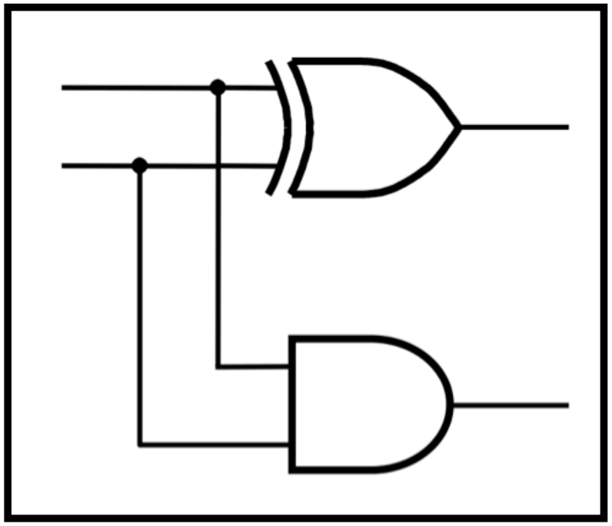


CprE 281: Digital Logic

Instructor: Alexander Stoytchev

http://www.ece.iastate.edu/~alexs/classes/

Signed Numbers

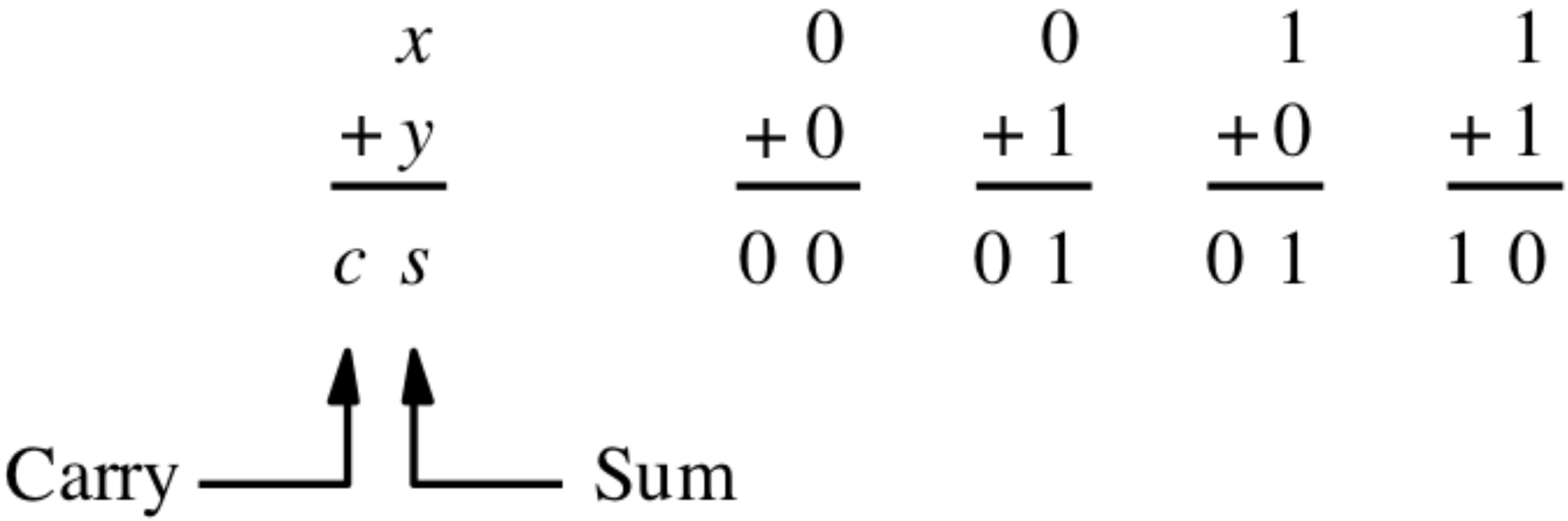
CprE 281: Digital Logic
Iowa State University, Ames, IA

Signed Integer Numbers

CprE 281: Digital Logic
Iowa State University, Ames, IA

Quick Review

Adding two bits
(there are four possible cases)



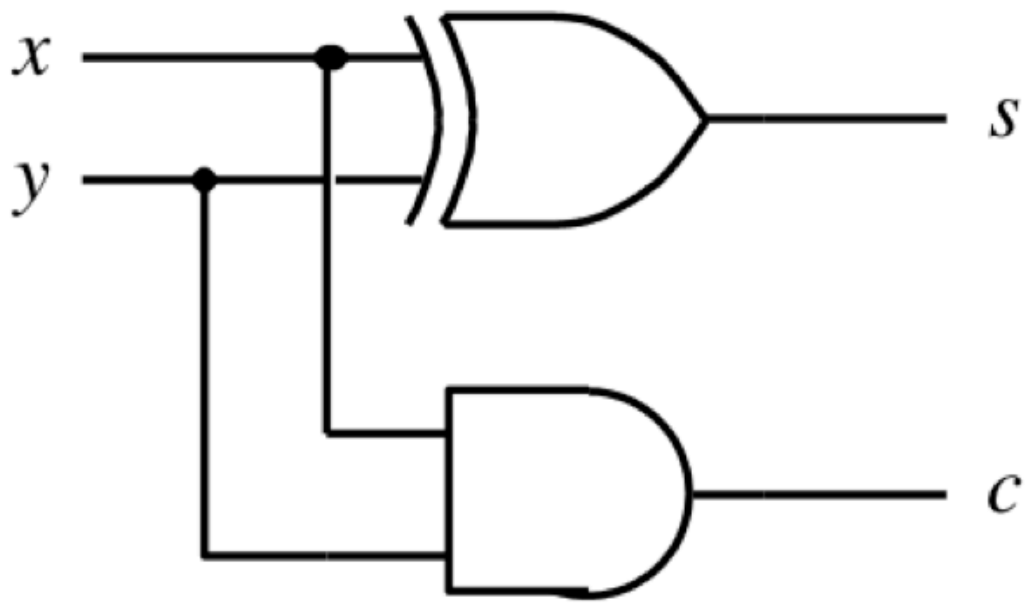
[Figure 3.1a from the textbook]

Adding two bits
(the truth table)

x	y	Carry c	Sum s
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

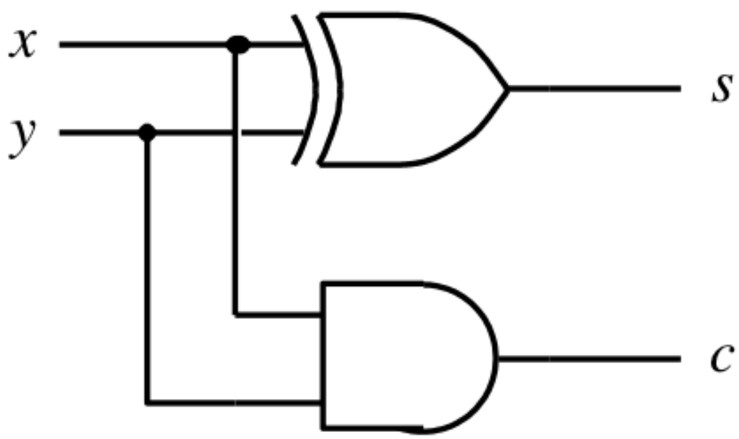
[Figure 3.1b from the textbook]

Adding two bits
(the logic circuit)

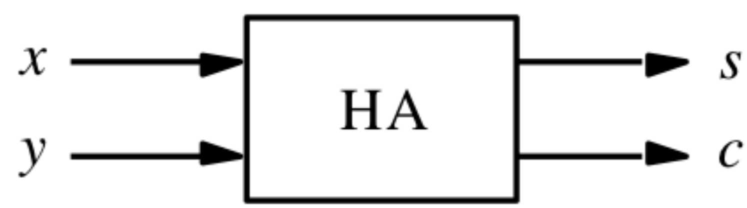


[Figure 3.1c from the textbook]

The Half-Adder



(c) Circuit



(d) Graphical symbol

[Figure 3.1c-d from the textbook]

Addition of multibit numbers

Generated carries →	1 1 1 0		
$X = x_4x_3x_2x_1x_0$	0 1 1 1 1	$(15)_{10}$	
$+\,Y = y_4y_3y_2y_1y_0$	+ 0 1 0 1 0	+ $(10)_{10}$	
$S = s_4s_3s_2s_1s_0$	1 1 0 0 1	$(25)_{10}$	

…	c_{i+1}	c_i	…
…	…	x_i	…
…	…	y_i	…
…	…	s_i	…

Bit position i

[Figure 3.2 from the textbook]

Analogy with addition in base 10

$$
\begin{array}{rccc}
 & x_2 & x_1 & x_0 \\
+ & y_2 & y_1 & y_0 \\
\hline
 & s_2 & s_1 & s_0
\end{array}
$$

Analogy with addition in base 10

$$
\begin{array}{r}
3\ 8\ 9 \\
+\ \ 1\ 5\ 7 \\
\hline
5\ 4\ 6
\end{array}
$$

Analogy with addition in base 10

carry	0	1	1	0
		3	8	9
+		1	5	7
		5	4	6

Analogy with addition in base 10

$$
\begin{array}{rcccc}
 & c_3 & c_2 & c_1 & c_0 \\
 & & x_2 & x_1 & x_0 \\
+ & & y_2 & y_1 & y_0 \\
\hline
 & & s_2 & s_1 & s_0
\end{array}
$$

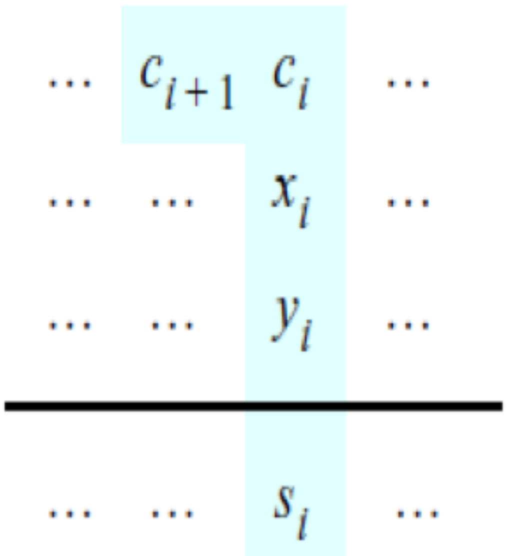
Another example in base 10

$$
\begin{array}{r}
938 \\
+\ 214 \\
\hline
1152
\end{array}
$$

Another example in base 10

carry		1	0	1	0
			9	3	8
	+		2	1	4
			1	5	2

Problem Statement and Truth Table



[Figure 3.2b from the textbook]

c_i	x_i	y_i	c_{i+1}	s_i
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

[Figure 3.3a from the textbook]

Let's fill-in the two K-maps

c_i	x_i	y_i	c_{i+1}	s_i
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

$c_i \backslash x_i y_i$	00	01	11	10
0				
1				

$s_i =$

$c_i \backslash x_i y_i$	00	01	11	10
0				
1				

$c_{i+1} =$

[Figure 3.3a-b from the textbook]

Let's fill-in the two K-maps

c_i	x_i	y_i	c_{i+1}	s_i
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

c_i \ $x_i y_i$	00	01	11	10
0		1		1
1	1		1	

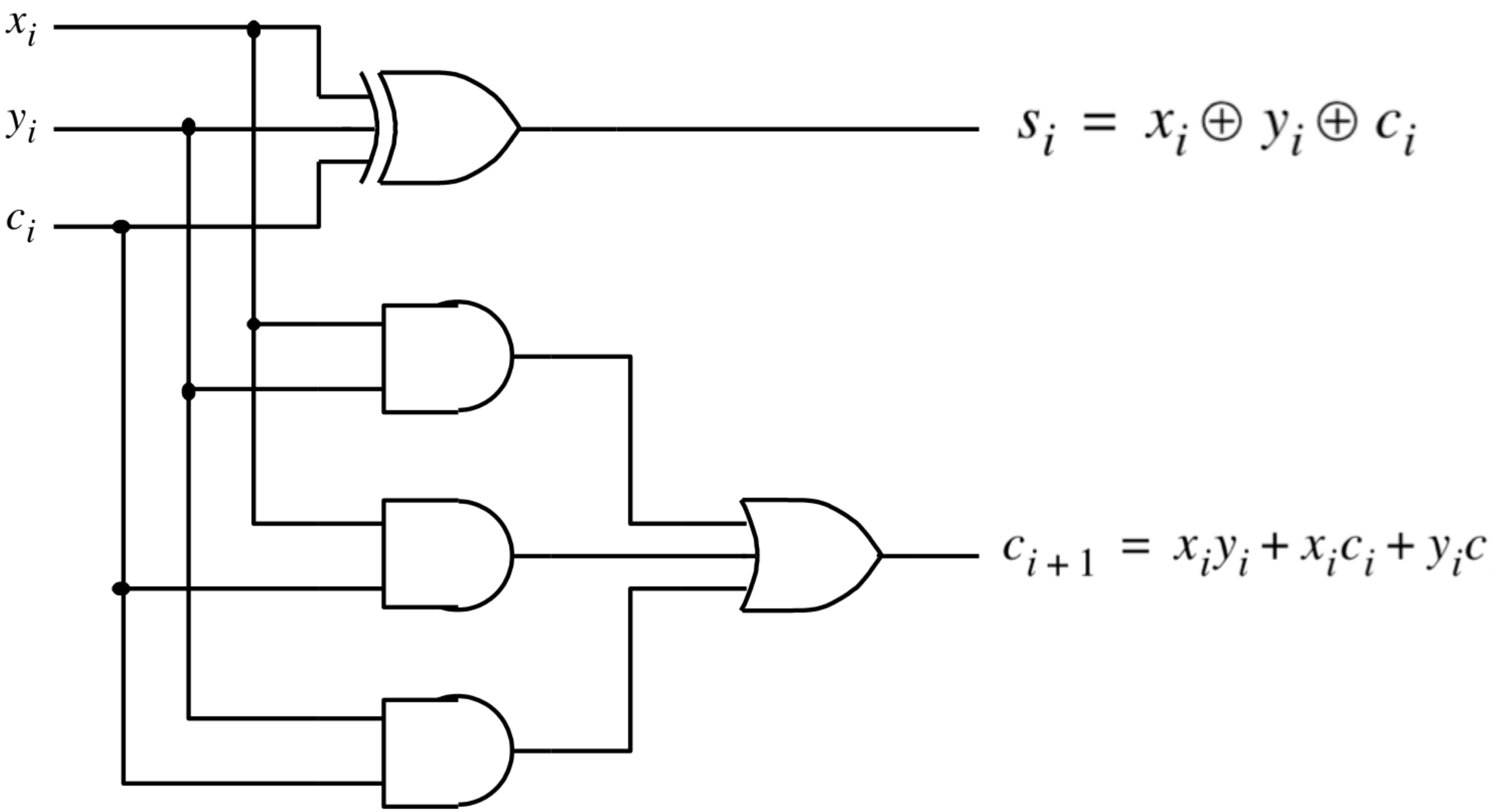
$$s_i = x_i \oplus y_i \oplus c_i$$

c_i \ $x_i y_i$	00	01	11	10
0			1	
1		1	1	1

$$c_{i+1} = x_i y_i + x_i c_i + y_i c_i$$

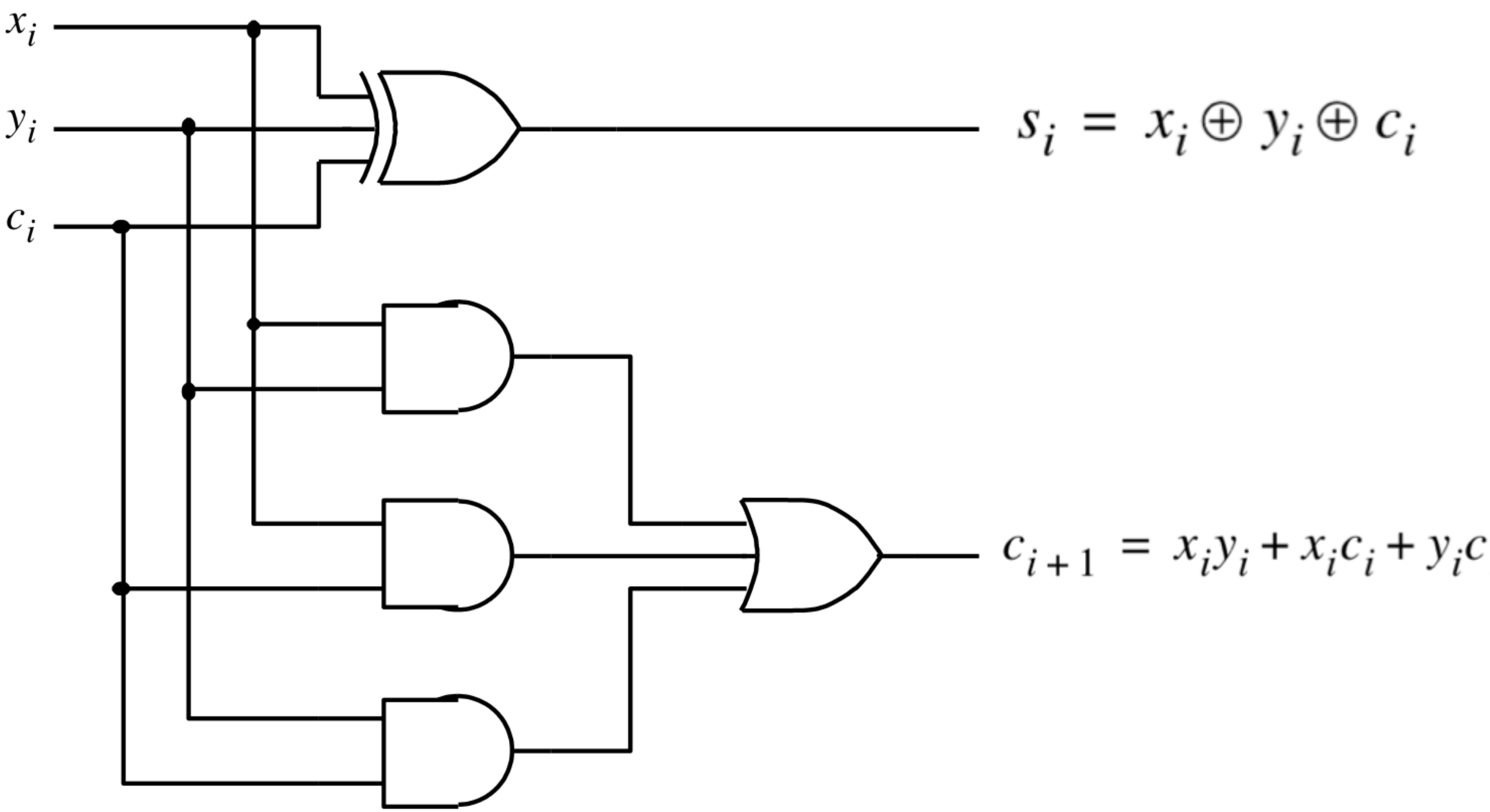
[Figure 3.3a-b from the textbook]

The circuit for the two expressions



[Figure 3.3c from the textbook]

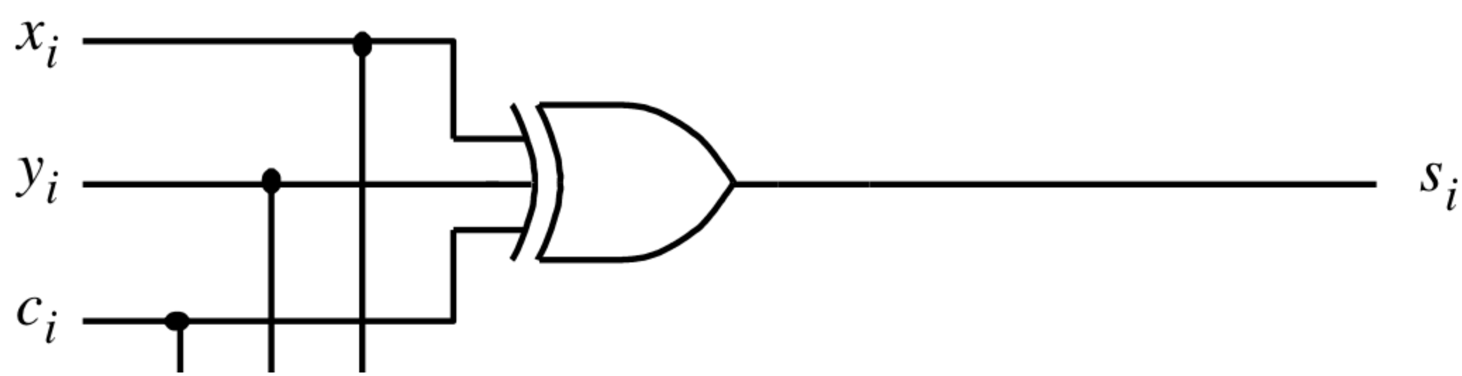
This is called the Full-Adder

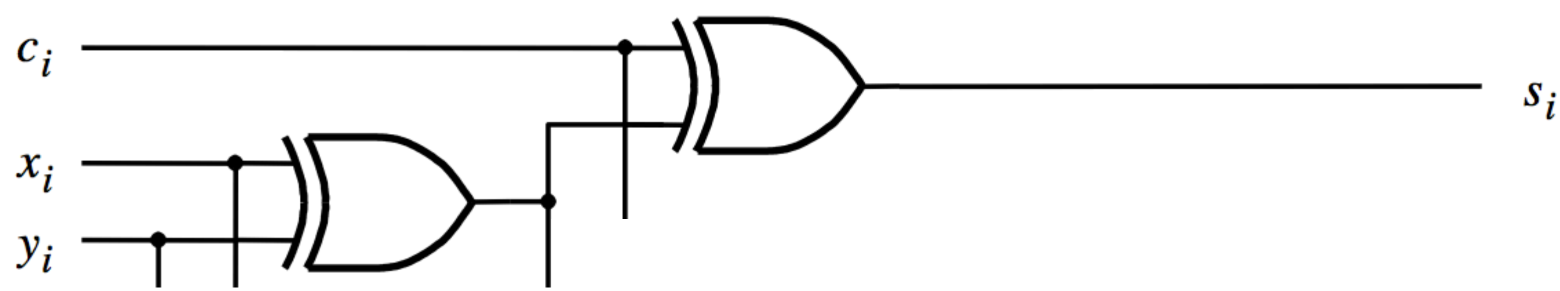


[Figure 3.3c from the textbook]

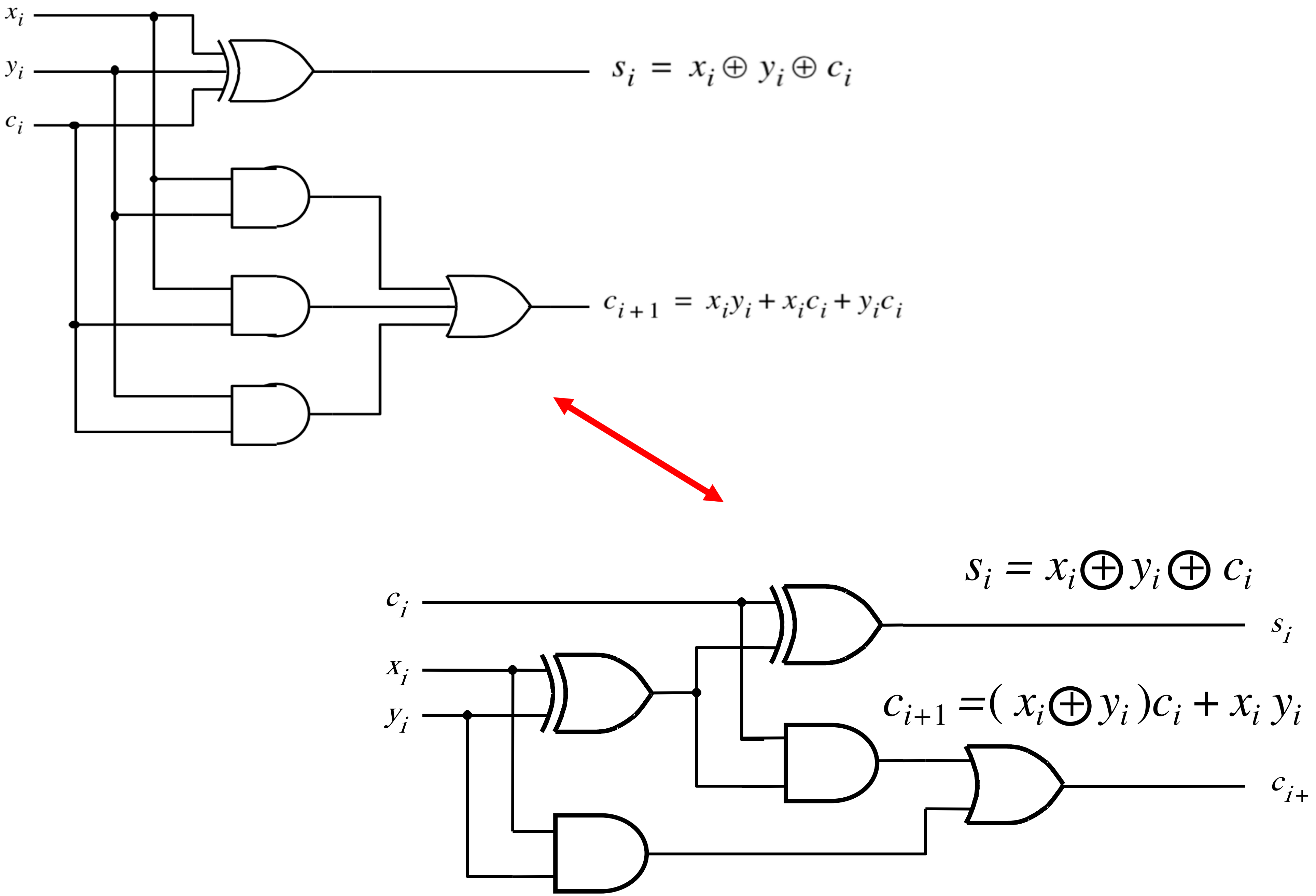
XOR Magic
(s_i can be implemented in two different ways)

$$s_i = x_i \oplus y_i \oplus c_i$$

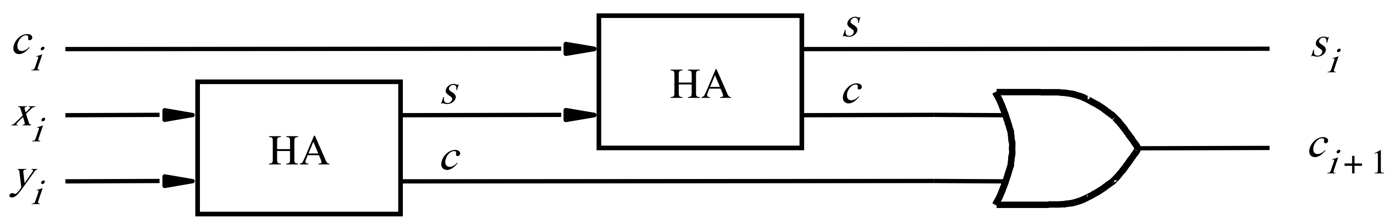




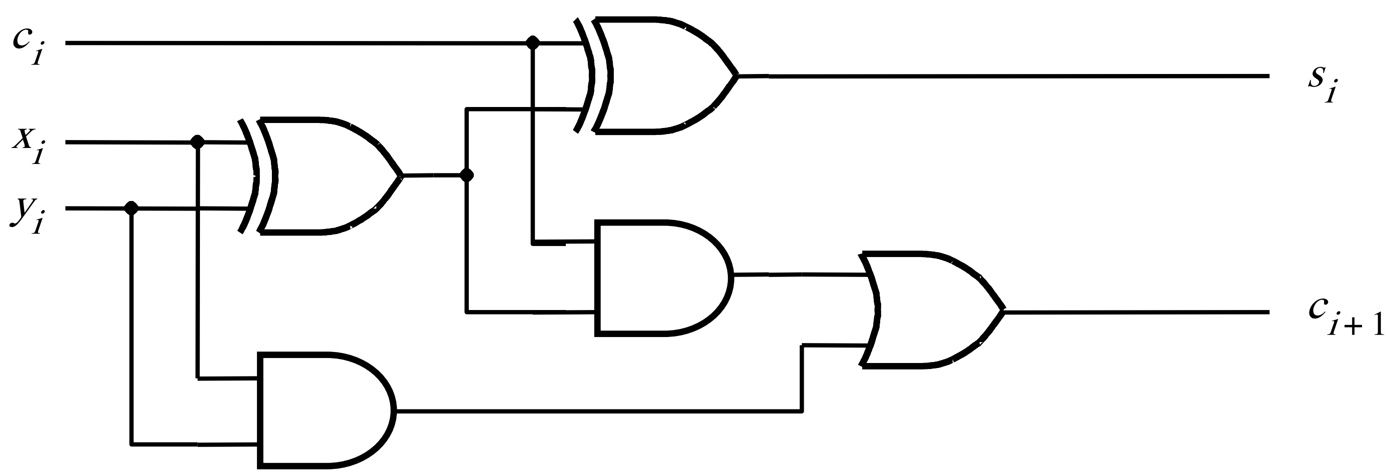
These two circuits are equivalent



A decomposed implementation of the full-adder circuit



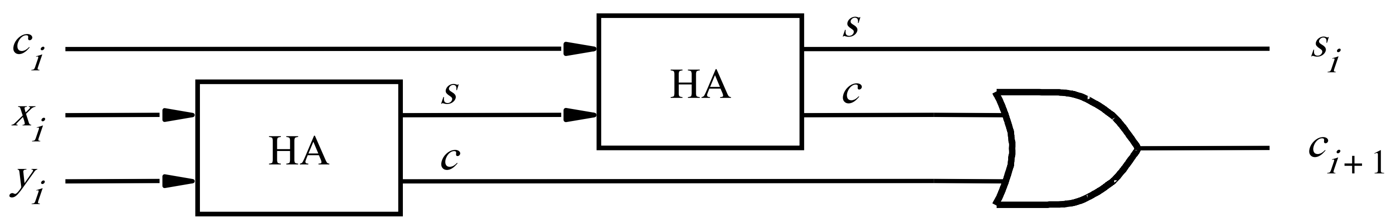
(a) Block diagram



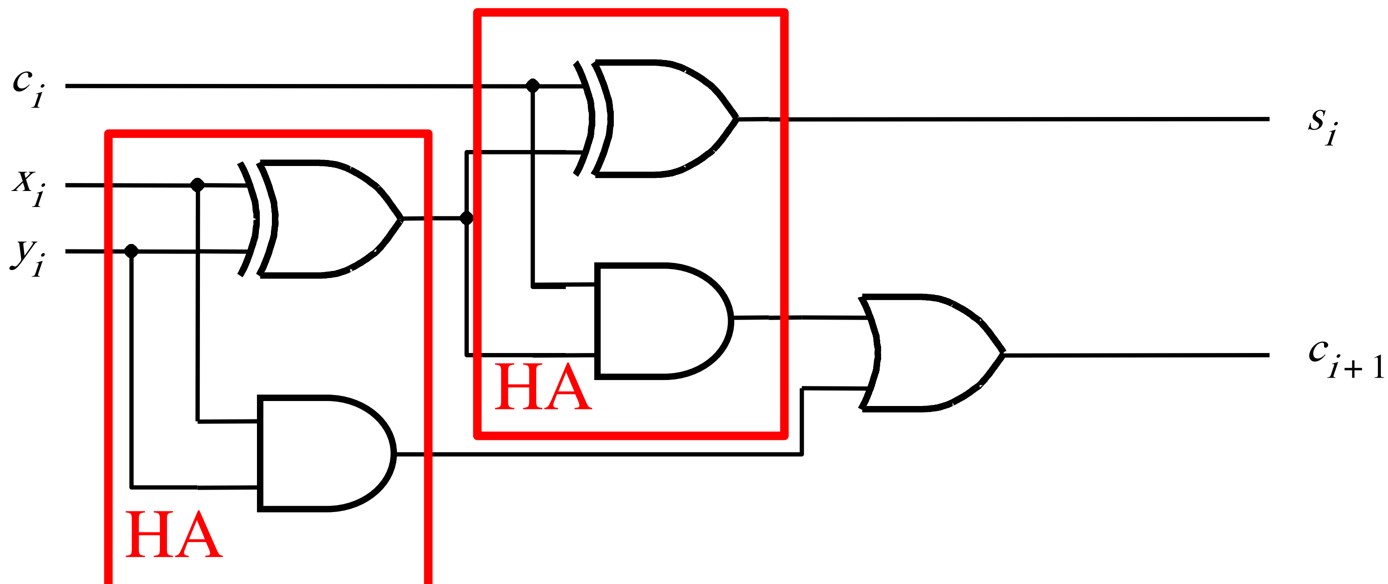
(b) Detailed diagram

[Figure 3.4 from the textbook]

