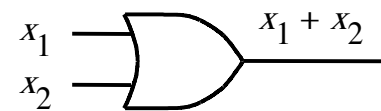
x_1	x_2	f
0	0	1
0	1	0
1	0	0
1	1	0

NOR followed by NOT = OR



x_1	x_2	f
0	0	1
0	1	0
1	0	0
1	1	0

f
0
1
1
1



x_1	x_2	f
0	0	0
0	1	1
1	0	1
1	1	1

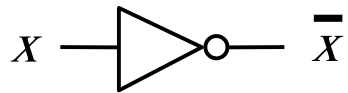
Why do we need two more gates?

Why do we need two more gates?

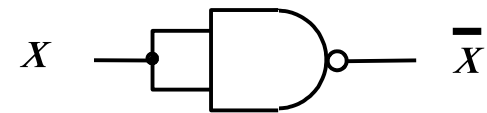
They can be implemented with fewer transistors.

**They are simpler to implement,
but are they also useful?**

Building a NOT Gate with NAND

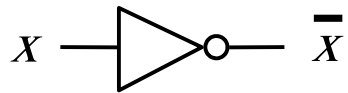


x	\bar{x}
0	1
1	0

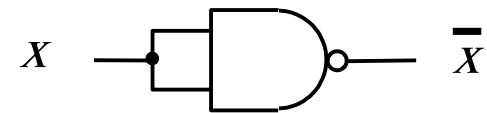


x	x	f
0	0	1
0	1	1
1	0	1
1	1	0

Building a NOT Gate with NAND



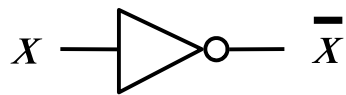
x	\bar{x}
0	1
1	0



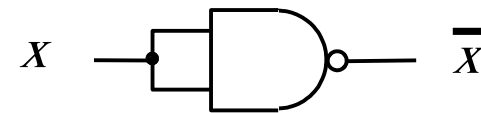
x	x	f
0	0	1
1	1	0

impossible combinations

Building a NOT Gate with NAND



x	\bar{x}
0	1
1	0

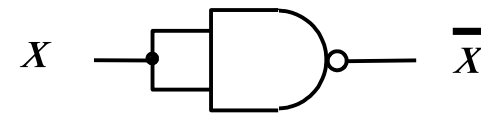
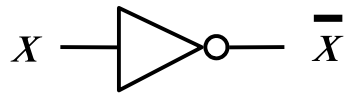


x	x	f
0	0	1
1	1	0

impossible combinations

Thus, the two truth tables are equal!

Building a NOT Gate with NAND



$$\overline{X \cdot X} = \bar{X}$$

Another way to think about this.

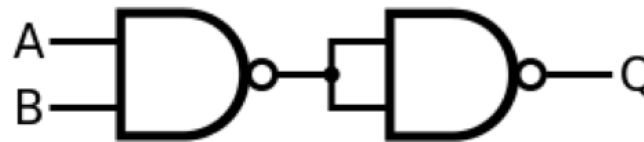
Building an AND gate with NAND gates

Desired AND Gate



$$Q = A \text{ AND } B$$

NAND Construction



$$= \text{NOT}(\text{NOT}(A \text{ AND } B))$$

Truth Table

Input A	Input B	Output Q
0	0	0
0	1	0
1	0	0
1	1	1

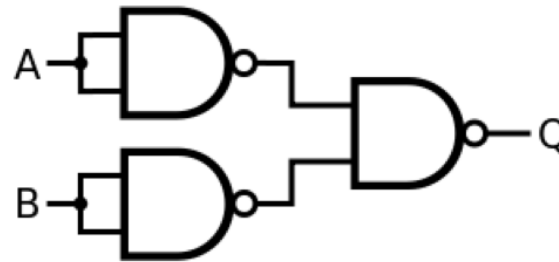
Building an OR gate with NAND gates

Desired OR Gate



$$Q = A \text{ OR } B$$

NAND Construction



$$= \text{NOT} [\text{NOT} (A \text{ AND } A) \text{ AND } \text{NOT} (B \text{ AND } B)]$$

Truth Table

Input A	Input B	Output Q
0	0	0
0	1	1
1	0	1
1	1	1

Implications

Implications

**Any Boolean function can be implemented
with only NAND gates!**

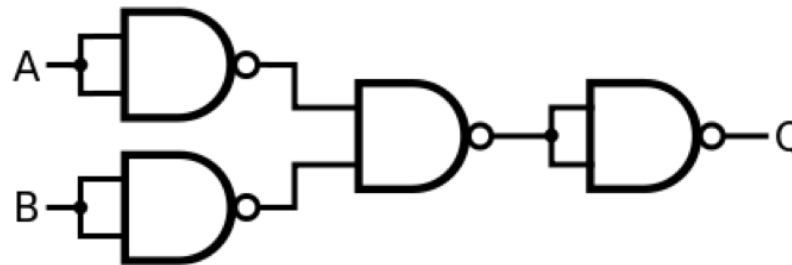
NOR gate with NAND gates

Desired NOR Gate



$$Q = \text{NOT}(A \text{ OR } B)$$

NAND Construction



$$= \text{NOT}\{ \text{NOT}[\text{NOT}(A \text{ AND } A) \text{ AND } \text{NOT}(B \text{ AND } B)] \}$$

Truth Table

Input A	Input B	Output Q
0	0	1
0	1	0
1	0	0
1	1	0

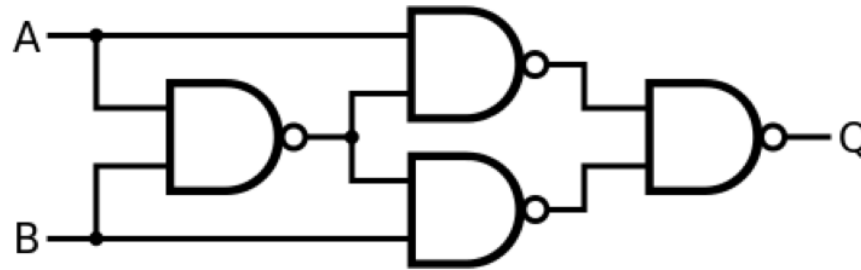
XOR gate with NAND gates

Desired XOR Gate



$$Q = A \text{ XOR } B$$

NAND Construction



$$= \text{NOT} [\text{NOT} \{ A \text{ AND } \text{NOT} (A \text{ AND } B) \} \text{ AND } \text{NOT} \{ B \text{ AND } \text{NOT} (A \text{ AND } B) \}]$$

Truth Table

Input A	Input B	Output Q
0	0	0
0	1	1
1	0	1
1	1	0

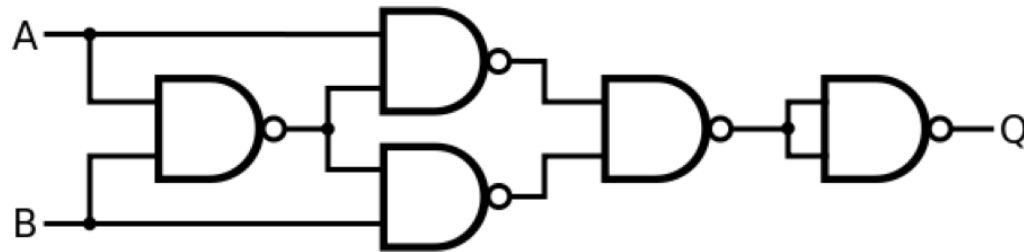
XNOR gate with NAND gates

Desired XNOR Gate



$$Q = \text{NOT}(A \text{ XOR } B)$$

NAND Construction

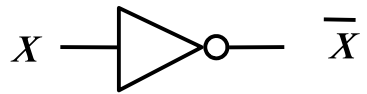


$$= \text{NOT}[\text{NOT}[\text{NOT}\{A \text{ AND } \text{NOT}(A \text{ AND } B)\} \text{ AND } \text{NOT}\{B \text{ AND } \text{NOT}(A \text{ AND } B)\}]]$$

Truth Table

Input A	Input B	Output Q
0	0	1
0	1	0
1	0	0
1	1	1

Building a NOT Gate with NOR

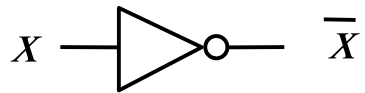


x	\bar{x}
0	1
1	0



x	x	f
0	0	1
0	1	0
1	0	0
1	1	0

Building a NOT Gate with NOR



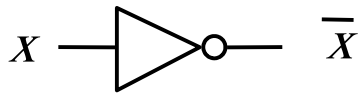
x	\bar{x}
0	1
1	0



x	x	f
0	0	1
[Redacted]		
1	1	0

impossible combinations

Building a NOT Gate with NOR



x	\bar{x}
0	1
1	0

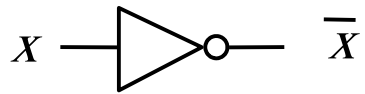


x	x	f
0	0	1
[Redacted]		
1	1	0

impossible combinations

Thus, the two truth tables are equal!

Building a NOT Gate with NOR



$$\overline{X + X} = \bar{X}$$

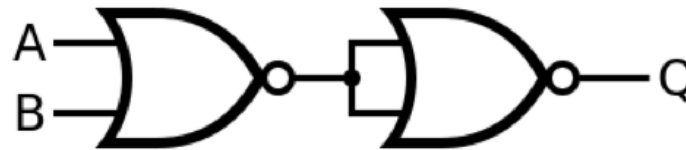
Another way to think about this.

Building an OR gate with NOR gates

Desired Gate



NOR Construction

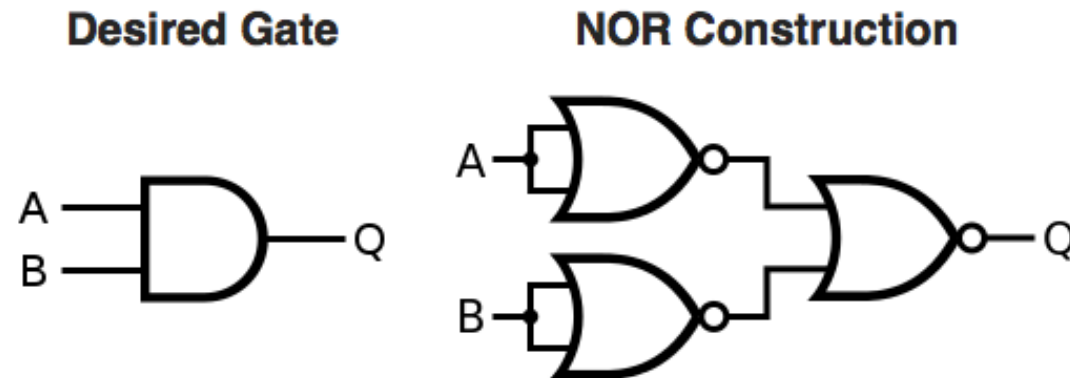


Truth Table

Input A	Input B	Output Q
0	0	0
0	1	1
1	0	1
1	1	1

Let's build an AND gate with NOR gates

Let's build an AND gate with NOR gates



Truth Table

Input A	Input B	Output Q
0	0	0
0	1	0
1	0	0
1	1	1

Implications

Implications

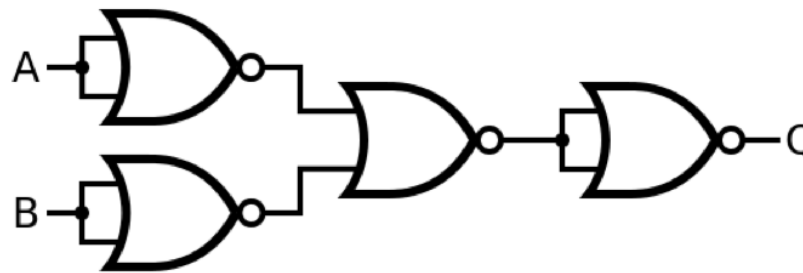
**Any Boolean function can be implemented
with only NOR gates!**

