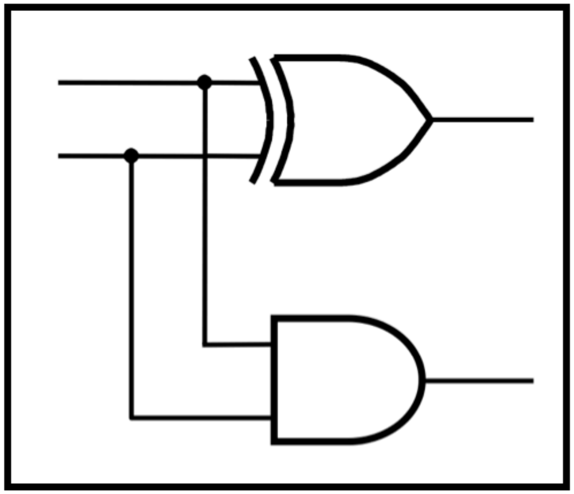


CprE 281: Digital Logic

Instructor: Alexander Stoytchev

http://www.ece.iastate.edu/~alexs/classes/

Fast Adders

CprE 281: Digital Logic
Iowa State University, Ames, IA

Administrative Stuff

- HW5 is out
- It is due on Monday Oct 3 @ 4pm.
- Please write clearly on the first page (in block capital letters) the following three things:
  - Your First and Last Name
  - Your Student ID Number
  - Your Lab Section Letter
- Also, please staple all of your pages together.

Administrative Stuff

- **Labs Next Week**
- **Mini-Project**
- **This one is worth 3% of your grade.**
- **Make sure to get all the points.**
- **http://www.ece.iastate.edu/~alexs/classes/2016_Fall_281/labs/Project-Mini/**

Quick Review

The problems in which row are easier to calculate?

$$\begin{array}{r} 82 \\ -\ 61 \\ \hline ?? \end{array} \qquad \begin{array}{r} 48 \\ -\ 26 \\ \hline ?? \end{array} \qquad \begin{array}{r} 32 \\ -\ 11 \\ \hline ?? \end{array}$$

$$\begin{array}{r} 82 \\ -\ 64 \\ \hline ?? \end{array} \qquad \begin{array}{r} 48 \\ -\ 29 \\ \hline ?? \end{array} \qquad \begin{array}{r} 32 \\ -\ 13 \\ \hline ?? \end{array}$$

The problems in which row are easier to calculate?

$$\begin{array}{r} 82 \\ -\ 61 \\ \hline 21 \end{array} \qquad \begin{array}{r} 48 \\ -\ 26 \\ \hline 22 \end{array} \qquad \begin{array}{r} 32 \\ -\ 11 \\ \hline 21 \end{array}$$

Why?

$$\begin{array}{r} 82 \\ -\ 64 \\ \hline 18 \end{array} \qquad \begin{array}{r} 48 \\ -\ 29 \\ \hline 19 \end{array} \qquad \begin{array}{r} 32 \\ -\ 13 \\ \hline 19 \end{array}$$

Another Way to Do Subtraction

$$82 - 64 = 82 + 100 - 100 - 64$$

Another Way to Do Subtraction

$$82 – 64 = 82 + 100 – 100 - 64$$

$$= 82 + (100 – 64) - 100$$

Another Way to Do Subtraction

$$82 – 64 = 82 + 100 – 100 - 64$$

$$= 82 + (100 – 64) - 100$$

$$= 82 + (99 + 1 – 64) - 100$$

Another Way to Do Subtraction

$$82 - 64 = 82 + 100 - 100 - 64$$

$$= 82 + (100 - 64) - 100$$

$$= 82 + (99 + 1 - 64) - 100$$

$$= 82 + (99 - 64) + 1 - 100$$

Another Way to Do Subtraction

$$82 – 64 = 82 + 100 – 100 - 64$$

$$= 82 + (100 – 64) - 100$$

$$= 82 + (99 + 1 – 64) - 100$$

Does not require borrows

$$= 82 + (99 – 64) +1 - 100$$

9's Complement
(subtract each digit from 9)

$$\begin{array}{r} 99 \\ -\ 64 \\ \hline 35 \end{array}$$

10's Complement
(subtract each digit from 9 and add 1 to the result)

$$\begin{array}{r} 99 \\ -\ 64 \\ \hline 35 \end{array} + 1 = 36$$

Another Way to Do Subtraction

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

$$= 82 + 35 + 1 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

10's complement

$$= 82 + 35 + 1 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

10's complement

$$= 82 + 35 + 1 - 100$$

$$= 82 + 36 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

10's complement

$$= 82 + 35 + 1 - 100$$

$$= 82 + 36 - 100$$ // Add the first two.

$$= 118 - 100$$

Another Way to Do Subtraction

9's complement

$$82 - 64 = 82 + (99 - 64) + 1 - 100$$

10's complement

$$= 82 + 35 + 1 - 100$$

$$= 82 + 36 - 100$$ // Add the first two.

$$= 118 - 100$$ // Just delete the leading 1.
// No need to subtract 100.

$$= 18$$

2's complement

Let K be the negative equivalent of an n-bit positive number P.

Then, in 2's complement representation K is obtained by subtracting P from 2^n , namely

$$K = 2^n - P$$

Deriving 2's complement

For a positive n-bit number P, let K_1 and K_2 denote its 1's and 2's complements, respectively.

$$K_1 = (2^n - 1) - P$$

$$K_2 = 2^n - P$$

Since $K_2 = K_1 + 1$, it is evident that in a logic circuit the 2's complement can computed by inverting all bits of P and then adding 1 to the resulting 1's-complement number.

Find the 2's complement of …

0 1 0 1

0 0 1 0

0 1 0 0

0 1 1 1

Find the 2's complement of …

0 1 0 1

1 0 1 0

0 0 1 0

1 1 0 1

0 1 0 0

1 0 1 1

0 1 1 1

1 0 0 0

Invert all bits.

Find the 2's complement of …

```
  0 1 0 1            0 0 1 0
  1 0 1 0            1 1 0 1
+       1          +       1
---------          ---------
  1 0 1 1            1 1 1 0


  0 1 0 0            0 1 1 1
  1 0 1 1            1 0 0 0
+       1          +       1
---------          ---------
  1 1 0 0            1 0 0 1
```

Then add 1.

Quick Way to find 2's complement

- Scan the binary number from right to left
- Copy all bits that are 0 from right to left
- Stop at the first 1
- Copy that 1 as well
- Invert all remaining bits

Find the 2's complement of …

0 1 0 1

0 0 1 0

0 1 0 0

0 1 1 1

Find the 2's complement of …

0 1 0 1

. . . .

0 0 1 0

. . . 0

0 1 0 0

. . 0 0

0 1 1 1

. . . .

Copy all bits that are 0 from right to left.

Find the 2's complement of …

0 1 0 1

. . . 1

0 0 1 0

. . 1 0

0 1 0 0

. 1 0 0

0 1 1 1

. . . 1

Stop at the first 1. Copy that 1 as well.

Find the 2's complement of …

0 1 0 1
1 0 1 1

0 0 1 0
1 1 1 0

0 1 0 0
1 1 0 0

0 1 1 1
1 0 0 1

Invert all remaining bits.

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

[Table 3.2 from the textbook]

Interpretation of four-bit signed integers

$b_3 b_2 b_1 b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

The top half is the same in all three representations.
It corresponds to the positive integers.

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

In all three representations the first bit represents the sign.
If that bit is 1, then the number is negative.

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

Notice that in this representation there are two zeros!

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

There are two zeros in this representation as well!

Interpretation of four-bit signed integers

$b_3b_2b_1b_0$	Sign and magnitude	1's complement	2's complement
0111	+7	+7	+7
0110	+6	+6	+6
0101	+5	+5	+5
0100	+4	+4	+4
0011	+3	+3	+3
0010	+2	+2	+2
0001	+1	+1	+1
0000	+0	+0	+0
1000	−0	−7	−8
1001	−1	−6	−7
1010	−2	−5	−6
1011	−3	−4	−5
1100	−4	−3	−4
1101	−5	−2	−3
1110	−6	−1	−2
1111	−7	−0	−1

In this representation there is one more negative number.