A decomposed implementation of the full-adder circuit



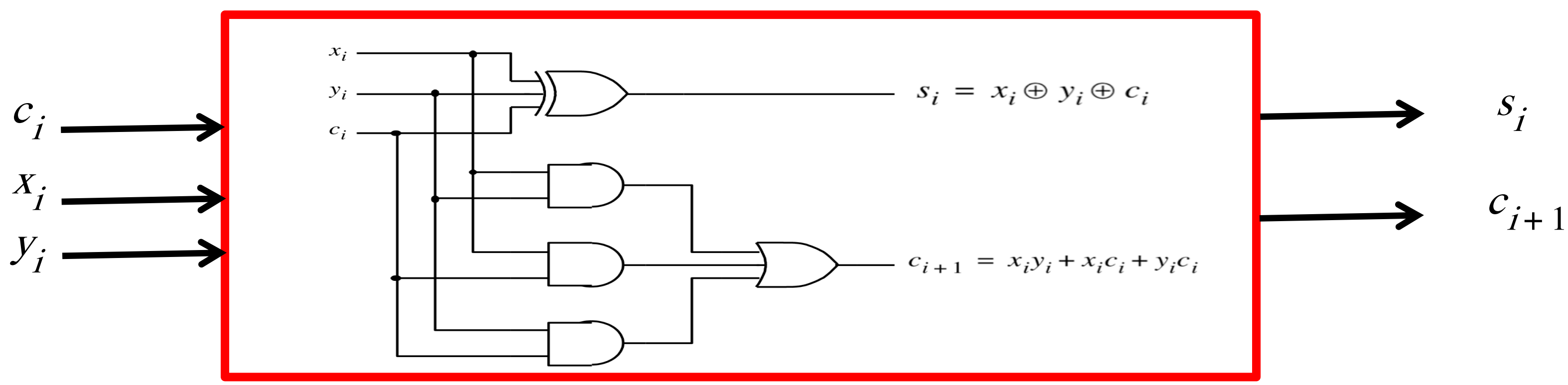
(a) Block diagram



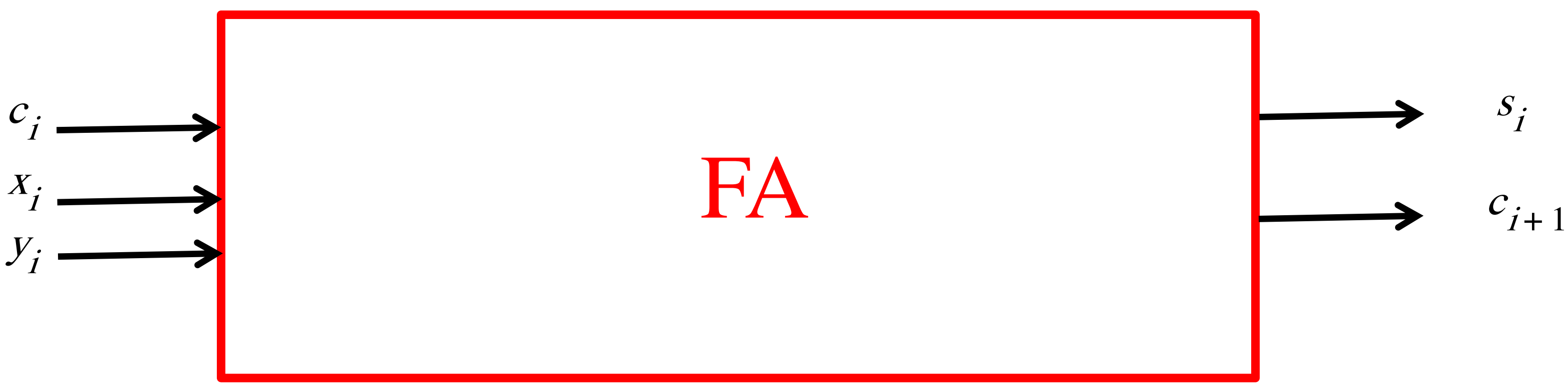
(b) Detailed diagram

[Figure 3.4 from the textbook]

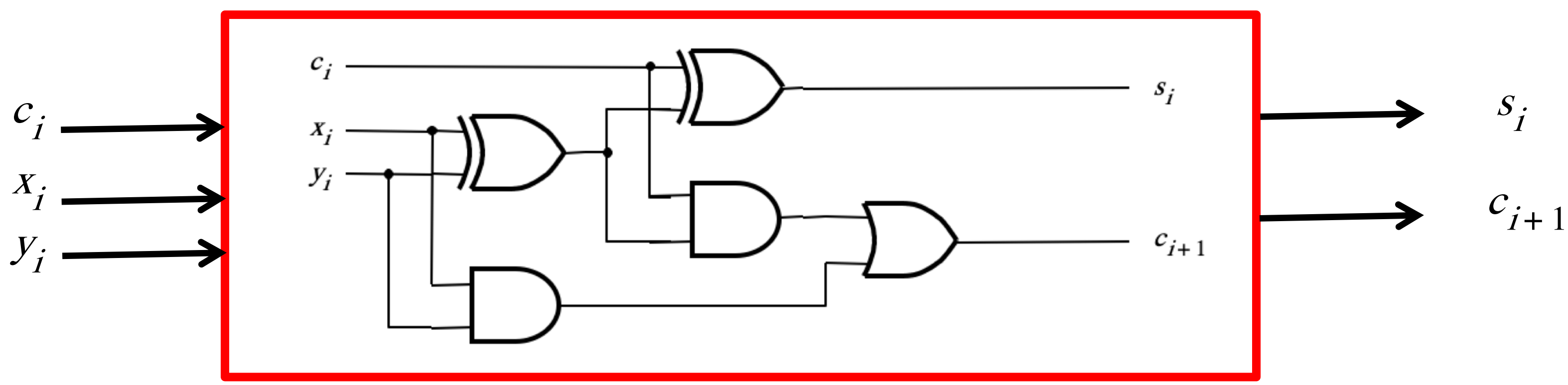
The Full-Adder Abstraction



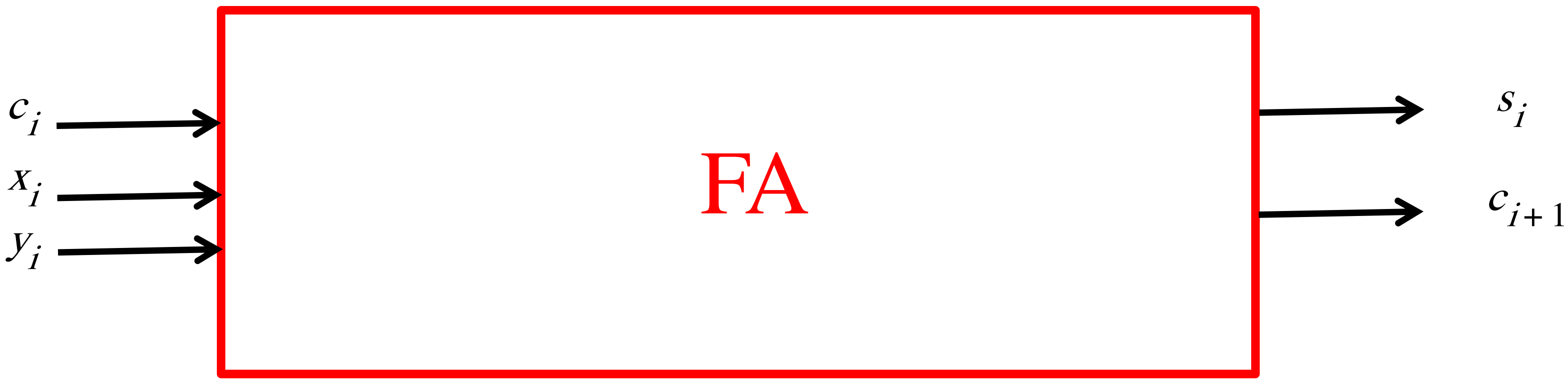
The Full-Adder Abstraction



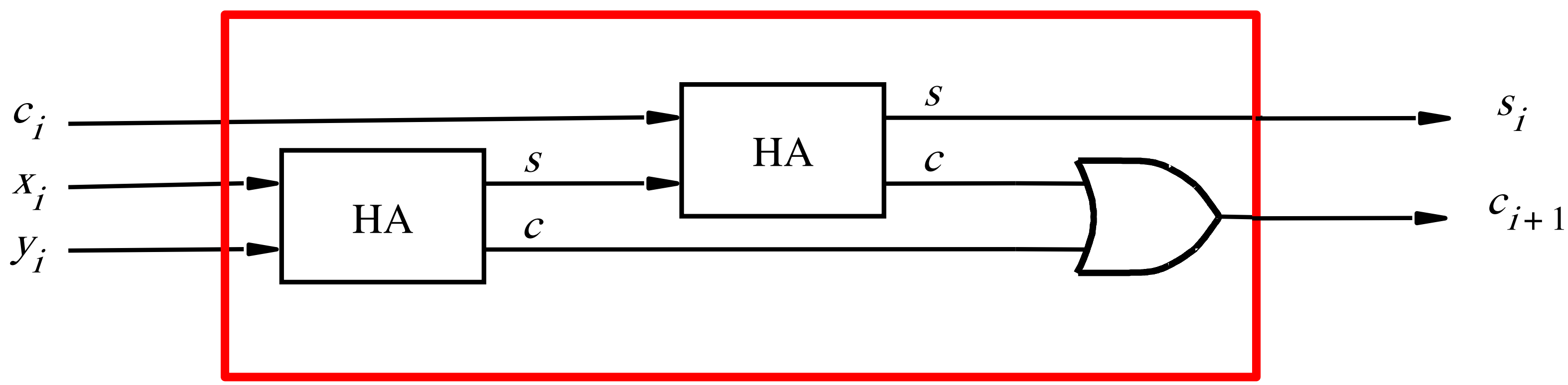
The Full-Adder Abstraction



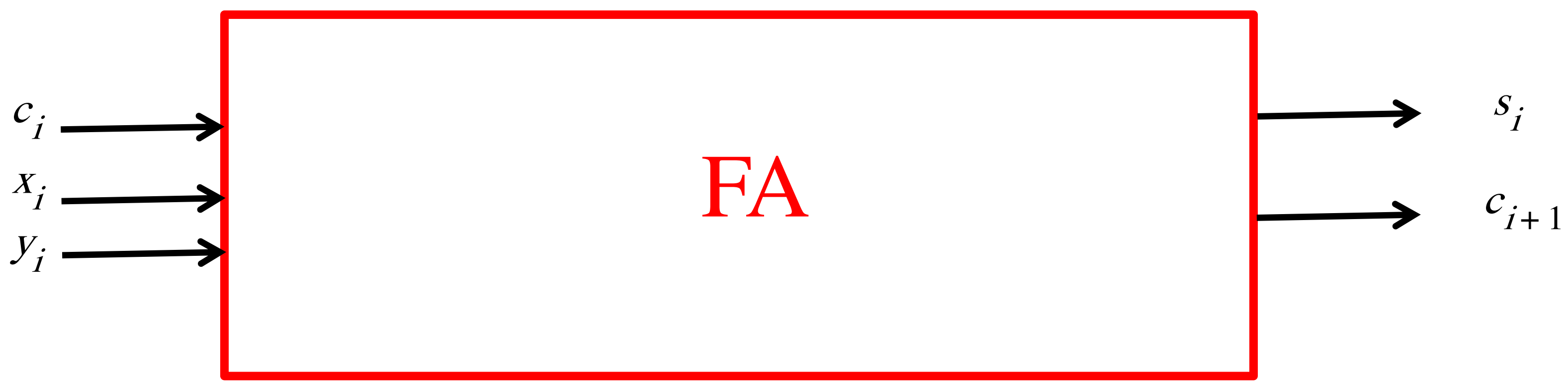
The Full-Adder Abstraction



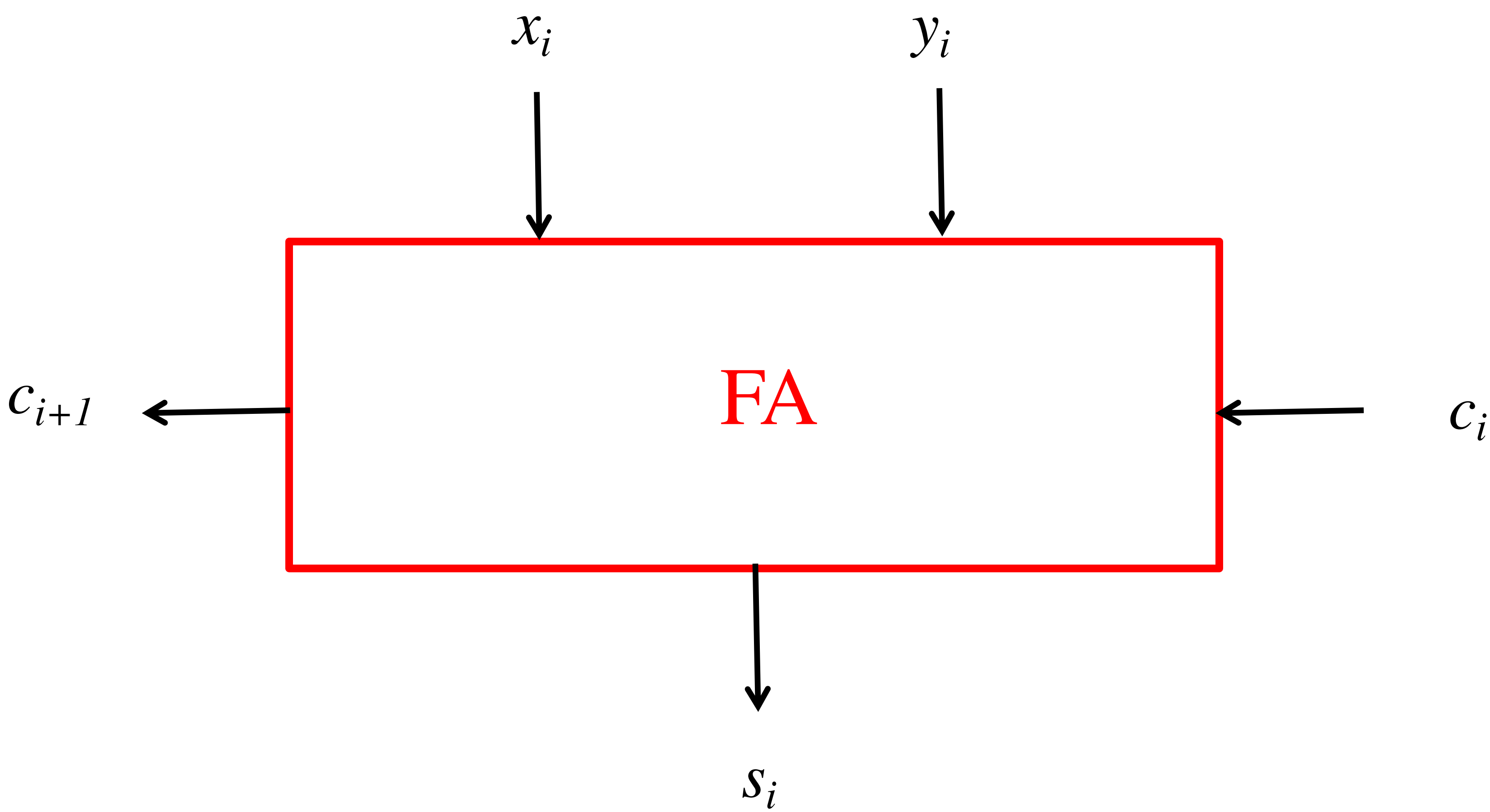
The Full-Adder Abstraction



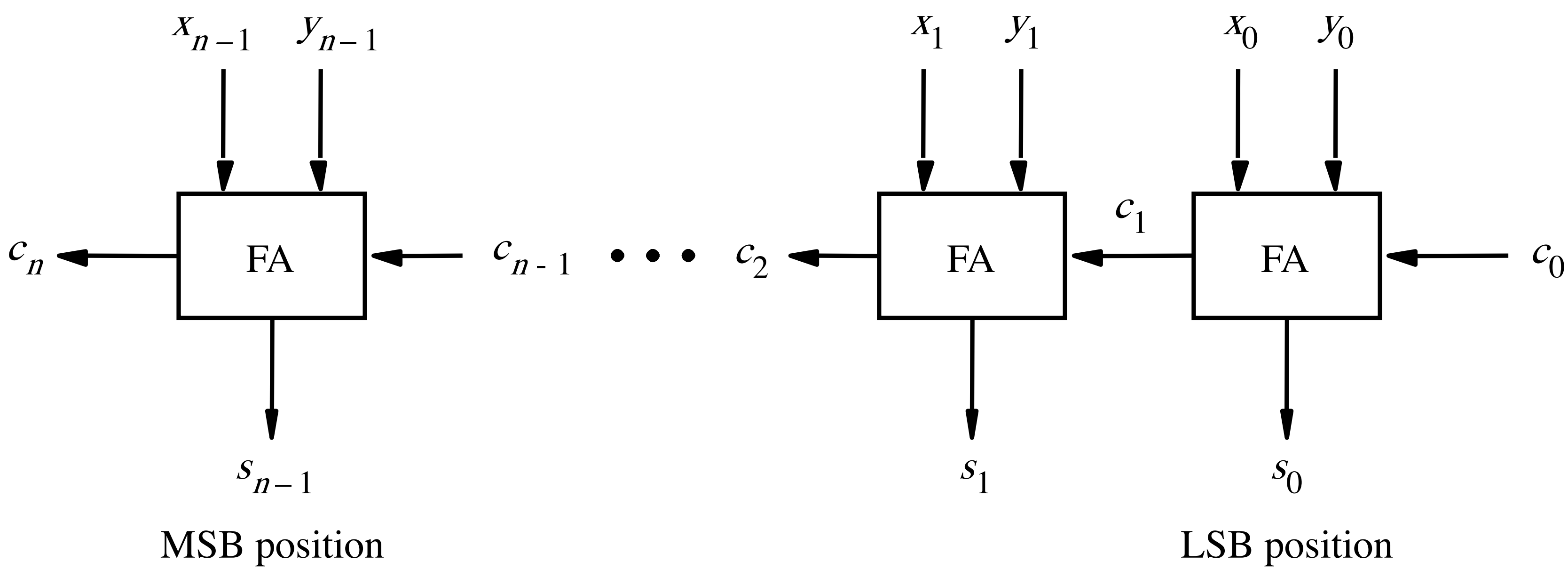
The Full-Adder Abstraction



We can place the arrows anywhere

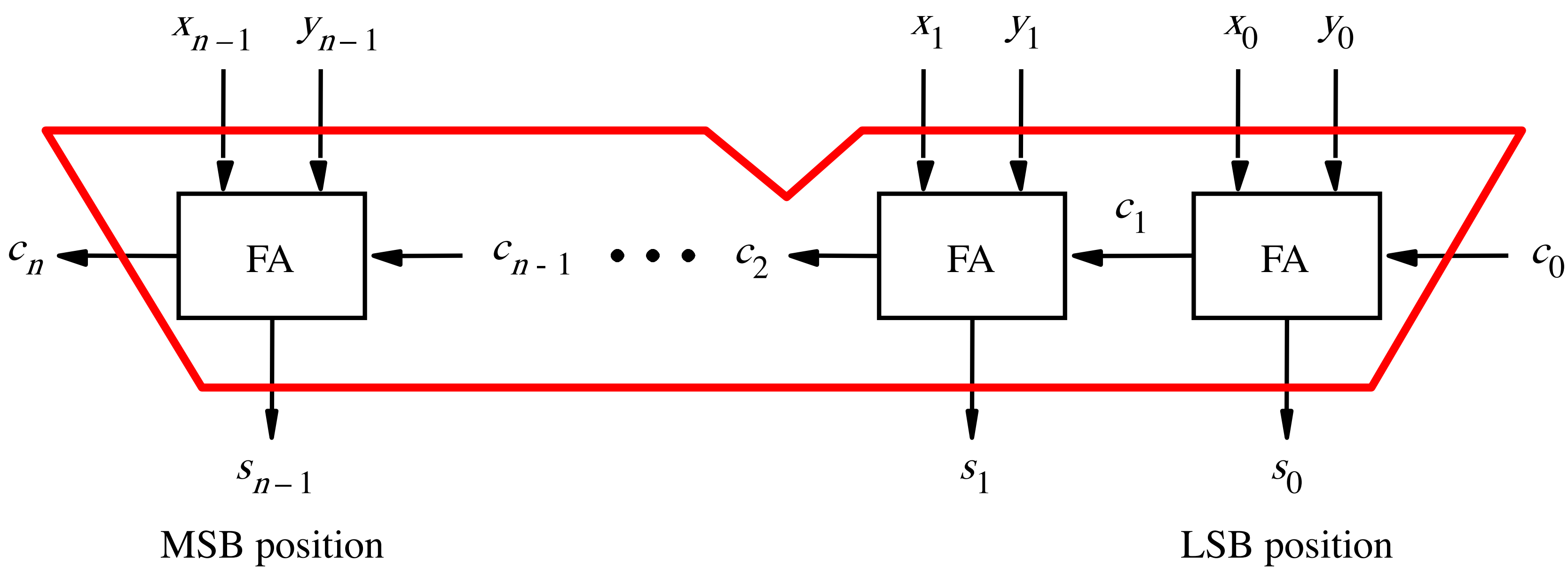


n-bit ripple-carry adder

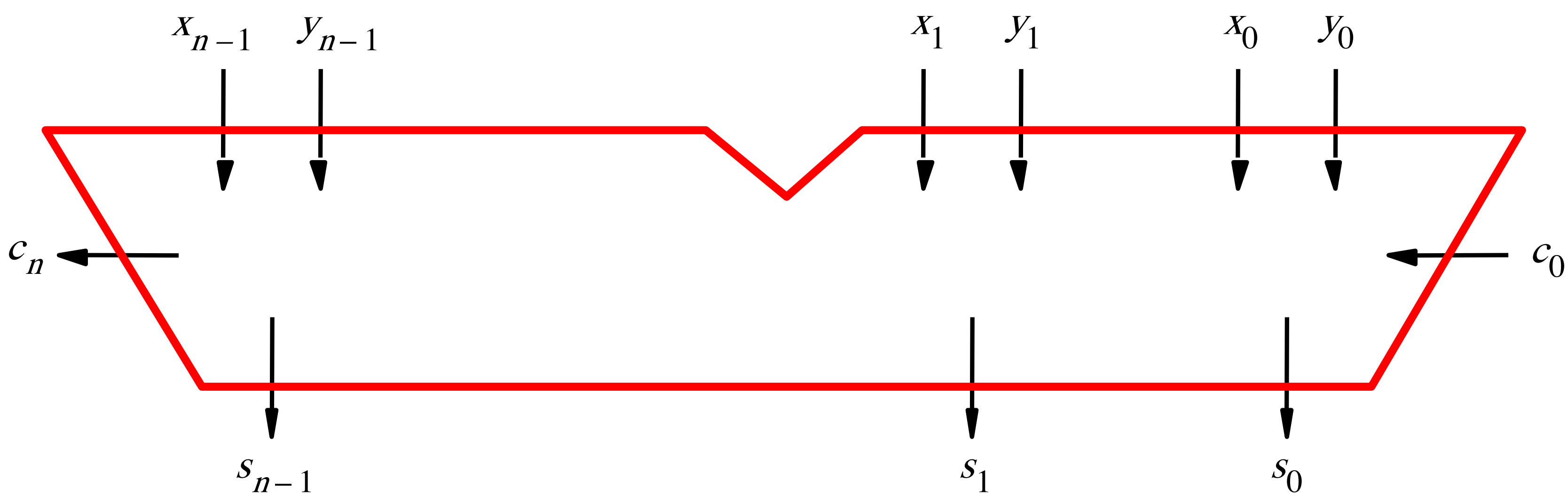


[Figure 3.5 from the textbook]

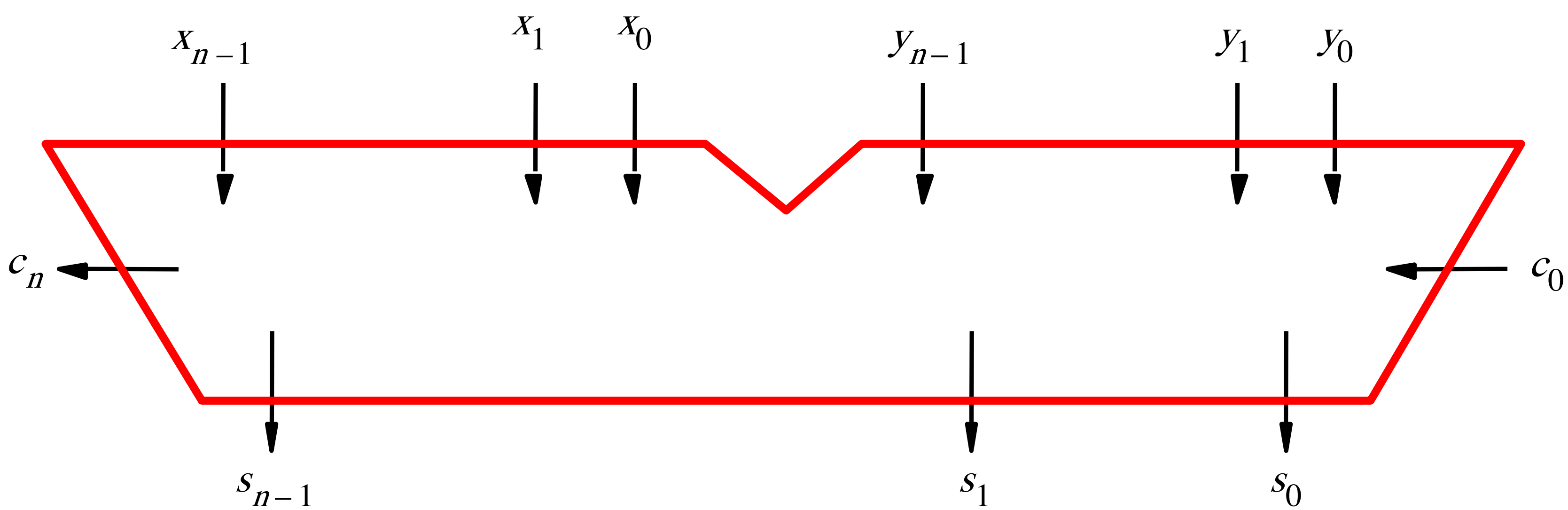
n-bit ripple-carry adder abstraction



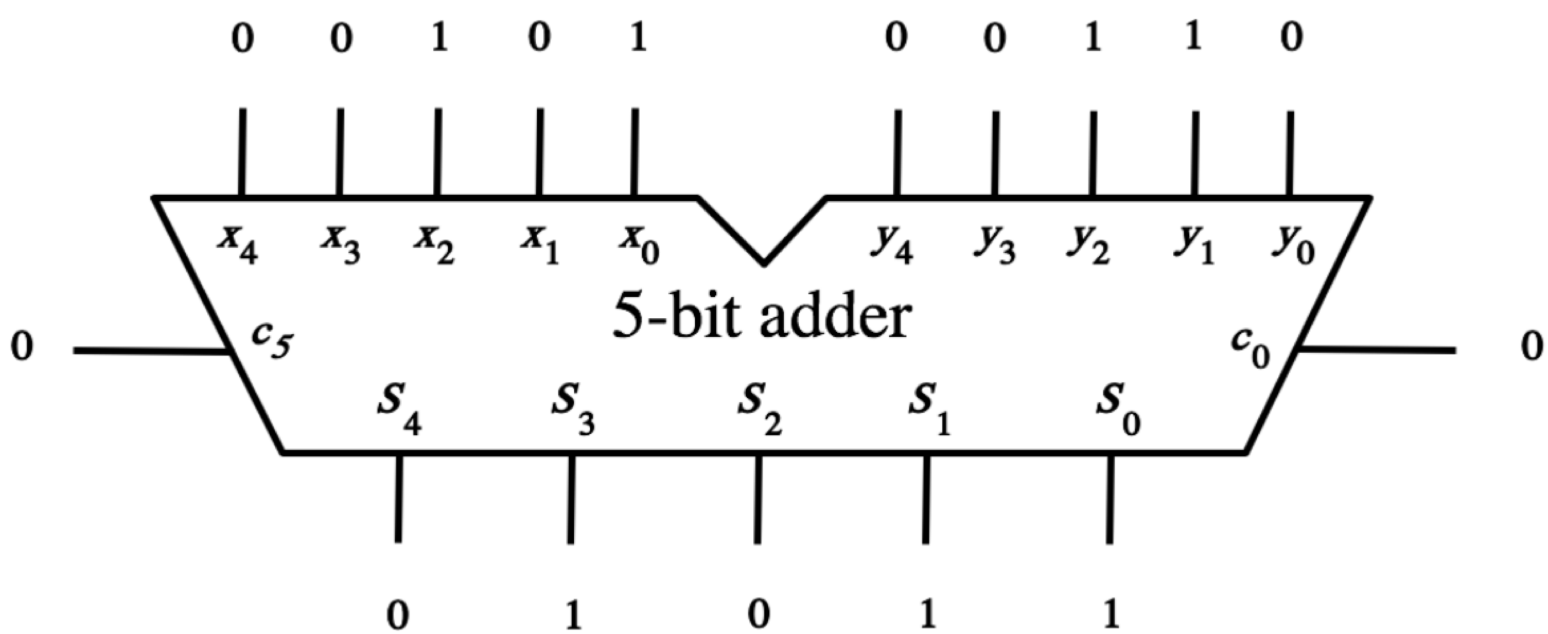
n-bit ripple-carry adder abstraction



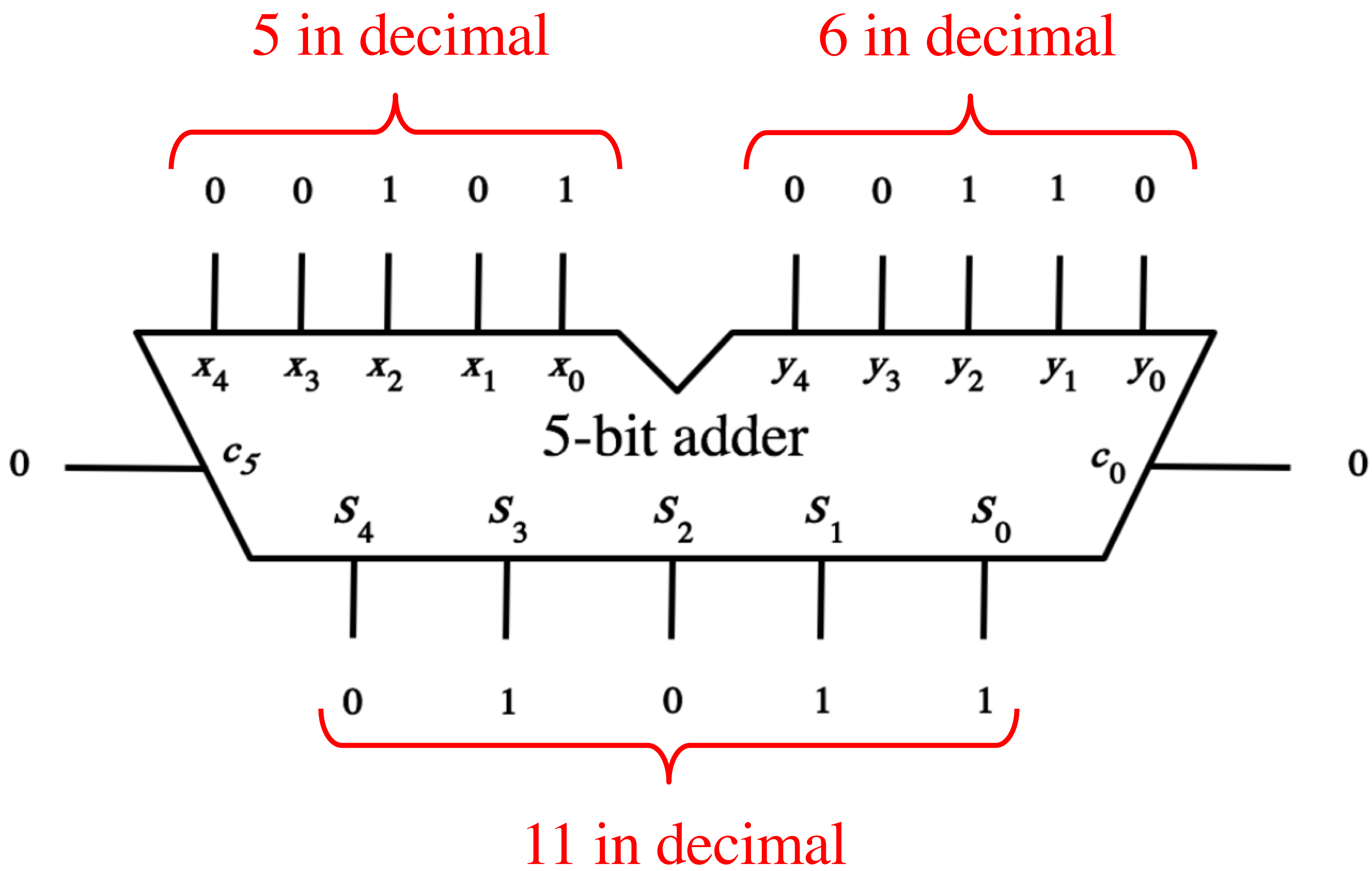
The x and y lines are typically grouped together for better visualization, but the underlying logic remains the same



Example:
Computing 5+6 using a 5-bit adder



Example:
Computing 5+6 using a 5-bit adder



Math Review: Subtraction

$$\begin{array}{r} 39 \\ -\ 15 \\ \hline ?? \end{array}$$

Math Review: Subtraction

$$\begin{array}{r} 39 \\ -\ 15 \\ \hline 24 \end{array}$$

Math Review: Subtraction

$$\begin{array}{r} 82 \\ -\ 61 \\ \hline ?? \end{array} \qquad \begin{array}{r} 48 \\ -\ 26 \\ \hline ?? \end{array} \qquad \begin{array}{r} 32 \\ -\ 11 \\ \hline ?? \end{array}$$

Math Review: Subtraction

$$\begin{array}{r} 82 \\ -\ 61 \\ \hline 21 \end{array} \qquad \begin{array}{r} 48 \\ -\ 26 \\ \hline 22 \end{array} \qquad \begin{array}{r} 32 \\ -\ 11 \\ \hline 21 \end{array}$$

Math Review: Subtraction

$$\begin{array}{r} 82 \\ -\ 64 \\ \hline ?? \end{array} \qquad \begin{array}{r} 48 \\ -\ 29 \\ \hline ?? \end{array} \qquad \begin{array}{r} 32 \\ -\ 13 \\ \hline ?? \end{array}$$

Math Review: Subtraction

$$\begin{array}{r} 82 \\ -\ 64 \\ \hline 18 \end{array} \qquad \begin{array}{r} 48 \\ -\ 29 \\ \hline 19 \end{array} \qquad \begin{array}{r} 32 \\ -\ 13 \\ \hline 19 \end{array}$$

The problems in which row are easier to calculate?

$$\begin{array}{r} 82 \\ -\ 61 \\ \hline ?? \end{array} \qquad \begin{array}{r} 48 \\ -\ 26 \\ \hline ?? \end{array} \qquad \begin{array}{r} 32 \\ -\ 11 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 82 \\ -\ 64 \\ \hline ?? \end{array} \qquad \begin{array}{r} 48 \\ -\ 29 \\ \hline ?? \end{array} \qquad \begin{array}{r} 32 \\ -\ 13 \\ \hline ?? \end{array}$$

The problems in which row are easier to calculate?

$$\begin{array}{r} 82 \\ -\ 61 \\ \hline 21 \end{array} \qquad \begin{array}{r} 48 \\ -\ 26 \\ \hline 22 \end{array} \qquad \begin{array}{r} 32 \\ -\ 11 \\ \hline 21 \end{array}$$

Why?

$$\begin{array}{r} 82 \\ -\ 64 \\ \hline 18 \end{array} \qquad \begin{array}{r} 48 \\ -\ 29 \\ \hline 19 \end{array} \qquad \begin{array}{r} 32 \\ -\ 13 \\ \hline 19 \end{array}$$

Another Way to Do Subtraction

$$82 - 64 = 82 + 100 - 100 - 64$$

Another Way to Do Subtraction

$$82 - 64 = 82 + 100 - 100 - 64$$

$$= 82 + (100 - 64) - 100$$

Another Way to Do Subtraction

$$82 – 64 = 82 + 100 – 100 - 64$$

$$= 82 + (100 – 64) - 100$$

$$= 82 + (99 + 1 – 64) - 100$$

Another Way to Do Subtraction

$$82 - 64 = 82 + 100 - 100 - 64$$

$$= 82 + (100 - 64) - 100$$

$$= 82 + (99 + 1 - 64) - 100$$

$$= 82 + (99 - 64) + 1 - 100$$

Another Way to Do Subtraction

$$82 - 64 = 82 + 100 - 100 - 64$$

$$= 82 + (100 - 64) - 100$$

$$= 82 + (99 + 1 - 64) - 100$$

Does not require borrows

$$= 82 + (99 - 64) + 1 - 100$$

9’s Complement
(subtract each digit from 9)

$$\begin{array}{r} 99 \\ -\ 64 \\ \hline 35 \end{array}$$

10’s Complement

(subtract each digit from 9 and add 1 to the result)

$$
\begin{array}{r@{\,}l}
 & 99 \\
- & 64 \\
\hline
 & 35 + 1 = 36
\end{array}
$$

Another Way to Do Subtraction

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

$$= 82 + 35 + 1 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

10's complement

$$= 82 + 35 + 1 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

10's complement

$$= 82 + 35 + 1 - 100$$

$$= 82 + 36 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

10's complement

$$= 82 + 35 + 1 - 100$$

$$= 82 + 36 - 100$$ // Add the first two.

$$= 118 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

10's complement

$$= 82 + 35 + 1 - 100$$

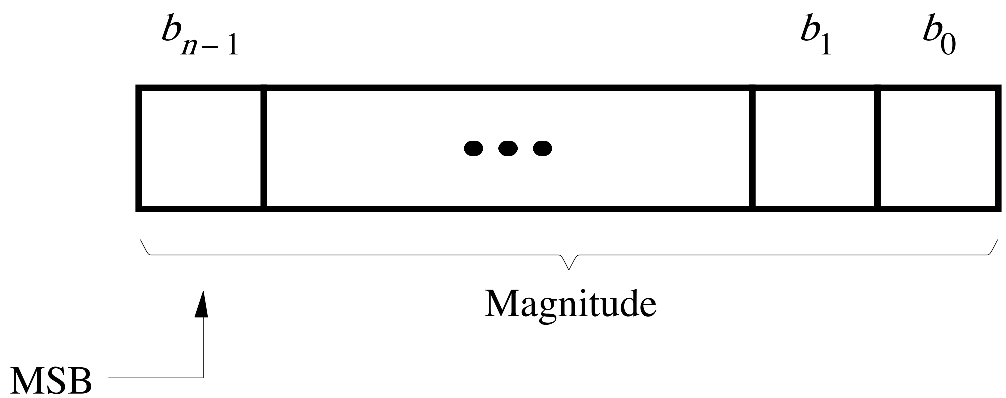
$$= 82 + 36 - 100$$ // Add the first two.

$$= 118 - 100$$ // Just delete the leading 1.