NAND gate with NOR gates

Desired Gate



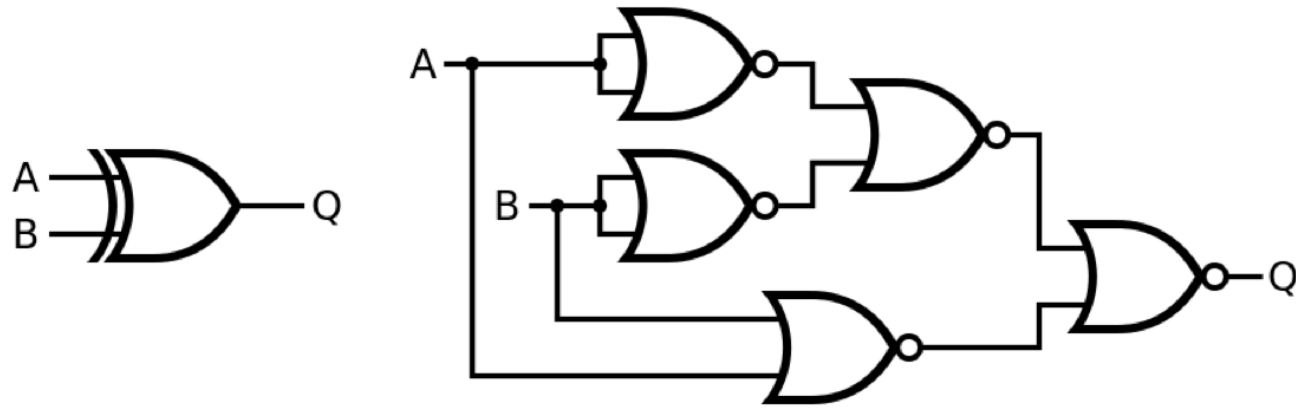
NOR Construction



Truth Table

Input A	Input B	Output Q
0	0	1
0	1	1
1	0	1
1	1	0

XOR gate with NOR gates



Truth Table

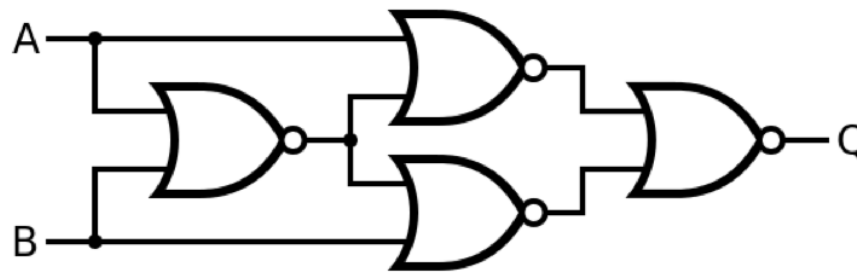
Input A	Input B	Output Q
0	0	0
0	1	1
1	0	1
1	1	0

XNOR gate with NOR gates

Desired XNOR Gate



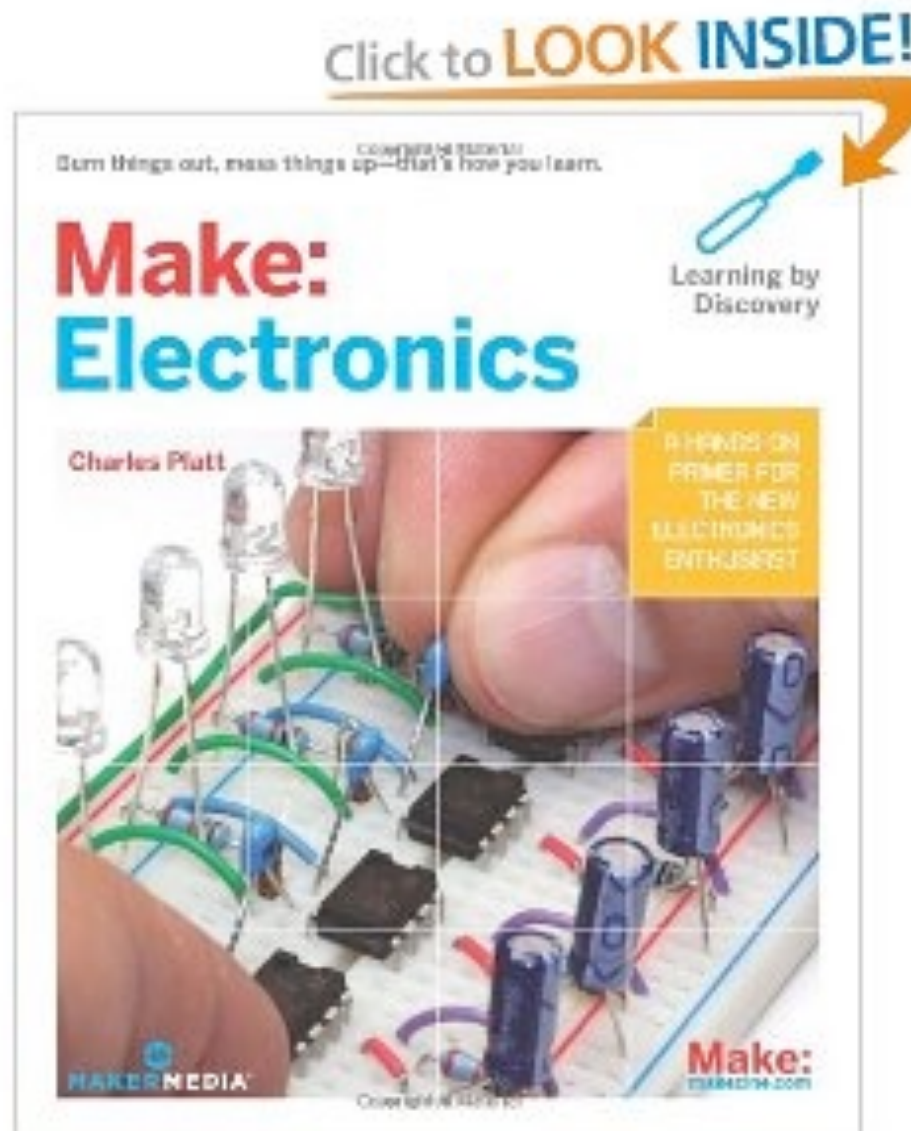
NOR Construction

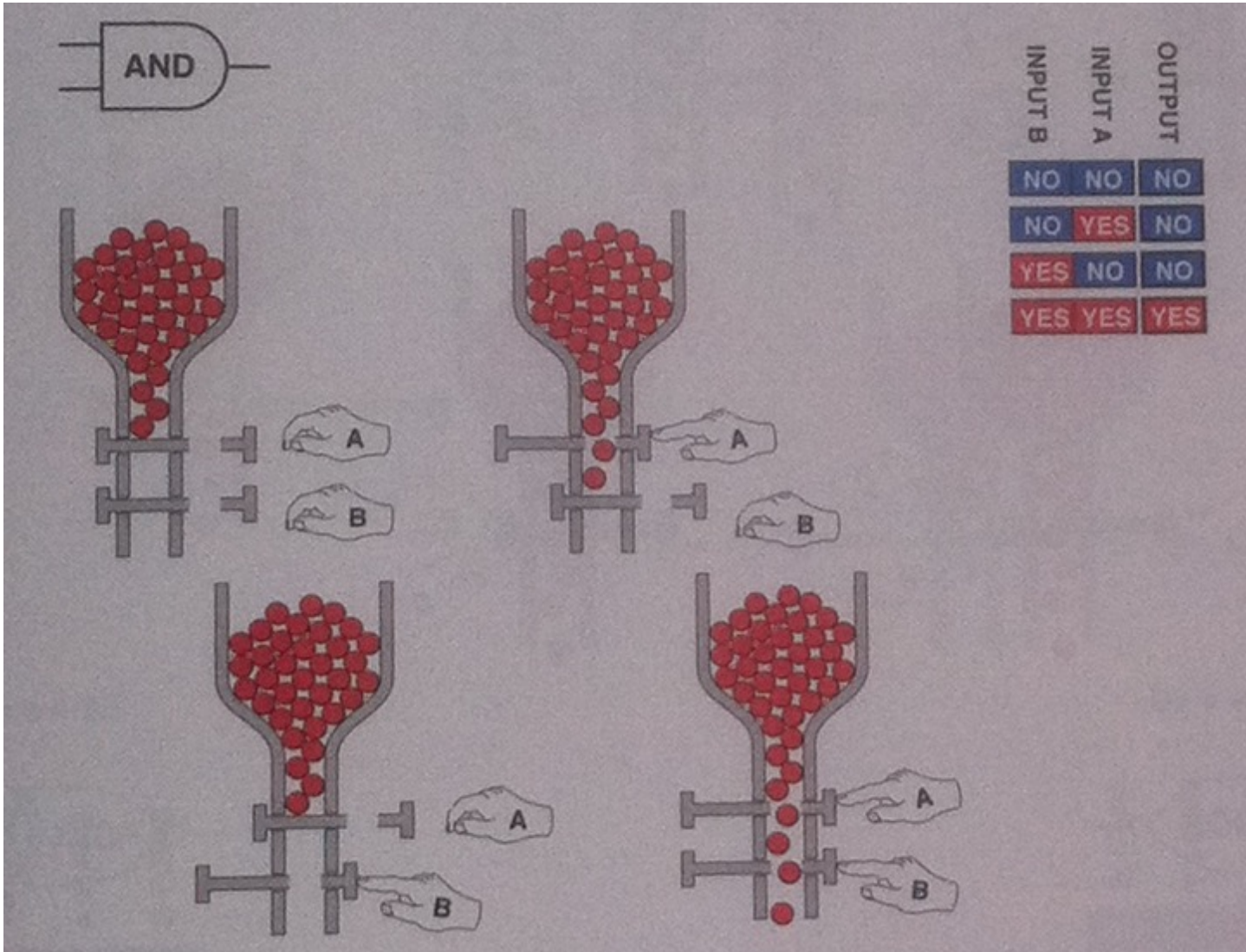


Truth Table

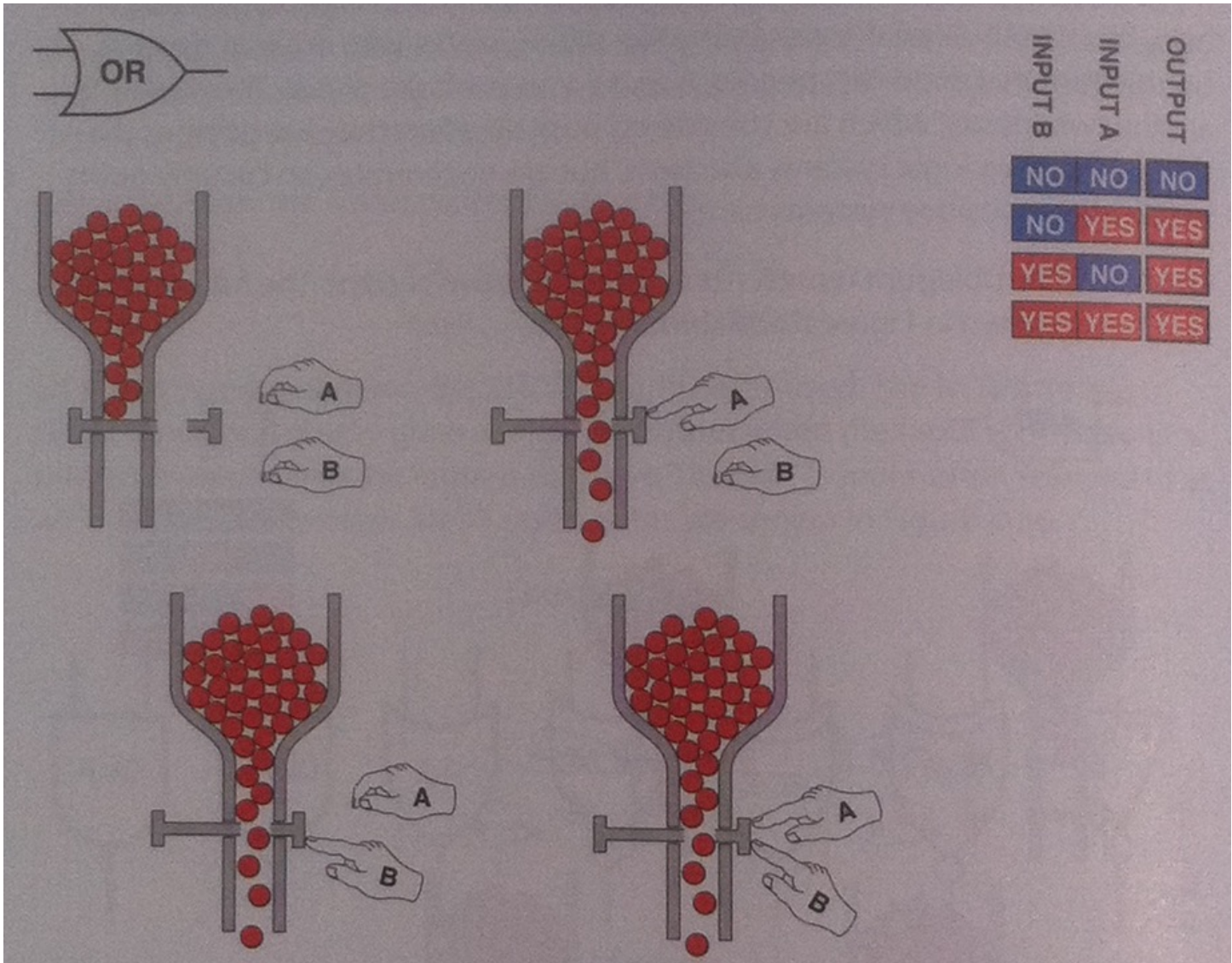
Input A	Input B	Output Q
0	0	1
0	1	0
1	0	0
1	1	1