Taking the 2's complement negates the number

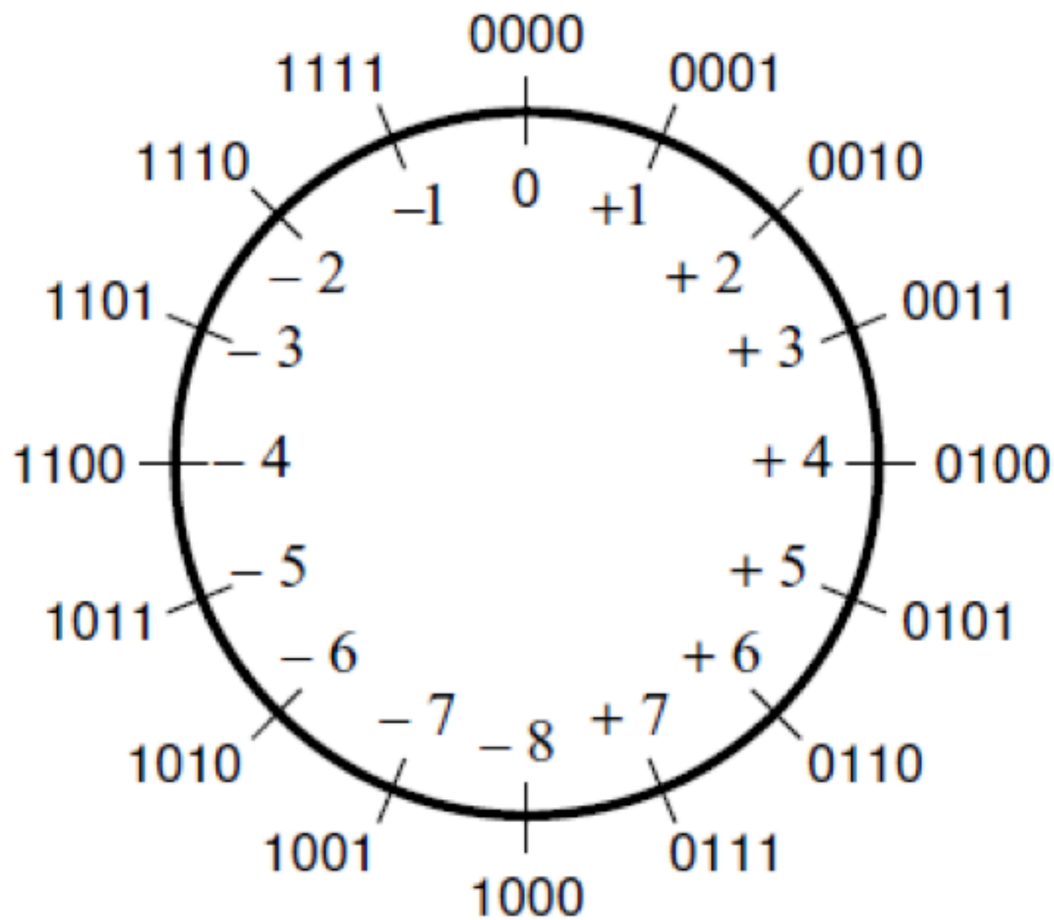
decimal	$b_3\,b_2\,b_1\,b_0$	take the 2's complement	$b_3\,b_2\,b_1\,b_0$	decimal
+7	0111	⟹	1001	−7
+6	0110	⟹	1010	−6
+5	0101	⟹	1011	−5
+4	0100	⟹	1100	−4
+3	0011	⟹	1101	−3
+2	0010	⟹	1110	−2
+1	0001	⟹	1111	−1
+0	0000	⟹	0000	+0
−8	1000	⟹	1000	−8
−7	1001	⟹	0111	+7
−6	1010	⟹	0110	+6
−5	1011	⟹	0101	+5
−4	1100	⟹	0100	+4
−3	1101	⟹	0011	+3
−2	1110	⟹	0010	+2
−1	1111	⟹	0001	+1

Taking the 2's complement negates the number

decimal	$b_3\, b_2\, b_1\, b_0$	take the 2's complement	$b_3\, b_2\, b_1\, b_0$	decimal
+7	0111	⟹	1001	−7
+6	0110	⟹	1010	−6
+5	0101	⟹	1011	−5
+4	0100	⟹	1100	−4
+3	0011	⟹	1101	−3
+2	0010	⟹	1110	−2
+1	0001	⟹	1111	−1
+0	0000	⟹	0000	+0
−8	1000	⟹	1000	−8
−7	1001	⟹	0111	+7
−6	1010	⟹	0110	+6
−5	1011	⟹	0101	+5
−4	1100	⟹	0100	+4
−3	1101	⟹	0011	+3
−2	1110	⟹	0010	+2
−1	1111	⟹	0001	+1

This is the only exception

The number circle for 2's complement



[Figure 3.11a from the textbook]

A) Example of 2's complement addition

(+ 5)		0 1 0 1
+ (+ 2)	+	0 0 1 0
(+ 7)		0 1 1 1

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.9 from the textbook]

B) Example of 2's complement addition

	(–5)	1 0 1 1
+	(+ 2)	+ 0 0 1 0
	(–3)	1 1 0 1

$b_3 b_2 b_1 b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.9 from the textbook]

C) Example of 2's complement addition

	(+ 5)	0 1 0 1
+	(–2)	+ 1 1 1 0
	(+ 3)	1 0 0 1 1

↑ ignore

$b_3b_2b_1b_0$	2's complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.9 from the textbook]

D) Example of 2’s complement addition

	(–5)	1 0 1 1
+	(–2)	+ 1 1 1 0
	(–7)	1 1 0 0 1

(the leading 1 of the result: ignore)

$b_3b_2b_1b_0$	2’s complement
0111	+7
0110	+6
0101	+5
0100	+4
0011	+3
0010	+2
0001	+1
0000	+0
1000	−8
1001	−7
1010	−6
1011	−5
1100	−4
1101	−3
1110	−2
1111	−1

[Figure 3.9 from the textbook]

Naming Ambiguity: 2's Complement

2's complement has two different meanings:

- **representation for signed integer numbers**

- **algorithm for computing the 2's complement (regardless of the representation of the number)**

Naming Ambiguity: 2's Complement

2's complement has two different meanings:

- **representation for signed integer numbers**

  in 2's complement

- **algorithm for computing the 2's complement (regardless of the representation of the number)**

  take the 2's complement

Example of 2’s complement <u>subtraction</u>

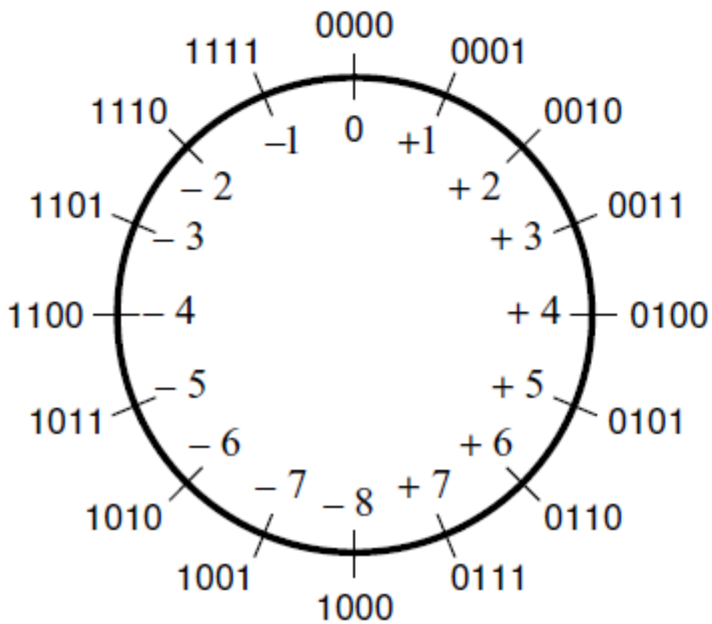
$$\begin{array}{r} (+5) \\ -\ (+2) \\ \hline (+3) \end{array} \qquad \begin{array}{r} 0\,1\,0\,1 \\ -\ 0\,0\,1\,0 \\ \hline \end{array} \quad \Longrightarrow \quad \begin{array}{r} 0\,1\,0\,1 \\ +\ 1\,1\,1\,0 \\ \hline 1\,0\,0\,1\,1 \end{array}$$

ignore

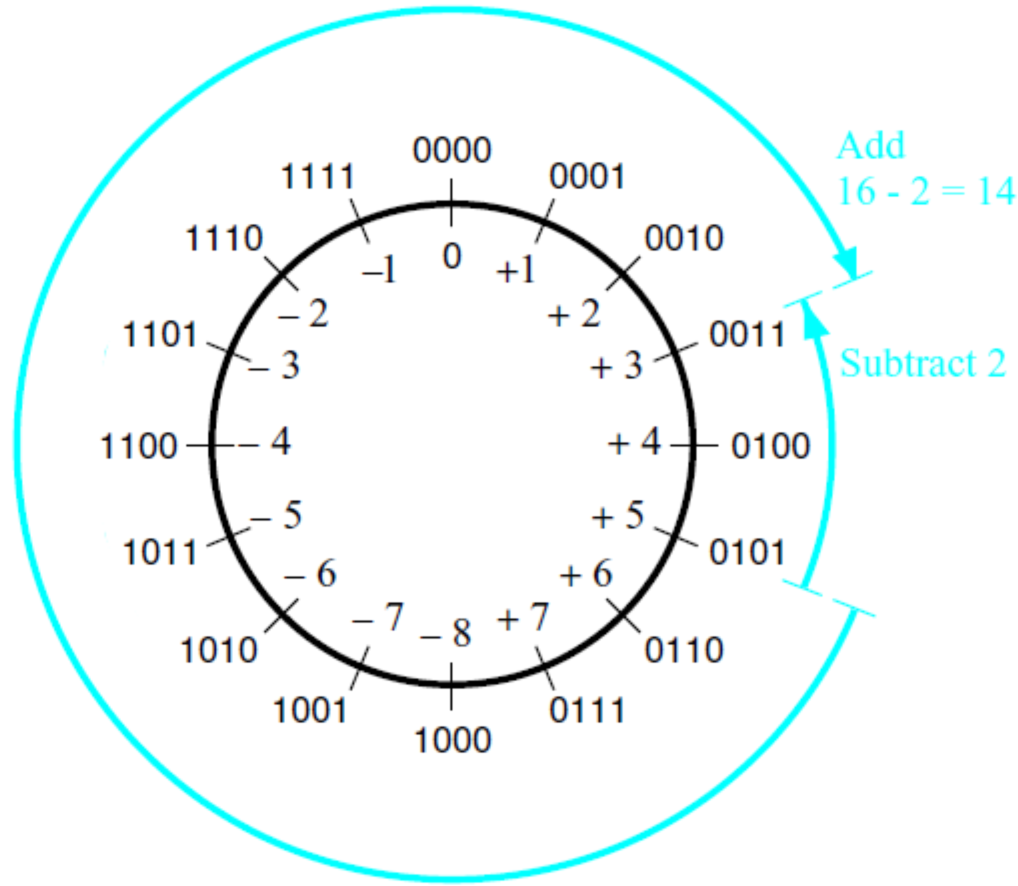
$\Longrightarrow$ means take the 2's complement