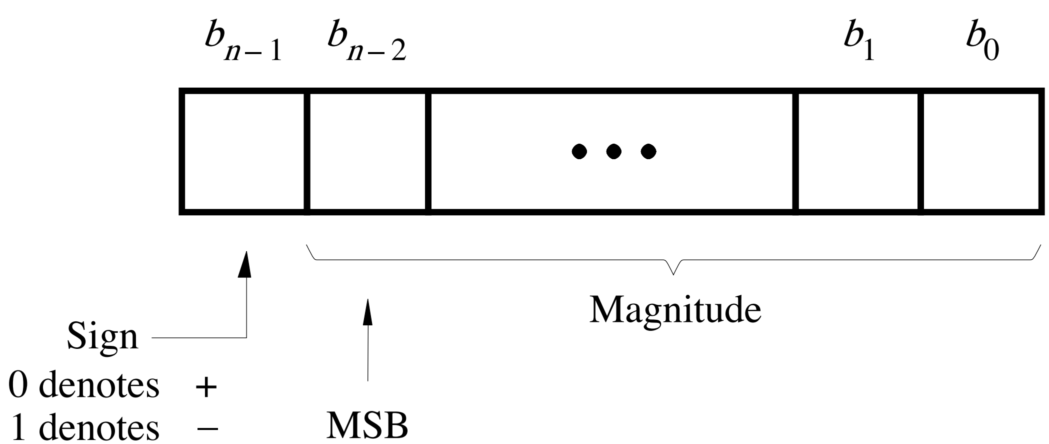
// No need to subtract 100.

$$= 18$$

Formats for representation of integers



(a) Unsigned number



(b) Signed number

[Figure 3.7 from the textbook]

Unsigned Representation

2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
0	0	1	0	1	1	0	0

This represents + 44.

Unsigned Representation

2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
1	0	1	0	1	1	0	0

This represents + 172.

Three Different Ways to Represent Negative Integer Numbers

- **Sign and magnitude**
- **1's complement**
- **2's complement**

Three Different Ways to Represent Negative Integer Numbers

- **Sign and magnitude**
- **1's complement**
- $\boxed{\textbf{2's complement}}$

only this method is used
in modern computers

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

[Table 3.2 from the textbook]

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

The top half is the same in all three representations.
It corresponds to the positive integers.

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

In all three representations the first bit represents the sign.
If that bit is 1, then the number is negative.

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

Notice that in this representation there are two zeros!

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

There are two zeros in this representation as well!

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

In this representation there is one more negative number.

Sign and Magnitude

Sign and Magnitude Representation
(using the left-most bit as the sign)

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0
0	0	1	0	1	1	0	0

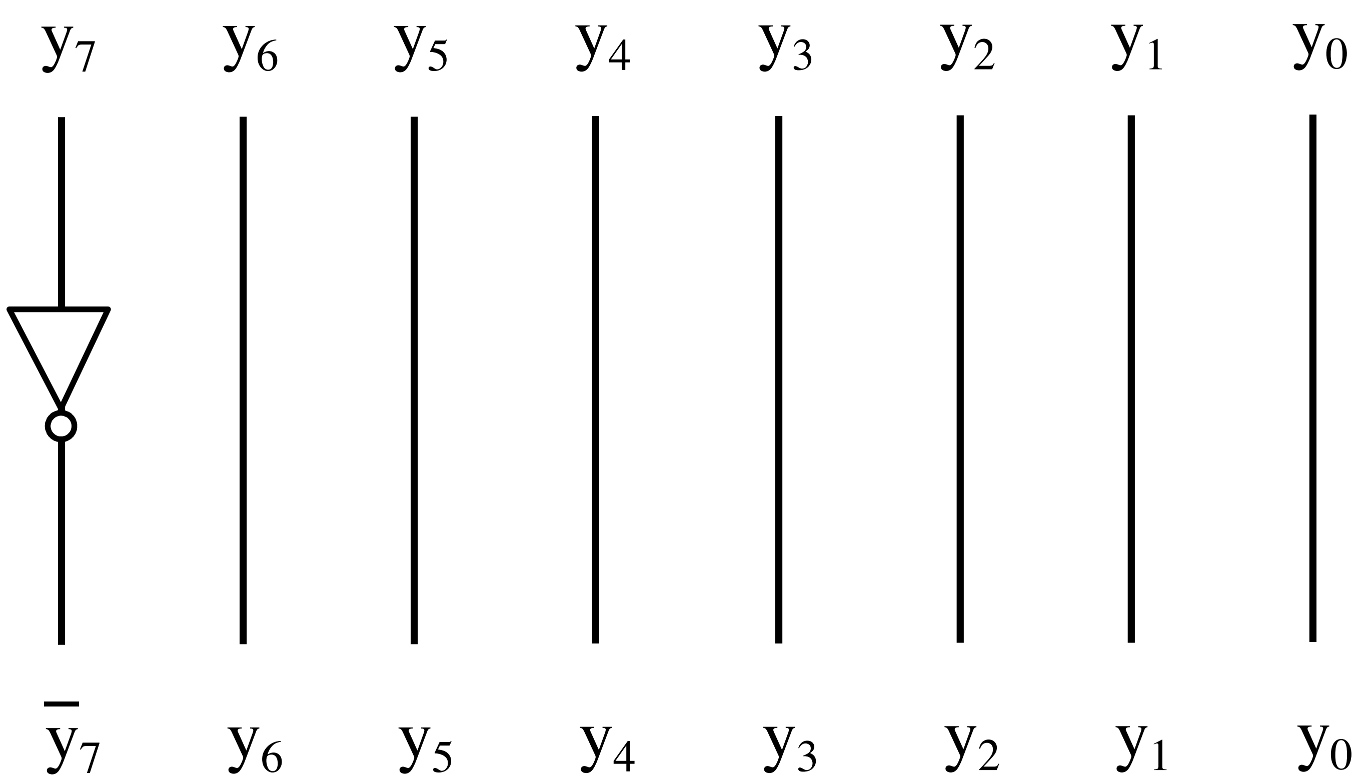
This represents + 44.

Sign and Magnitude Representation
(using the left-most bit as the sign)

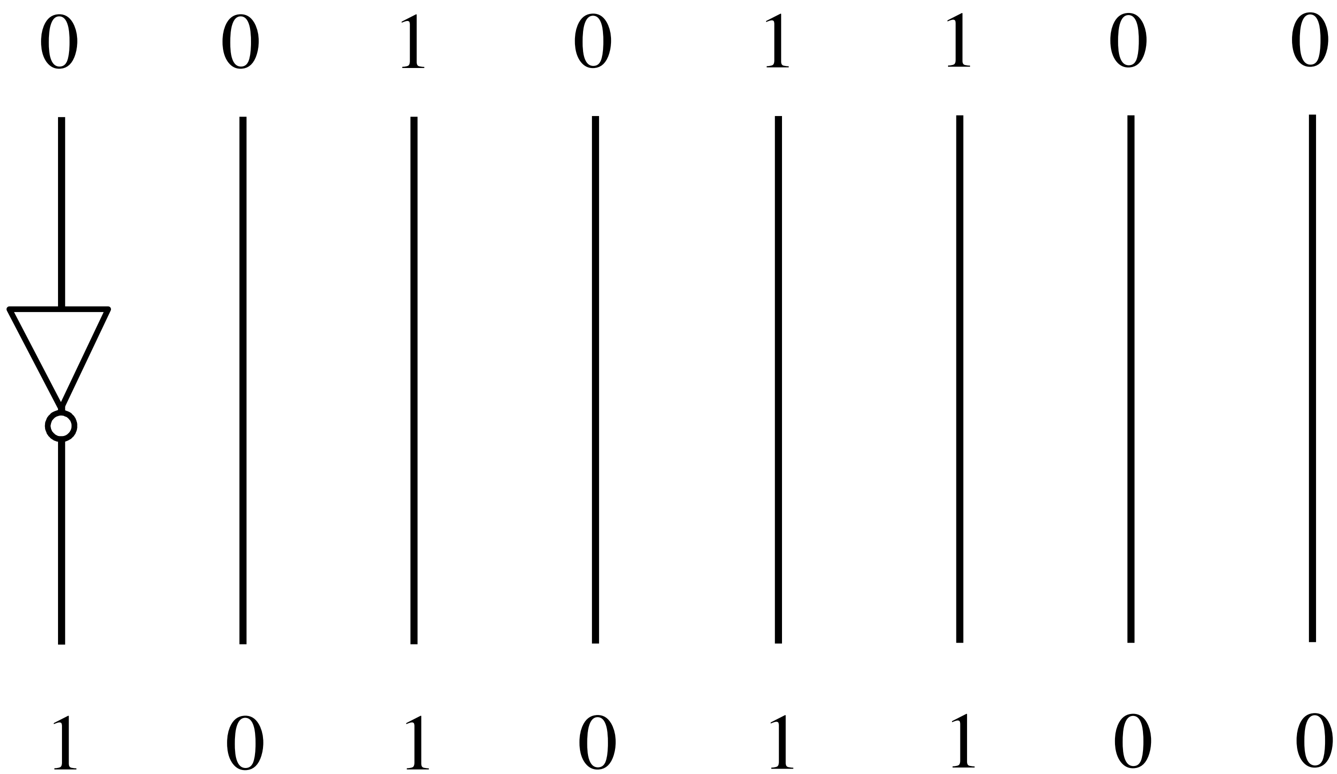
sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0
1	0	1	0	1	1	0	0

This represents – 44.

Circuit for negating a number stored in sign and magnitude representation



Circuit for negating a number stored in sign and magnitude representation



1’s Complement

1's complement
(subtract each digit from 1)

Let K be the negative equivalent of an n-bit positive number P.

Then, in 1's complement representation K is obtained by subtracting P from $2^n - 1$, namely

$$K = (2^n - 1) - P$$

This means that K can be obtained by inverting all bits of P.

1's complement
(subtract each digit from 1)

Let K be the negative equivalent of an 8-bit positive number P.

Then, in 1's complement representation K is obtained by subtracting P from $2^8 - 1$, namely

$$K = (2^8 - 1) - P = 255 - P$$

This means that K can be obtained by inverting all bits of P.

Provided that P is between 0 and 127, because the most significant bit must be zero to indicate that it is positive.

1's Complement Representation

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0
0	0	1	0	1	1	0	0

1's Complement Representation

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0
0	0	1	0	1	1	0	0

$+\,44$

$$2^5 + 2^3 + 2^2 = 44$$

1's Complement Representation

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0
0	0	1	0	1	1	0	0

+ 44

+ 44 in 1's complement representation

1's Complement Representation
(invert all the bits to negate the number)

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	0	1	1	- 44

1's Complement Representation
(invert all the bits to negate the number)

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

positive

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	0	1	1	- 44

negative

1's Complement Representation
(invert all the bits to negate the number)

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	0	1	1	- 44

$2^7 + 2^6 + 2^4 + 2^1 + 2^0 = 211$ (as unsigned)

1's Complement Representation
(invert all the bits to negate the number)

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	0	1	1	- 44

211 = 255 – 44 (as unsigned)

1's Complement Representation
(invert all the bits to negate the number)

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	0	1	1	- 44

- 44 in 1's complement representation

1's complement
(subtract each digit from 1)

$$
\begin{array}{r@{\quad}cccccccc}
 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
- & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 \\
\hline
 & 1 & 1 & 0 & 1 & 0 & 0 & 1 & 1
\end{array}
$$

1's complement
(subtract each digit from 1)

No need to borrow!

$$
\begin{array}{rcccccccc}
 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
- & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 \\
\hline
 & 1 & 1 & 0 & 1 & 0 & 0 & 1 & 1
\end{array}
$$

1’s complement
(subtract each digit from 1)

$$\begin{array}{rcl} & \boxed{1\quad 1\quad 1\quad 1\quad 1\quad 1\quad 1\quad 1} & 255 \\ - & \boxed{0\quad 0\quad 1\quad 0\quad 1\quad 1\quad 0\quad 0} & 44 \\ \hline & 1\quad 1\quad 0\quad 1\quad 0\quad 0\quad 1\quad 1 & \end{array}$$

1's complement
(subtract each digit from 1)

$$
\begin{array}{r}
1\ 1\ 1\ 1\ 1\ 1\ 1\ 1 \\
-\ \ 0\ 0\ 1\ 0\ 1\ 1\ 0\ 0 \\
\hline
\boxed{1\ 1\ 0\ 1\ 0\ 0\ 1\ 1}
\end{array}
\quad 211
$$

211 = 255 – 44 (as unsigned)

1's complement
(subtract each digit from 1)

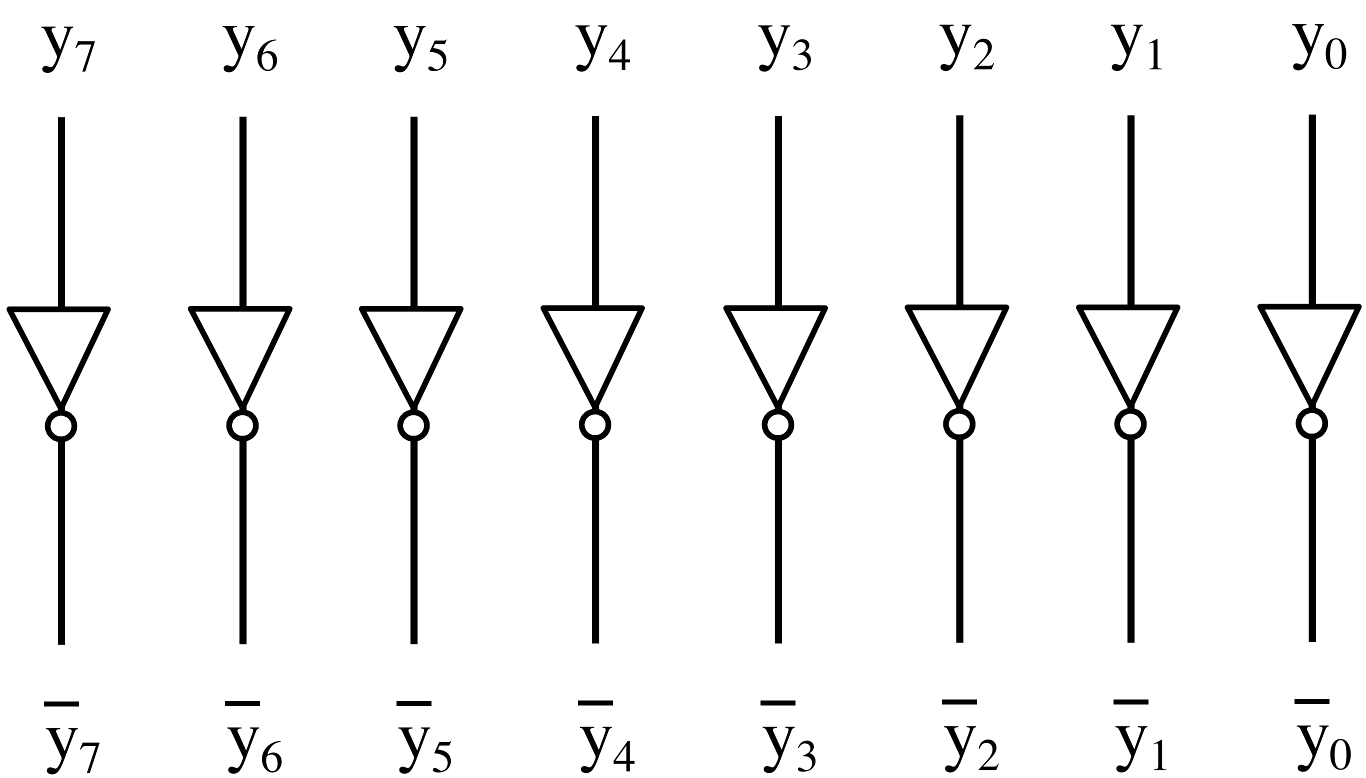
$$
\begin{array}{rcccccccc}
 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
- & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 \\
\hline
 & \boxed{1} & \boxed{1} & \boxed{0} & \boxed{1} & \boxed{0} & \boxed{0} & \boxed{1} & \boxed{1}
\end{array}
\quad - 44
$$

211 = 255 – 44 (as unsigned)

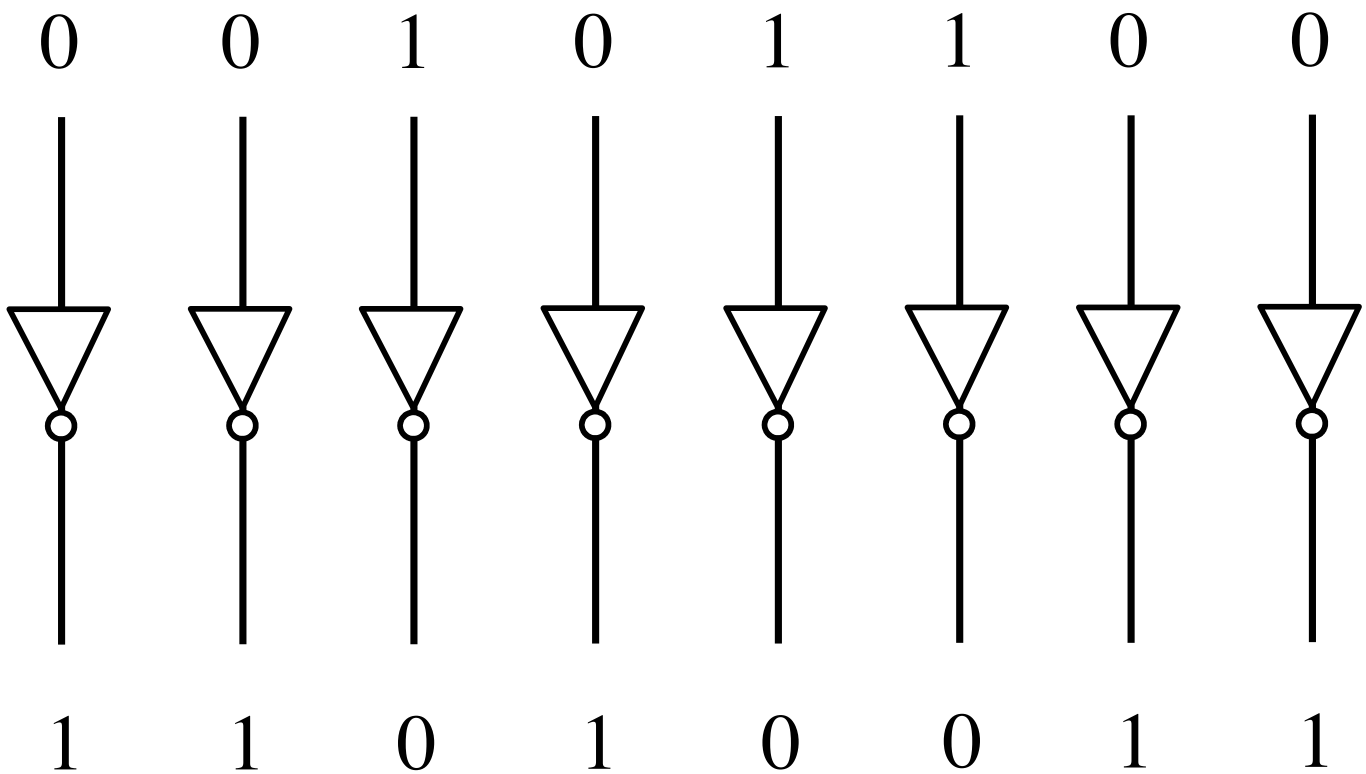
or

- 44 in 1's complement representation

Circuit for negating a number stored in 1's complement representation



Circuit for negating a number stored in 1's complement representation



This works in reverse too
(from negative to positive)

1's Complement Representation

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0
1	1	0	1	0	0	1	1

- 44

1's Complement Representation
(invert all the bits to negate the number)

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	0	1	1	- 44

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

1's Complement Representation
(invert all the bits to negate the number)

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	0	1	1	- 44

211 (as unsigned)

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

44 = 255 – 211 (as unsigned)

1’s Complement Representation
(invert all the bits to negate the number)

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	0	1	1	- 44

- 44 in 1’s complement representation

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

+ 44 in 1’s complement representation

Negate these numbers stored in 1's complement representation

0 1 0 1

1 0 1 1

1 1 1 0

0 1 1 1

Negate these numbers stored in 1's complement representation

0 1 0 1	1 0 1 1
1 0 1 0	0 1 0 0
1 1 1 0	0 1 1 1
0 0 0 1	1 0 0 0

Just flip 1's to 0's and vice versa.

Negate these numbers stored in 1's complement representation

0 1 0 1 = +5

1 0 1 0 = -5

1 0 1 1 = -4

0 1 0 0 = +4

1 1 1 0 = -1

0 0 0 1 = +1

0 1 1 1 = +7

1 0 0 0 = -7

Just flip 1's to 0's and vice versa.

Addition of two numbers stored in 1's complement representation

There are four cases to consider

- $(+5) \quad + \quad (+2)$
- $(-5) \quad + \quad (+2)$
- $(+5) \quad + \quad (-2)$
- $(-5) \quad + \quad (-2)$

There are four cases to consider

- `(+5)  +  (+2)      positive plus positive`
- `(–5)  +  (+2)      negative plus positive`
- `(+5)  +  (–2)      positive plus negative`
- `(–5)  +  (–2)      negative plus negative`

A) Example of 1's complement addition

$$\begin{array}{rr} (+5) & 0\,1\,0\,1 \\ +\,(+2) & +\,0\,0\,1\,0 \\ \hline (+7) & 0\,1\,1\,1 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

[Figure 3.8 from the textbook]

A) Example of 1's complement addition

$$\begin{array}{r} (+5) \\ +\,(+2) \\ \hline (+7) \end{array} \qquad \begin{array}{r} 0\,1\,0\,1 \\ +\,0\,0\,1\,0 \\ \hline 0\,1\,1\,1 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

B) Example of 1's complement addition

$$\begin{array}{r} (-5) \\ +\,(+2) \\ \hline (-3) \end{array} \qquad \begin{array}{r} 1\,0\,1\,0 \\ +\,0\,0\,1\,0 \\ \hline 1\,1\,0\,0 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

[Figure 3.8 from the textbook]

B) Example of 1's complement addition

$$\begin{array}{rr} (-\,5) & 1\,0\,1\,0 \\ +\,(+2) & +\,0\,0\,1\,0 \\ \hline (-\,3) & 1\,1\,0\,0 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

C) Example of 1's complement addition

$$
\begin{array}{r}
(+5) \\
+\,(-2) \\
\hline
(+3)
\end{array}
\qquad
\begin{array}{r}
0\,1\,0\,1 \\
+\,1\,1\,0\,1 \\
\hline
1\,0\,0\,1\,0
\end{array}
$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

[Figure 3.8 from the textbook]

C) Example of 1's complement addition

$$\begin{array}{r} (+5) \\ +(-2) \\ \hline (+3) \end{array} \qquad \begin{array}{r} 0\,1\,0\,1 \\ +\,1\,1\,0\,1 \\ \hline 1\,0\,0\,1\,0 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

C) Example of 1's complement addition

	$(+5)$		$0\ 1\ 0\ 1$
$+$	(-2)	$+$	$1\ 1\ 0\ 1$
	$(+3)$	1	$0\ 0\ 1\ 0$

But this is 2!

$b_3 b_2 b_1 b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

C) Example of 1's complement addition

$$\begin{array}{r} (+5) \\ +(-2) \\ \hline (+3) \end{array} \qquad \begin{array}{r} 0\,1\,0\,1 \\ +\,1\,1\,0\,1 \\ \hline 1\,0\,0\,1\,0 \\ 1 \\ \hline 0\,0\,1\,1 \end{array}$$

We need to perform one more addition to get the result.

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

C) Example of 1's complement addition

$$\begin{array}{rr} (+5) & 0\,1\,0\,1 \\ +(-2) & +\,1\,1\,0\,1 \\ \hline (+3) & 1\;\boxed{0\,0\,1\,0} \\ & 1 \\ \hline & 0\,0\,1\,1 \end{array}$$

We need to perform one more addition to get the result.

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

D) Example of 1's complement addition

$$\begin{array}{r} (-5) \\ +\ (-2) \\ \hline (-7) \end{array} \qquad \begin{array}{r} 1\,0\,1\,0 \\ +\ 1\,1\,0\,1 \\ \hline 1\,0\,1\,1\,1 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

[Figure 3.8 from the textbook]

D) Example of 1's complement addition

$$\begin{array}{r} (-5) \\ +\,(-2) \\ \hline (-7) \end{array} \qquad \begin{array}{r} 1\,0\,1\,0 \\ +\,1\,1\,0\,1 \\ \hline 1\;0\,1\,1\,1 \end{array}$$

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

D) Example of 1's complement addition

$$\begin{array}{r} (-5) \\ +\ (-2) \\ \hline (-7) \end{array} \qquad \begin{array}{r} 1\,0\,1\,0 \\ +\ 1\,1\,0\,1 \\ \hline 1\ 0\,1\,1\,1 \end{array}$$

But this is +7!

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

D) Example of 1's complement addition

$$\begin{array}{r} (-5) \\ +\,(-2) \\ \hline (-7) \end{array} \qquad \begin{array}{r} 1\,0\,1\,0 \\ +\,1\,1\,0\,1 \\ \hline 1\,0\,1\,1\,1 \\ 1 \\ \hline 1\,0\,0\,0 \end{array}$$

We need to perform one more addition to get the result.

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

D) Example of 1's complement addition

$$\begin{array}{r} (-5) \\ +\,(-2) \\ \hline (-7) \end{array} \qquad \begin{array}{r} 1\,0\,1\,0 \\ +\;1\,1\,0\,1 \\ \hline 1\;0\,1\,1\,1 \\ 1 \\ \hline 1\,0\,0\,0 \end{array}$$

We need to perform one more addition to get the result.

$b_3b_2b_1b_0$	1's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−7
1001	−6
1010	−5
1011	−4
1100	−3
1101	−2
1110	−1
1111	−0

Implications for arithmetic operations in 1's complement representation

- We could do addition in 1's complement, but the circuit will need to handle these exceptions.
- In some cases it will run faster that others, thus creating uncertainties in the timing.
- Therefore, 1's complement is not used in practice to do arithmetic operations.
- But it may show up as an intermediary step in doing 2's complement operations.

2’s Complement

2's complement
(subtract each digit from 1 and add 1 to the result)

Let K be the negative equivalent of an n-bit positive number P.

Then, in 2's complement representation K is obtained by subtracting P from 2^n , namely

$$K = 2^n - P$$

2’s complement
(subtract each digit from 1 and add 1 to the result)

Let K be the negative equivalent of an 8-bit positive number P.

Then, in 2’s complement representation K is obtained by subtracting P from 2^8, namely

$$K = 2^8 - P = 256 - P$$

2's Complement Representation

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

2's Complement Representation

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	1	0	0	- 44

2’s Complement Representation

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44
positive								

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	1	0	0	- 44
negative								

2's Complement Representation

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
0	0	1	0	1	1	0	0	+ 44

sign	2^6	2^5	2^4	2^3	2^2	2^1	2^0	
1	1	0	1	0	1	0	0	- 44

212 = 256 - 44

Deriving 2's complement

For a positive n-bit number P, let K_1 and K_2 denote its 1's and 2's complements, respectively.

$$K_1 = (2^n - 1) - P$$

$$K_2 = 2^n - P$$

Since $K_2 = K_1 + 1$, it is evident that in a logic circuit the 2's complement can be computed by inverting all bits of P and then adding 1 to the resulting 1's-complement number.

Deriving 2's complement

For a positive 8-bit number P, let K_1 and K_2 denote its 1's and 2's complements, respectively.

$$K_1 = (2^n - 1) - P = 255 - P$$

$$K_2 = 2^n - P = 256 - P$$

Since $K_2 = K_1 + 1$, it is evident that in a logic circuit the 2's complement can be computed by inverting all bits of P and then adding 1 to the resulting 1's-complement number.

Negate these numbers stored in 2's complement representation

0 1 0 1

1 1 1 0

1 1 0 0

0 1 1 1

Negate these numbers stored in 2's complement representation

0 1 0 1

1 0 1 0

1 1 1 0

0 0 0 1

1 1 0 0

0 0 1 1

0 1 1 1

1 0 0 0

Invert all bits…

Negate these numbers stored in 2's complement representation

$$\begin{array}{r} 0\,1\,0\,1 \\ 1\,0\,1\,0 \\ +\quad\quad 1 \\ \hline 1\,0\,1\,1 \end{array}$$

$$\begin{array}{r} 1\,1\,1\,0 \\ 0\,0\,0\,1 \\ +\quad\quad 1 \\ \hline 0\,0\,1\,0 \end{array}$$

$$\begin{array}{r} 1\,1\,0\,0 \\ 0\,0\,1\,1 \\ +\quad\quad 1 \\ \hline 0\,1\,0\,0 \end{array}$$

$$\begin{array}{r} 0\,1\,1\,1 \\ 1\,0\,0\,0 \\ +\quad\quad 1 \\ \hline 1\,0\,0\,1 \end{array}$$

.. then add 1.