The following examples came from this book

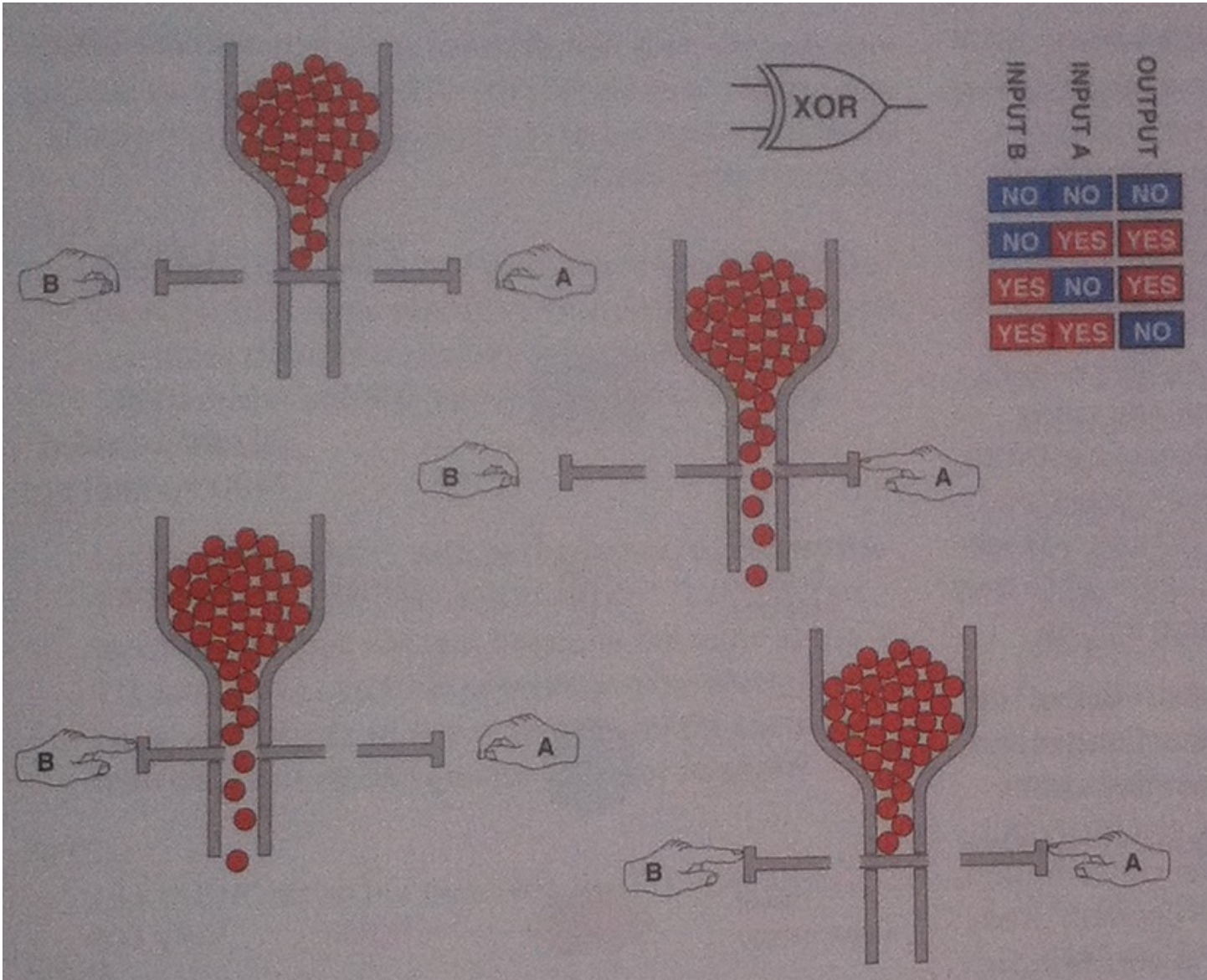




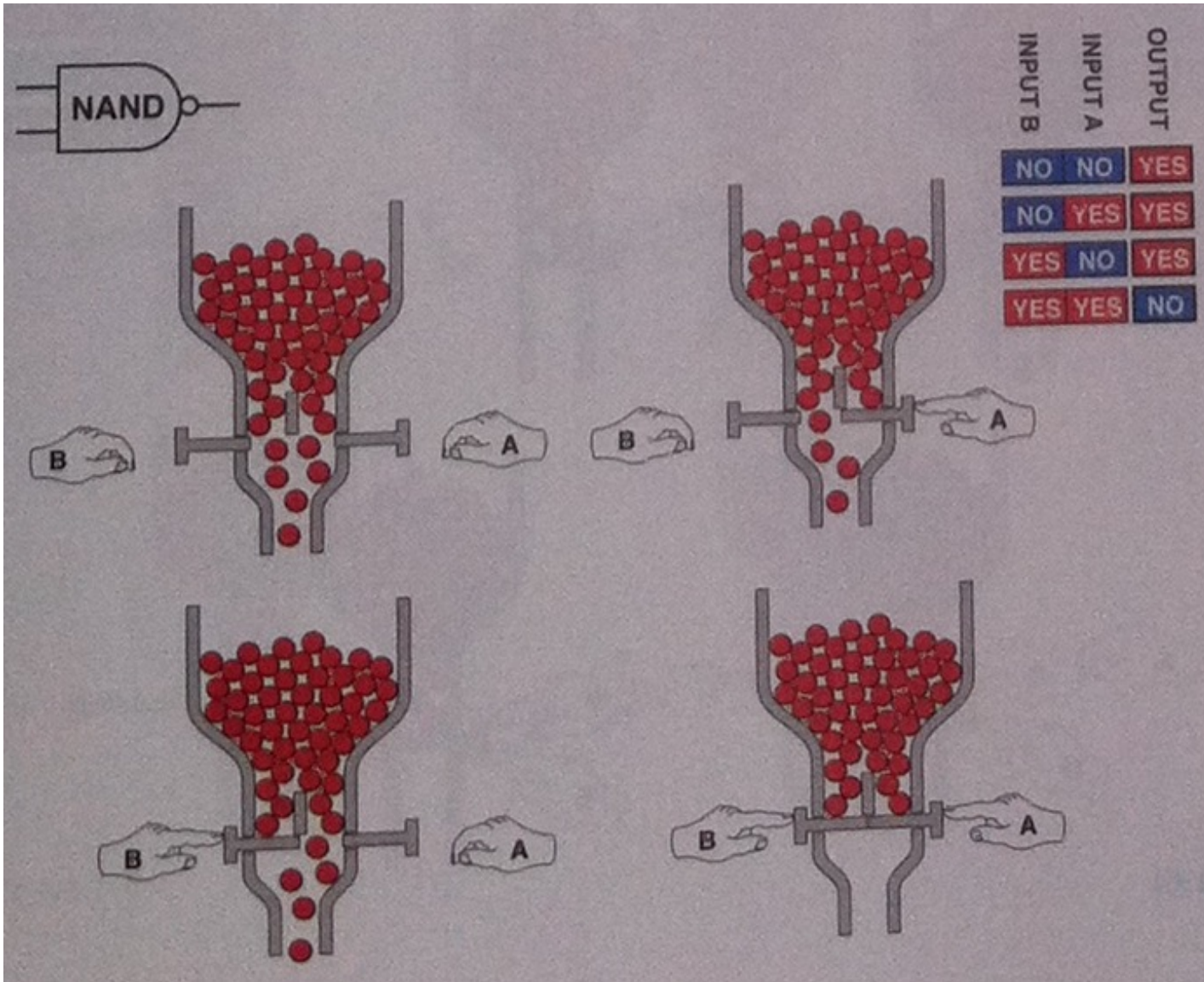
[Platt 2009]



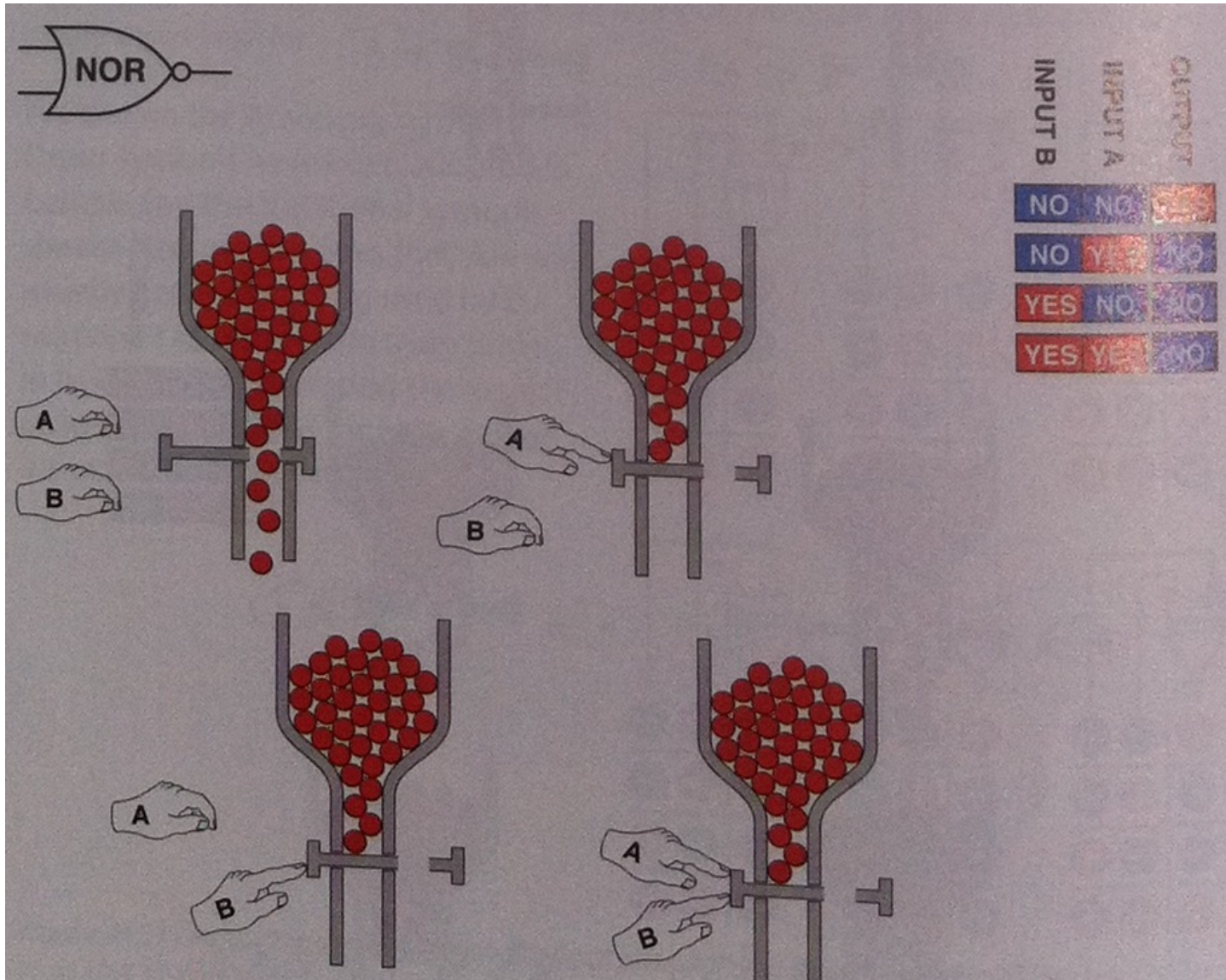
[Platt 2009]