[Figure 3.10 from the textbook]

Graphical interpretation of four-bit 2's complement numbers



(a) The number circle



(b) Subtracting 2 by adding its 2's complement

[Figure 3.11 from the textbook]

Example of 2's complement subtraction

$$\begin{array}{r} (-5) \\ -\ (+2) \\ \hline (-7) \end{array} \qquad \begin{array}{r} 1\,0\,1\,1 \\ -\ 0\,0\,1\,0 \\ \hline \end{array} \qquad \Longrightarrow \qquad \begin{array}{r} 1\,0\,1\,1 \\ +\ 1\,1\,1\,0 \\ \hline 1\,1\,0\,0\,1 \end{array}$$

↑ ignore

[Figure 3.10 from the textbook]

Example of 2' s complement subtraction

(+ 5)	0 1 0 1		0 1 0 1
− (−2)	− 1 1 1 0	⟹	+ 0 0 1 0
(+ 7)			0 1 1 1

[Figure 3.10 from the textbook]

Example of 2's complement subtraction

$$
\begin{array}{r}
(-5) \\
-\ (-2) \\
\hline
(-3)
\end{array}
\qquad
\begin{array}{r}
1\,0\,1\,1 \\
-\ 1\,1\,1\,0 \\
\hline
\end{array}
\qquad \Longrightarrow \qquad
\begin{array}{r}
1\,0\,1\,1 \\
+\ 0\,0\,1\,0 \\
\hline
1\,1\,0\,1
\end{array}
$$

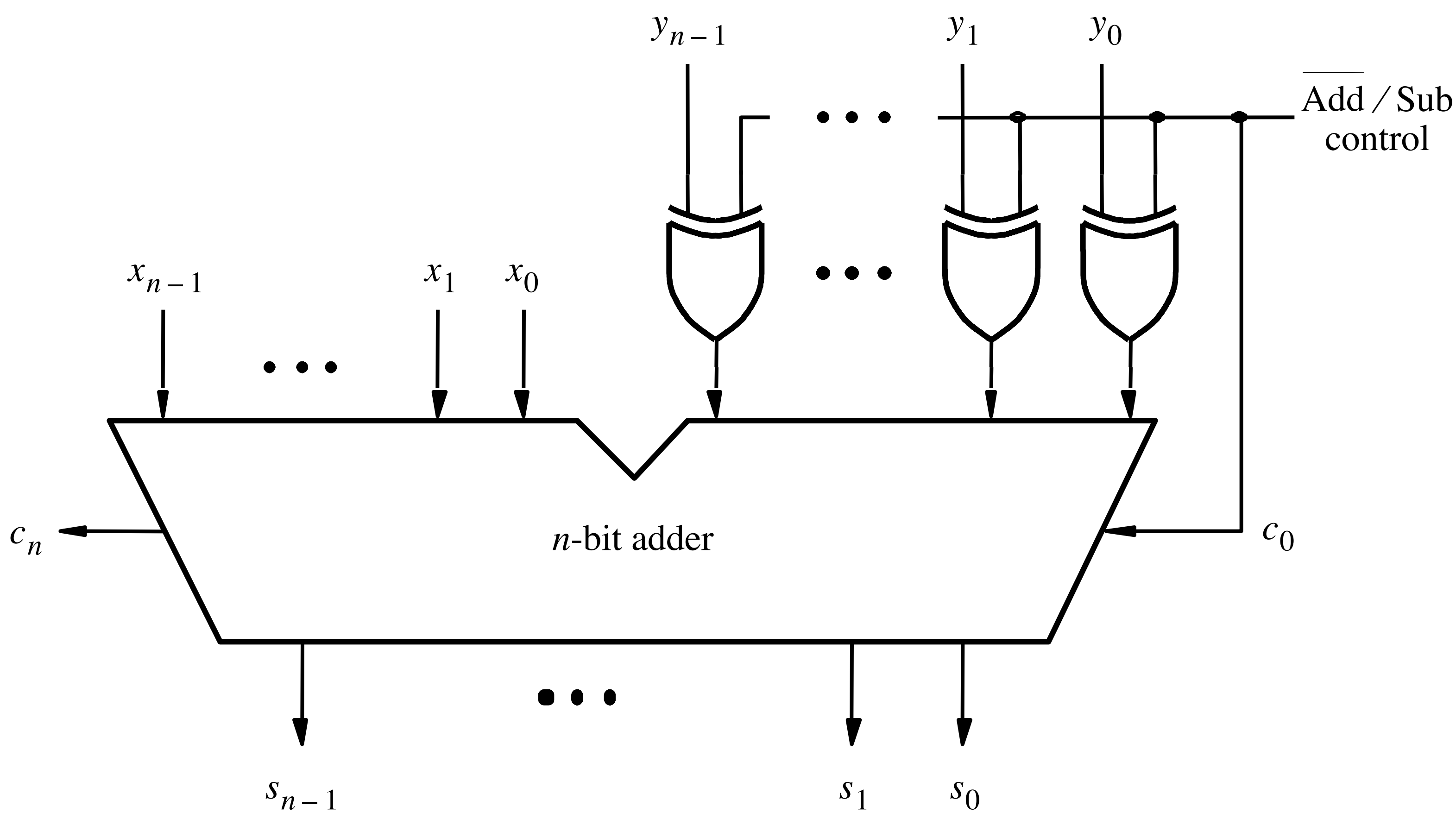
[Figure 3.10 from the textbook]

Take Home Message

- Subtraction can be performed by simply adding the 2's complement of the second number, regardless of the signs of the two numbers.

- Thus, the same adder circuit can be used to perform both addition and subtraction !!!

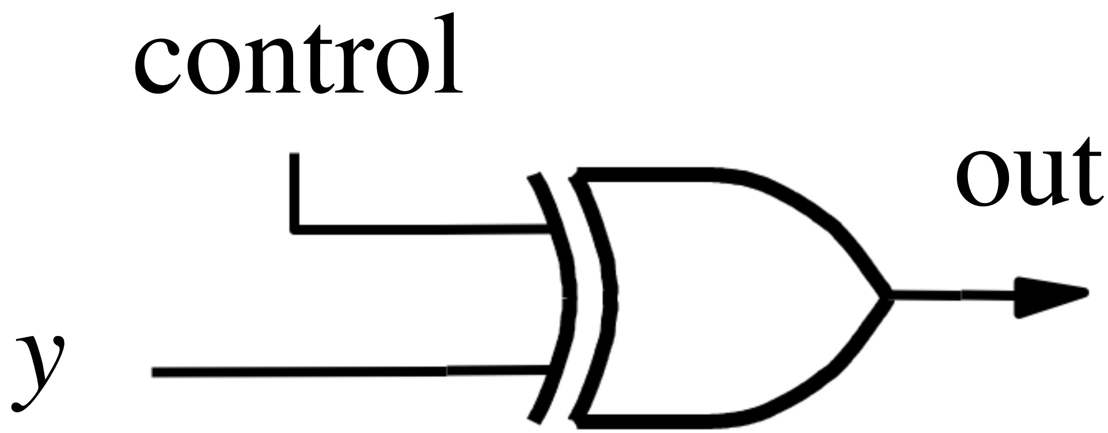
Adder/subtractor unit



[Figure 3.12 from the textbook]