Negate these numbers stored in 2's complement representation

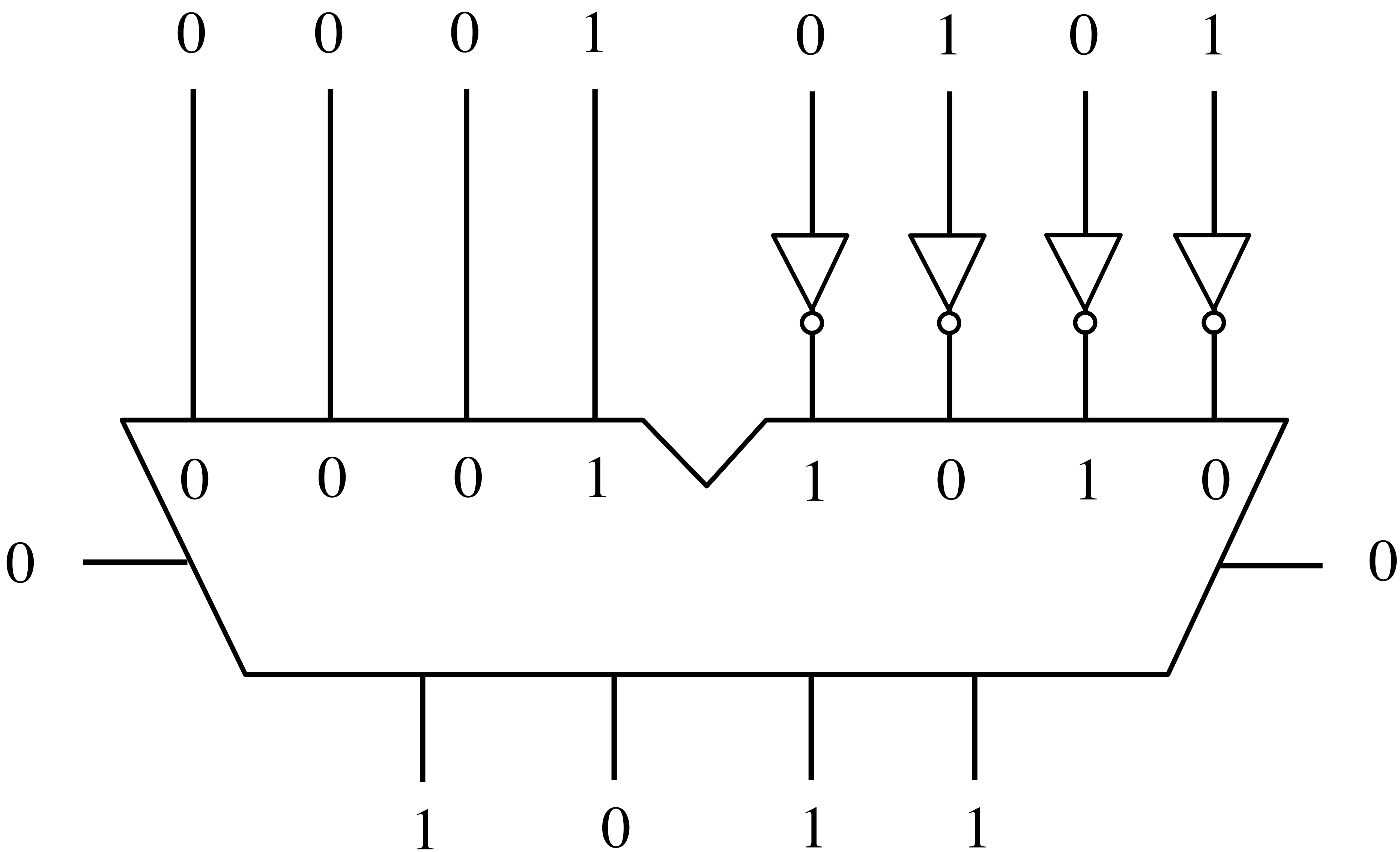
$$0\,1\,0\,1 = +5$$
$$\begin{array}{r} 1\,0\,1\,0 \\ +\quad 1 \\ \hline 1\,0\,1\,1 \end{array} = -5$$

$$1\,1\,1\,0 = -2$$
$$\begin{array}{r} 0\,0\,0\,1 \\ +\quad 1 \\ \hline 0\,0\,1\,0 \end{array} = +2$$

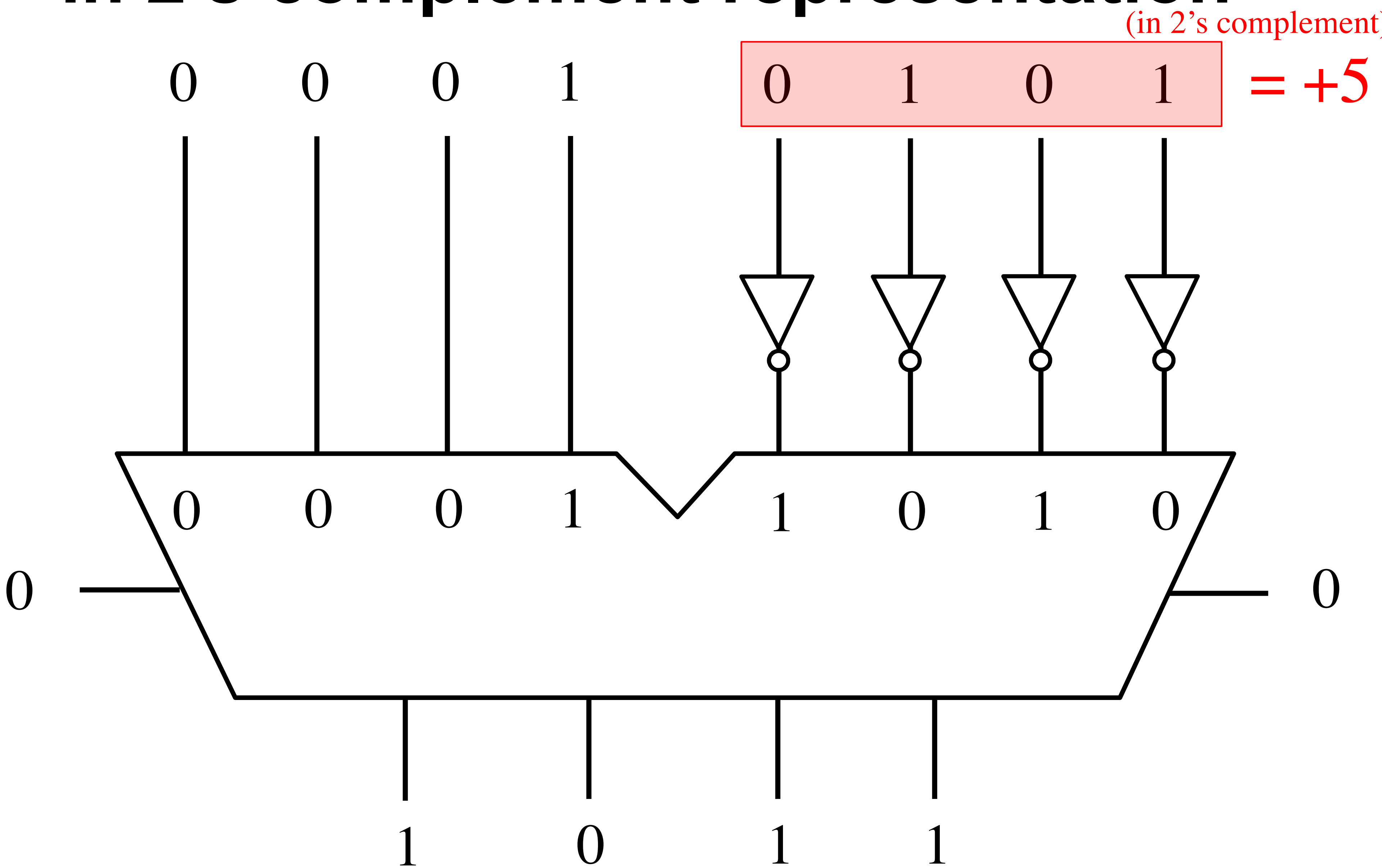
$$1\,1\,0\,0 = -4$$
$$\begin{array}{r} 0\,0\,1\,1 \\ +\quad 1 \\ \hline 0\,1\,0\,0 \end{array} = +4$$

$$0\,1\,1\,1 = +7$$
$$\begin{array}{r} 1\,0\,0\,0 \\ +\quad 1 \\ \hline 1\,0\,0\,1 \end{array} = -7$$

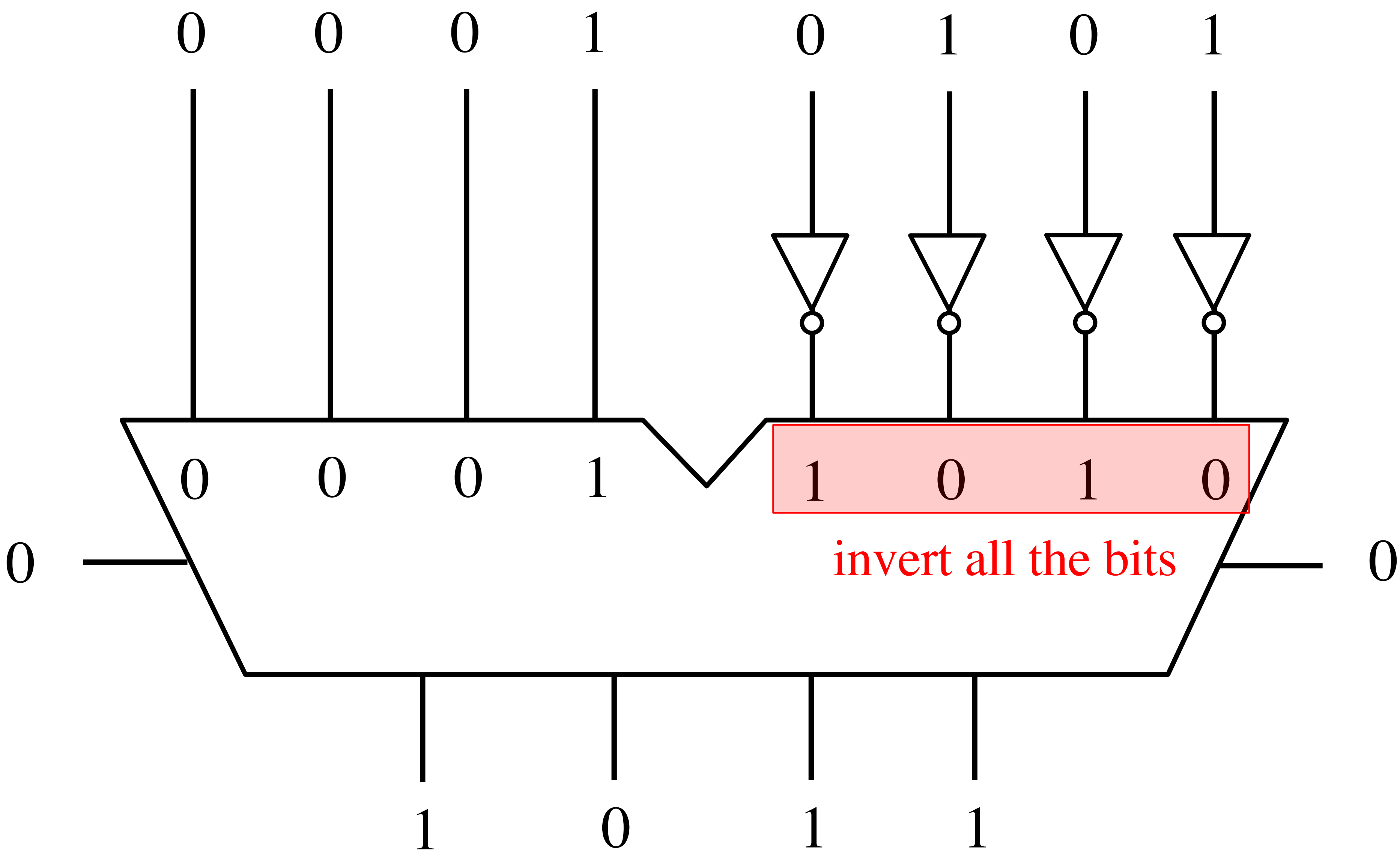
Circuit #1 for negating a number stored in 2's complement representation



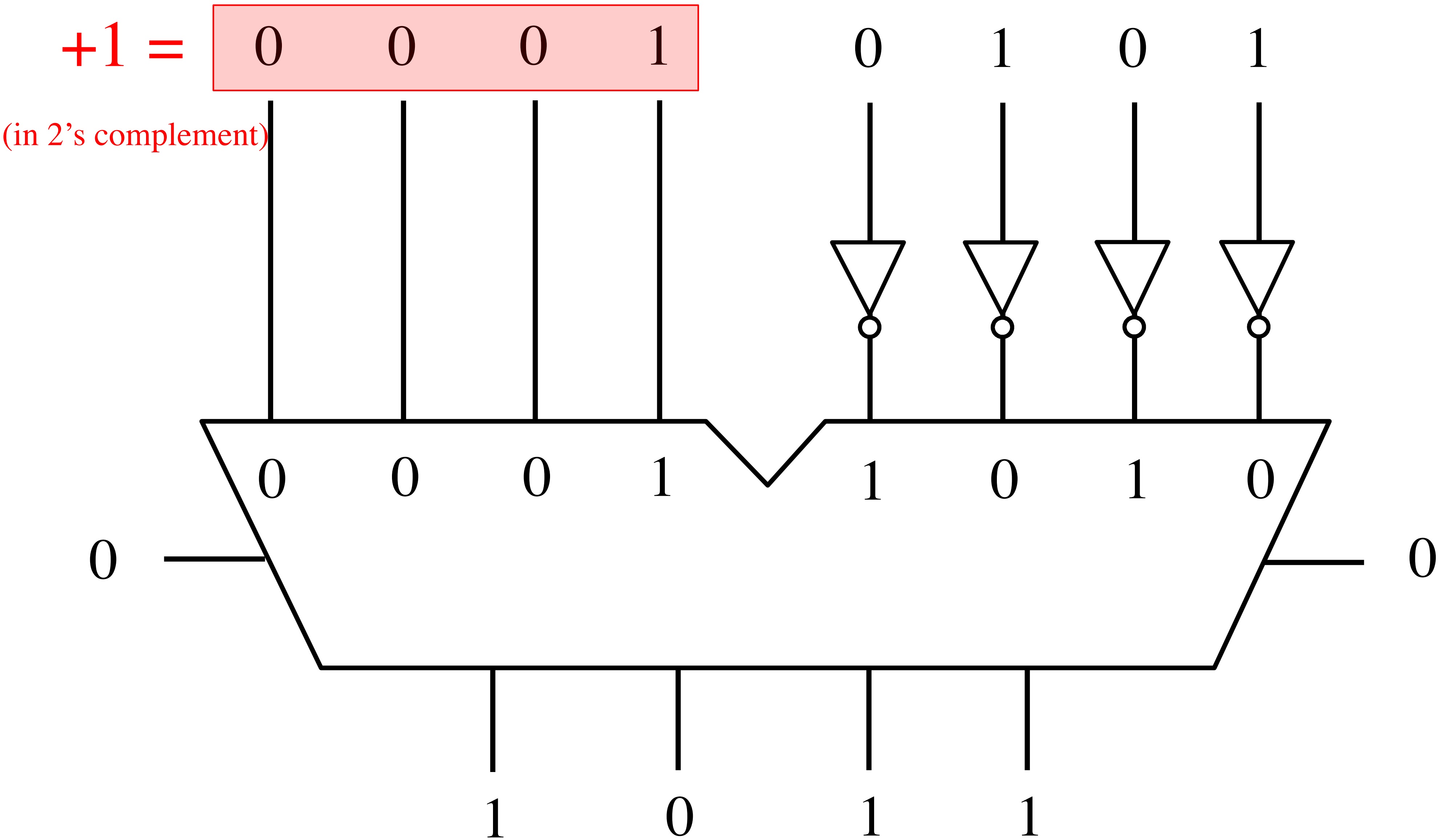
Circuit #1 for negating a number stored in 2's complement representation



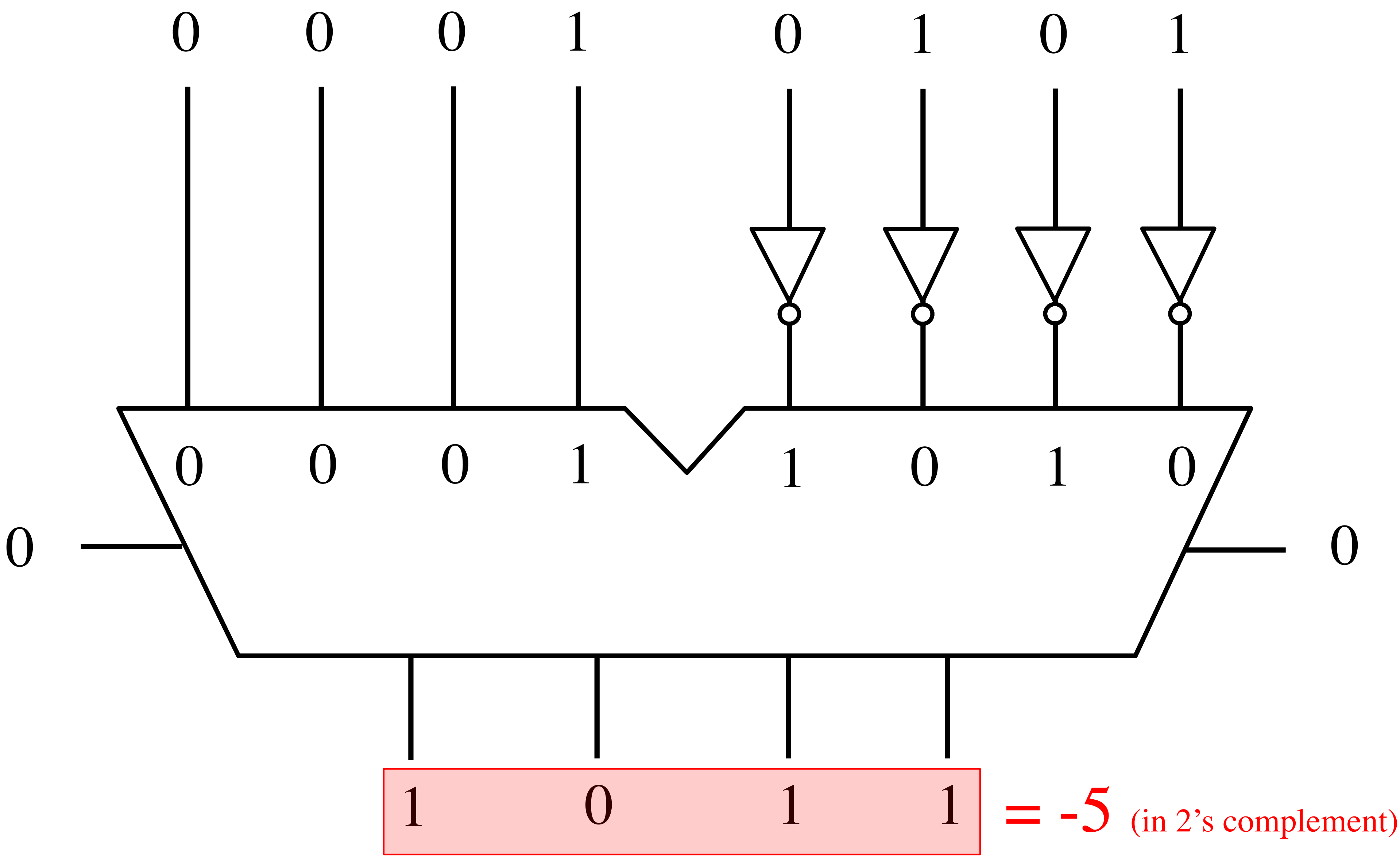
Circuit #1 for negating a number stored in 2's complement representation



Circuit #1 for negating a number stored in 2's complement representation

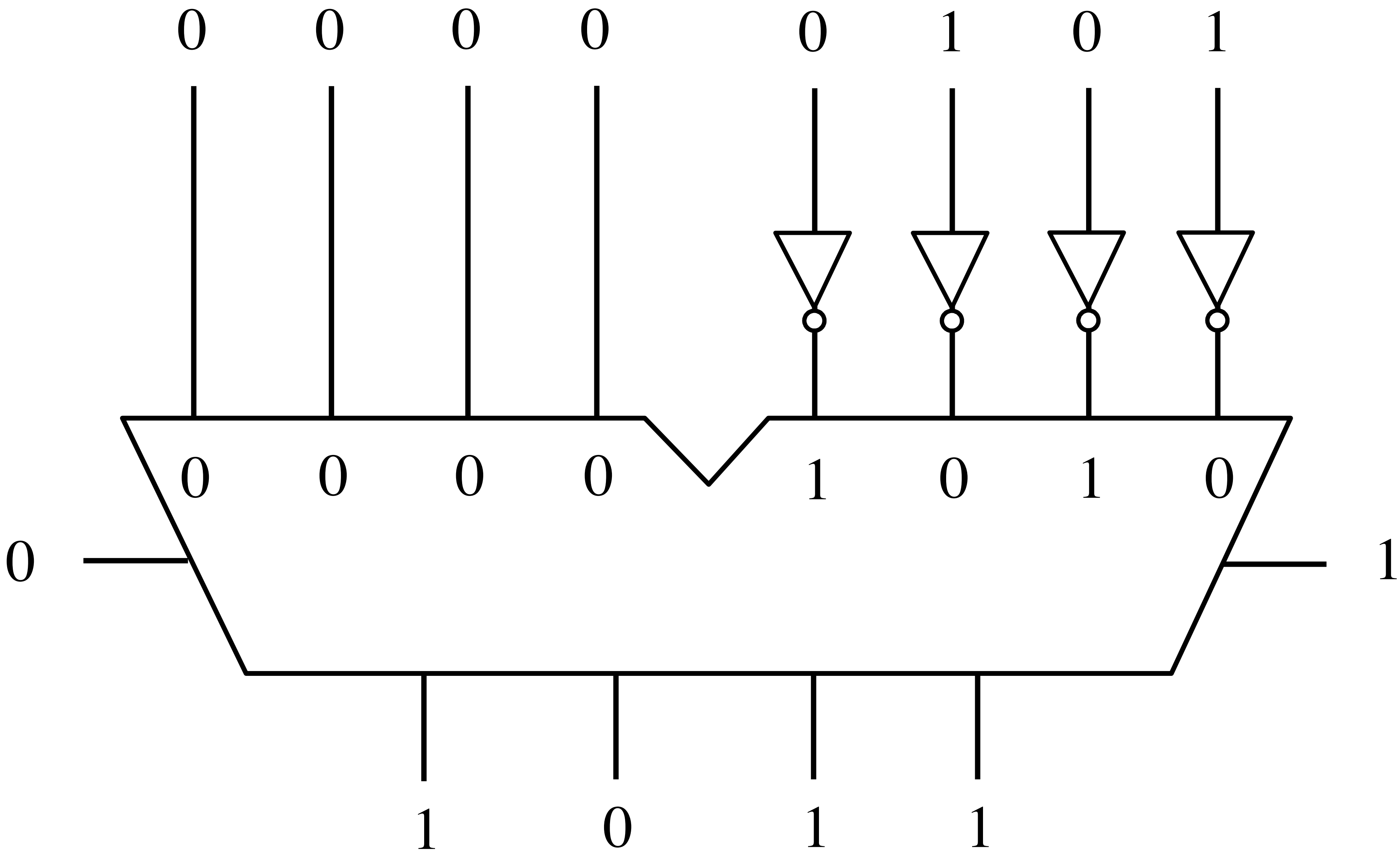


Circuit #1 for negating a number stored in 2's complement representation

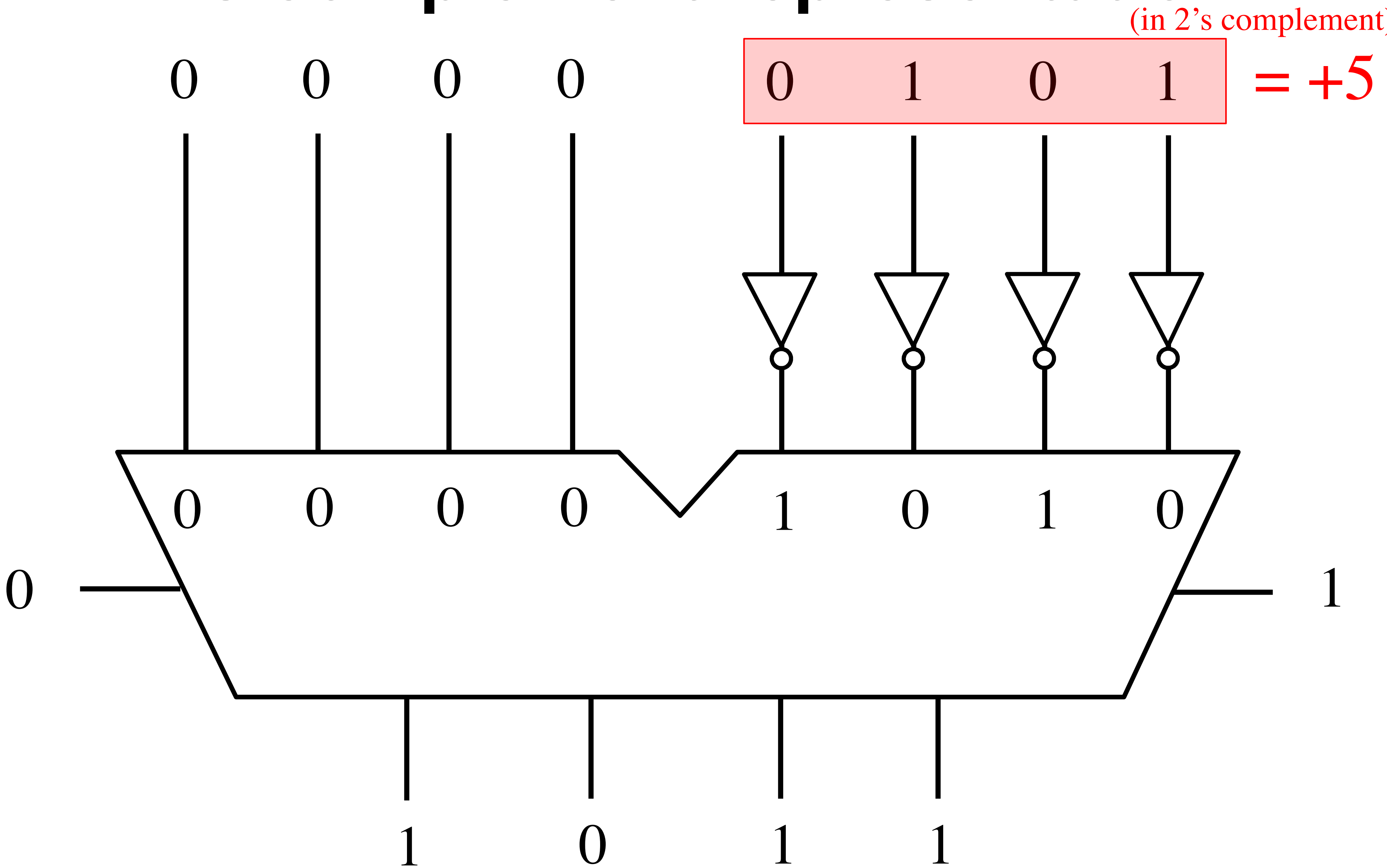


Alternative Circuit

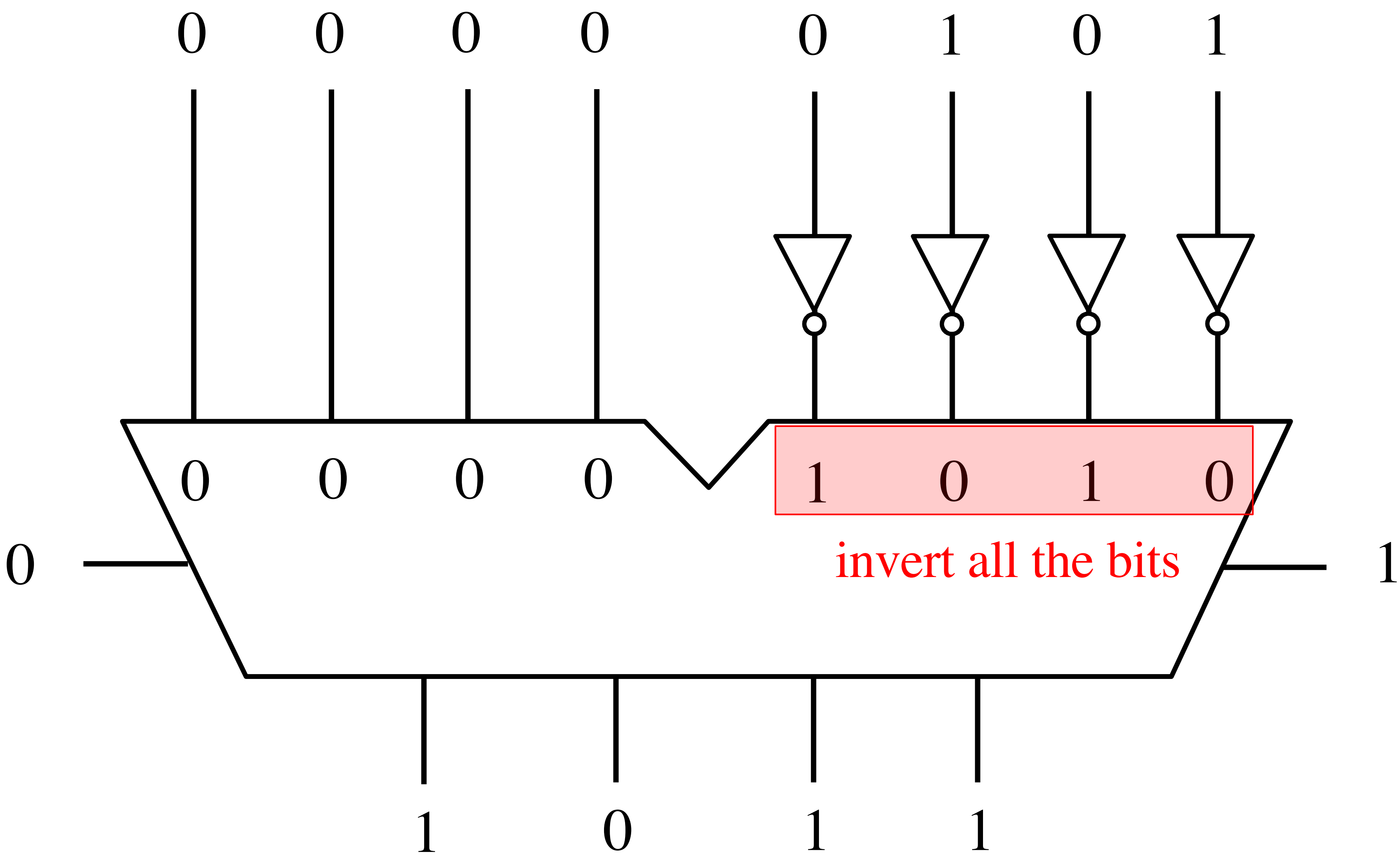
Circuit #2 for negating a number stored in 2's complement representation



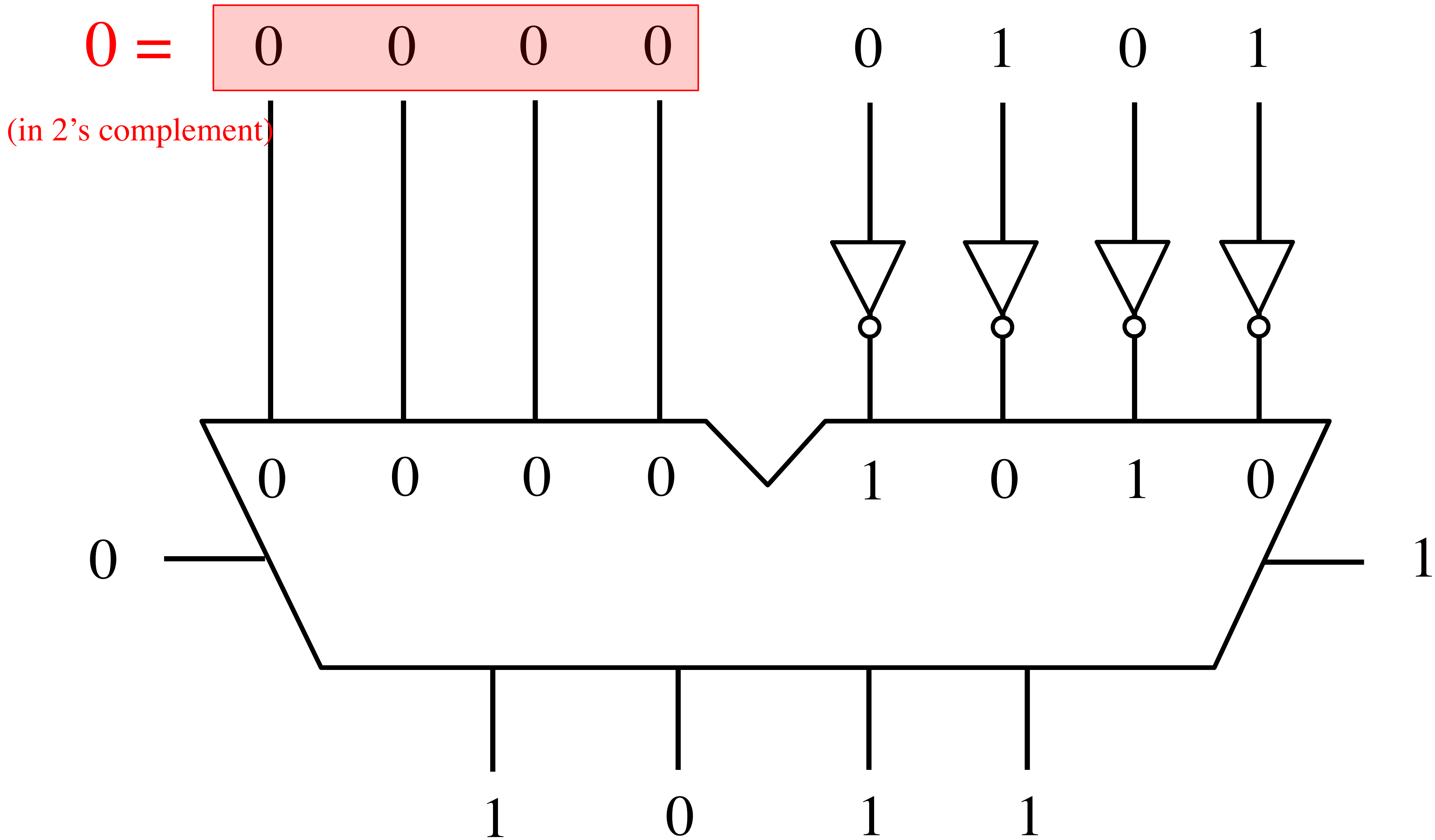
Circuit #2 for negating a number stored in 2's complement representation