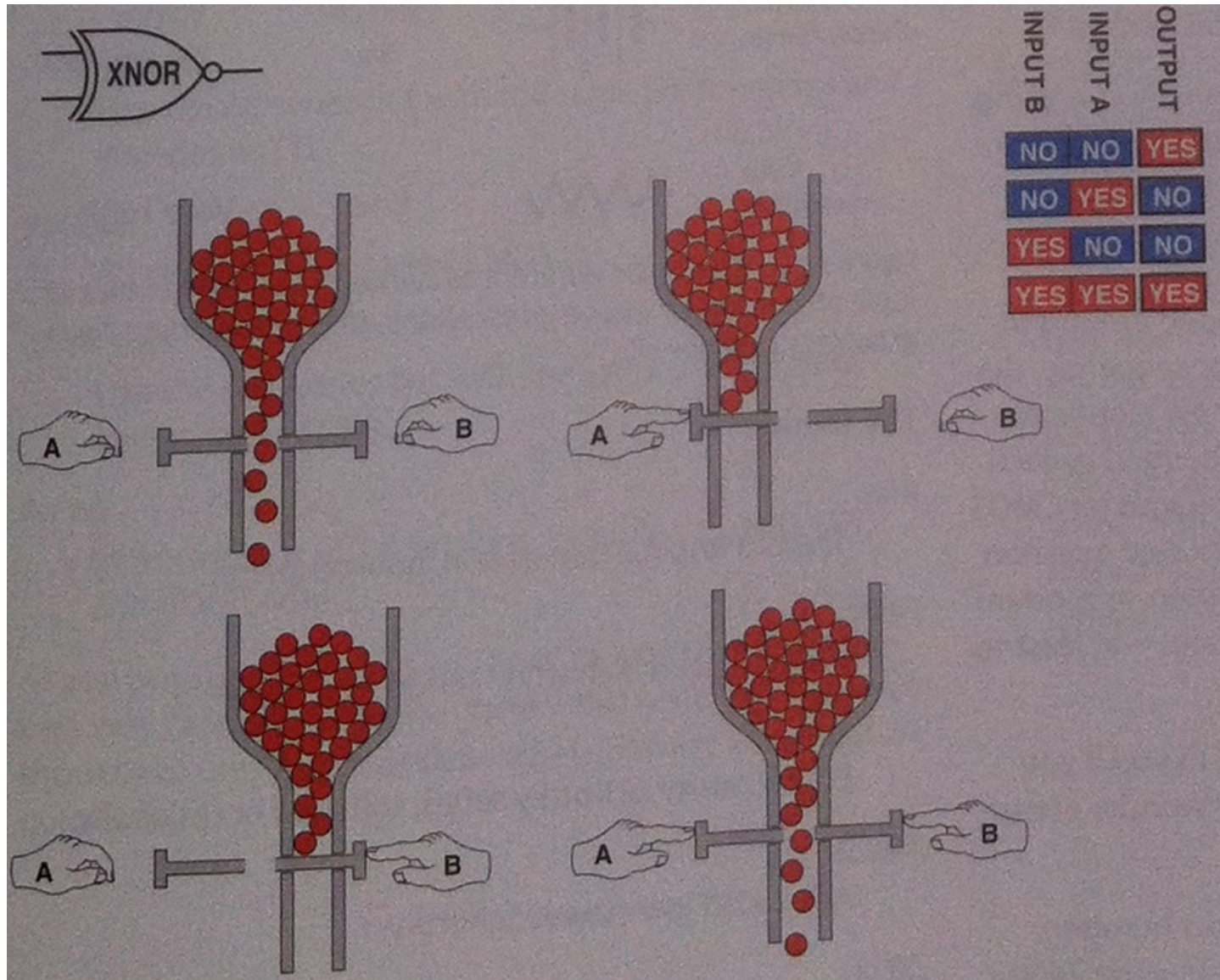
[Platt 2009]



[Platt 2009]



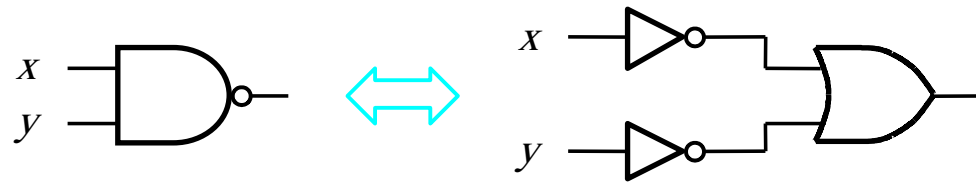
[Platt 2009]



[Platt 2009]

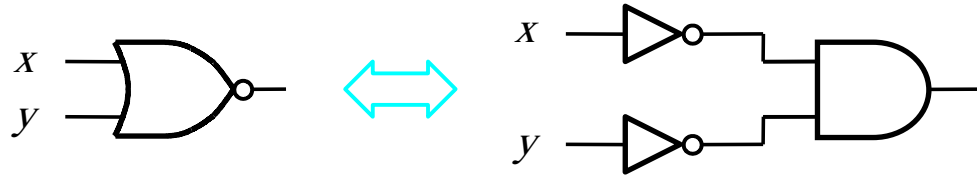
DeMorgan's Theorem Revisited

DeMorgan's theorem (in terms of logic gates)



$$\overline{x \cdot y} = \overline{x} + \overline{y}$$

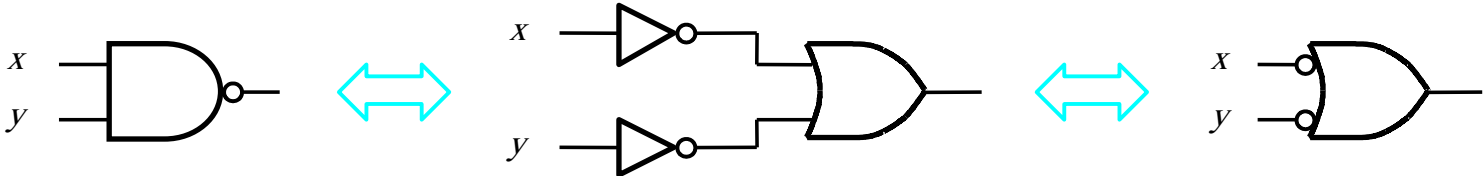
The other DeMorgan's theorem (in terms of logic gates)



$$\overline{x + y} = \overline{x} \cdot \overline{y}$$

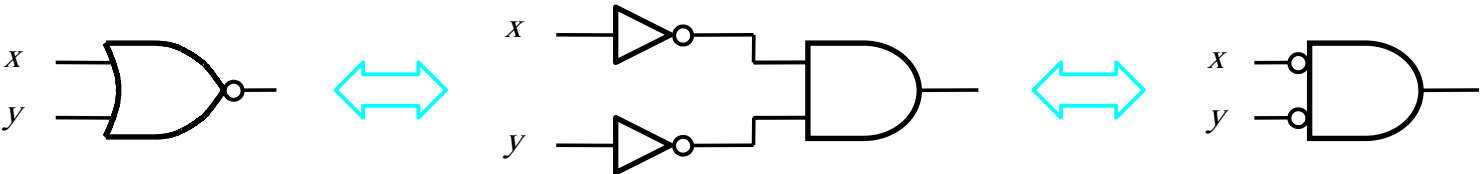
Shortcut Notation

DeMorgan's theorem in terms of logic gates



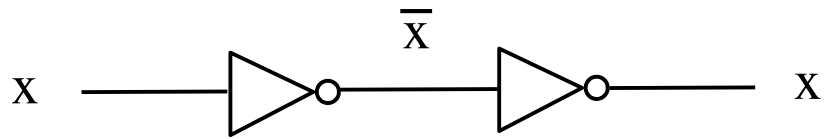
(Theorem 15.a) $\overline{x \cdot y} = \overline{x} + \overline{y}$

DeMorgan's theorem in terms of logic gates

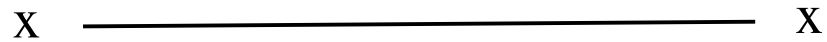
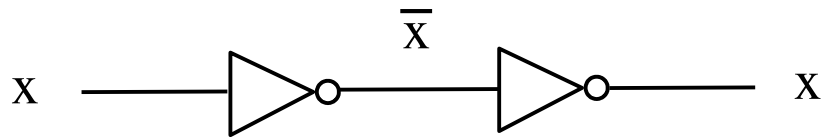


(Theorem 15.b) $\overline{x + y} = \overline{x} \overline{y}$

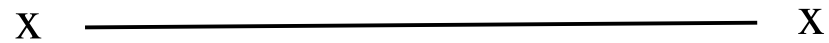
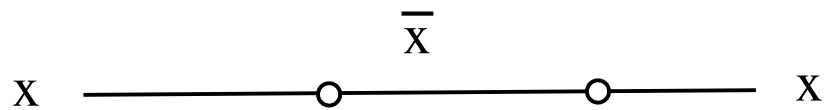
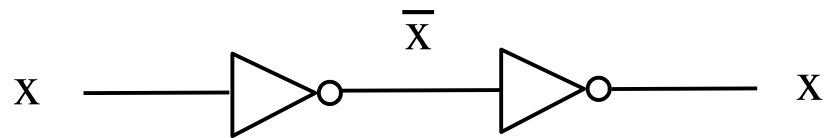
Two NOTs in a row



Two NOTs in a row

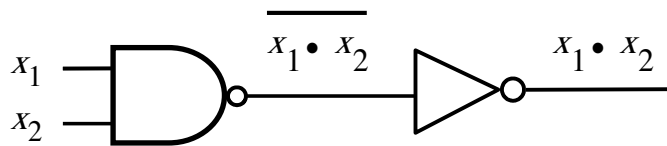


Two NOTs in a row



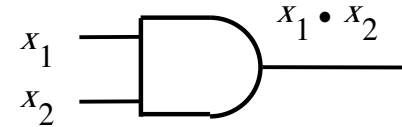
NAND-NAND Implementation of Sum-of-Products Expressions

NAND followed by NOT = AND



x_1	x_2	f
0	0	1
0	1	1
1	0	1
1	1	0

f
0
0
0
1



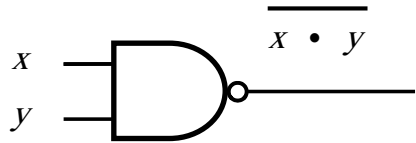
x_1	x_2	f
0	0	0
0	1	0
1	0	0
1	1	1

DeMorgan's Theorem

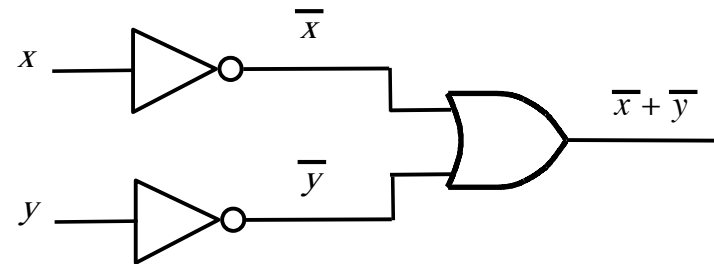
$$15a. \quad \overline{x \cdot y} = \bar{x} + \bar{y}$$

DeMorgan's Theorem

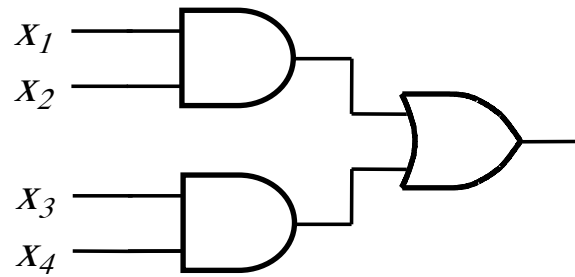
15a. $\overline{\mathbf{x} \cdot \mathbf{y}} = \overline{\mathbf{x}} + \overline{\mathbf{y}}$