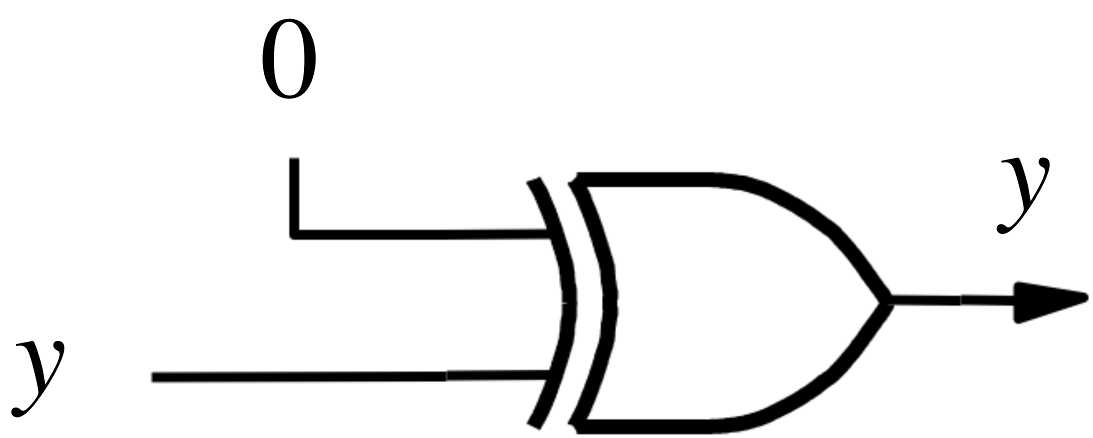
XOR Tricks

control	y	out
0	0	0
0	1	1
1	0	1
1	1	0



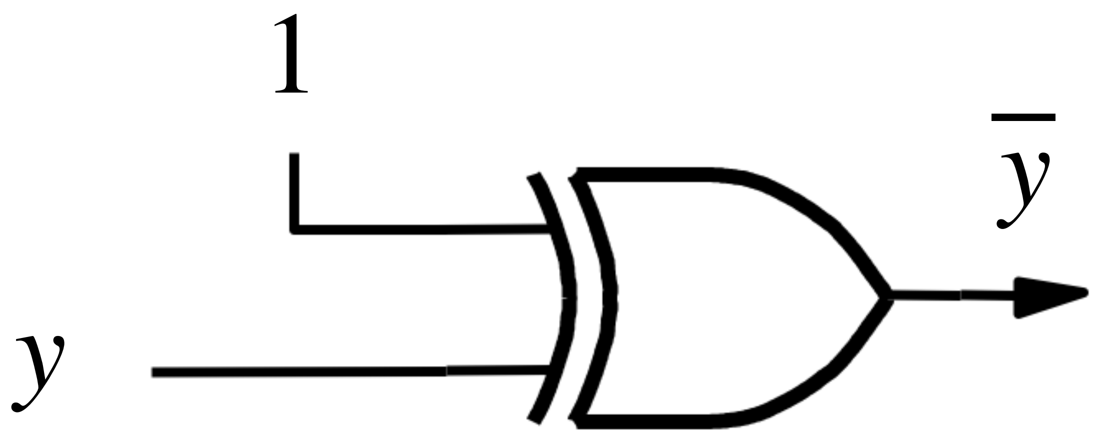
XOR as a repeater

control	*y*	out
0	0	0
0	1	1

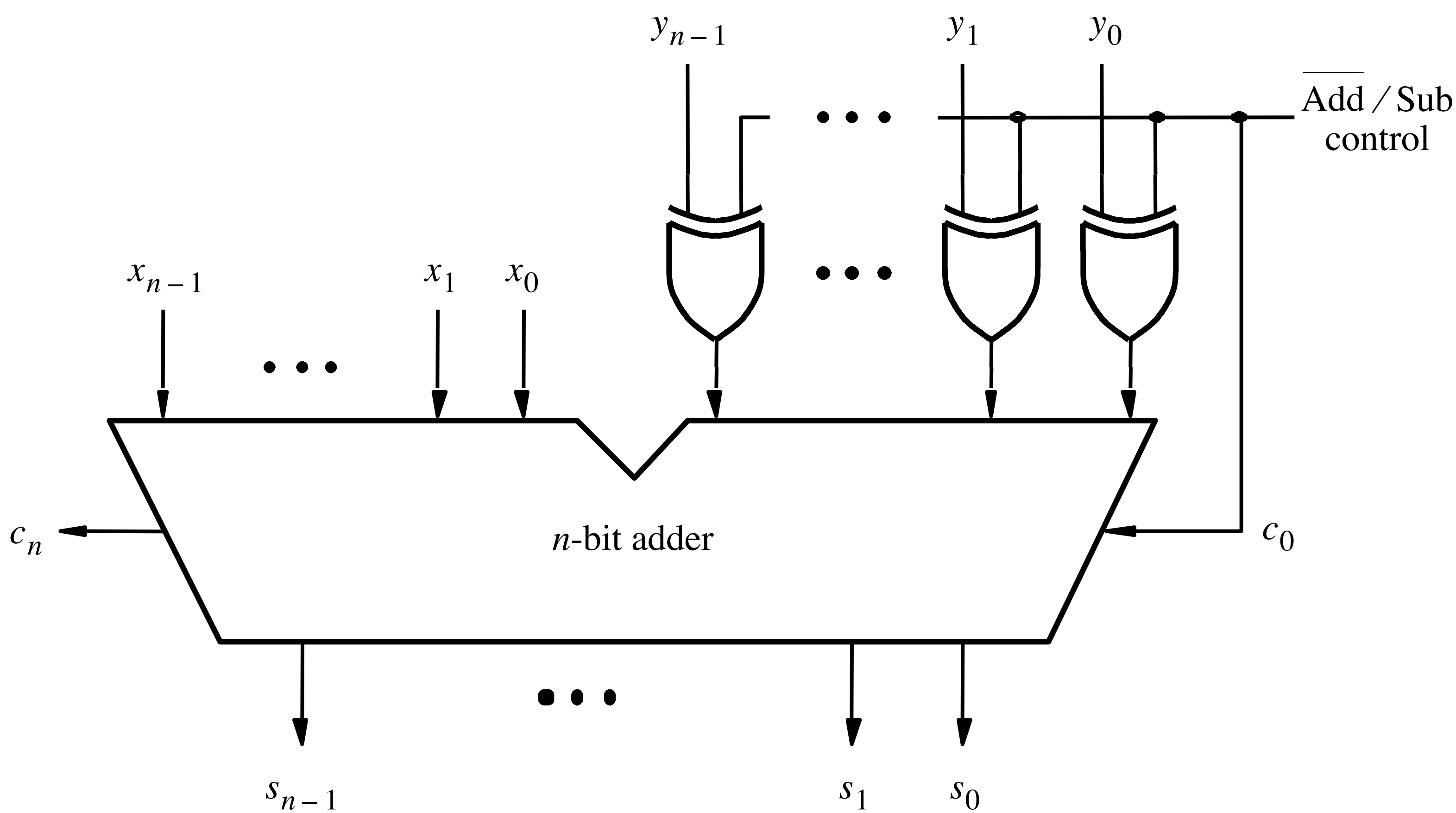


XOR as an inverter

control	y	out