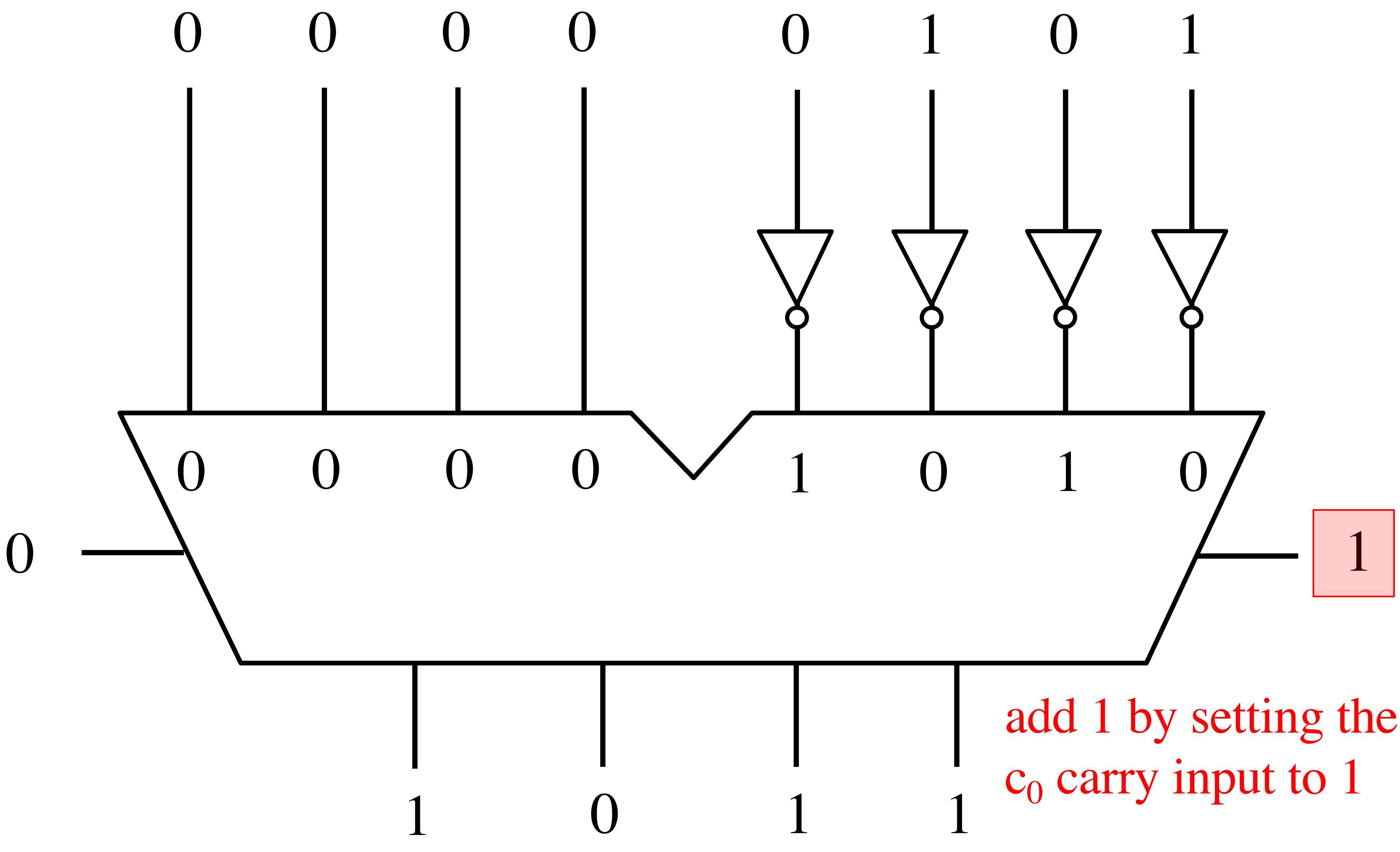
Circuit #2 for negating a number stored in 2's complement representation



Circuit #2 for negating a number stored in 2's complement representation

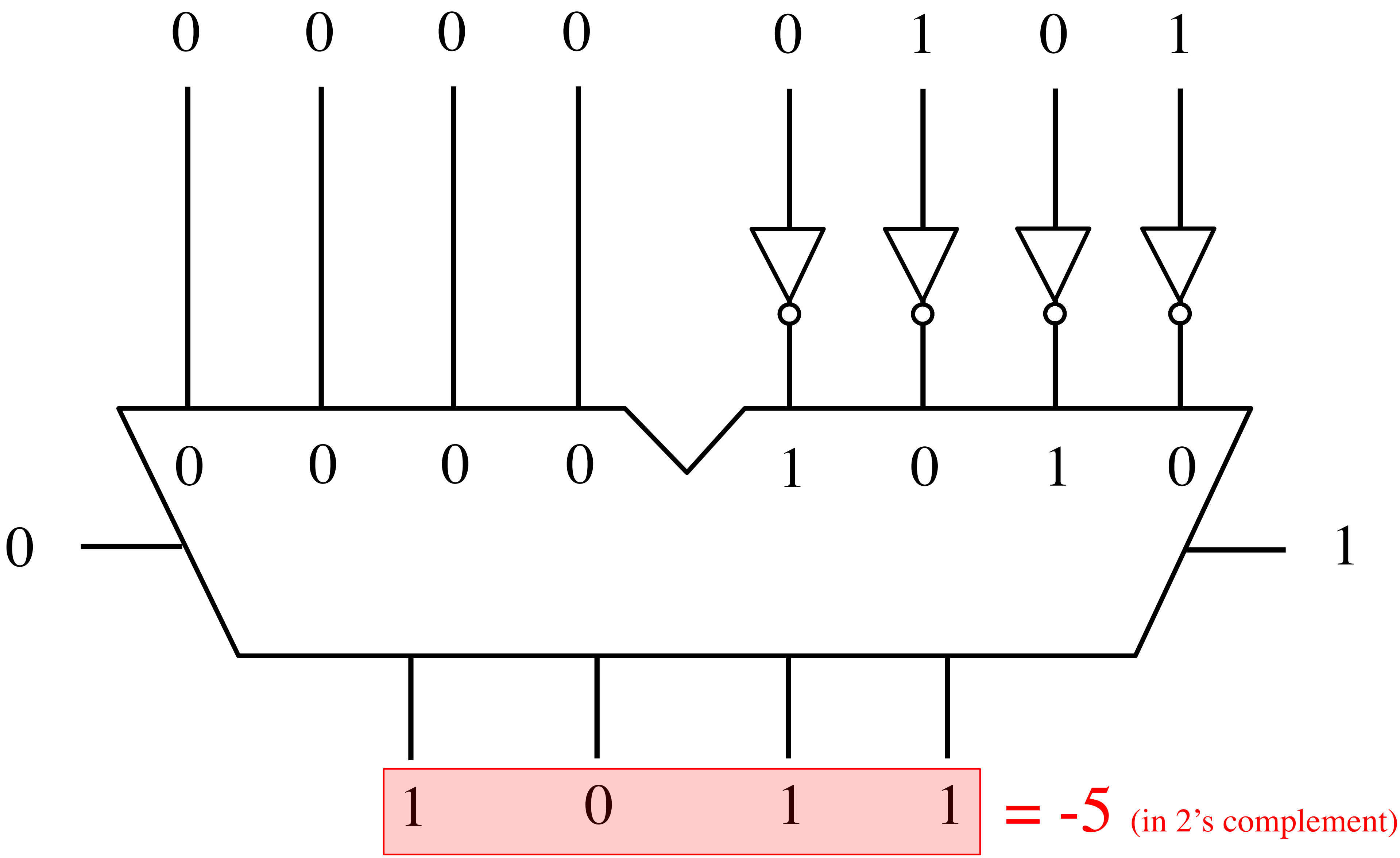


Circuit #2 for negating a number stored in 2's complement representation



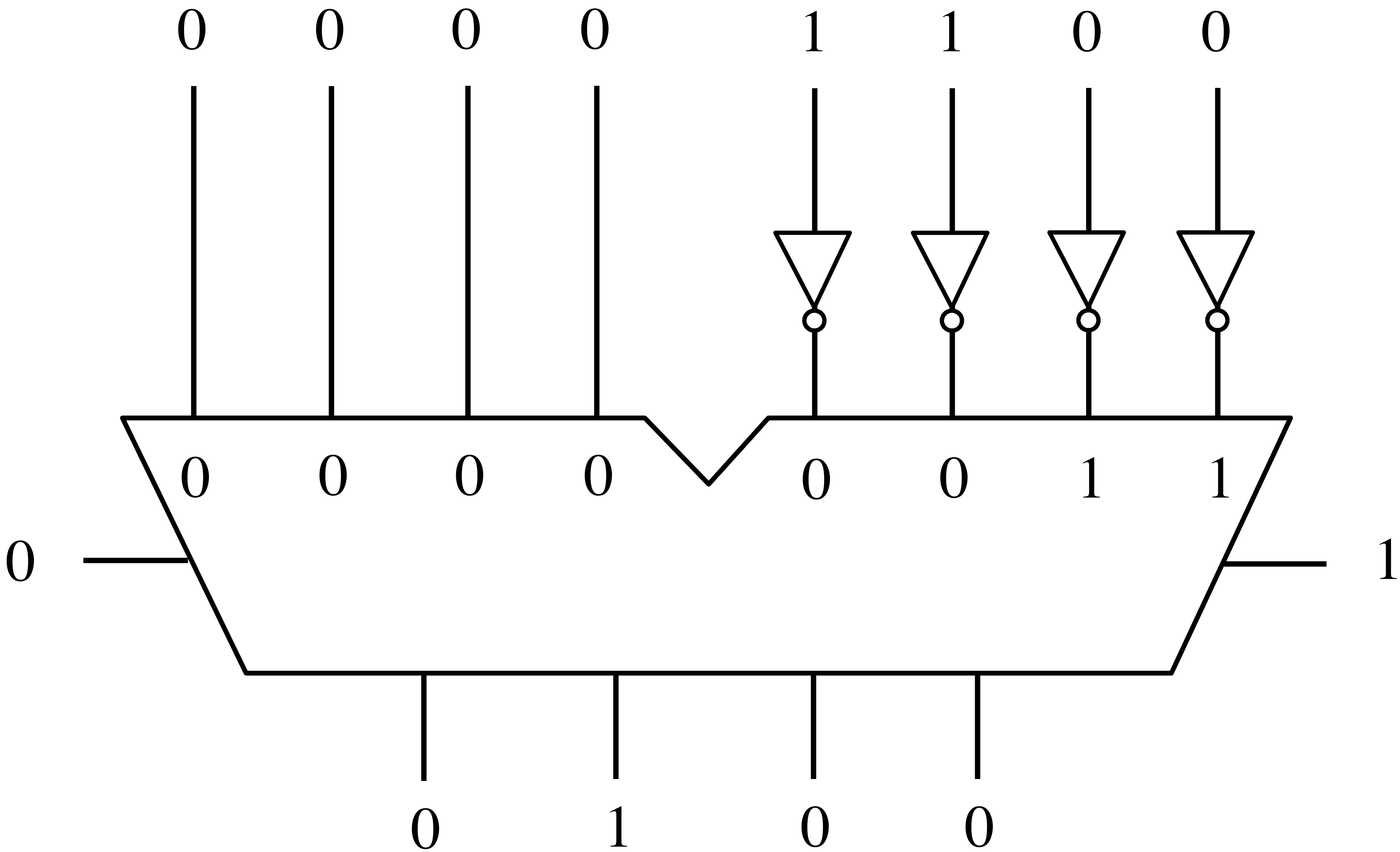
add 1 by setting the c_0 carry input to 1

Circuit #2 for negating a number stored in 2's complement representation

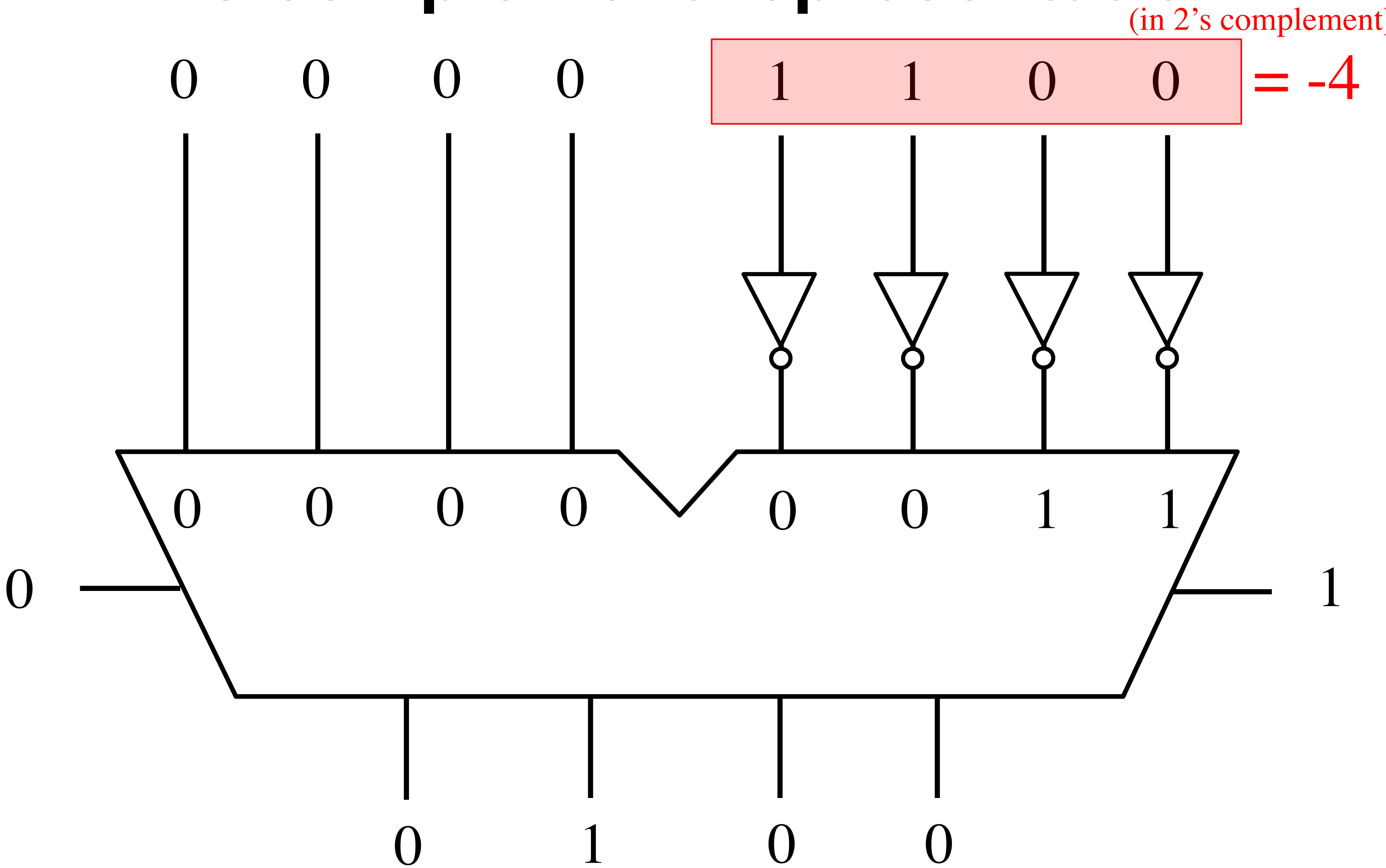


This also works for negating a negative number, thus making it positive

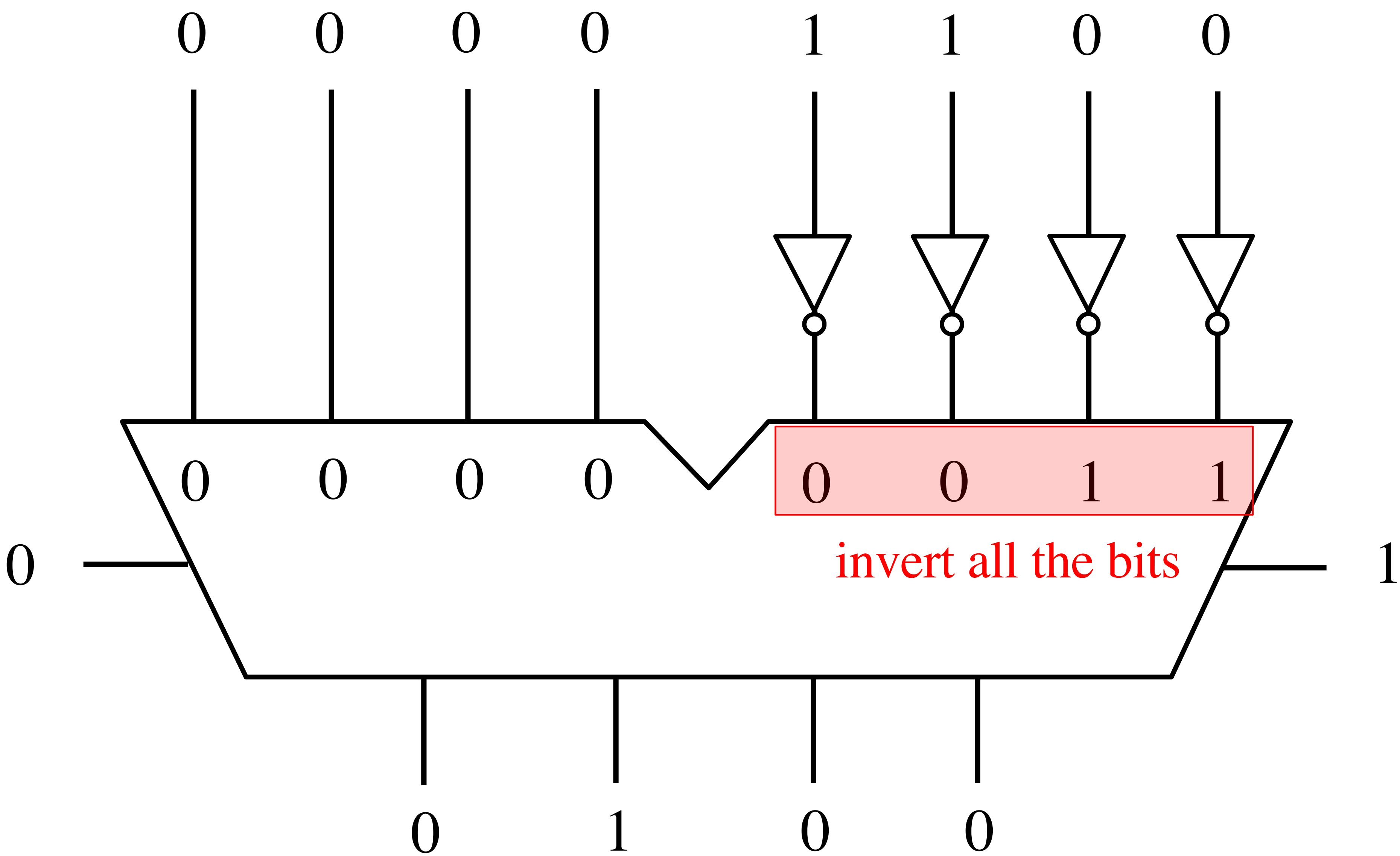
Circuit #2 for negating a number stored in 2's complement representation



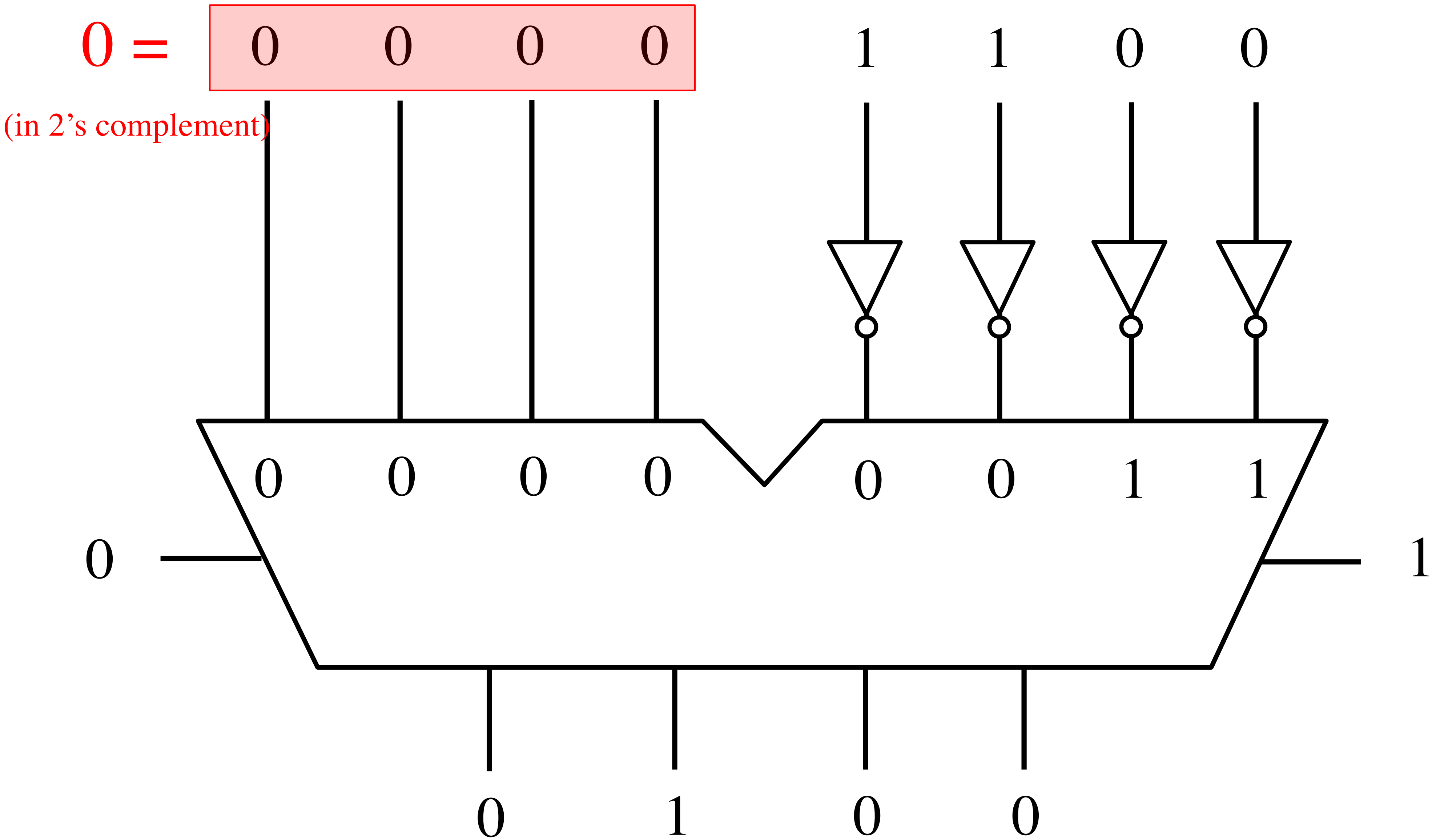
Circuit #2 for negating a number stored in 2's complement representation



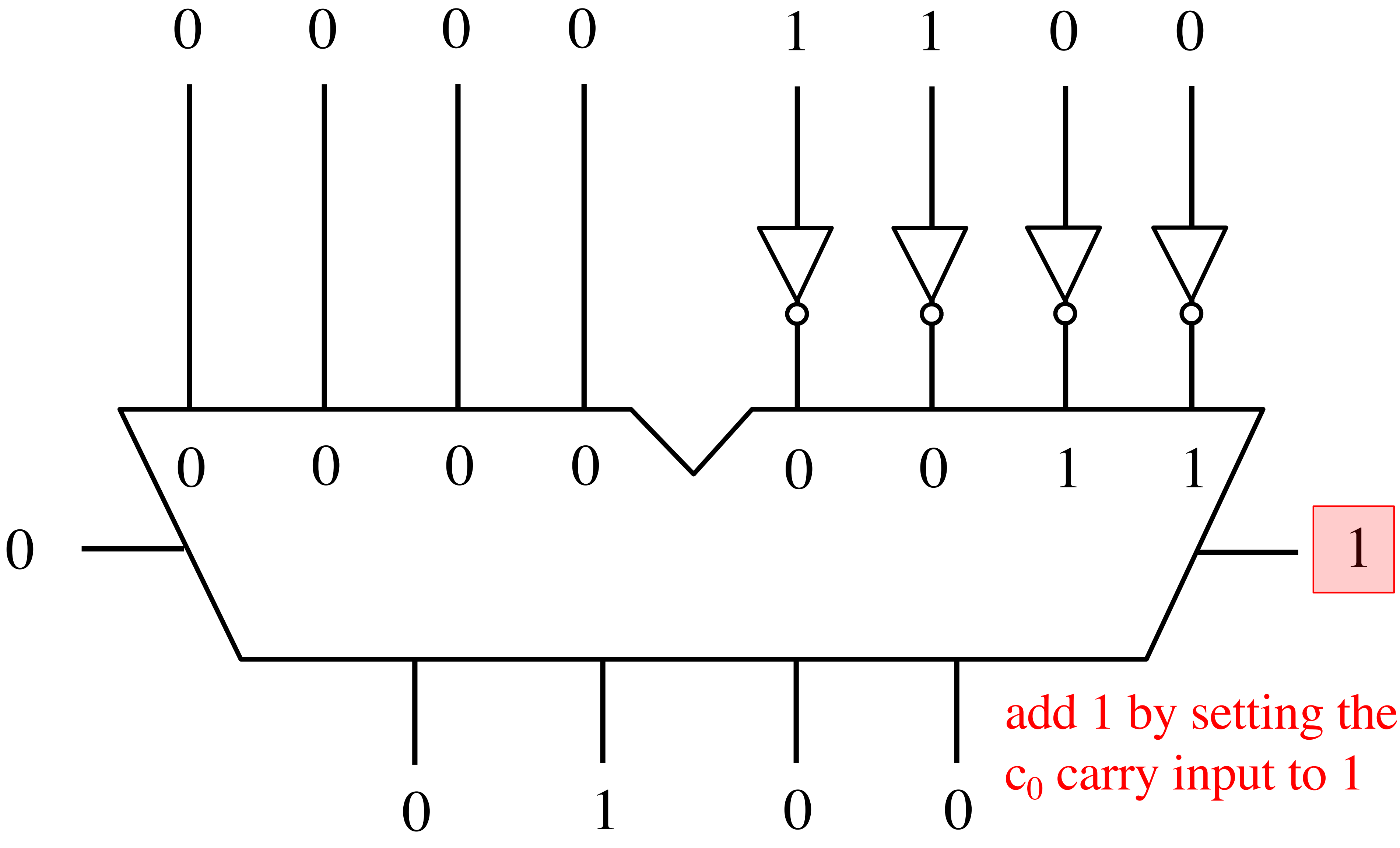
Circuit #2 for negating a number stored in 2's complement representation



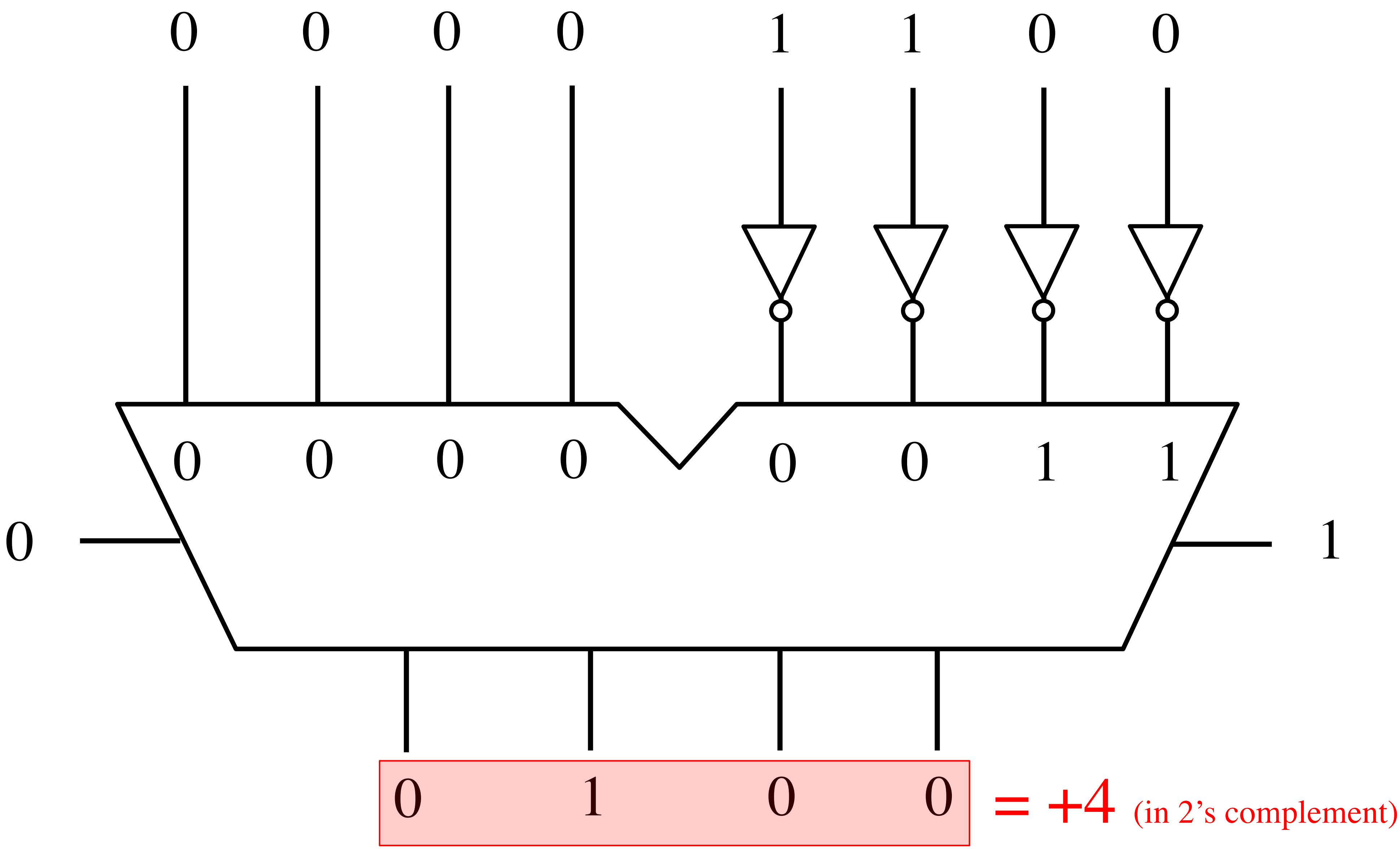
Circuit #2 for negating a number stored in 2's complement representation



Circuit #2 for negating a number stored in 2's complement representation



Circuit #2 for negating a number stored in 2's complement representation



Quick way (for a human) negate a number stored in 2's complement

- Scan the binary number from right to left
- Copy all bits that are 0 from right to left
- Stop at the first 1
- Copy that 1 as well
- Invert all remaining bits

Negate these numbers stored in 2’s complement representation

0 1 0 1

1 1 1 0

1 1 0 0

0 1 1 1

Negate these numbers stored in 2's complement representation

0 1 0 1	1 1 1 0
. . . .	. . . 0
1 1 0 0	0 1 1 1
. . 0 0	. . . .

Copy all bits that are 0 from right to left.

Negate these numbers stored in 2’s complement representation

0 1 0 1	1 1 1 0
. . . 1	. . 1 0
1 1 0 0	0 1 1 1
. 1 0 0	. . . 1

Stop at the first 1. Copy that 1 as well.

Negate these numbers stored in 2's complement representation

0 1 0 1	1 1 1 0
1 0 1 1	0 0 1 0
1 1 0 0	0 1 1 1
0 1 0 0	1 0 0 1

Invert all remaining bits.

Negate these numbers stored in 2's complement representation

0 1 0 1 = +5

1 0 1 1 = -5

1 1 1 0 = -2

0 0 1 0 = +2

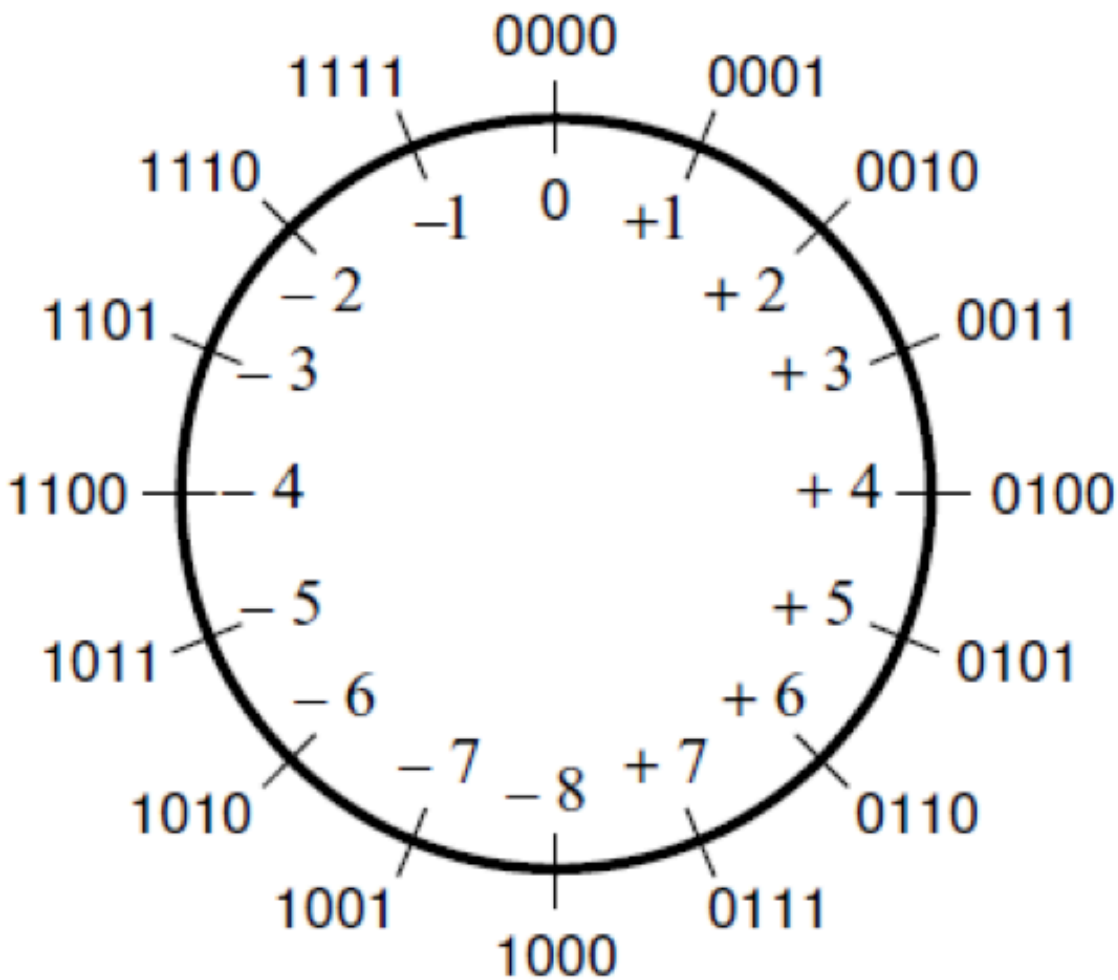
1 1 0 0 = -4

0 1 0 0 = +4

0 1 1 1 = +7

1 0 0 1 = -7

The number circle for 2's complement



[Figure 3.11a from the textbook]

Addition of two numbers stored in 2’s complement representation

There are four cases to consider

- $(+5) \quad + \quad (+2)$
- $(-5) \quad + \quad (+2)$
- $(+5) \quad + \quad (-2)$
- $(-5) \quad + \quad (-2)$

There are four cases to consider

- `(+5) + (+2)` positive plus positive
- `(–5) + (+2)` negative plus positive
- `(+5) + (–2)` positive plus negative
- `(–5) + (–2)` negative plus negative

A) Example of 2's complement addition

$$
\begin{array}{rr}
(+5) & 0\,1\,0\,1 \\
+\ (+2) & +\ 0\,0\,1\,0 \\
\hline
(+7) & 0\,1\,1\,1
\end{array}
$$

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.9 from the textbook]

B) Example of 2's complement addition

(−5)		1 0 1 1
+ (+ 2)	+	0 0 1 0
(−3)		1 1 0 1

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.9 from the textbook]

C) Example of 2’s complement addition

$$\begin{array}{r} (+5) \\ +\ (-2) \\ \hline (+3) \end{array} \qquad\qquad \begin{array}{r} 0\ 1\ 0\ 1 \\ +\ 1\ 1\ 1\ 0 \\ \hline 1\ \ 0\ 0\ 1\ 1 \end{array}$$

ignore

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.9 from the textbook]

D) Example of 2' s complement addition

$$
\begin{array}{rr}
(-5) & 1011 \\
+\ (-2) & +\ 1110 \\
\hline
(-7) & 1\ 1001
\end{array}
$$

ignore

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.9 from the textbook]

Naming Ambiguity: 2's Complement

2's complement has two different meanings:

- **representation for signed integer numbers**

- **algorithm for computing the 2's complement (regardless of the representation of the number)**

Naming Ambiguity: 2's Complement

2's complement has two different meanings:

- **representation for signed integer numbers**

  in 2's complement

- **algorithm for computing the 2's complement (regardless of the representation of the number)**

  take the 2's complement (or negate)

Subtraction of two numbers stored in 2’s complement representation

There are four cases to consider

- $(+5) - (+2)$
- $(-5) - (+2)$
- $(+5) - (-2)$
- $(-5) - (-2)$

There are four cases to consider

- `(+5) – (+2)` **positive minus positive**
- `(–5) – (+2)` **negative minus positive**
- `(+5) – (–2)` **positive minus negative**
- `(–5) – (–2)` **negative minus negative**

There are four cases to consider

- $(+5) - (+2)$
- $(-5) - (+2)$
- $(+5) - (-2)$
- $(-5) - (-2)$

There are four cases to consider

- $(+5) - (+2) = (+5) + (-2)$
- $(-5) - (+2) = (-5) + (-2)$
- $(+5) - (-2) = (+5) + (+2)$
- $(-5) - (-2) = (-5) + (+2)$

There are four cases to consider

- $(+5) - (+2) = (+5) + (-2)$
- $(-5) - (+2) = (-5) + (-2)$
- $(+5) - (-2) = (+5) + (+2)$
- $(-5) - (-2) = (-5) + (+2)$

We can change subtraction into addition ...