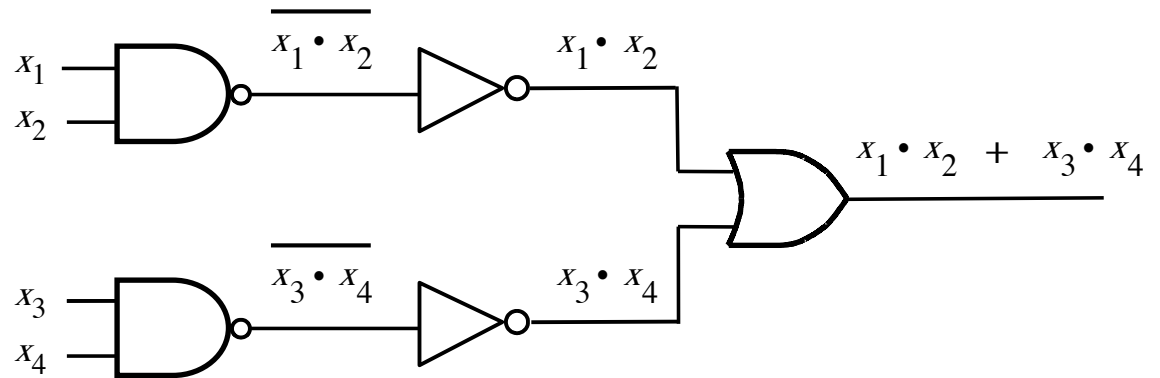
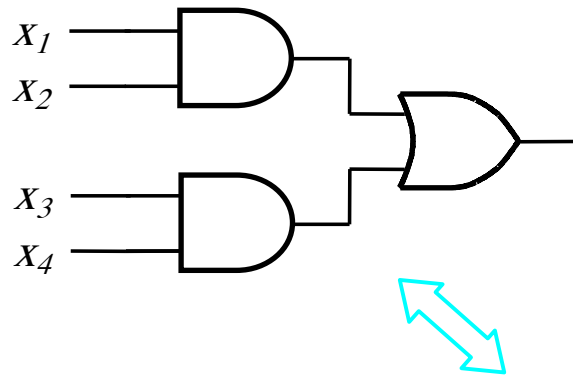
=



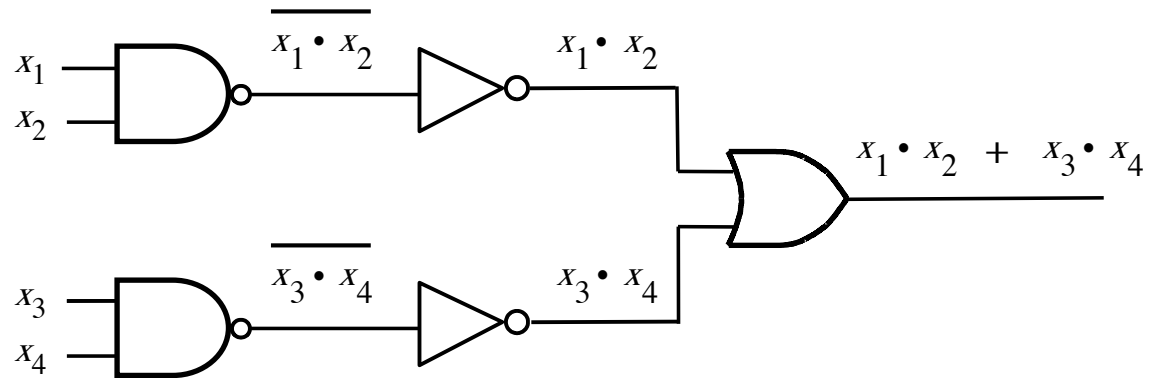
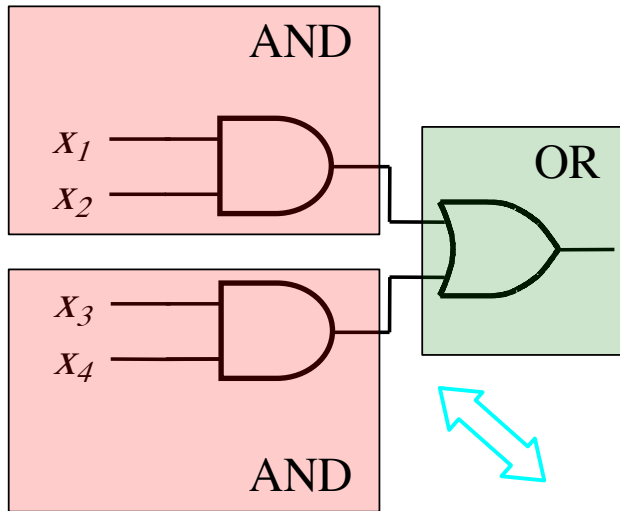
Sum-Of-Products



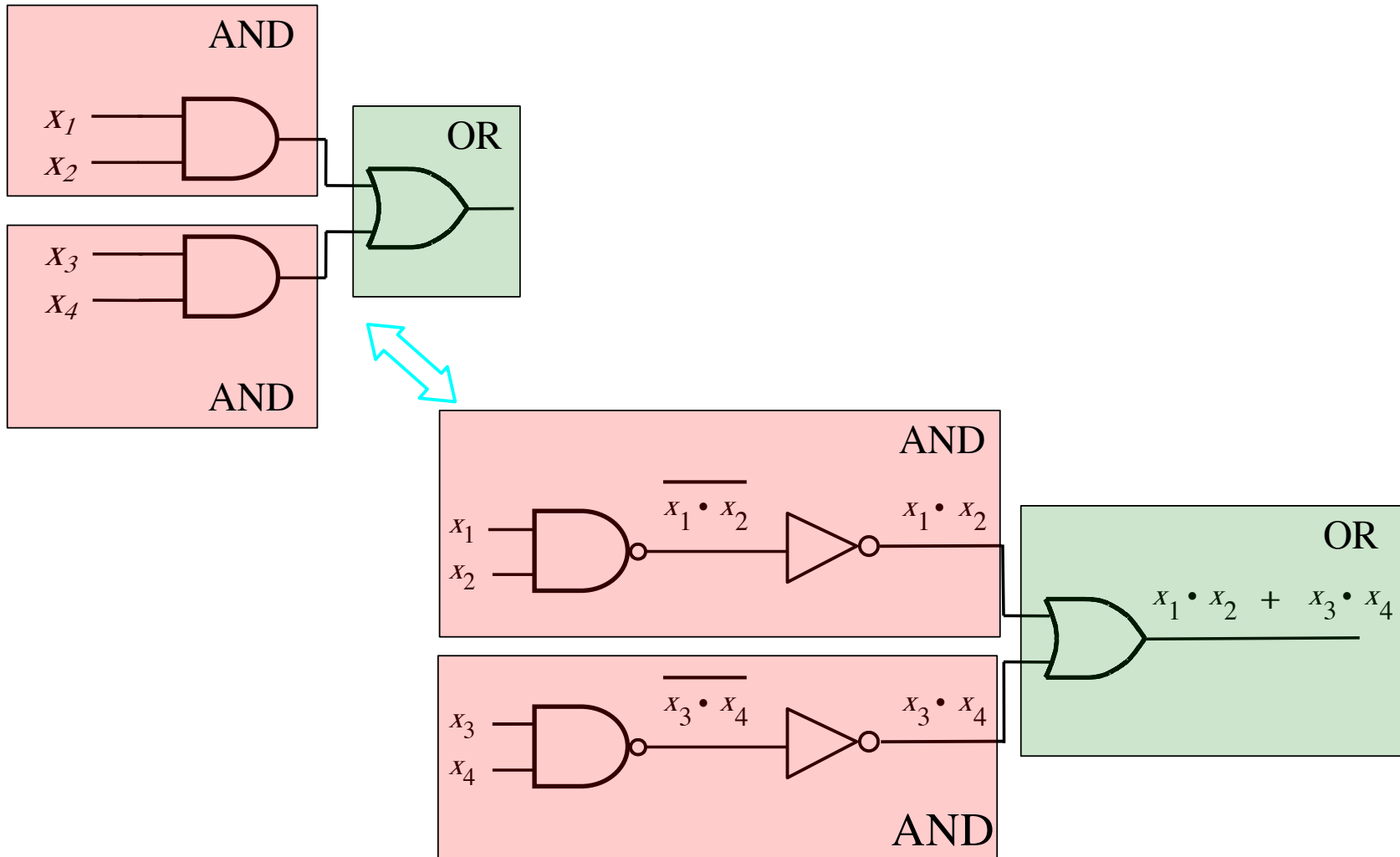
Sum-Of-Products



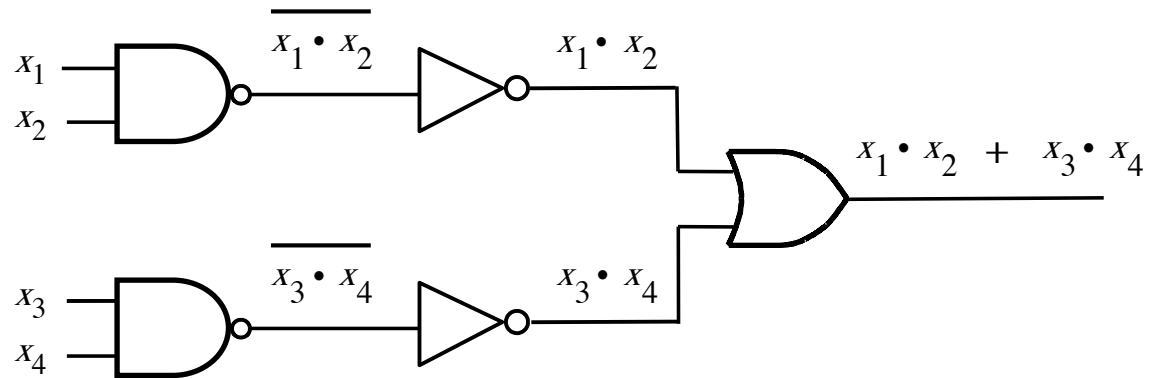
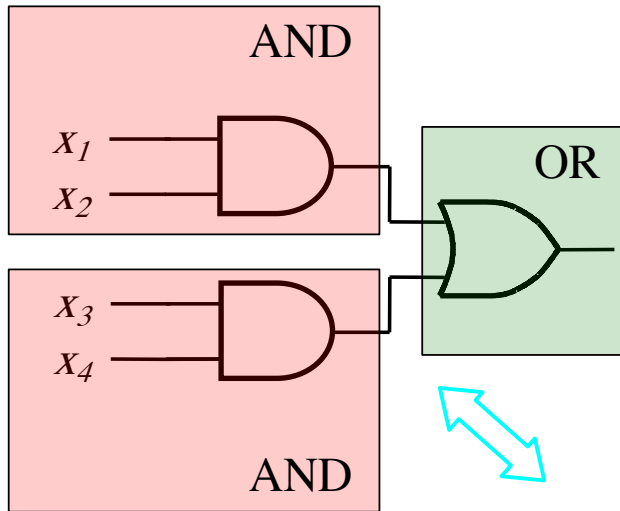
Sum-Of-Products