1	0	1
1	1	0

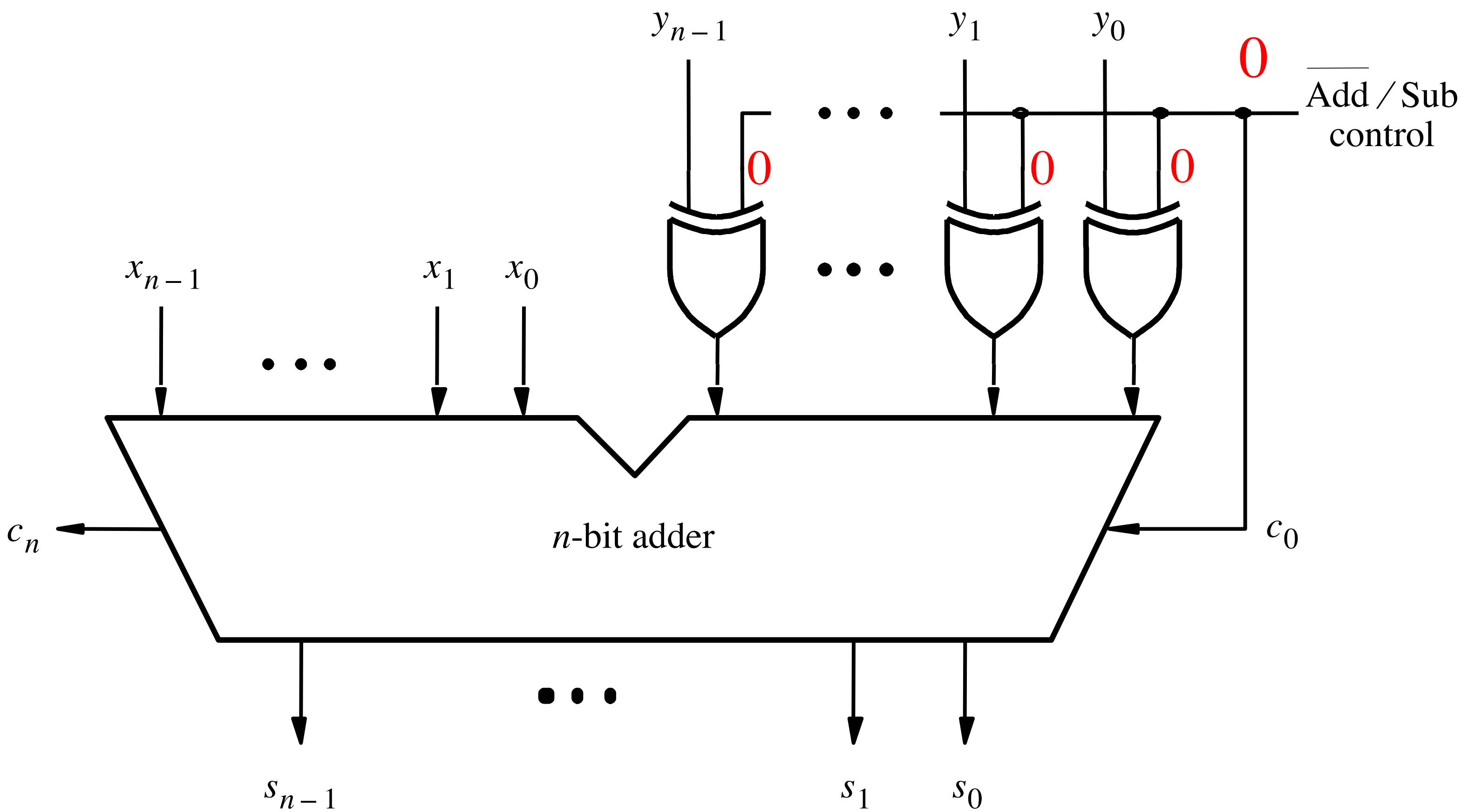


Addition: when control = 0



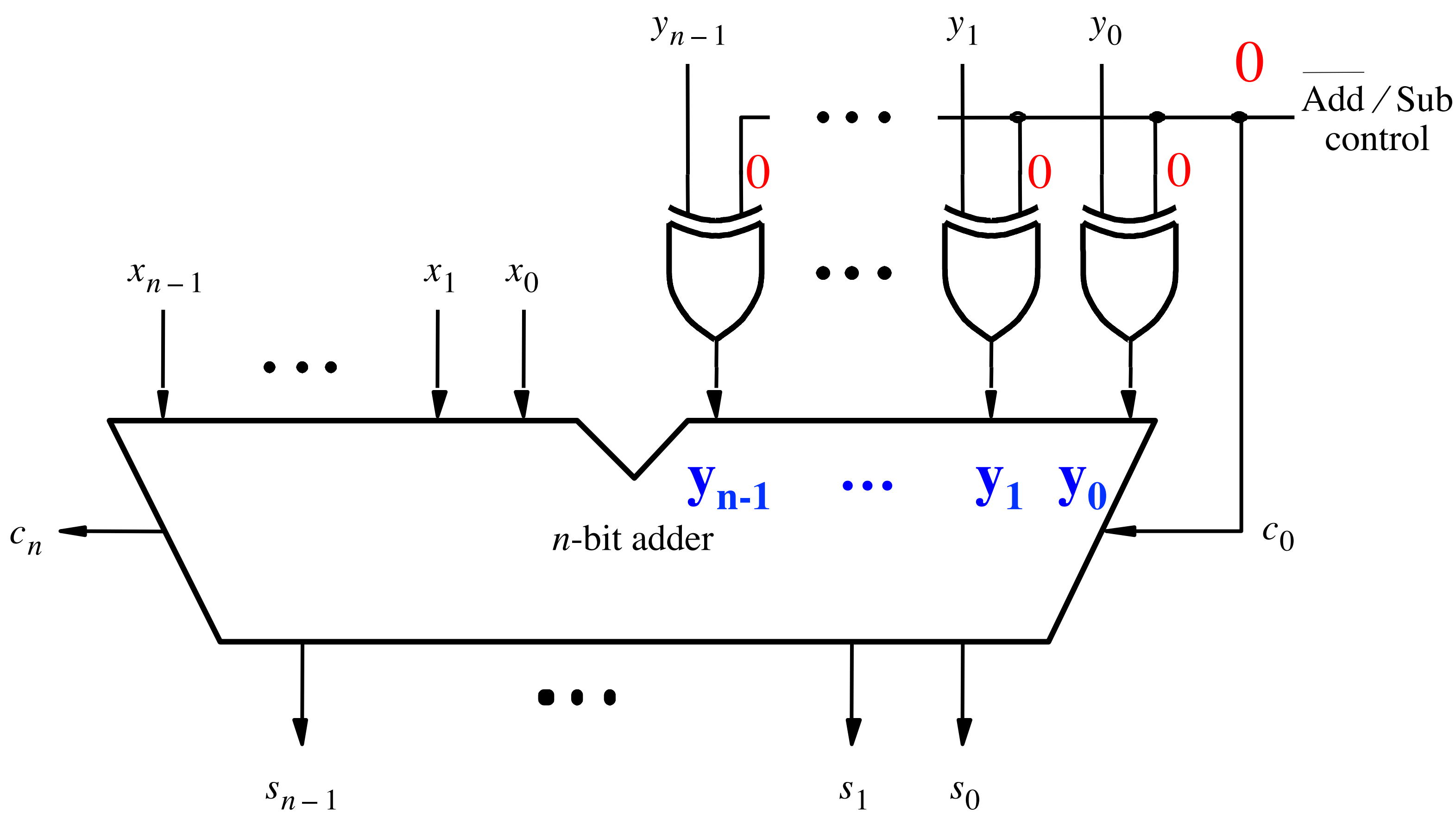
[Figure 3.12 from the textbook]

Addition: when control = 0



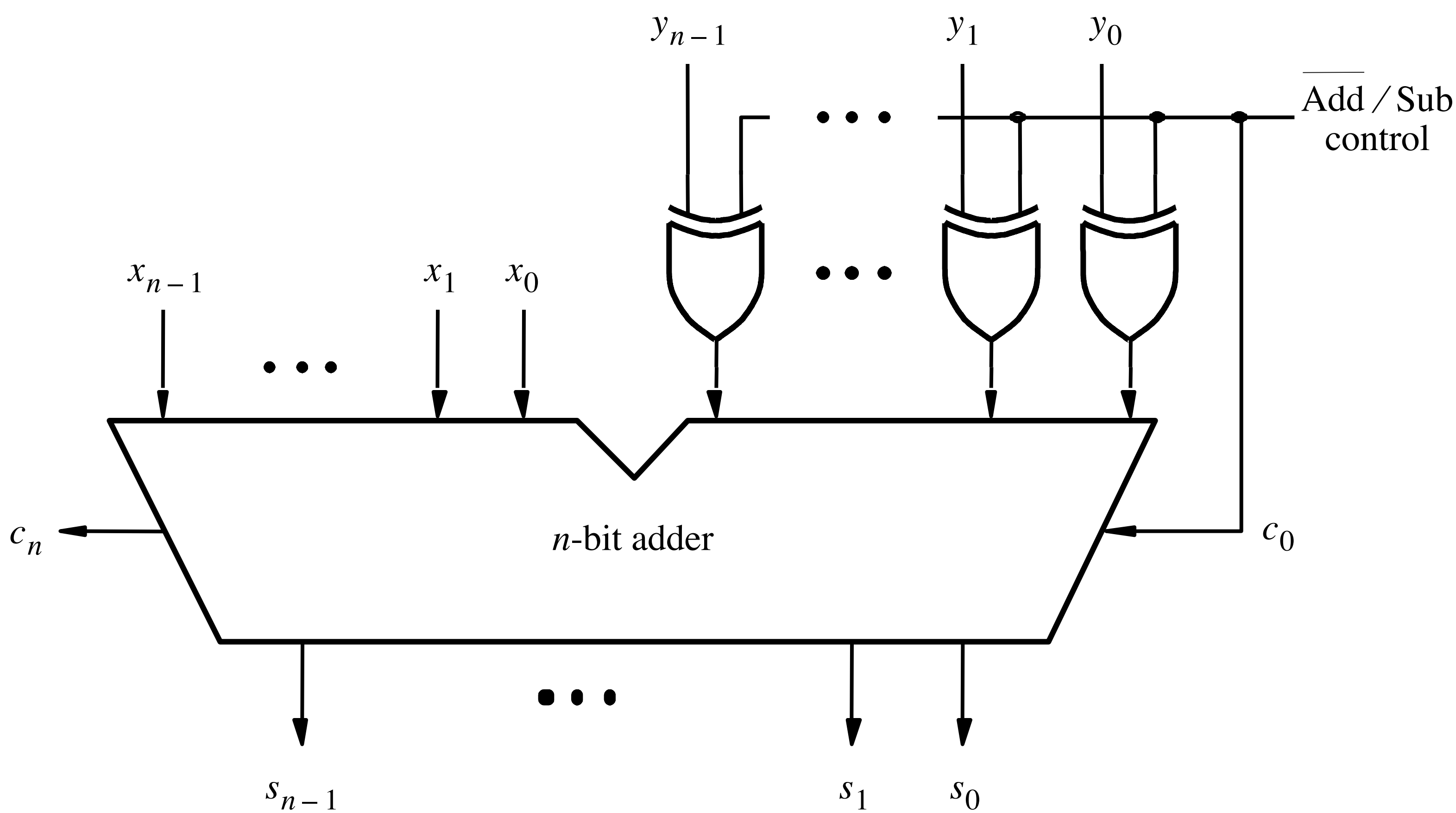
[Figure 3.12 from the textbook]