There are four cases to consider

- $(+5) \quad - \quad (+2) \quad = \quad (+5) \quad + \quad (-2)$
- $(-5) \quad - \quad (+2) \quad = \quad (-5) \quad + \quad (-2)$
- $(+5) \quad - \quad (-2) \quad = \quad (+5) \quad + \quad (+2)$
- $(-5) \quad - \quad (-2) \quad = \quad (-5) \quad + \quad (+2)$

… if we negate the second number.

There are four cases to consider

- $(+5) - (+2) = (+5) + (-2)$
- $(-5) - (+2) = (-5) + (-2)$
- $(+5) - (-2) = (+5) + (+2)$
- $(-5) - (-2) = (-5) + (+2)$

These are the four addition cases
(arranged in a shuffled order)

Example of 2's complement <u>subtraction</u>

$$\begin{array}{r} (+5) \\ -\ (+2) \\ \hline (+3) \end{array} \qquad \begin{array}{r} 0\,1\,0\,1 \\ -\ 0\,0\,1\,0 \\ \hline \end{array} \quad \Longrightarrow \quad \begin{array}{r} 0\,1\,0\,1 \\ +\ 1\,1\,1\,0 \\ \hline 1\,0\,0\,1\,1 \end{array}$$

ignore

$\Longrightarrow$ means take the 2's complement (or negate)

[Figure 3.10 from the textbook]

Example of 2' s complement <u>subtraction</u>

$$
\begin{array}{rr}
 & (+5) \\
- & (+2) \\
\hline
 & (+3)
\end{array}
\qquad
\begin{array}{rr}
 & 0\,1\,0\,1 \\
- & 0\,0\,1\,0 \\
\hline
\end{array}
\quad \Longrightarrow \quad
\begin{array}{rr}
 & 0\,1\,0\,1 \\
+ & 1\,1\,1\,0 \\
\hline
1 & 0\,0\,1\,1
\end{array}
$$

ignore

Notice that the minus changes to a plus.

$\Longrightarrow$ means take the 2's complement (or negate)

[Figure 3.10 from the textbook]

Example of 2' s complement subtraction

$$\begin{array}{r} (+\,5) \\ -\,(+\,2) \\ \hline (+\,3) \end{array} \qquad \begin{array}{r} 0\;1\;0\;1 \\ -\;0\;0\;1\;0 \\ \hline \end{array} \quad \Longrightarrow \quad \begin{array}{r} 0\;1\;0\;1 \\ +\;1\;1\;1\;0 \\ \hline 1\;0\;0\;1\;1 \end{array}$$

ignore

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.10 from the textbook]

Example of 2's complement subtraction

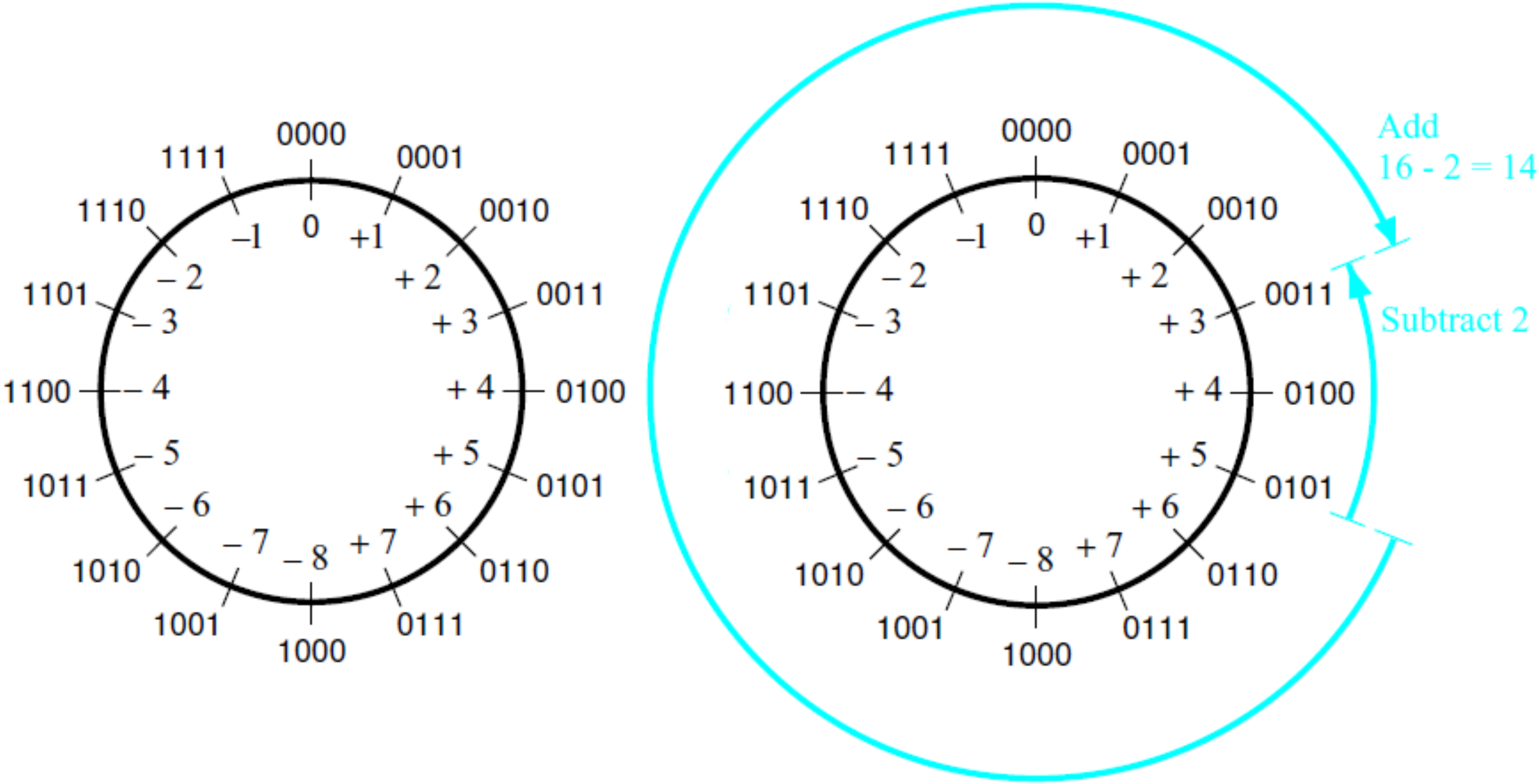
$$\begin{array}{r} (+5) \\ -\,(+2) \\ \hline (+3) \end{array} \qquad \begin{array}{r} 0\,1\,0\,1 \\ -\;0\,0\,1\,0 \\ \hline \end{array} \qquad \Longrightarrow \qquad \begin{array}{r} 0\,1\,0\,1 \\ +\;1\,1\,1\,0 \\ \hline 1\,0\,0\,1\,1 \end{array}$$

ignore

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.10 from the textbook]

Graphical interpretation of four-bit 2's complement numbers



(a) The number circle (b) Subtracting 2 by adding its 2's complement

[Figure 3.11 from the textbook]

Example of 2' s complement subtraction

$$\begin{array}{r} (-5) \\ -\ (+2) \\ \hline (-7) \end{array} \qquad \begin{array}{r} 1\,0\,1\,1 \\ -\ 0\,0\,1\,0 \\ \hline \end{array} \qquad \Longrightarrow \qquad \begin{array}{r} 1\,0\,1\,1 \\ +\ 1\,1\,1\,0 \\ \hline 1\ 1\,0\,0\,1 \end{array}$$

ignore

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.10 from the textbook]

Example of 2’s complement subtraction

$$\begin{array}{r} (+5) \\ -\ (-2) \\ \hline (+7) \end{array} \qquad \begin{array}{r} 0\,1\,0\,1 \\ -\ 1\,1\,1\,0 \\ \hline \\ \end{array} \qquad \Longrightarrow \qquad \begin{array}{r} 0\,1\,0\,1 \\ +\ 0\,0\,1\,0 \\ \hline 0\,1\,1\,1 \end{array}$$

$b_3b_2b_1b_0$	2’s complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.10 from the textbook]

Example of 2’s complement subtraction

$$\begin{array}{r} (-5) \\ -\ (-2) \\ \hline (-3) \end{array} \qquad \begin{array}{r} 1\,0\,1\,1 \\ -\ 1\,1\,1\,0 \\ \hline \end{array} \qquad \Longrightarrow \qquad \begin{array}{r} 1\,0\,1\,1 \\ +\ 0\,0\,1\,0 \\ \hline 1\,1\,0\,1 \end{array}$$

$b_3b_2b_1b_0$	2’s complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.10 from the textbook]

Taking the 2’s complement negates the number

decimal	$b_3\, b_2\, b_1\, b_0$	take the 2's complement	$b_3\, b_2\, b_1\, b_0$	decimal
+7	0111	$\Longrightarrow$	1001	−7
+6	0110	$\Longrightarrow$	1010	−6
+5	0101	$\Longrightarrow$	1011	−5
+4	0100	$\Longrightarrow$	1100	−4
+3	0011	$\Longrightarrow$	1101	−3
+2	0010	$\Longrightarrow$	1110	−2
+1	0001	$\Longrightarrow$	1111	−1
+0	0000	$\Longrightarrow$	0000	+0
−8	1000	$\Longrightarrow$	1000	−8
−7	1001	$\Longrightarrow$	0111	+7
−6	1010	$\Longrightarrow$	0110	+6
−5	1011	$\Longrightarrow$	0101	+5
−4	1100	$\Longrightarrow$	0100	+4
−3	1101	$\Longrightarrow$	0011	+3
−2	1110	$\Longrightarrow$	0010	+2
−1	1111	$\Longrightarrow$	0001	+1

Taking the 2’s complement negates the number

decimal	$b_3\,b_2\,b_1\,b_0$	take the 2's complement	$b_3\,b_2\,b_1\,b_0$	decimal
+7	0111	$\Longrightarrow$	1001	−7
+6	0110	$\Longrightarrow$	1010	−6
+5	0101	$\Longrightarrow$	1011	−5
+4	0100	$\Longrightarrow$	1100	−4
+3	0011	$\Longrightarrow$	1101	−3
+2	0010	$\Longrightarrow$	1110	−2
+1	0001	$\Longrightarrow$	1111	−1
+0	0000	$\Longrightarrow$	0000	+0
−8	1000	$\Longrightarrow$	1000	−8
−7	1001	$\Longrightarrow$	0111	+7
−6	1010	$\Longrightarrow$	0110	+6
−5	1011	$\Longrightarrow$	0101	+5
−4	1100	$\Longrightarrow$	0100	+4
−3	1101	$\Longrightarrow$	0011	+3
−2	1110	$\Longrightarrow$	0010	+2
−1	1111	$\Longrightarrow$	0001	+1

This is the only exception

Taking the 2's complement negates the number

decimal	$b_3\,b_2\,b_1\,b_0$	take the 2's complement	$b_3\,b_2\,b_1\,b_0$	decimal
+7	0111	$\Longrightarrow$	1001	−7
+6	0110	$\Longrightarrow$	1010	−6
+5	0101	$\Longrightarrow$	1011	−5
+4	0100	$\Longrightarrow$	1100	−4
+3	0011	$\Longrightarrow$	1101	−3
+2	0010	$\Longrightarrow$	1110	−2
+1	0001	$\Longrightarrow$	1111	−1
+0	0000	$\Longrightarrow$	0000	+0
−8	1000	$\Longrightarrow$	1000	−8
−7	1001	$\Longrightarrow$	0111	+7
−6	1010	$\Longrightarrow$	0110	+6
−5	1011	$\Longrightarrow$	0101	+5
−4	1100	$\Longrightarrow$	0100	+4
−3	1101	$\Longrightarrow$	0011	+3
−2	1110	$\Longrightarrow$	0010	+2
−1	1111	$\Longrightarrow$	0001	+1

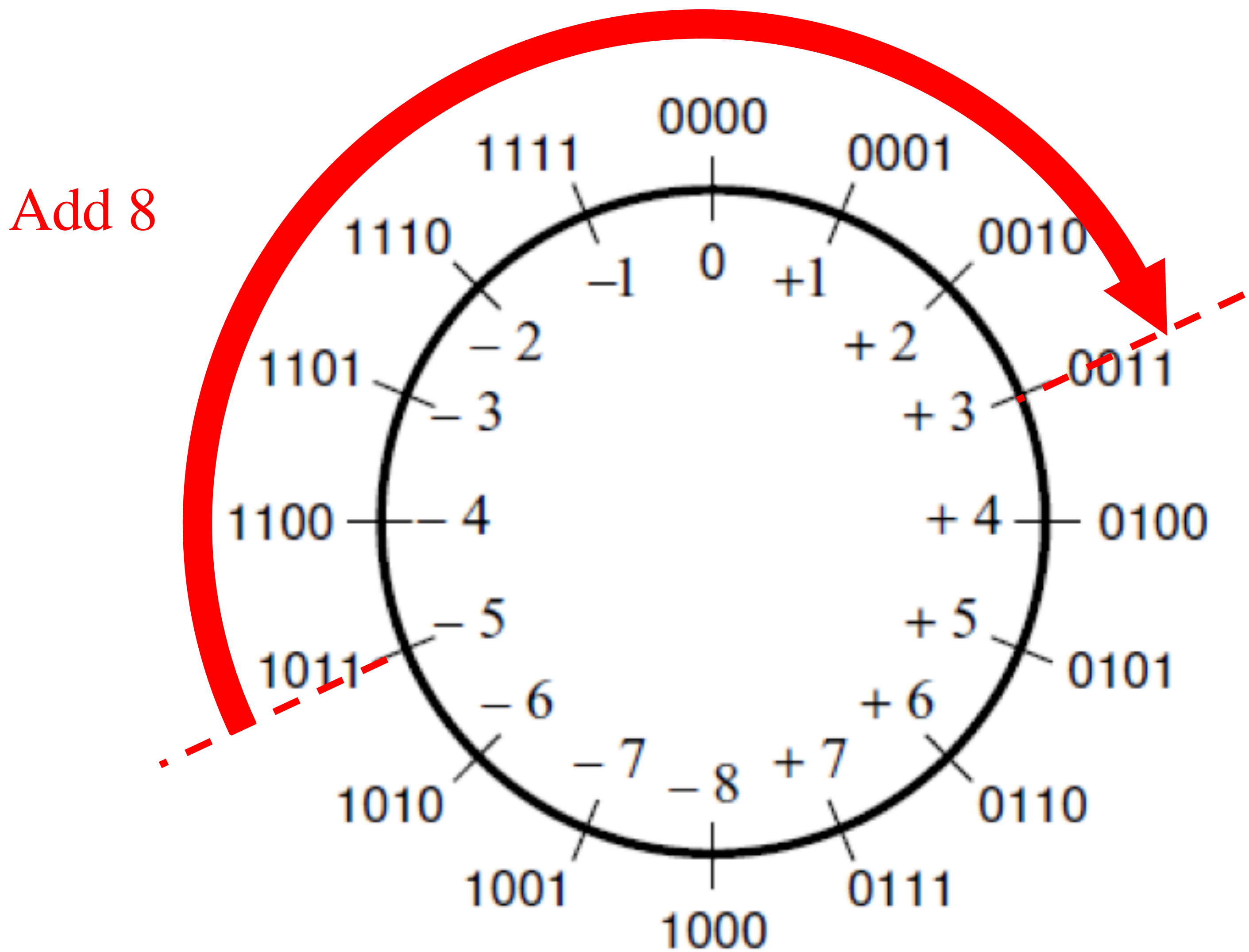
And this one too.

But that exception does not matter

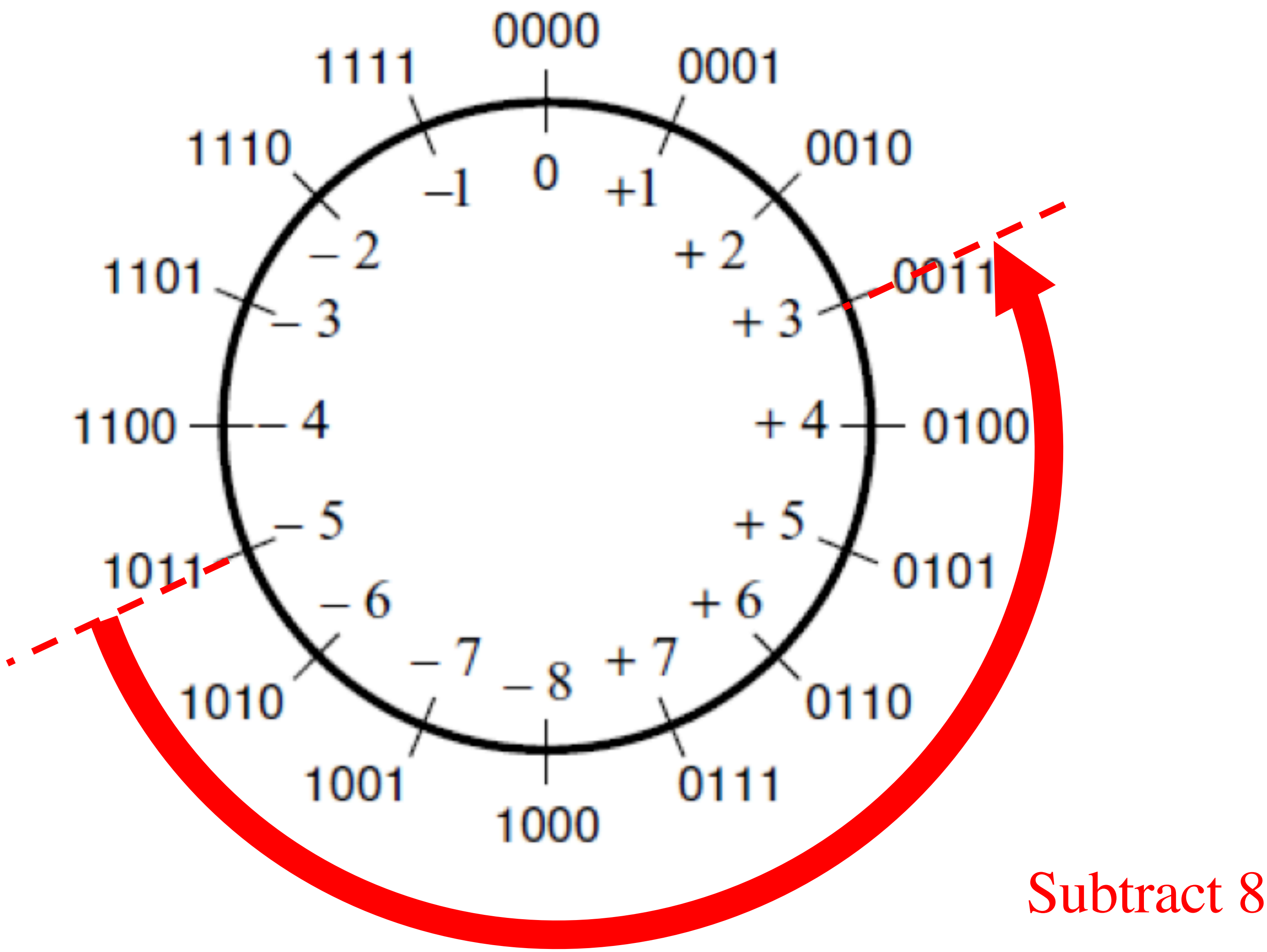
$$\begin{array}{r} (-5) \\ -\ (-8) \\ \hline (+3) \end{array} \qquad \begin{array}{r} 1\,0\,1\,1 \\ -\ 1\,0\,0\,0 \\ \hline \end{array} \qquad \Longrightarrow \qquad \begin{array}{r} 1\,0\,1\,1 \\ +\ 1\,0\,0\,0 \\ \hline 1\,0\,0\,1\,1 \end{array}$$

ignore

But that exception does not matter



But that exception does not matter



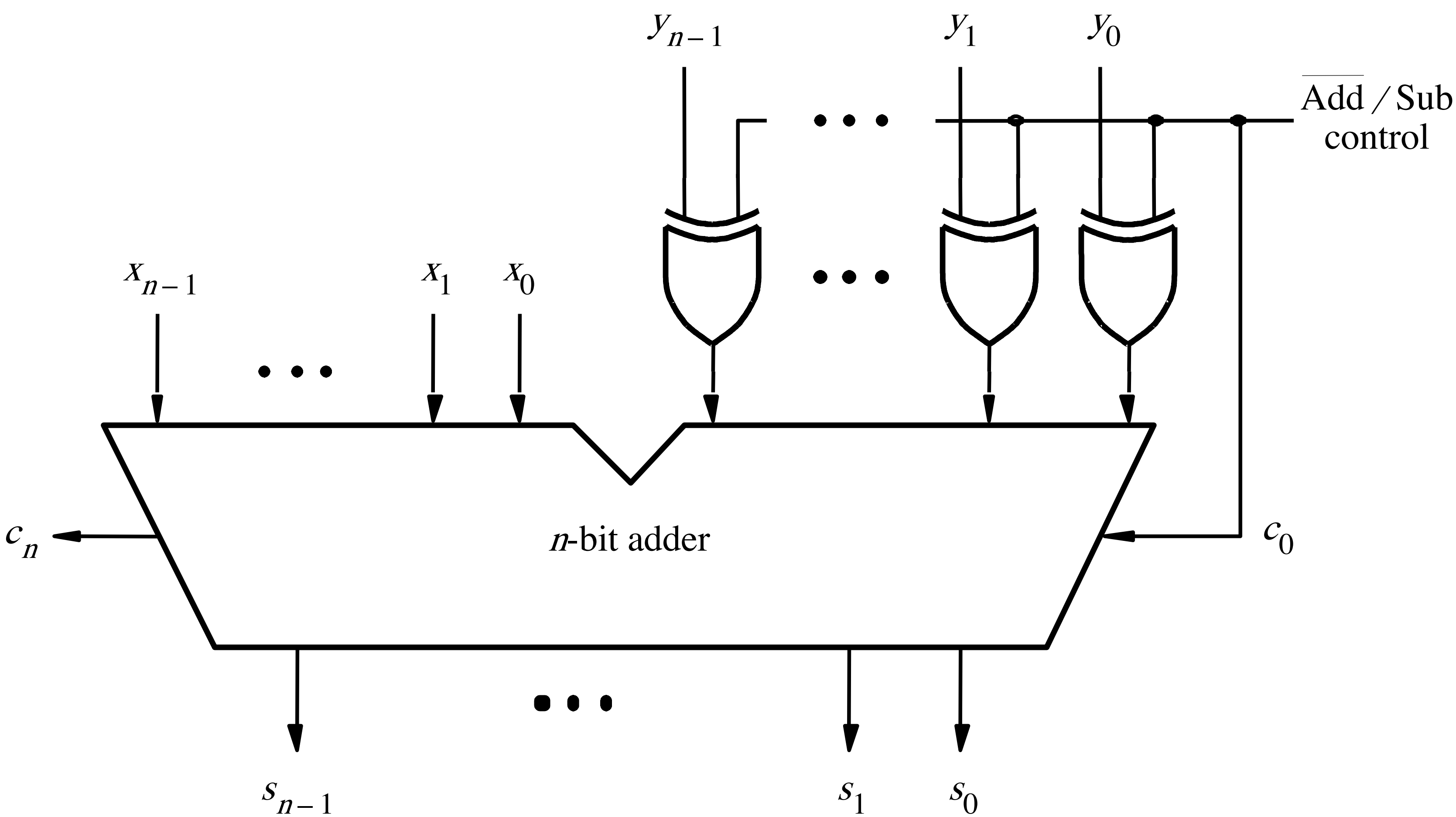
Take-Home Message

Take-Home Message

- **Subtraction can be performed by simply adding the 2's complement of the second number, regardless of the signs of the two numbers.**

- **Thus, the same adder circuit can be used to perform <u>both addition and subtraction</u> !!!**

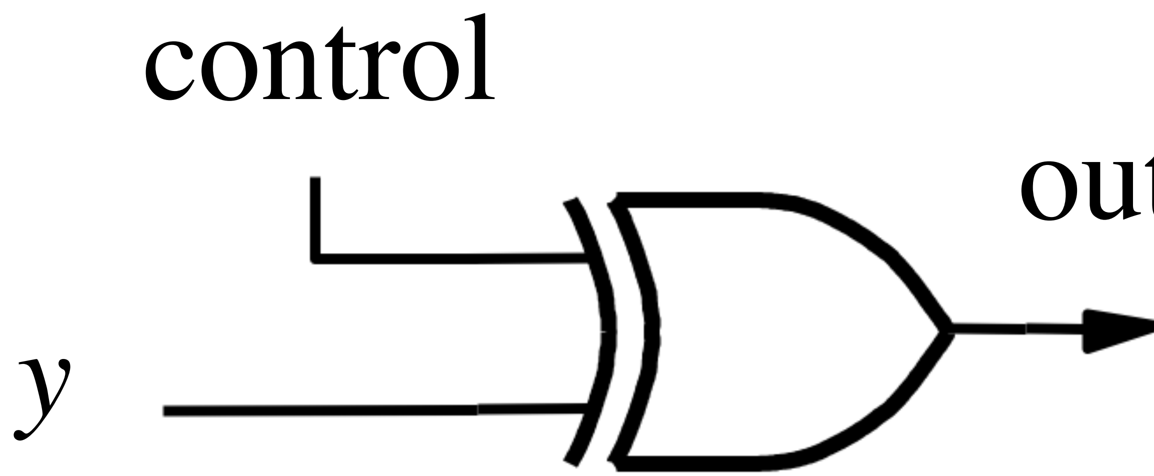
Adder/subtractor unit



[Figure 3.12 from the textbook]