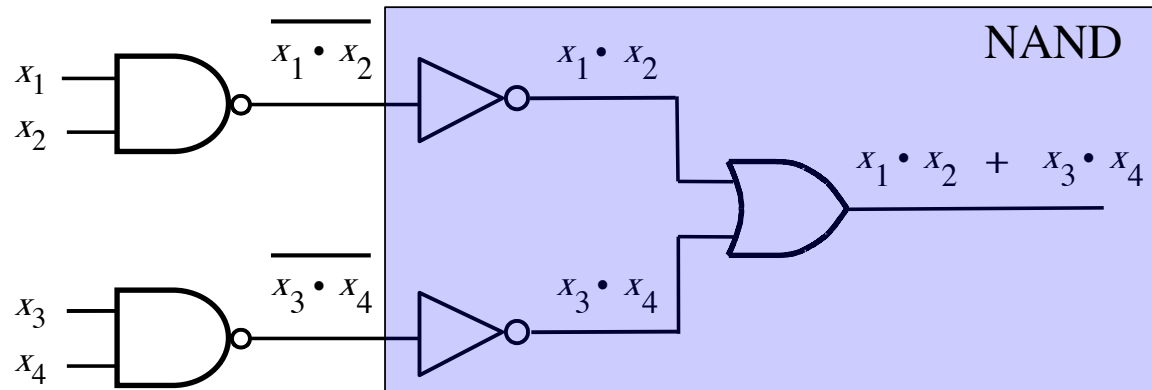
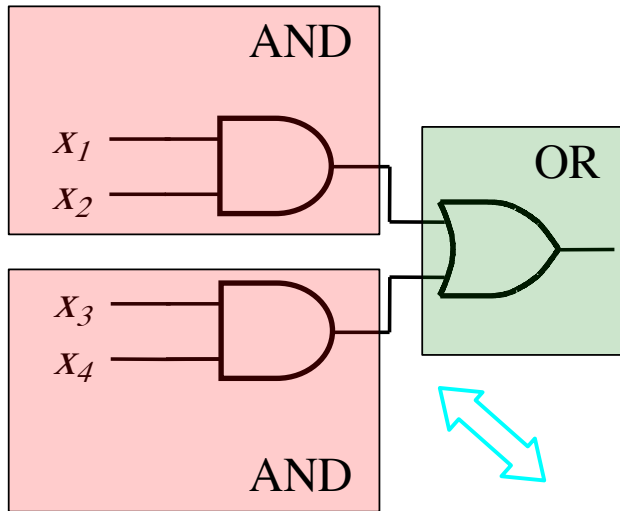
Sum-Of-Products



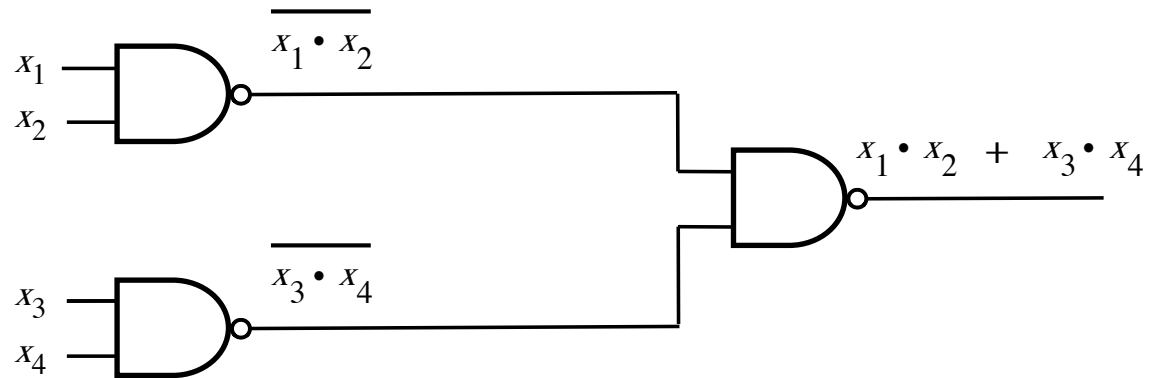
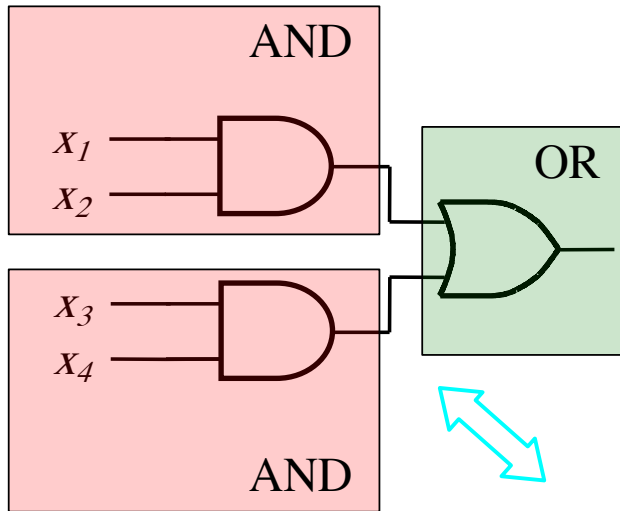
Sum-Of-Products



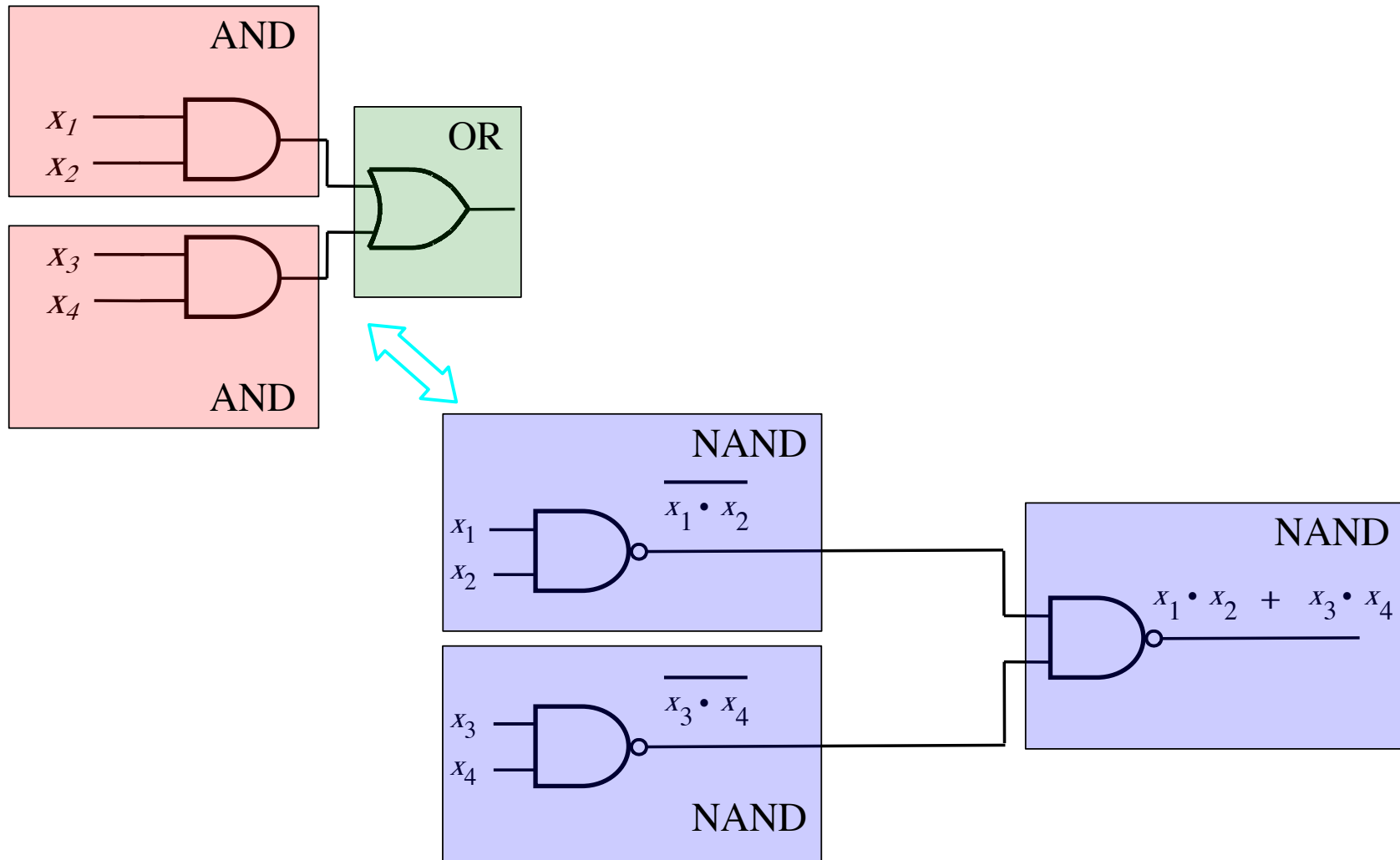
Sum-Of-Products



Sum-Of-Products

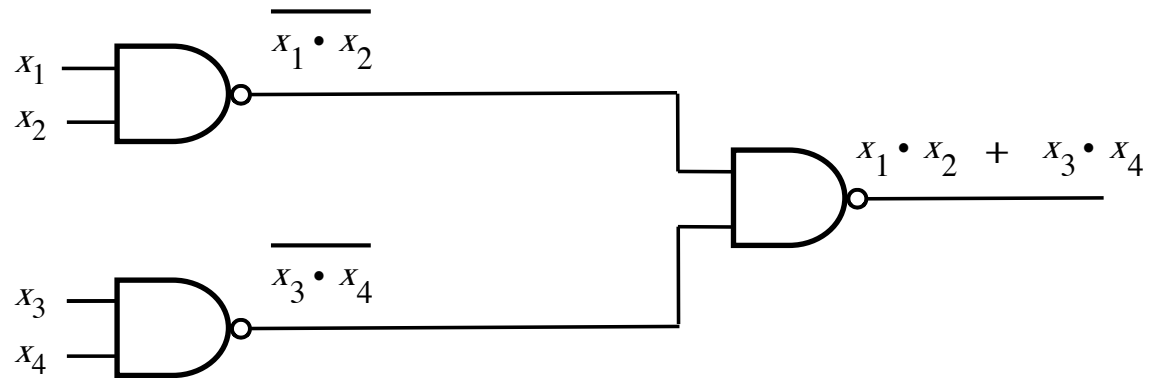
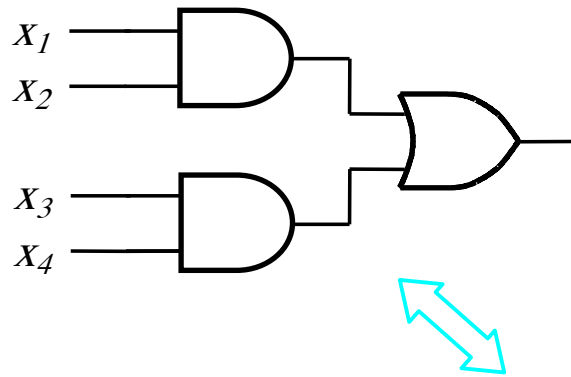


Sum-Of-Products



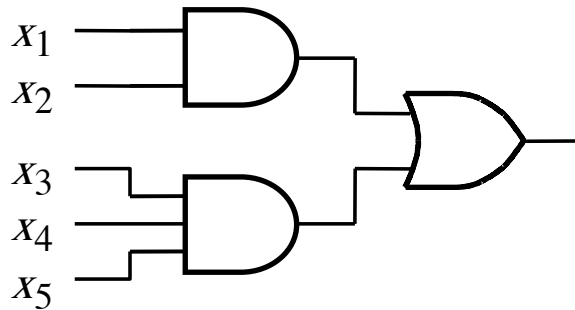
This circuit uses only NANDs

Sum-Of-Products

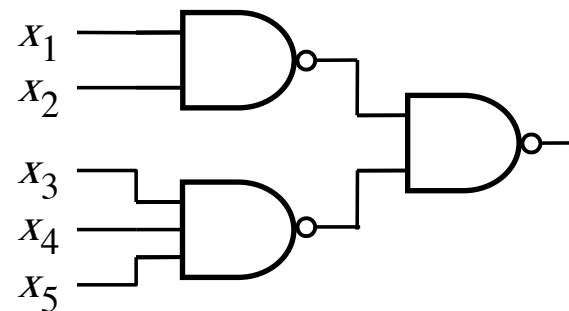
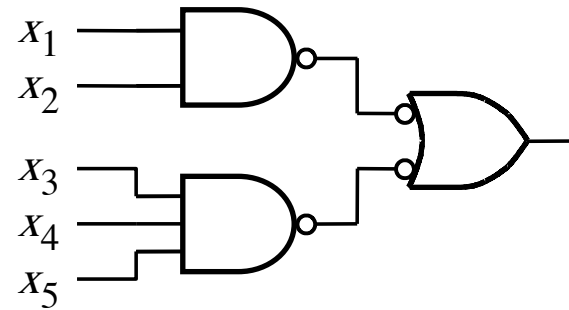
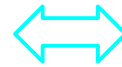


This circuit uses only NANDs

Another SOP Example



This circuit uses ANDs & OR

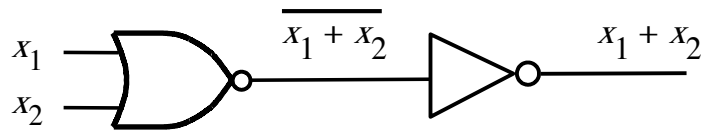


This circuit uses only NANDs

[Figure 2.27 from the textbook]