Addition: when control = 0



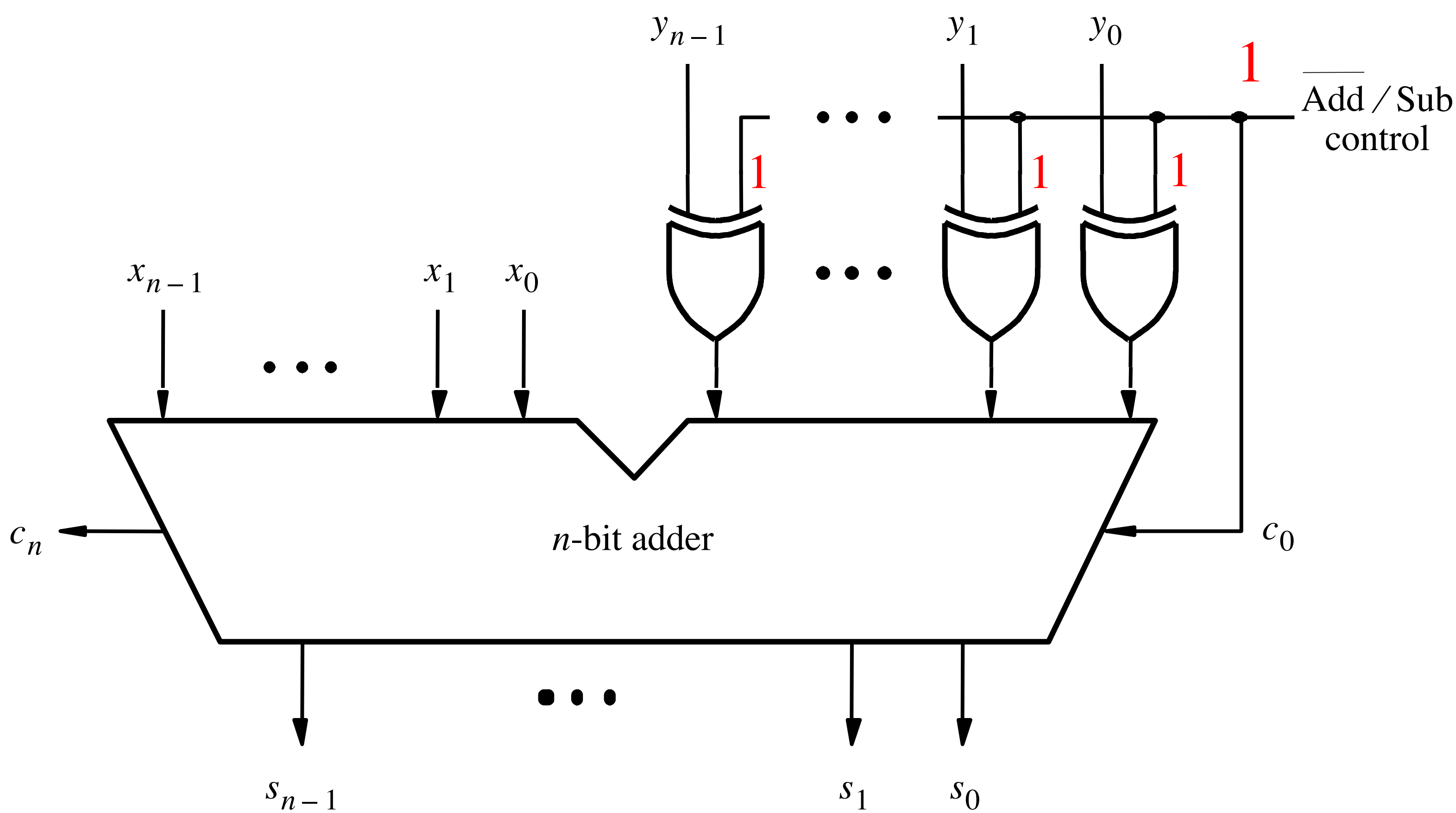
[Figure 3.12 from the textbook]

Subtraction: when control = 1



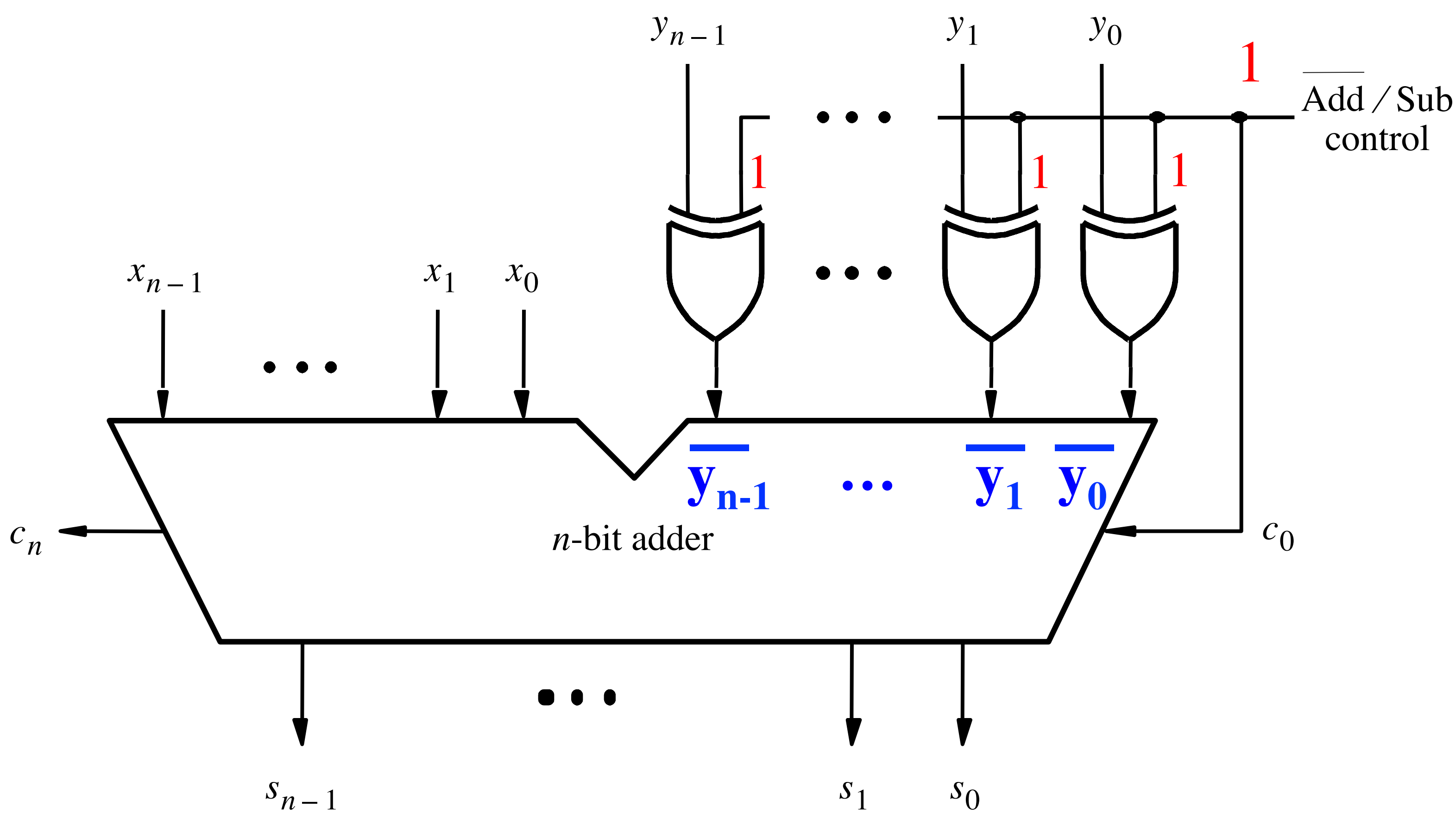
[Figure 3.12 from the textbook]