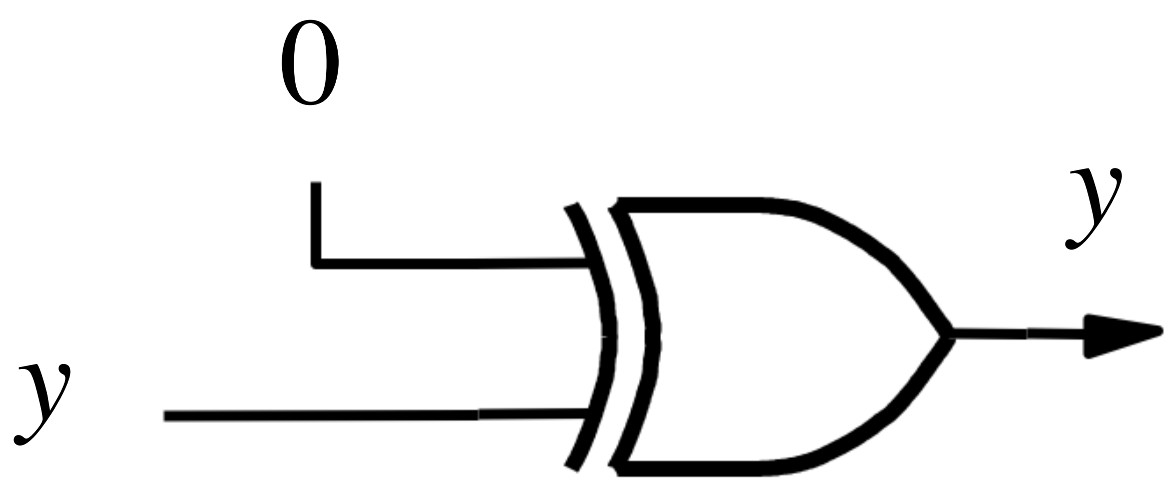
XOR Tricks

control	y	out
0	0	0
0	1	1
1	0	1
1	1	0



XOR as a repeater

control	y	out
0	0	0
0	1	1

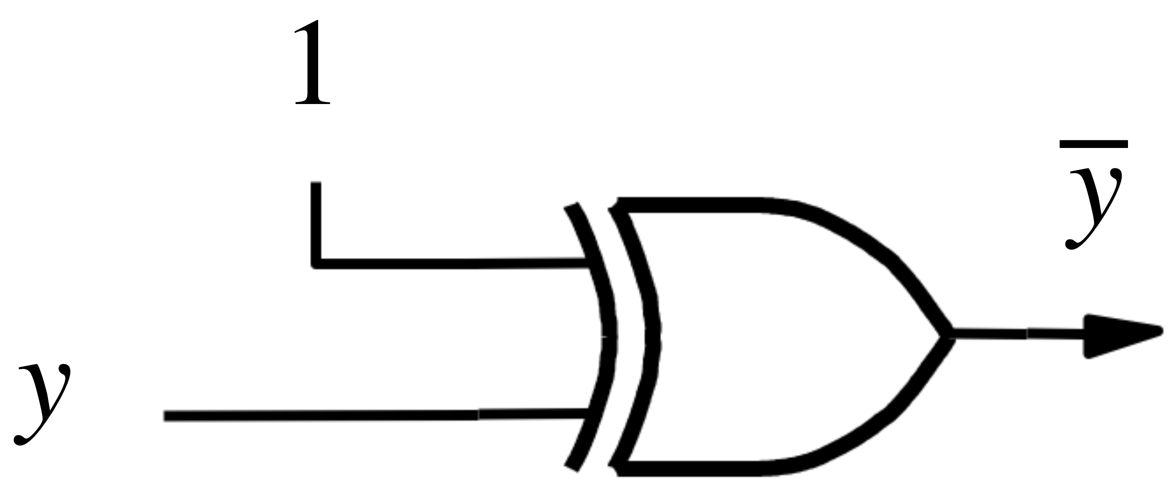


XOR as a repeater

control	y	out
0	0	0
0	1	1

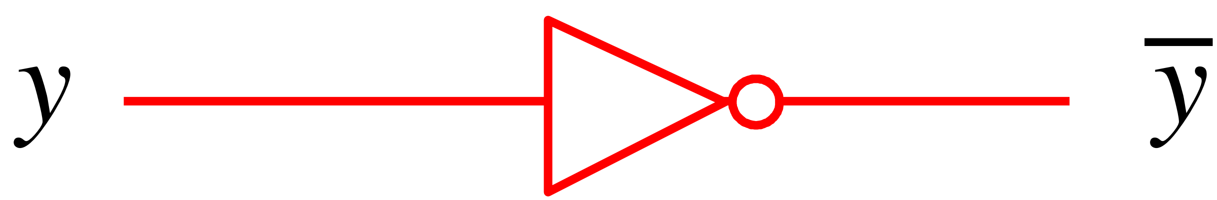
XOR as an inverter

control	y	out
1	0	1
1	1	0

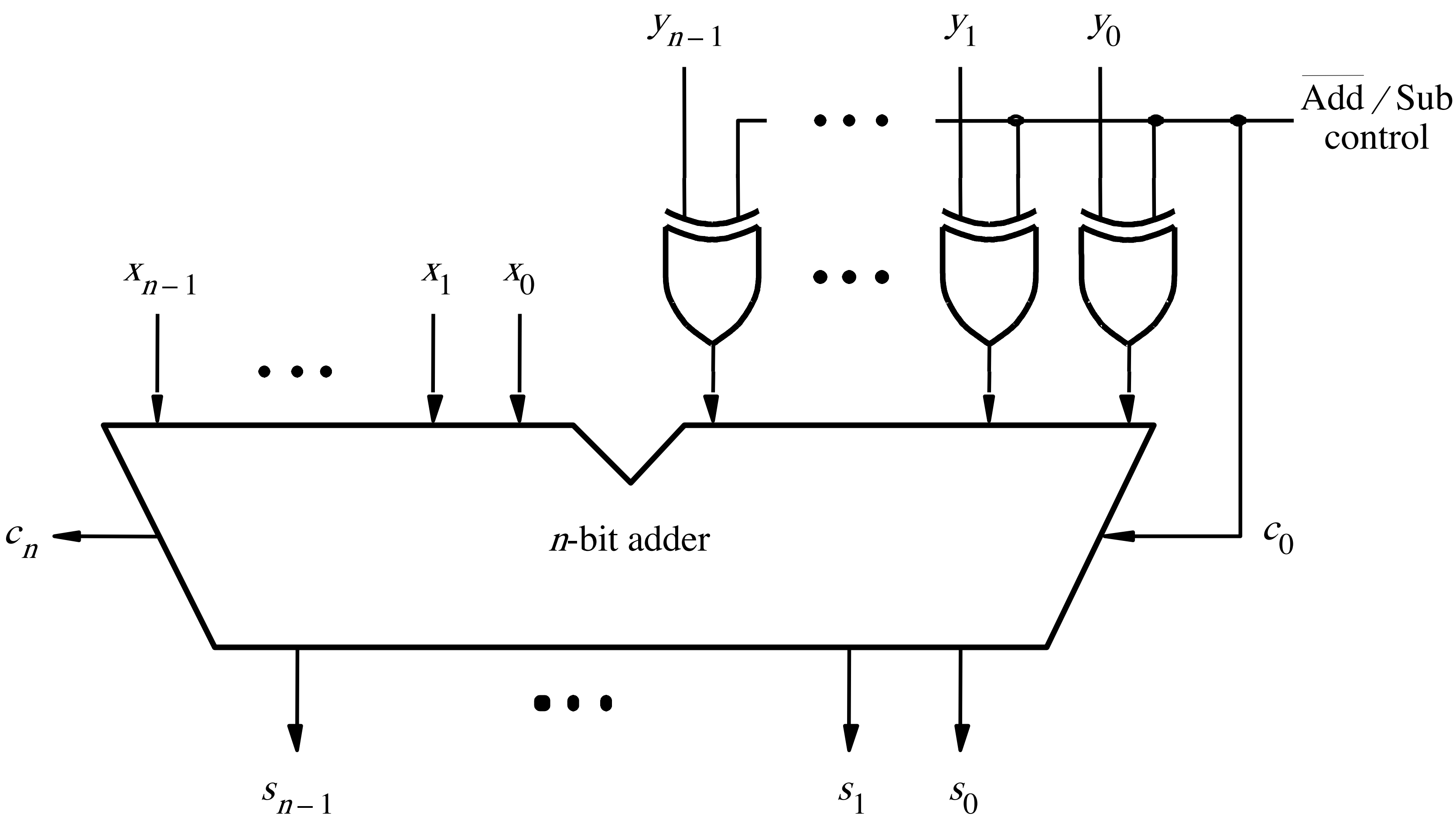


XOR as an inverter

control	y	out
1	0	1
1	1	0

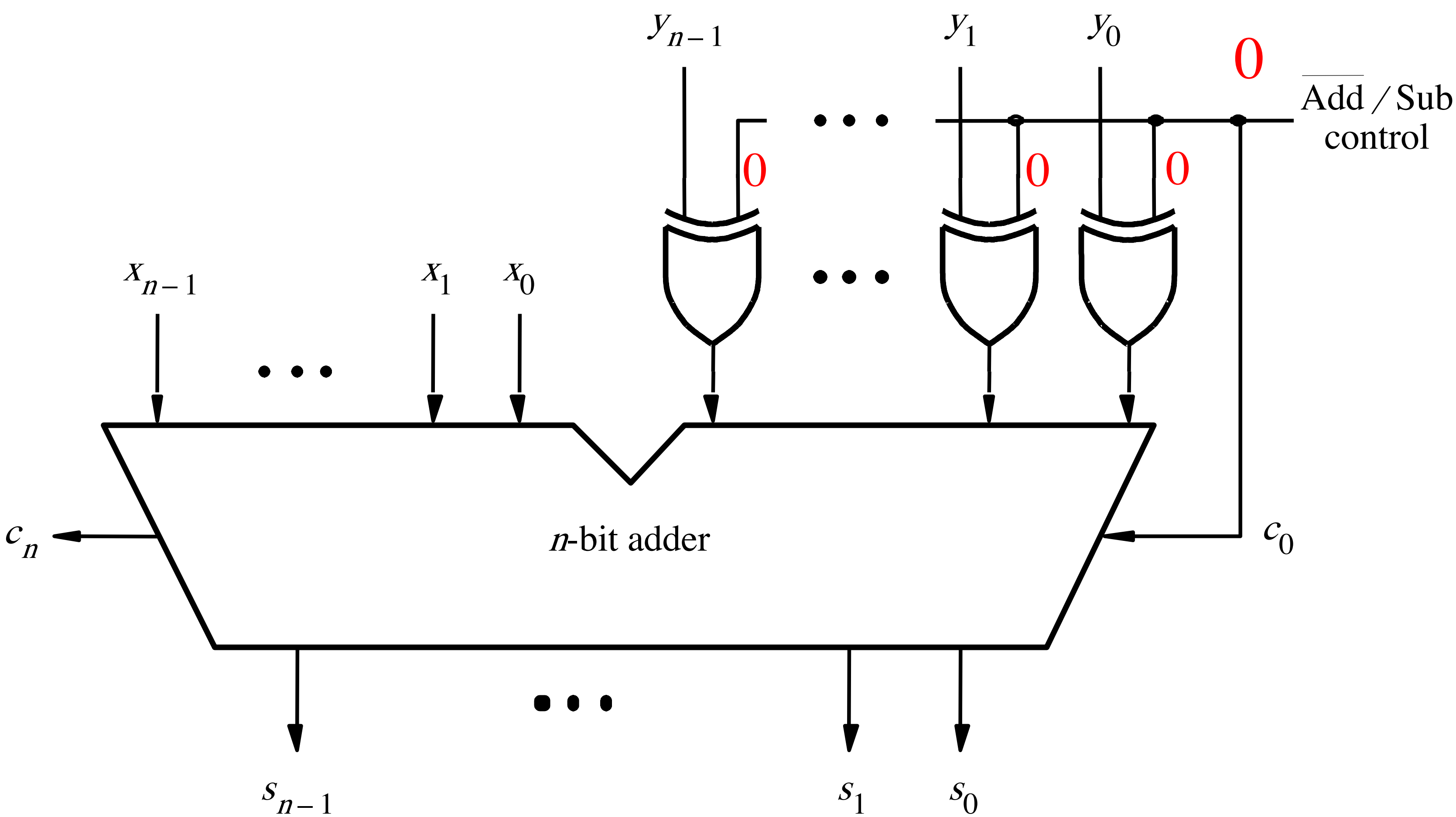


Addition: when control = 0



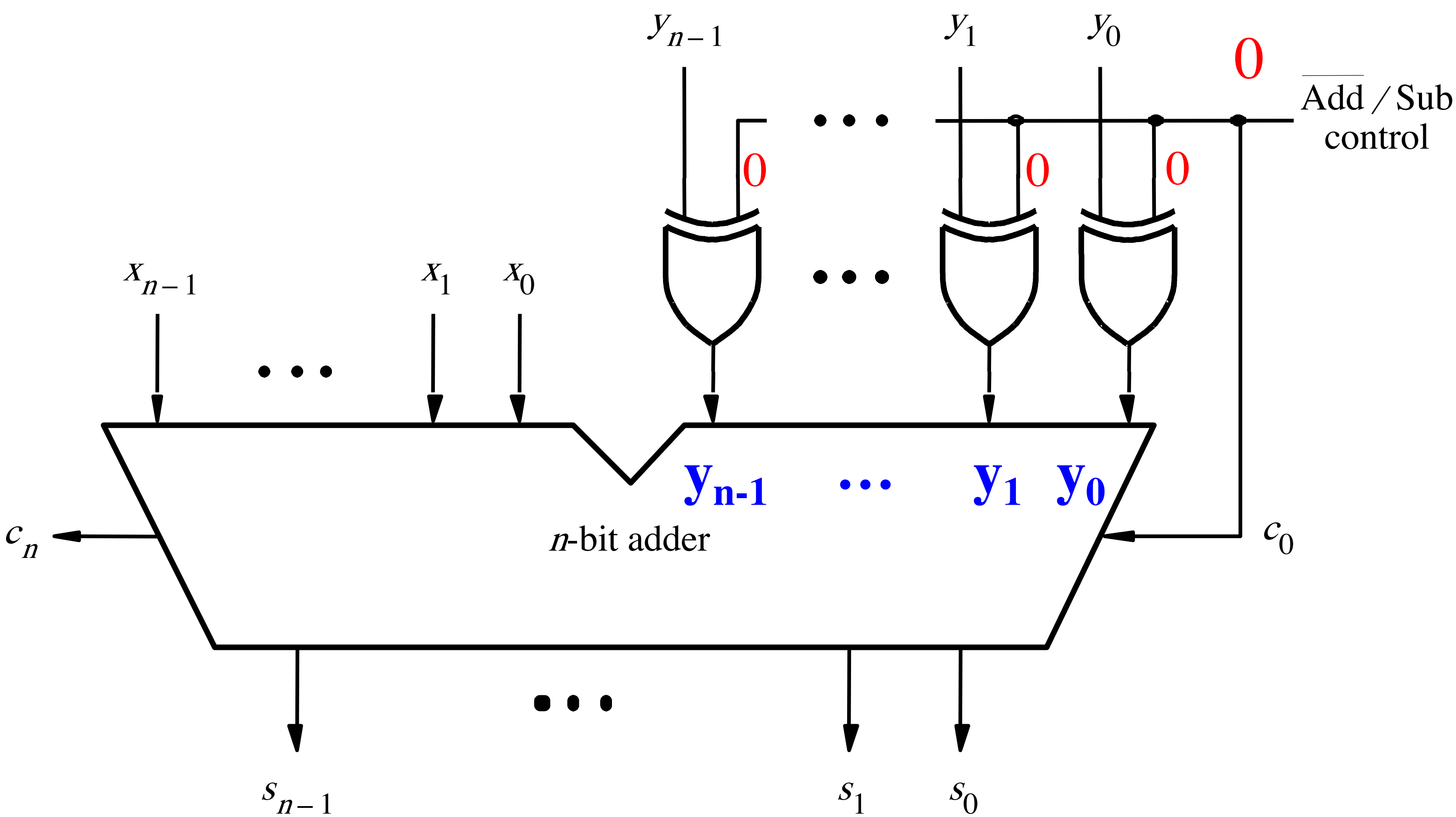
[Figure 3.12 from the textbook]

Addition: when control = 0



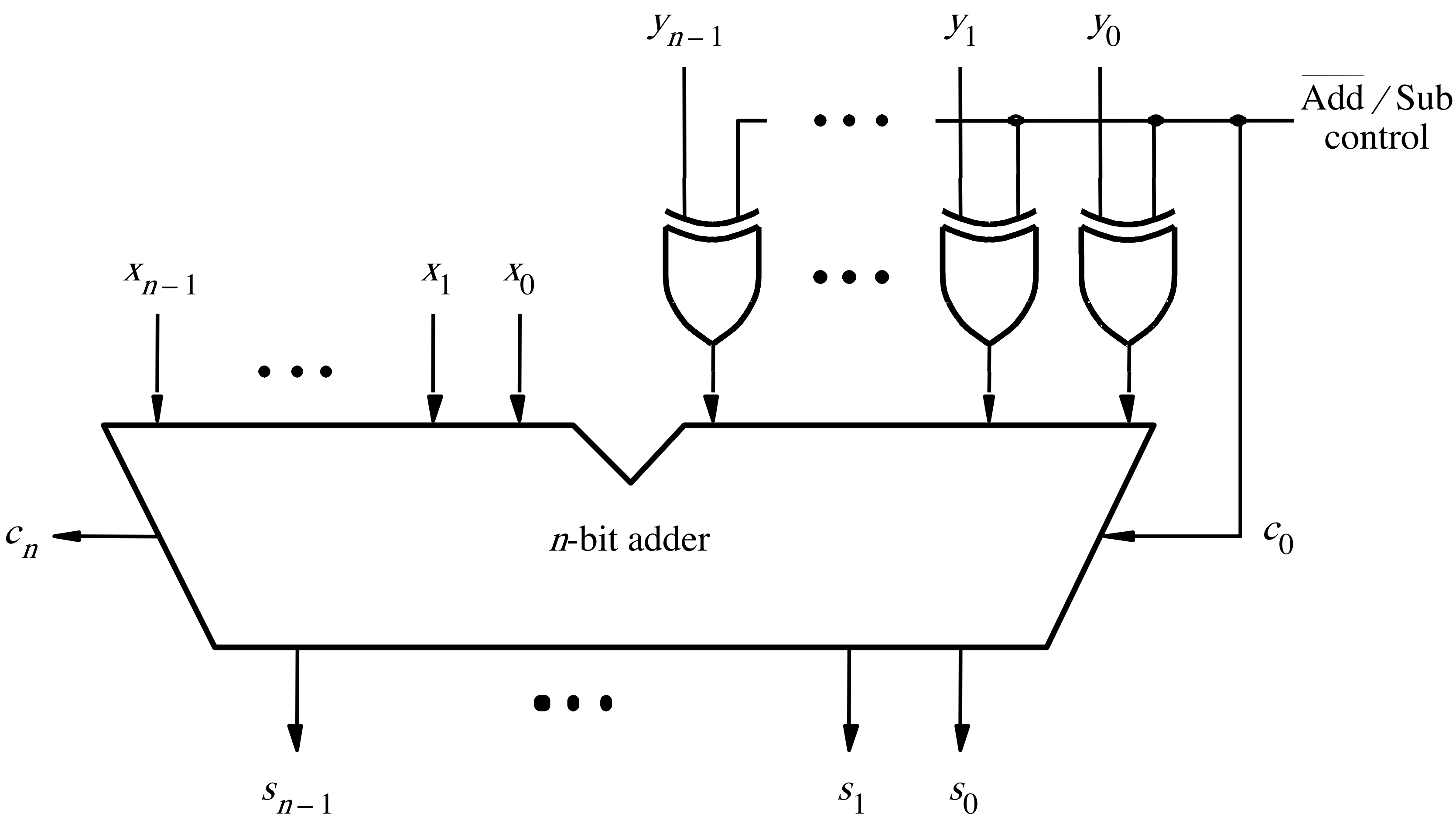
[Figure 3.12 from the textbook]

Addition: when control = 0



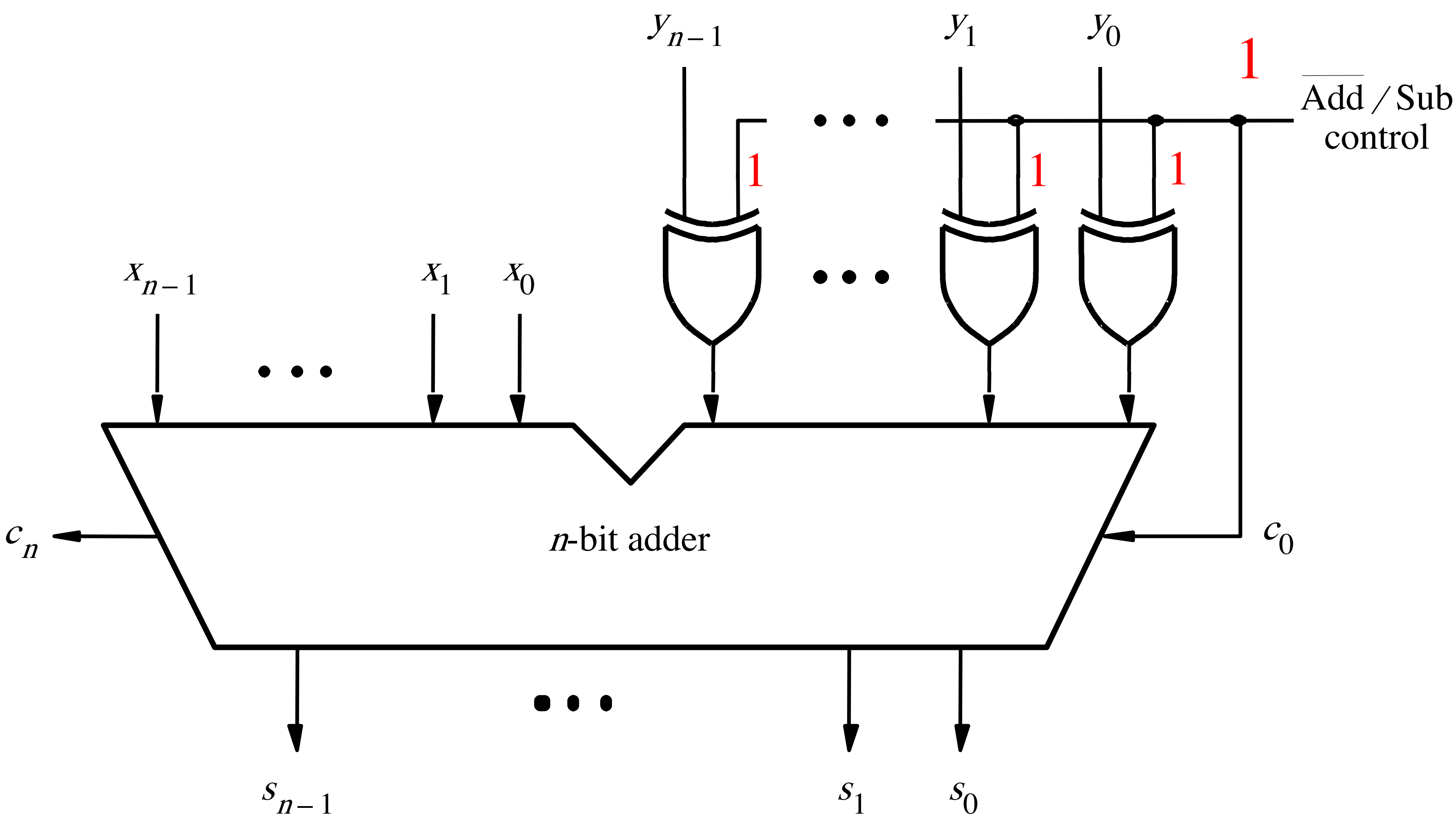
[Figure 3.12 from the textbook]

Subtraction: when control = 1



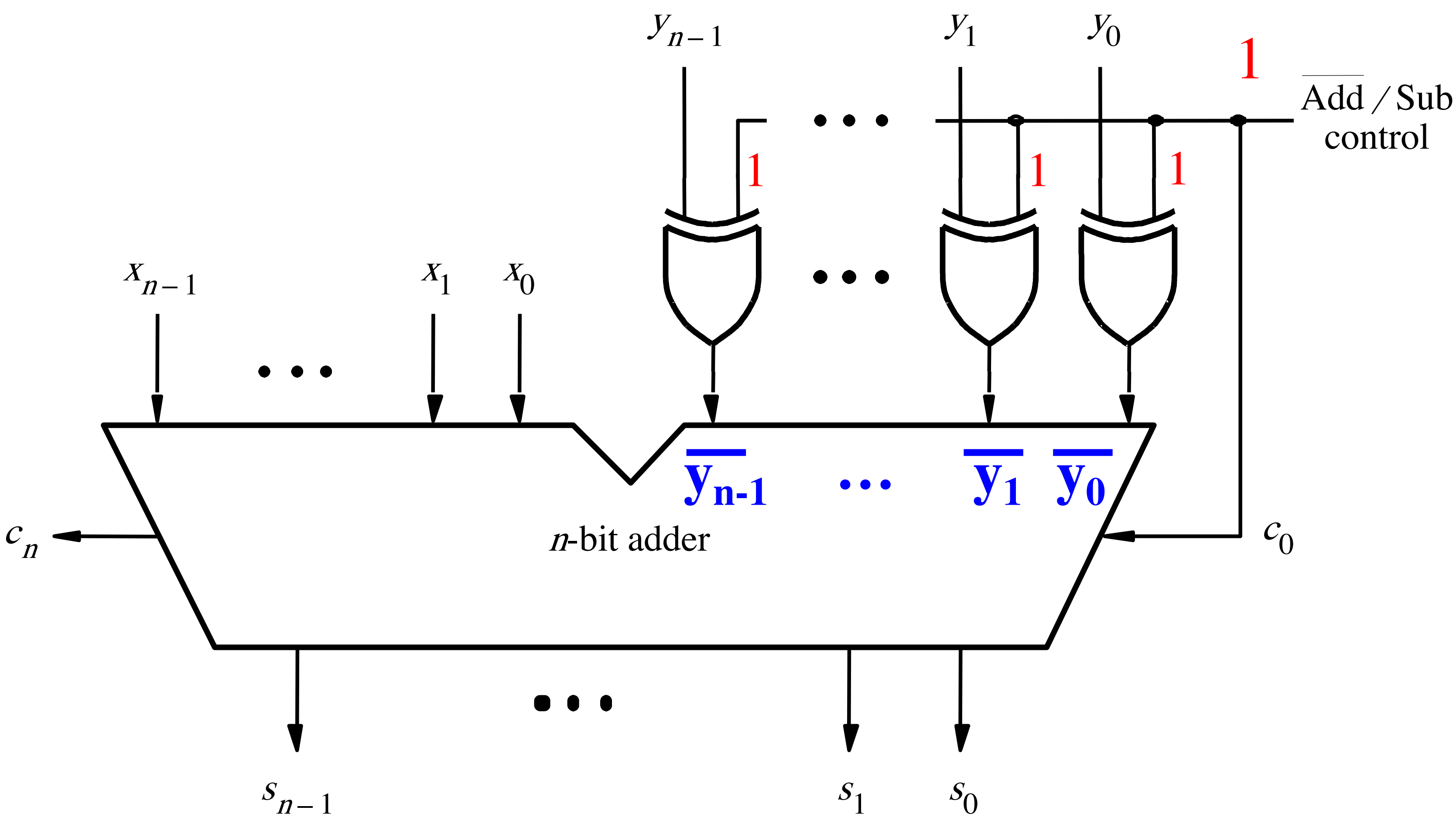
[Figure 3.12 from the textbook]

Subtraction: when control = 1



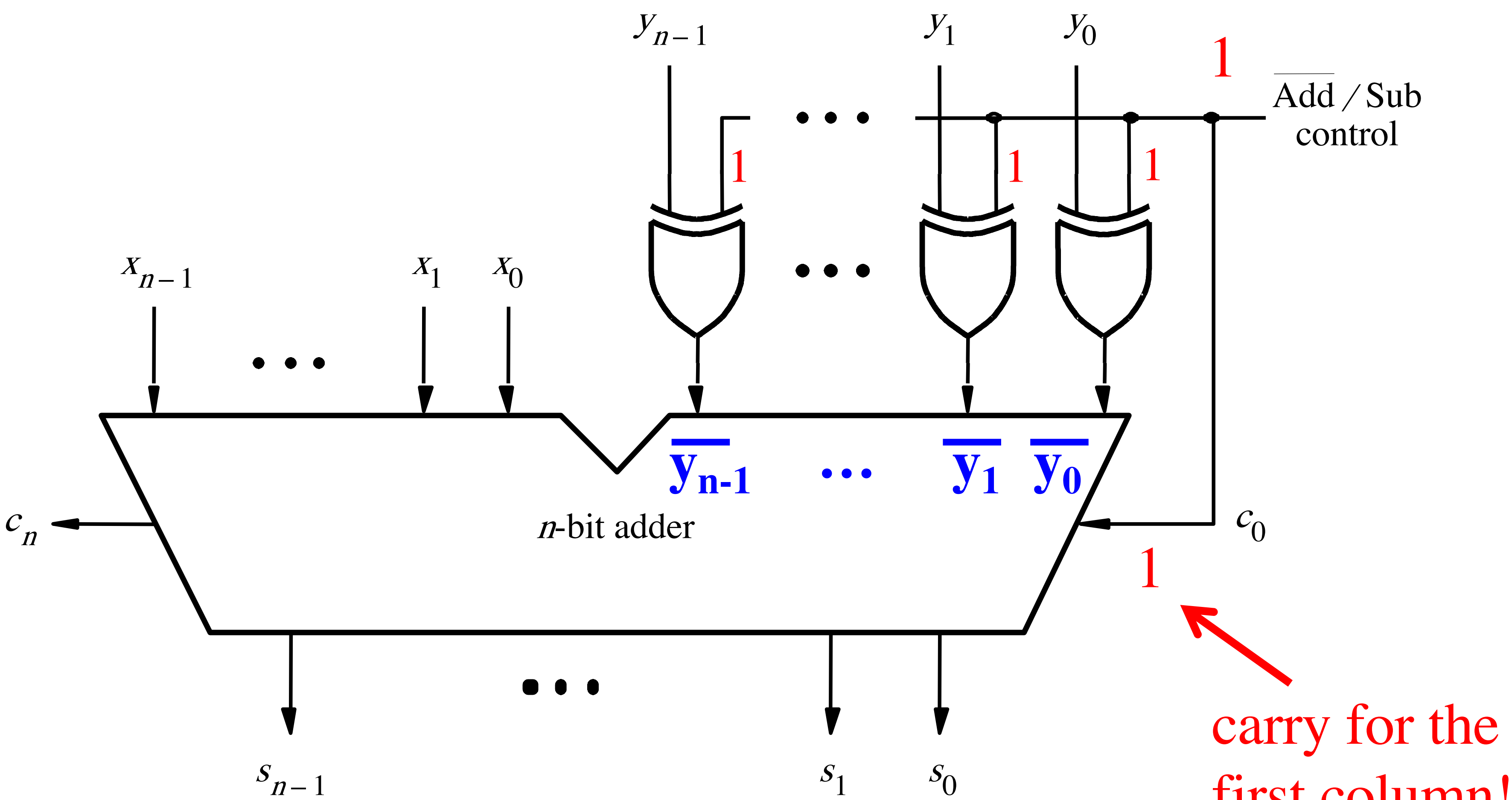
[Figure 3.12 from the textbook]

Subtraction: when control = 1



[Figure 3.12 from the textbook]

Subtraction: when control = 1



[Figure 3.12 from the textbook]

Detecting Overflow

Examples of determination of overflow

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{r} 0111 \\ +\quad 0010 \\ \hline 1001 \end{array} \qquad\qquad \begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{r} 1001 \\ +\quad 0010 \\ \hline 1011 \end{array}$$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{r} 0111 \\ +\quad 1110 \\ \hline 10101 \end{array} \qquad\qquad \begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{r} 1001 \\ +\quad 1110 \\ \hline 10111 \end{array}$$

[Figure 3.13 from the textbook]

Examples of determination of overflow

$$\begin{array}{rr} & \\ (+7) & \\ +\ (+2) & + \\ \hline (+9) & \end{array} \quad \begin{array}{r} 0\,1\,1\,0\,0 \\ 0\,1\,1\,1 \\ 0\,0\,1\,0 \\ \hline 1\,0\,0\,1 \end{array} \qquad\qquad \begin{array}{rr} & \\ (-7) & \\ +\ (+2) & + \\ \hline (-5) & \end{array} \quad \begin{array}{r} 0\,0\,0\,0\,0 \\ 1\,0\,0\,1 \\ 0\,0\,1\,0 \\ \hline 1\,0\,1\,1 \end{array}$$

$$\begin{array}{rr} & \\ (+7) & \\ +\ (-2) & + \\ \hline (+5) & \end{array} \quad \begin{array}{r} 1\,1\,1\,0\,0 \\ 0\,1\,1\,1 \\ 1\,1\,1\,0 \\ \hline 1\,0\,1\,0\,1 \end{array} \qquad\qquad \begin{array}{rr} & \\ (-7) & \\ +\ (-2) & + \\ \hline (-9) & \end{array} \quad \begin{array}{r} 1\,0\,0\,0\,0 \\ 1\,0\,0\,1 \\ 1\,1\,1\,0 \\ \hline 1\,0\,1\,1\,1 \end{array}$$

Include the carry bits: $c_4\ c_3\ c_2\ c_1\ c_0$

Examples of determination of overflow

$$
\begin{array}{r}
(+7) \\
+\ (+2) \\
\hline
(+9)
\end{array}
\qquad
\begin{array}{rr}
 & \boxed{0\ 1}\ 1\ 0\ 0 \\
 & 0\ 1\ 1\ 1 \\
+ & 0\ 0\ 1\ 0 \\
\hline
 & 1\ 0\ 0\ 1
\end{array}
\qquad\qquad
\begin{array}{r}
(-7) \\
+\ (+2) \\
\hline
(-5)
\end{array}
\qquad
\begin{array}{rr}
 & \boxed{0\ 0}\ 0\ 0\ 0 \\
 & 1\ 0\ 0\ 1 \\
+ & 0\ 0\ 1\ 0 \\
\hline
 & 1\ 0\ 1\ 1
\end{array}
$$

$$
\begin{array}{r}
(+7) \\
+\ (-2) \\
\hline
(+5)
\end{array}
\qquad
\begin{array}{rr}
 & \boxed{1\ 1}\ 1\ 0\ 0 \\
 & 0\ 1\ 1\ 1 \\
+ & 1\ 1\ 1\ 0 \\
\hline
 & 1\ \ 0\ 1\ 0\ 1
\end{array}
\qquad\qquad
\begin{array}{r}
(-7) \\
+\ (-2) \\
\hline
(-9)
\end{array}
\qquad
\begin{array}{rr}
 & \boxed{1\ 0}\ 0\ 0\ 0 \\
 & 1\ 0\ 0\ 1 \\
+ & 1\ 1\ 1\ 0 \\
\hline
 & 1\ 0\ 1\ 1\ 1
\end{array}
$$

Include the carry bits: $\boxed{c_4\ c_3}\ c_2\ c_1\ c_0$

Examples of determination of overflow

$c_4 = 0$
$c_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{cr} & \boxed{0\ 1}\,1\ 0\ 0 \\ & 0\ 1\ 1\ 1 \\ + & 0\ 0\ 1\ 0 \\ \hline & 1\ 0\ 0\ 1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{cr} & \boxed{0\ 0}\,0\ 0\ 0 \\ & 1\ 0\ 0\ 1 \\ + & 0\ 0\ 1\ 0 \\ \hline & 1\ 0\ 1\ 1 \end{array}$$

$c_4 = 0$
$c_3 = 0$

$c_4 = 1$
$c_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{cr} & \boxed{1\ 1}\,1\ 0\ 0 \\ & 0\ 1\ 1\ 1 \\ + & 1\ 1\ 1\ 0 \\ \hline & 1\ 0\ 1\ 0\ 1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{cr} & \boxed{1\ 0}\,0\ 0\ 0 \\ & 1\ 0\ 0\ 1 \\ + & 1\ 1\ 1\ 0 \\ \hline & 1\ 0\ 1\ 1\ 1 \end{array}$$

$c_4 = 1$
$c_3 = 0$

Include the carry bits: $\boxed{c_4\ c_3}\, c_2\ c_1\ c_0$

Examples of determination of overflow

$$c_4 = 0 \qquad c_3 = 1$$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{rr} & \boxed{0\ 1}\,1\ 0\ 0 \\ & 0\ 1\ 1\ 1 \\ + & 0\ 0\ 1\ 0 \\ \hline & 1\ 0\ 0\ 1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{rr} & \boxed{0\ 0}\,0\ 0\ 0 \\ & 1\ 0\ 0\ 1 \\ + & 0\ 0\ 1\ 0 \\ \hline & 1\ 0\ 1\ 1 \end{array}$$

$$c_4 = 0 \qquad c_3 = 0$$

$$c_4 = 1 \qquad c_3 = 1$$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{rr} & \boxed{1\ 1}\,1\ 0\ 0 \\ & 0\ 1\ 1\ 1 \\ + & 1\ 1\ 1\ 0 \\ \hline & 1\ 0\ 1\ 0\ 1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{rr} & \boxed{1\ 0}\,0\ 0\ 0 \\ & 1\ 0\ 0\ 1 \\ + & 1\ 1\ 1\ 0 \\ \hline & 1\ 0\ 1\ 1\ 1 \end{array}$$

$$c_4 = 1 \qquad c_3 = 0$$

Overflow occurs only in these two cases.