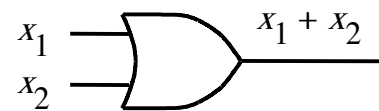
NOR-NOR Implementation of Product-of-Sums Expressions

NOR followed by NOT = OR



x_1	x_2	f
0	0	1
0	1	0
1	0	0
1	1	0

f
0
1
1
1



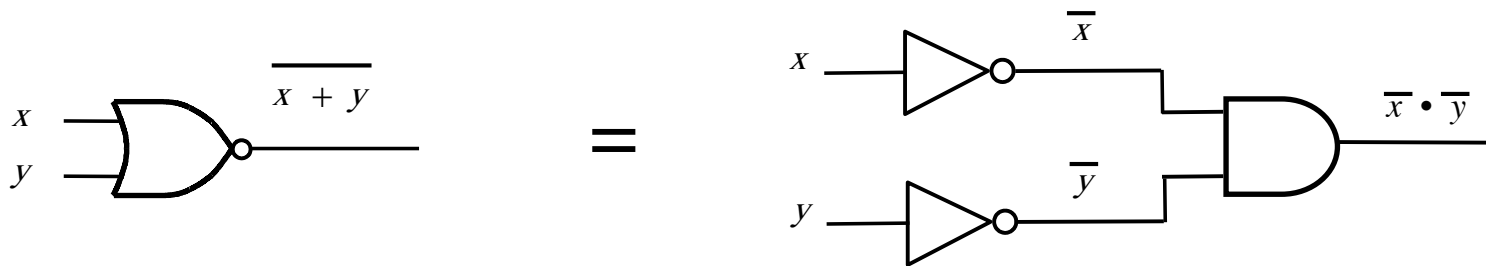
x_1	x_2	f
0	0	0
0	1	1
1	0	1
1	1	1

DeMorgan's Theorem

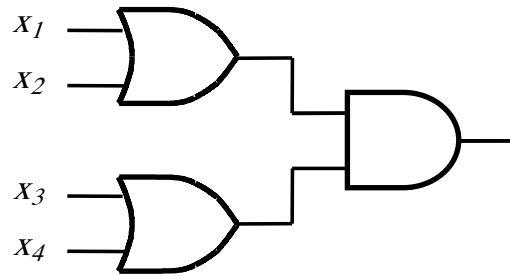
$$15b. \quad \overline{x + y} = \bar{x} \cdot \bar{y}$$

DeMorgan's Theorem

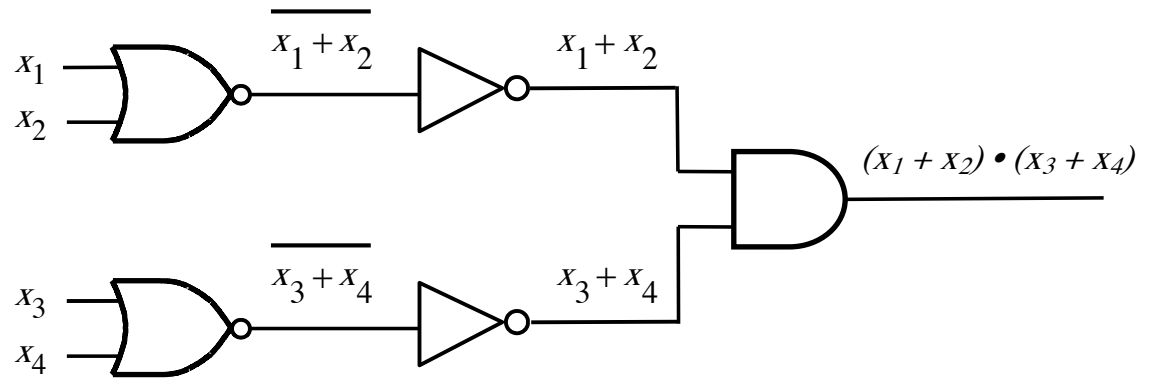
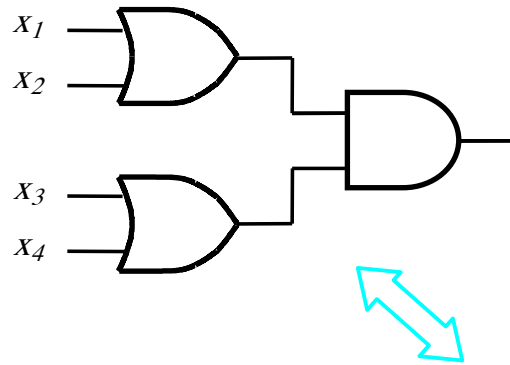
15b. $\overline{x + y} = \bar{x} \cdot \bar{y}$



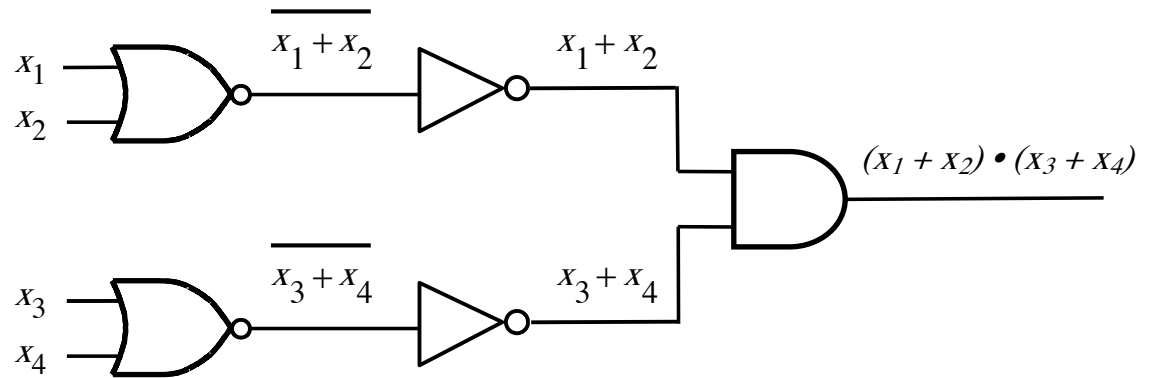
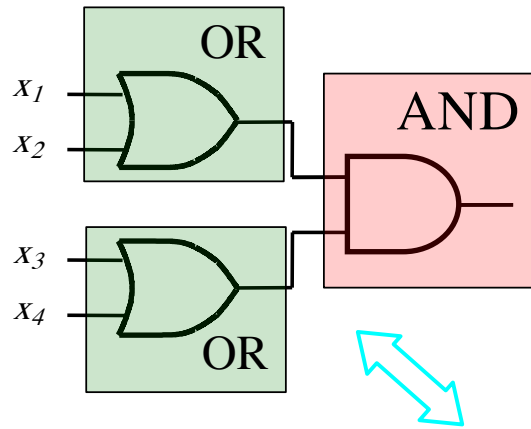
Product-Of-Sums



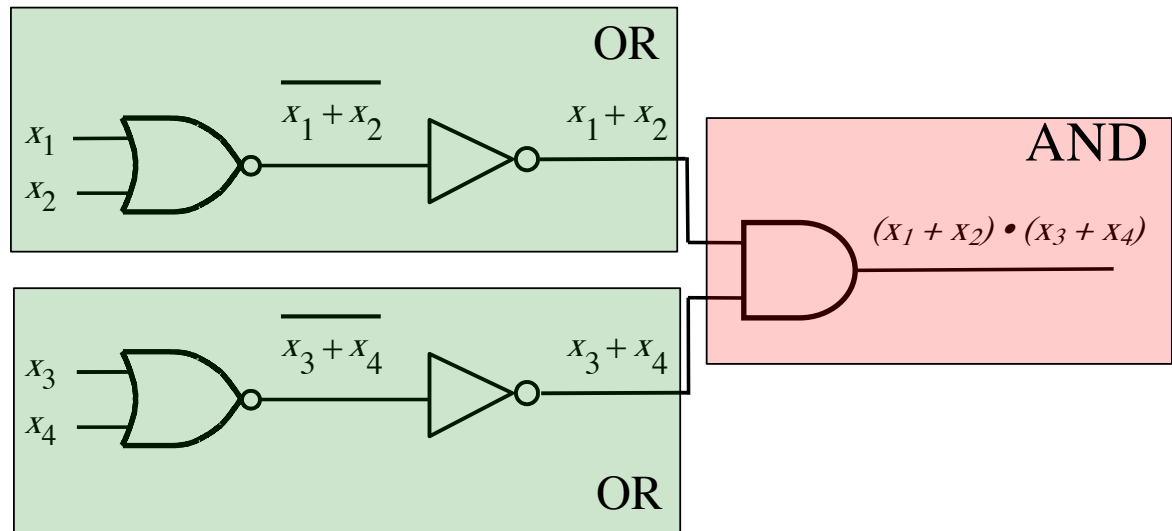
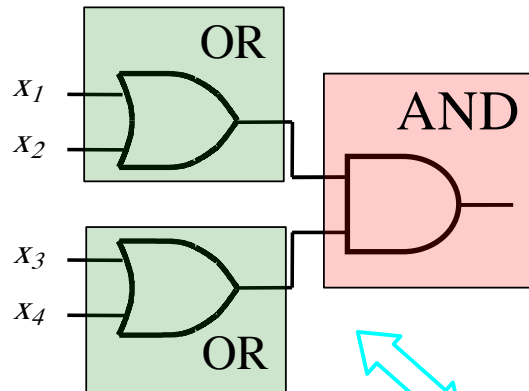
Product-Of-Sums



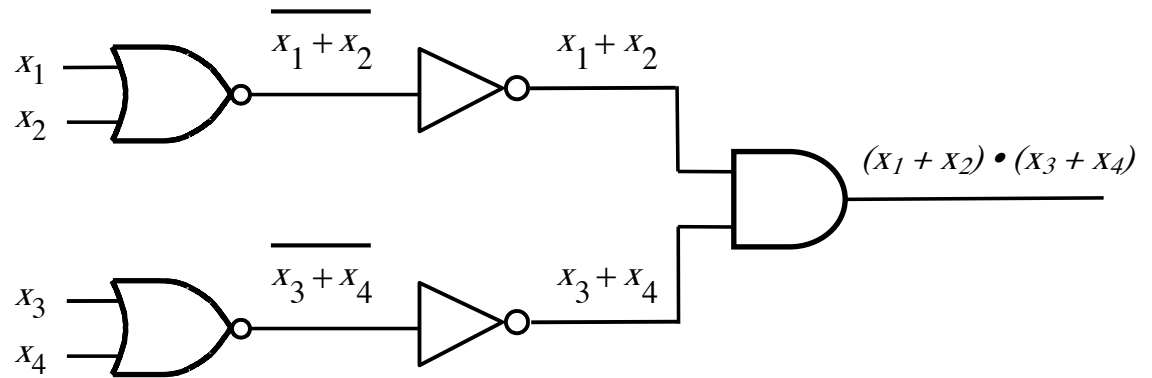
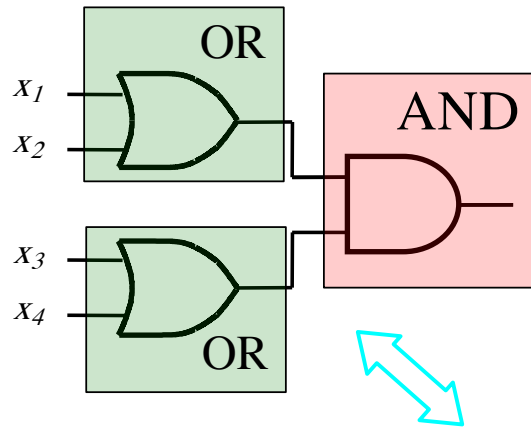
Product-Of-Sums