Subtraction: when control = 1



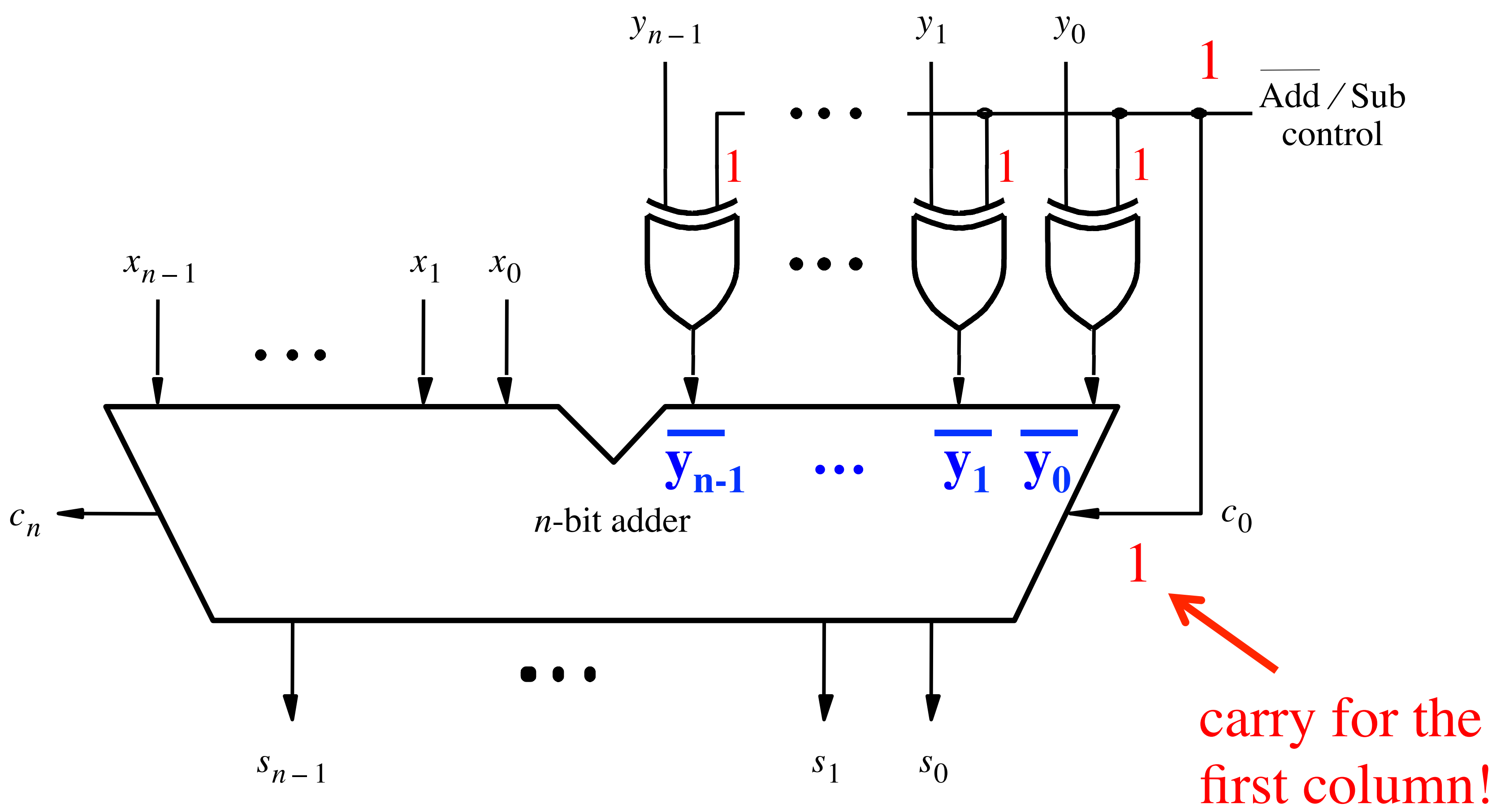
[Figure 3.12 from the textbook]

Subtraction: when control = 1



[Figure 3.12 from the textbook]

Subtraction: when control = 1



[Figure 3.12 from the textbook]

Examples of determination of overflow

(+7)	0 1 1 1	(−7)	1 0 0 1
+ (+2)	+ 0 0 1 0	+ (+2)	+ 0 0 1 0
(+9)	1 0 0 1	(−5)	1 0 1 1

(+7)	0 1 1 1	(−7)	1 0 0 1
+ (−2)	+ 1 1 1 0	+ (−2)	+ 1 1 1 0
(+5)	1 0 1 0 1	(−9)	1 0 1 1 1

[Figure 3.13 from the textbook]

Examples of determination of overflow

Decimal	Binary
	0 1 1 0 0 (carries)
(+7)	0 1 1 1
+ (+2)	+ 0 0 1 0
(+9)	1 0 0 1

Decimal	Binary
	0 0 0 0 0 (carries)
(−7)	1 0 0 1
+ (+2)	+ 0 0 1 0
(−5)	1 0 1 1

Decimal	Binary
	1 1 1 0 0 (carries)
(+7)	0 1 1 1
+ (−2)	+ 1 1 1 0
(+5)	1 0 1 0 1

Decimal	Binary
	1 0 0 0 0 (carries)
(−7)	1 0 0 1
+ (−2)	+ 1 1 1 0
(−9)	1 0 1 1 1

Include the carry bits: $c_4\ c_3\ c_2\ c_1\ c_0$

Examples of determination of overflow

		[0 1] 1 0 0
(+7)		0 1 1 1
+ (+2)	+	0 0 1 0
(+9)		1 0 0 1

		[0 0] 0 0 0
(–7)		1 0 0 1
+ (+2)	+	0 0 1 0
(–5)		1 0 1 1

		[1 1] 1 0 0
(+7)		0 1 1 1
+ (–2)	+	1 1 1 0
(+5)		1 0 1 0 1

		[1 0] 0 0 0
(–7)		1 0 0 1
+ (–2)	+	1 1 1 0
(–9)		1 0 1 1 1

Include the carry bits: [c_4 c_3] c_2 c_1 c_0

Examples of determination of overflow

$c_4 = 0$
$c_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{r} \boxed{0\ 1}\,1\ 0\ 0 \\ 0\ 1\ 1\ 1 \\ +\quad 0\ 0\ 1\ 0 \\ \hline 1\ 0\ 0\ 1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{r} \boxed{0\ 0}\,0\ 0\ 0 \\ 1\ 0\ 0\ 1 \\ +\quad 0\ 0\ 1\ 0 \\ \hline 1\ 0\ 1\ 1 \end{array}$$

$c_4 = 0$
$c_3 = 0$

$c_4 = 1$
$c_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{r} \boxed{1\ 1}\,1\ 0\ 0 \\ 0\ 1\ 1\ 1 \\ +\quad 1\ 1\ 1\ 0 \\ \hline 1\ 0\ 1\ 0\ 1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{r} \boxed{1\ 0}\,0\ 0\ 0 \\ 1\ 0\ 0\ 1 \\ +\quad 1\ 1\ 1\ 0 \\ \hline 1\ 0\ 1\ 1\ 1 \end{array}$$

$c_4 = 1$
$c_3 = 0$

Include the carry bits: $\boxed{c_4\ c_3}\,c_2\ c_1\ c_0$

Examples of determination of overflow

$c_4 = 0$
$c_3 = 1$

Carries: 0 1 1 0 0

$$\begin{array}{rr} & (+7) \\ + & (+2) \\ \hline & (+9) \end{array} \qquad \begin{array}{rr} & 0111 \\ + & 0010 \\ \hline & 1001 \end{array}$$

Carries: 0 0 0 0 0

$$\begin{array}{rr} & (-7) \\ + & (+2) \\ \hline & (-5) \end{array} \qquad \begin{array}{rr} & 1001 \\ + & 0010 \\ \hline & 1011 \end{array}$$

$c_4 = 0$
$c_3 = 0$

$c_4 = 1$
$c_3 = 1$

Carries: 1 1 1 0 0

$$\begin{array}{rr} & (+7) \\ + & (-2) \\ \hline & (+5) \end{array} \qquad \begin{array}{rr} & 0111 \\ + & 1110 \\ \hline & 10101 \end{array}$$

Carries: 1 0 0 0 0

$$\begin{array}{rr} & (-7) \\ + & (-2) \\ \hline & (-9) \end{array} \qquad \begin{array}{rr} & 1001 \\ + & 1110 \\ \hline & 10111 \end{array}$$

$c_4 = 1$
$c_3 = 0$

Overflow occurs only in these two cases.

Examples of determination of overflow

$c_4 = 0$
$c_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{r} 0\,1\,1\,0\,0 \\ 0\,1\,1\,1 \\ +\quad 0\,0\,1\,0 \\ \hline 1\,0\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{r} 0\,0\,0\,0\,0 \\ 1\,0\,0\,1 \\ +\quad 0\,0\,1\,0 \\ \hline 1\,0\,1\,1 \end{array}$$

$c_4 = 0$
$c_3 = 0$

$c_4 = 1$
$c_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{r} 1\,1\,1\,0\,0 \\ 0\,1\,1\,1 \\ +\quad 1\,1\,1\,0 \\ \hline 1\,0\,1\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{r} 1\,0\,0\,0\,0 \\ 1\,0\,0\,1 \\ +\quad 1\,1\,1\,0 \\ \hline 1\,0\,1\,1\,1 \end{array}$$

$c_4 = 1$
$c_3 = 0$

$$\text{Overflow} = c_3\overline{c_4} + \overline{c_3}c_4$$

Examples of determination of overflow

$c_4 = 0$
$c_3 = 1$

$$\begin{array}{r} (+7) \\ +\,(+2) \\ \hline (+9) \end{array} \qquad \begin{array}{r} \boxed{0\,1}\,1\,0\,0 \\ 0\,1\,1\,1 \\ +\quad 0\,0\,1\,0 \\ \hline 1\,0\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\,(+2) \\ \hline (-5) \end{array} \qquad \begin{array}{r} \boxed{0\,0}\,0\,0\,0 \\ 1\,0\,0\,1 \\ +\quad 0\,0\,1\,0 \\ \hline 1\,0\,1\,1 \end{array}$$