Examples of determination of overflow

$c_4 = 0$
$c_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{r} \boxed{0\ 1}\,1\ 0\ 0 \\ 0\ 1\ 1\ 1 \\ +\quad 0\ 0\ 1\ 0 \\ \hline 1\ 0\ 0\ 1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{r} \boxed{0\ 0}\,0\ 0\ 0 \\ 1\ 0\ 0\ 1 \\ +\quad 0\ 0\ 1\ 0 \\ \hline 1\ 0\ 1\ 1 \end{array}$$

$c_4 = 0$
$c_3 = 0$

$c_4 = 1$
$c_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{r} \boxed{1\ 1}\,1\ 0\ 0 \\ 0\ 1\ 1\ 1 \\ +\quad 1\ 1\ 1\ 0 \\ \hline 1\ 0\ 1\ 0\ 1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{r} \boxed{1\ 0}\,0\ 0\ 0 \\ 1\ 0\ 0\ 1 \\ +\quad 1\ 1\ 1\ 0 \\ \hline 1\ 0\ 1\ 1\ 1 \end{array}$$

$c_4 = 1$
$c_3 = 0$

$$\text{Overflow} = c_3\overline{c}_4 + \overline{c}_3 c_4$$

Examples of determination of overflow

$c_4 = 0$
$c_3 = 1$

$$\begin{array}{r} \\ (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{rr} & \boxed{0\ 1}\,1\ 0\ 0 \\ & 0\ 1\ 1\ 1 \\ + & 0\ 0\ 1\ 0 \\ \hline & 1\ 0\ 0\ 1 \end{array}$$

$$\begin{array}{r} \\ (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{rr} & \boxed{0\ 0}\,0\ 0\ 0 \\ & 1\ 0\ 0\ 1 \\ + & 0\ 0\ 1\ 0 \\ \hline & 1\ 0\ 1\ 1 \end{array}$$

$c_4 = 0$
$c_3 = 0$

$c_4 = 1$
$c_3 = 1$

$$\begin{array}{r} \\ (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{rr} & \boxed{1\ 1}\,1\ 0\ 0 \\ & 0\ 1\ 1\ 1 \\ + & 1\ 1\ 1\ 0 \\ \hline & 1\ 0\ 1\ 0\ 1 \end{array}$$

$$\begin{array}{r} \\ (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{rr} & \boxed{1\ 0}\,0\ 0\ 0 \\ & 1\ 0\ 0\ 1 \\ + & 1\ 1\ 1\ 0 \\ \hline & 1\ 0\ 1\ 1\ 1 \end{array}$$

$c_4 = 1$
$c_3 = 0$

$$\text{Overflow} = \underbrace{c_3\overline{c}_4 + \overline{c}_3 c_4}_{\text{XOR}}$$

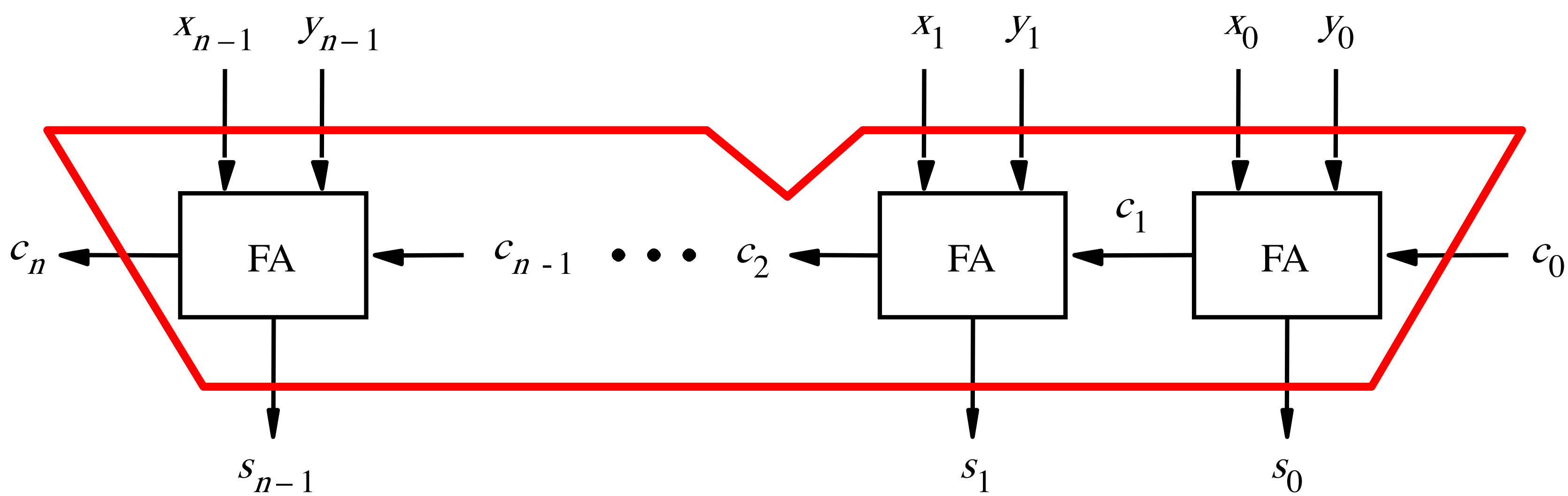
Calculating overflow for 4-bit numbers with only three significant bits

$$\begin{aligned}\text{Overflow} &= c_3\overline{c}_4 + \overline{c}_3 c_4 \\ &= c_3 \oplus c_4\end{aligned}$$

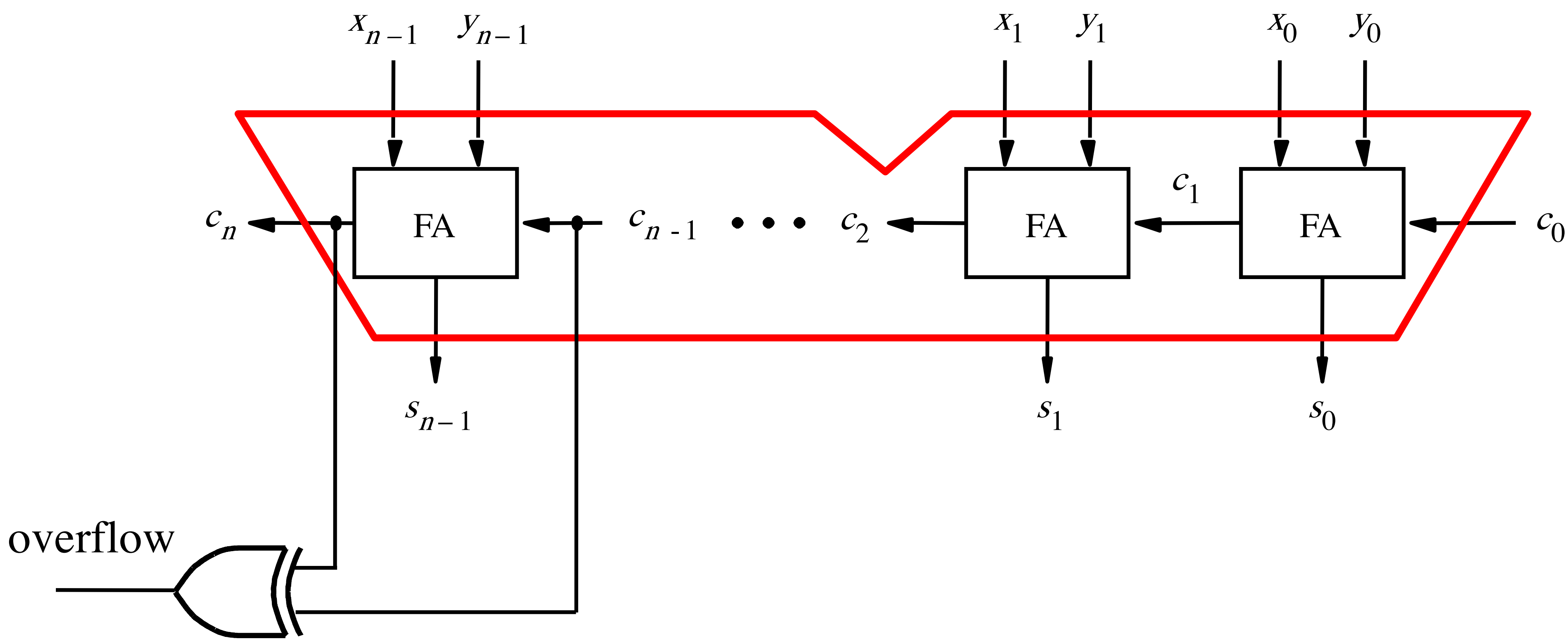
Calculating overflow for n-bit numbers with only n-1 significant bits

$$\text{Overflow} = c_{n-1} \oplus c_n$$

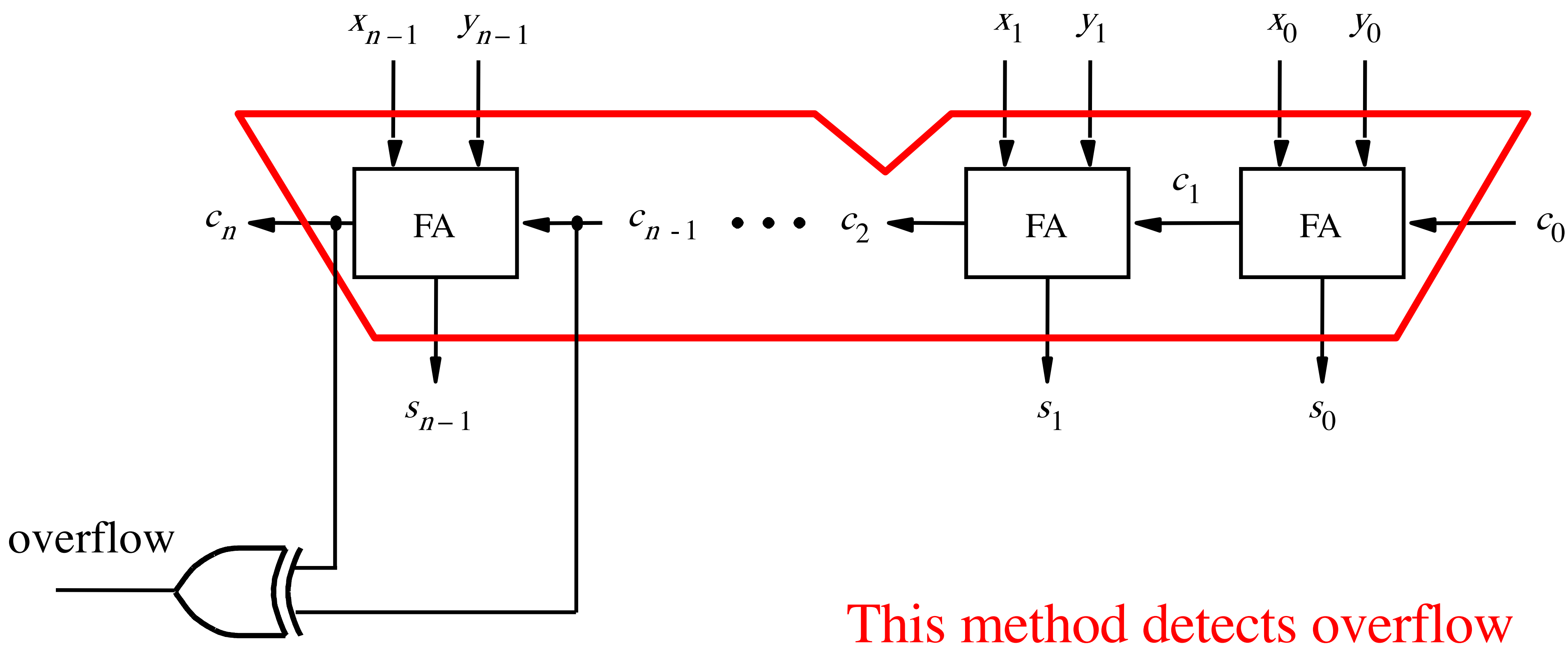
Detecting Overflow



Detecting Overflow (with one extra XOR)



Detecting Overflow (with one extra XOR)



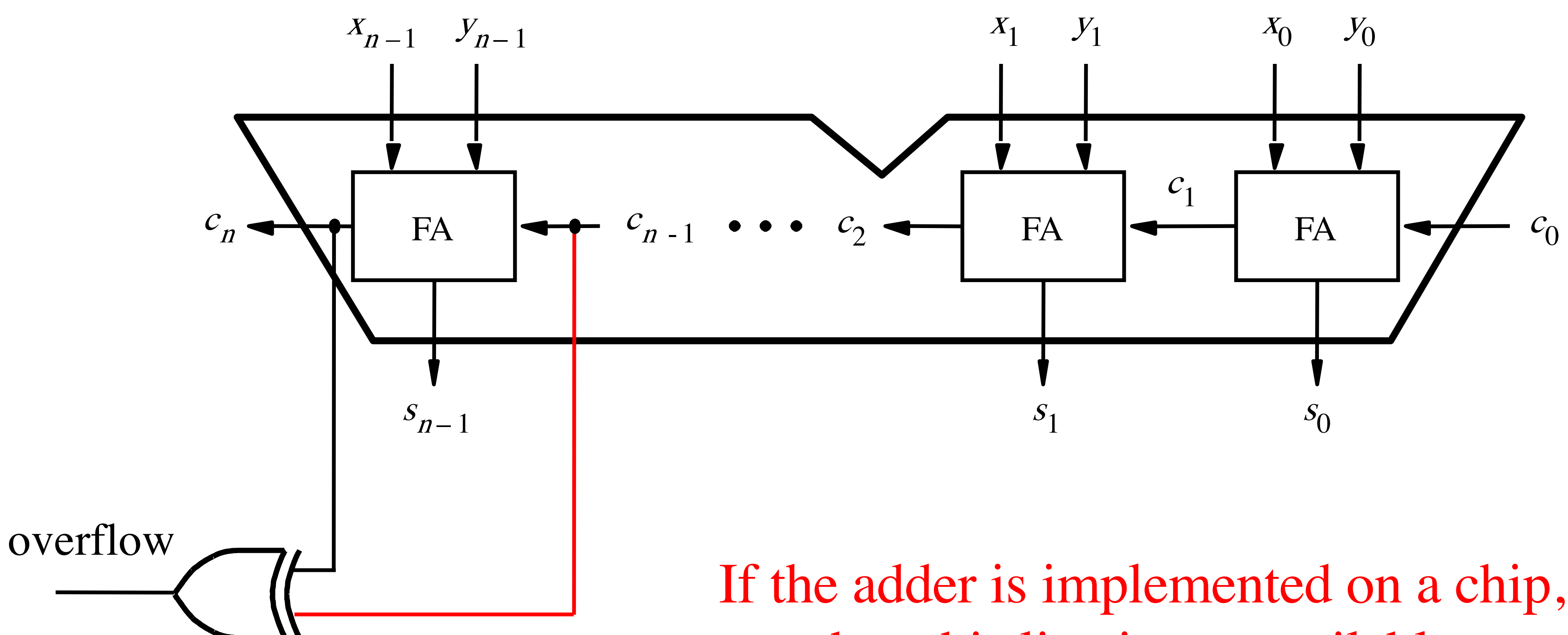
This method detects overflow for both addition and subtraction.

Detecting Overflow
(alternative method)

Detecting Overflow
(alternative method)

Used if you don't have access to the internal carries of the adder.

Detecting Overflow
(with one extra XOR)



If the adder is implemented on a chip,
then this line is not available.
So the first method can't be used.

Another way to look at the overflow issue

$$
\begin{array}{rl}
 & X = x_3\ x_2\ x_1\ x_0 \\
+ & Y = y_3\ y_2\ y_1\ y_0 \\
\hline
 & S = s_3\ s_2\ s_1\ s_0
\end{array}
$$

Another way to look at the overflow issue

$$
\begin{array}{rcccc}
 & X= & \boxed{x_3} & x_2 & x_1 & x_0 \\
+ & Y= & \boxed{y_3} & y_2 & y_1 & y_0 \\
\hline
 & S= & \boxed{s_3} & s_2 & s_1 & s_0
\end{array}
$$

If both numbers that we are adding have the same sign but the sum does not, then we have an overflow.

Examples of determination of overflow

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{r} 0\,1\,1\,1 \\ +\quad 0\,0\,1\,0 \\ \hline 1\,0\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{r} 1\,0\,0\,1 \\ +\quad 0\,0\,1\,0 \\ \hline 1\,0\,1\,1 \end{array}$$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{r} 0\,1\,1\,1 \\ +\quad 1\,1\,1\,0 \\ \hline 1\,0\,1\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{r} 1\,0\,0\,1 \\ +\quad 1\,1\,1\,0 \\ \hline 1\,0\,1\,1\,1 \end{array}$$

Examples of determination of overflow

$$
\begin{array}{rr}
 & (+7) \\
+ & (+2) \\
\hline
 & (+9)
\end{array}
\qquad
\begin{array}{rcccc}
 & 0 & 1 & 1 & 1 \\
+ & 0 & 0 & 1 & 0 \\
\hline
 & 1 & 0 & 0 & 1
\end{array}
\qquad\qquad
\begin{array}{rr}
 & (-7) \\
+ & (+2) \\
\hline
 & (-5)
\end{array}
\qquad
\begin{array}{rcccc}
 & 1 & 0 & 0 & 1 \\
+ & 0 & 0 & 1 & 0 \\
\hline
 & 1 & 0 & 1 & 1
\end{array}
$$

$$
\begin{array}{rr}
 & (+7) \\
+ & (-2) \\
\hline
 & (+5)
\end{array}
\qquad
\begin{array}{rccccc}
 & & 0 & 1 & 1 & 1 \\
+ & & 1 & 1 & 1 & 0 \\
\hline
 & 1 & 0 & 1 & 0 & 1
\end{array}
\qquad\qquad
\begin{array}{rr}
 & (-7) \\
+ & (-2) \\
\hline
 & (-9)
\end{array}
\qquad
\begin{array}{rccccc}
 & & 1 & 0 & 0 & 1 \\
+ & & 1 & 1 & 1 & 0 \\
\hline
 & 1 & 0 & 1 & 1 & 1
\end{array}
$$

Examples of determination of overflow

$x_3 = 0$
$y_3 = 0$
$s_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{rr} & 0\,1\,1\,1 \\ + & 0\,0\,1\,0 \\ \hline & 1\,0\,0\,1 \end{array}$$

$x_3 = 1$
$y_3 = 0$
$s_3 = 1$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{rr} & 1\,0\,0\,1 \\ + & 0\,0\,1\,0 \\ \hline & 1\,0\,1\,1 \end{array}$$

$x_3 = 0$
$y_3 = 1$
$s_3 = 0$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{rr} & 0\,1\,1\,1 \\ + & 1\,1\,1\,0 \\ \hline & 1\,0\,1\,0\,1 \end{array}$$

$x_3 = 1$
$y_3 = 1$
$s_3 = 0$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{rr} & 1\,0\,0\,1 \\ + & 1\,1\,1\,0 \\ \hline & 1\,0\,1\,1\,1 \end{array}$$

Examples of determination of overflow

$x_3 = 0$, $y_3 = 0$, $s_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{rr} & 0\,1\,1\,1 \\ + & 0\,0\,1\,0 \\ \hline & 1\,0\,0\,1 \end{array}$$

$x_3 = 1$, $y_3 = 0$, $s_3 = 1$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{rr} & 1\,0\,0\,1 \\ + & 0\,0\,1\,0 \\ \hline & 1\,0\,1\,1 \end{array}$$

$x_3 = 0$, $y_3 = 1$, $s_3 = 0$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{rr} & 0\,1\,1\,1 \\ + & 1\,1\,1\,0 \\ \hline & 1\,0\,1\,0\,1 \end{array}$$

$x_3 = 1$, $y_3 = 1$, $s_3 = 0$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{rr} & 1\,0\,0\,1 \\ + & 1\,1\,1\,0 \\ \hline & 1\,0\,1\,1\,1 \end{array}$$

In 2's complement, both +9 and -9 are not representable with 4 bits.

Examples of determination of overflow

$x_3 = 0$
$y_3 = 0$
$s_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{cr} & 0\,1\,1\,1 \\ + & 0\,0\,1\,0 \\ \hline & 1\,0\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{cr} & 1\,0\,0\,1 \\ + & 0\,0\,1\,0 \\ \hline & 1\,0\,1\,1 \end{array}$$

$x_3 = 1$
$y_3 = 0$
$s_3 = 1$

$x_3 = 0$
$y_3 = 1$
$s_3 = 0$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{cr} & 0\,1\,1\,1 \\ + & 1\,1\,1\,0 \\ \hline & 1\,0\,1\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{cr} & 1\,0\,0\,1 \\ + & 1\,1\,1\,0 \\ \hline & 1\,0\,1\,1\,1 \end{array}$$

$x_3 = 1$
$y_3 = 1$
$s_3 = 0$

Overflow occurs only in these two cases.

Examples of determination of overflow

$x_3 = 0$, $y_3 = 0$, $s_3 = 1$

$$\begin{array}{r} (+7) \\ +\,(+2) \\ \hline (+9) \end{array} \qquad \begin{array}{rr} & 0\,1\,1\,1 \\ + & 0\,0\,1\,0 \\ \hline & 1\,0\,0\,1 \end{array}$$

$x_3 = 1$, $y_3 = 0$, $s_3 = 1$

$$\begin{array}{r} (-7) \\ +\,(+2) \\ \hline (-5) \end{array} \qquad \begin{array}{rr} & 1\,0\,0\,1 \\ + & 0\,0\,1\,0 \\ \hline & 1\,0\,1\,1 \end{array}$$

$x_3 = 0$, $y_3 = 1$, $s_3 = 0$

$$\begin{array}{r} (+7) \\ +\,(-2) \\ \hline (+5) \end{array} \qquad \begin{array}{rr} & 0\,1\,1\,1 \\ + & 1\,1\,1\,0 \\ \hline & 1\,0\,1\,0\,1 \end{array}$$

$x_3 = 1$, $y_3 = 1$, $s_3 = 0$

$$\begin{array}{r} (-7) \\ +\,(-2) \\ \hline (-9) \end{array} \qquad \begin{array}{rr} & 1\,0\,0\,1 \\ + & 1\,1\,1\,0 \\ \hline & 1\,0\,1\,1\,1 \end{array}$$

$$\text{Overflow} = \overline{x}_3\,\overline{y}_3\,s_3 + x_3\,y_3\,\overline{s}_3$$

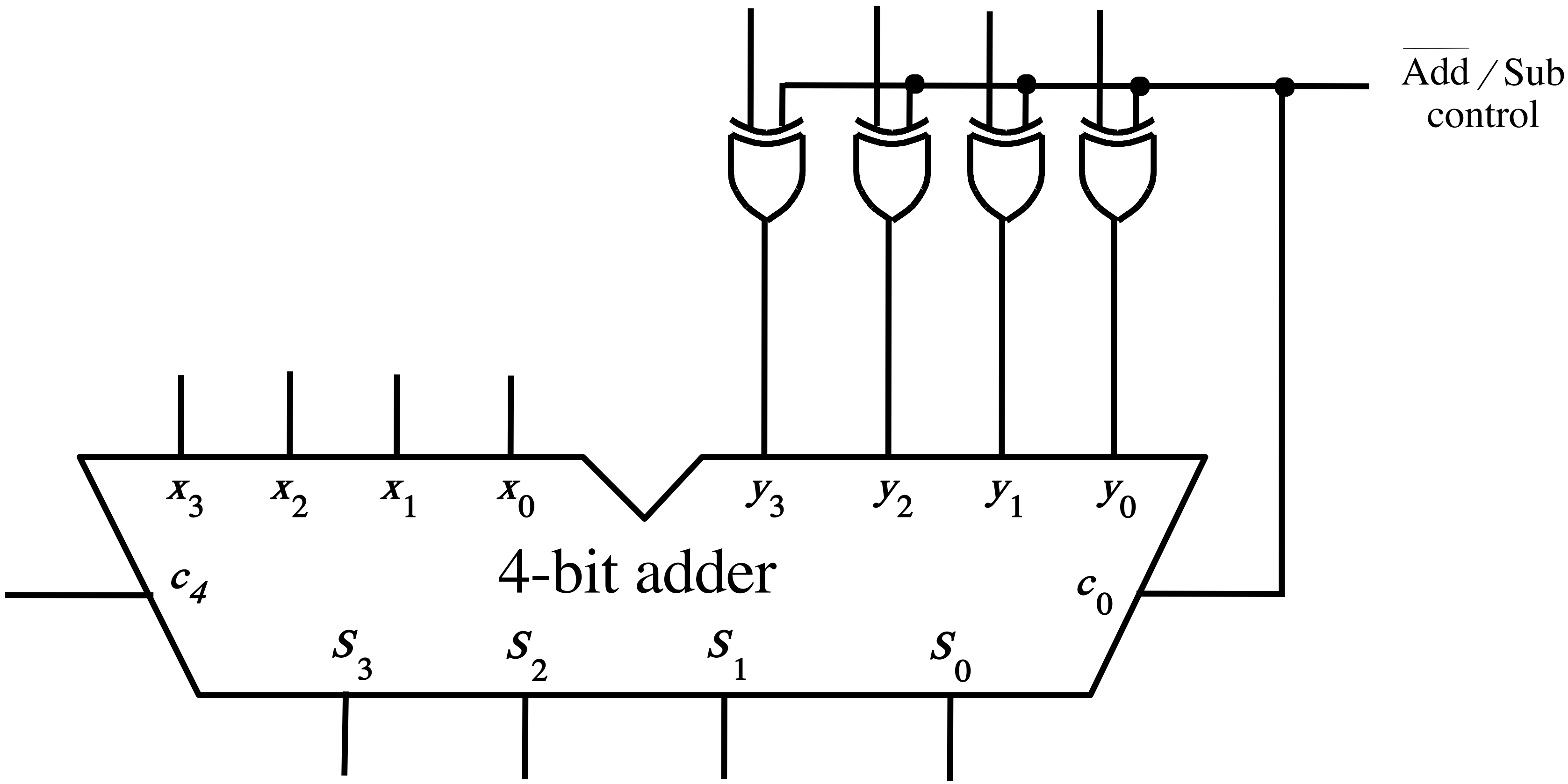
Another way to look at the overflow issue

$$
\begin{array}{rcccc}
 & X= & \boxed{x_3} & x_2 & x_1 & x_0 \\
+ & Y= & \boxed{y_3} & y_2 & y_1 & y_0 \\
\hline
 & S= & \boxed{s_3} & s_2 & s_1 & s_0
\end{array}
$$

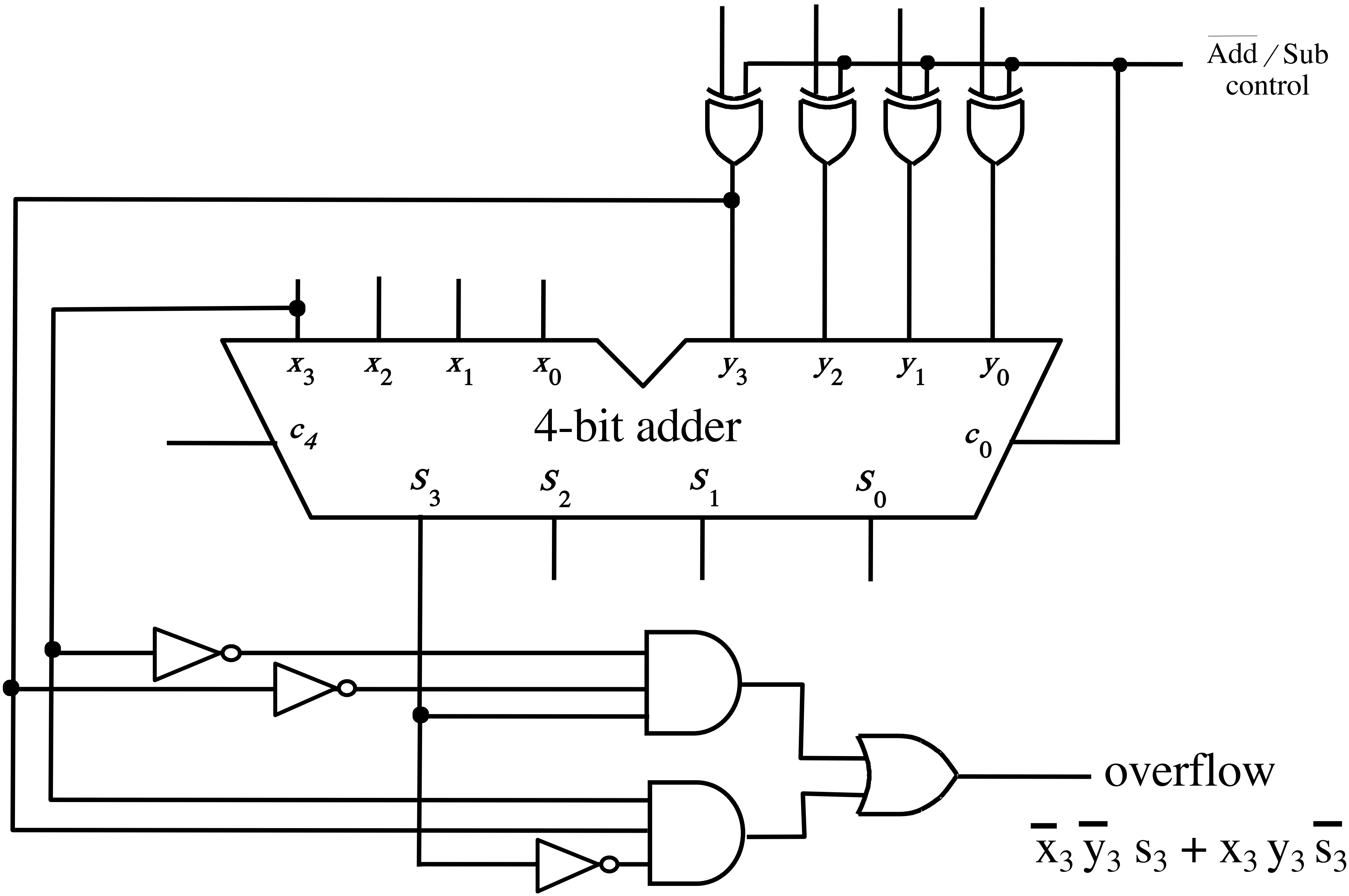
If both numbers that we are adding have the same sign but the sum does not, then we have an overflow.

$$\text{Overflow} = \overline{x}_3\,\overline{y}_3\,s_3 + x_3\,y_3\,\overline{s}_3$$

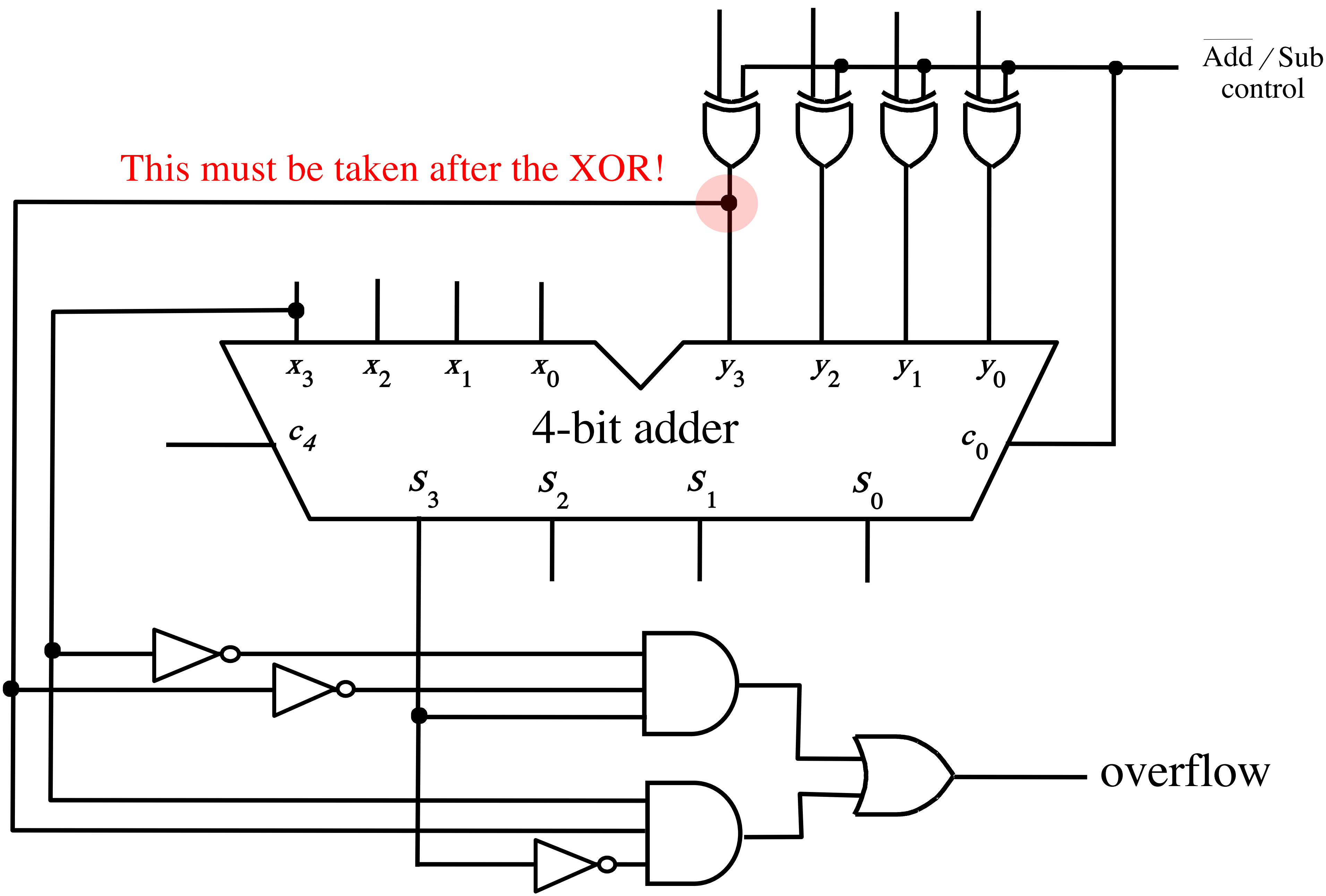
Overflow Detection



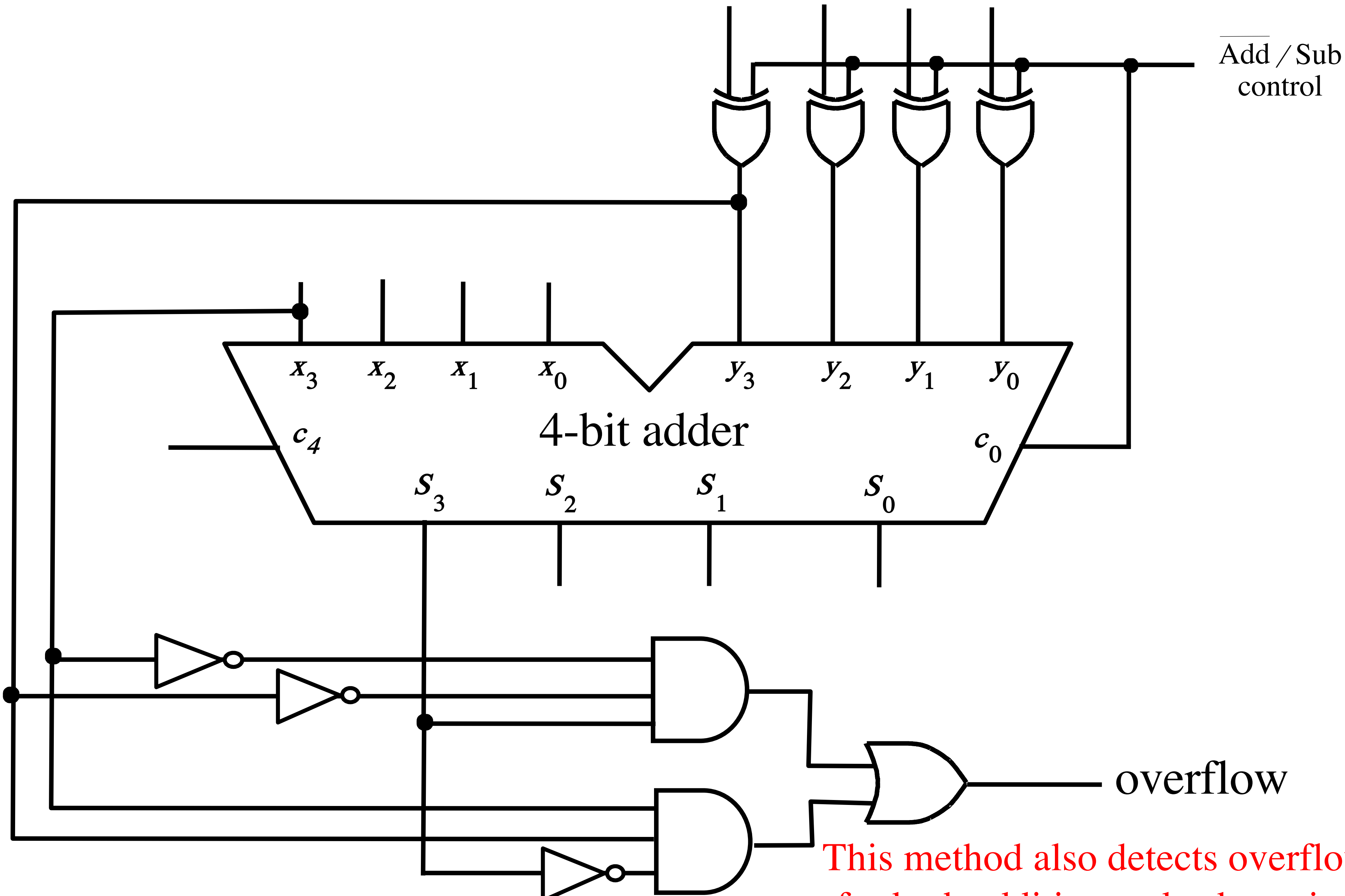
Overflow Detection



Overflow Detection



Overflow Detection



This method also detects overflow for both addition and subtraction.

Questions?

THE END