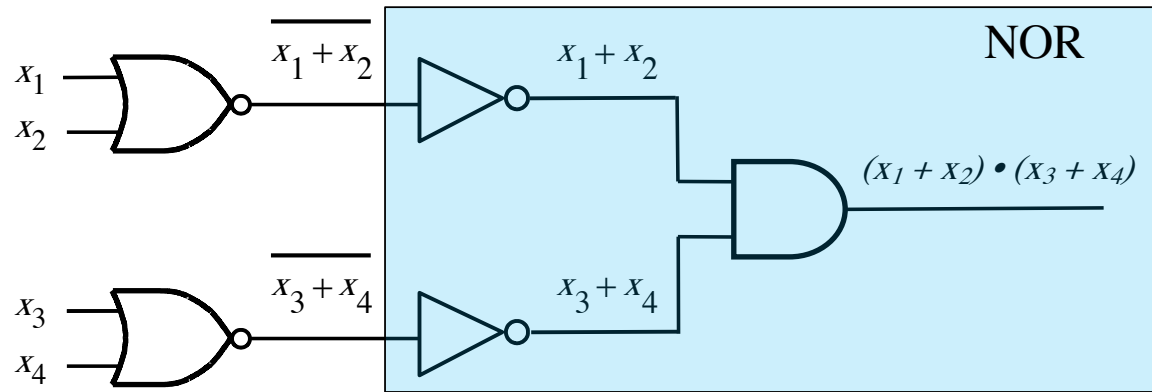
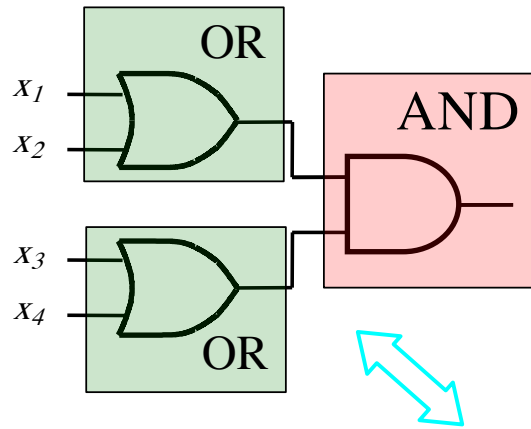
Product-Of-Sums



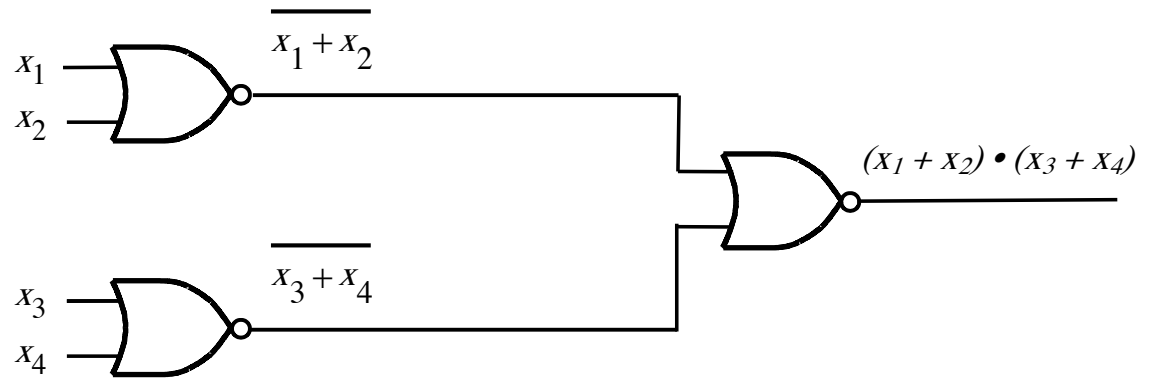
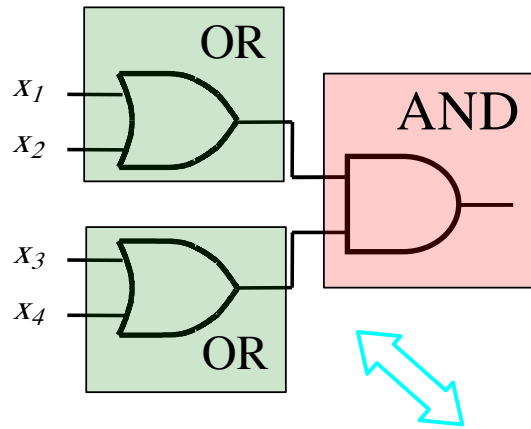
Product-Of-Sums



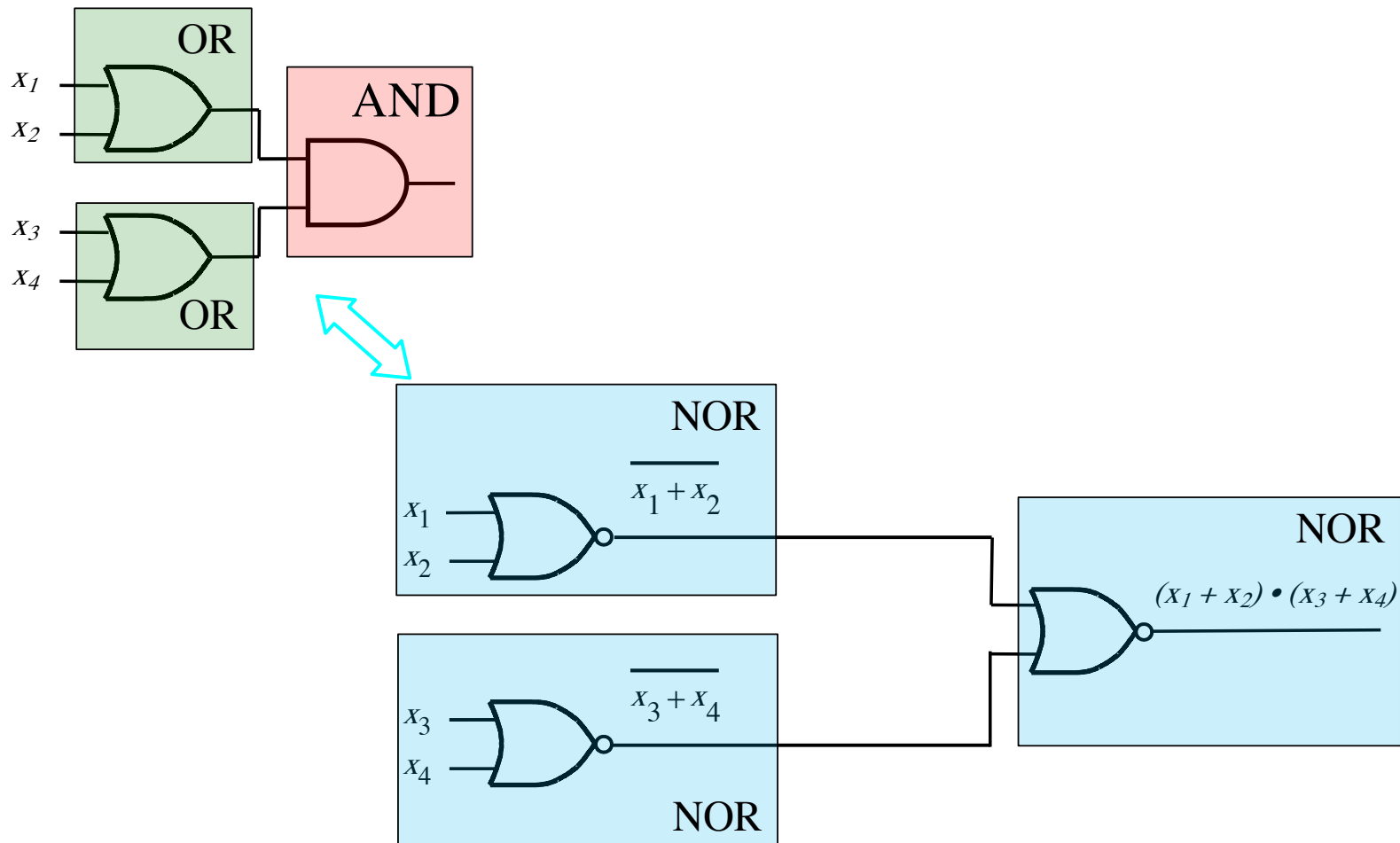
Product-Of-Sums



Product-Of-Sums

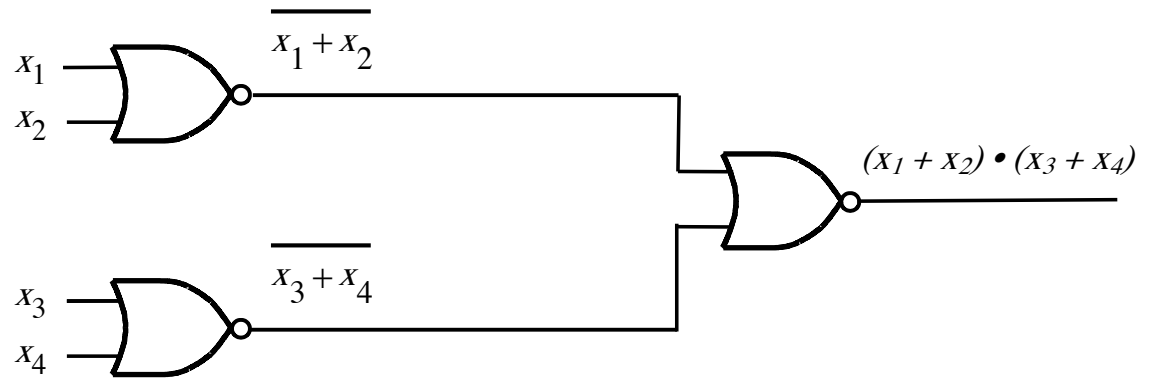
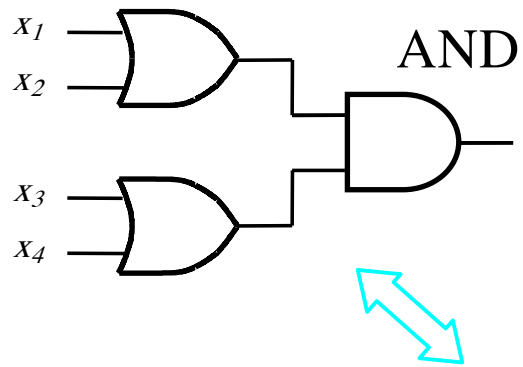


Product-Of-Sums



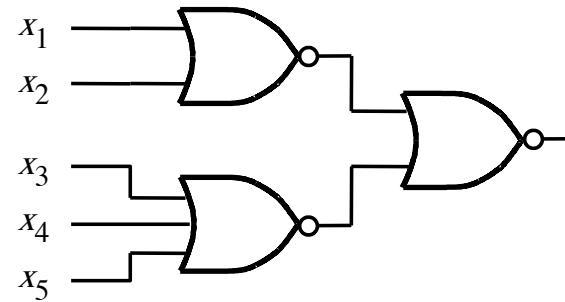
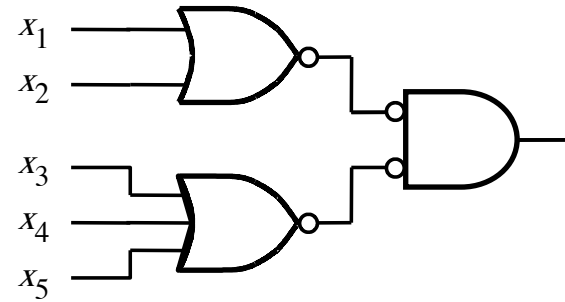
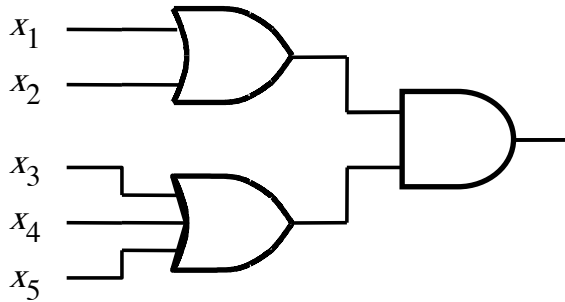
This circuit uses only NORs

Product-Of-Sums



This circuit uses only NORs

Another POS Example



This circuit uses ORs & AND

This circuit uses only NORs

[Figure 2.28 from the textbook]

Summary

- **Sum-of-Products (SOP) expressions are directly mappable to NAND-NAND implementation.**
- **Product-of-Sums (POS) expressions are directly mappable to NOR-NOR implementation.**
- **Going from SOP to NOR-NOR is not that easy.**
- **Similarly, converting from POS to NAND-NAND implementation requires extra work.**

Questions?

THE END