$c_4 = 0$
$c_3 = 0$

$c_4 = 1$
$c_3 = 1$

$$\begin{array}{r} (+7) \\ +\,(-2) \\ \hline (+5) \end{array} \qquad \begin{array}{r} \boxed{1\,1}\,1\,0\,0 \\ 0\,1\,1\,1 \\ +\quad 1\,1\,1\,0 \\ \hline 1\;0\,1\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\,(-2) \\ \hline (-9) \end{array} \qquad \begin{array}{r} \boxed{1\,0}\,0\,0\,0 \\ 1\,0\,0\,1 \\ +\quad 1\,1\,1\,0 \\ \hline 1\;0\,1\,1\,1 \end{array}$$

$c_4 = 1$
$c_3 = 0$

$$\text{Overflow} = \underbrace{c_3\overline{c_4} + \overline{c_3}c_4}_{\text{XOR}}$$

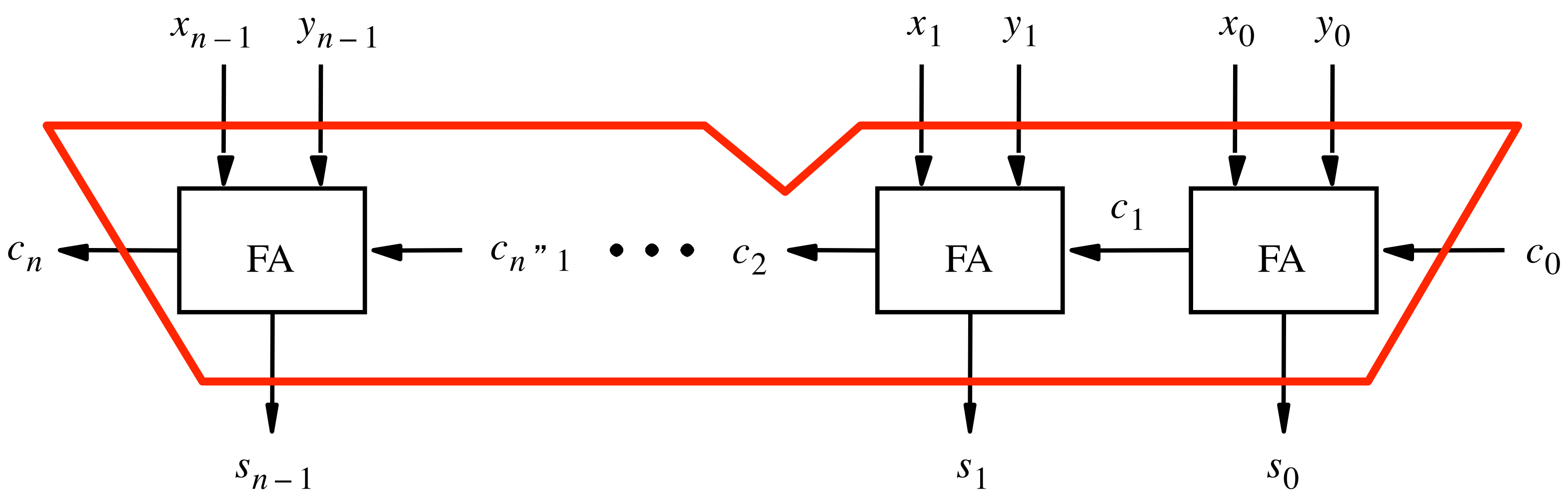
Calculating overflow for 4-bit numbers with only three significant bits

$$\begin{aligned}\text{Overflow} &= c_3\overline{c}_4 + \overline{c}_3 c_4 \\ &= c_3 \oplus c_4\end{aligned}$$

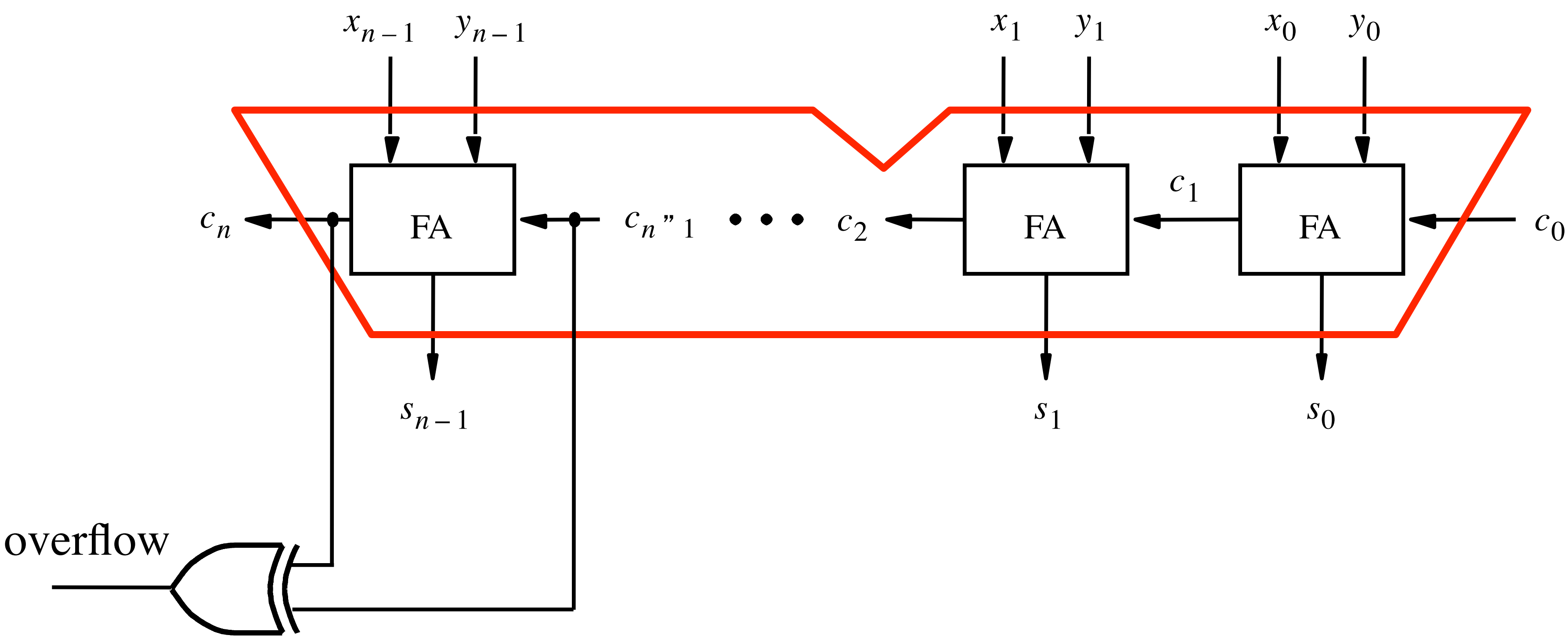
Calculating overflow for n-bit numbers with only n-1 significant bits

$$\text{Overflow} = c_{n-1} \oplus c_n$$

Detecting Overflow



Detecting Overflow (with one extra XOR)



Another way to look at the overflow issue

$$
\begin{array}{rl}
 & X = x_3\ x_2\ x_1\ x_0 \\
+ & Y = y_3\ y_2\ y_1\ y_0 \\
\hline
 & S = s_3\ s_2\ s_1\ s_0
\end{array}
$$

Another way to look at the overflow issue

$$
\begin{array}{rrrrrr}
 & X= & x_3 & x_2 & x_1 & x_0 \\
+ & Y= & y_3 & y_2 & y_1 & y_0 \\
\hline
 & S= & s_3 & s_2 & s_1 & s_0
\end{array}
$$

If both numbers that we are adding have the same sign but the sum does not, then we have an overflow.

Examples of determination of overflow

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{r} 0\,1\,1\,1 \\ +\quad 0\,0\,1\,0 \\ \hline 1\,0\,0\,1 \end{array} \qquad\qquad \begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{r} 1\,0\,0\,1 \\ +\quad 0\,0\,1\,0 \\ \hline 1\,0\,1\,1 \end{array}$$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{r} 0\,1\,1\,1 \\ +\quad 1\,1\,1\,0 \\ \hline 1\,0\,1\,0\,1 \end{array} \qquad\qquad \begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{r} 1\,0\,0\,1 \\ +\quad 1\,1\,1\,0 \\ \hline 1\,0\,1\,1\,1 \end{array}$$

Examples of determination of overflow

$$
\begin{array}{r}
(+7) \\
+\ (+2) \\
\hline
(+9)
\end{array}
\qquad
\begin{array}{rr}
 & 0\,1\,1\,1 \\
+ & 0\,0\,1\,0 \\
\hline
 & 1\,0\,0\,1
\end{array}
$$

$$
\begin{array}{r}
(-7) \\
+\ (+2) \\
\hline
(-5)
\end{array}
\qquad
\begin{array}{rr}
 & 1\,0\,0\,1 \\
+ & 0\,0\,1\,0 \\
\hline
 & 1\,0\,1\,1
\end{array}
$$

$$
\begin{array}{r}
(+7) \\
+\ (-2) \\
\hline
(+5)
\end{array}
\qquad
\begin{array}{rr}
 & 0\,1\,1\,1 \\
+ & 1\,1\,1\,0 \\
\hline
 & 1\,0\,1\,0\,1
\end{array}
$$

$$
\begin{array}{r}
(-7) \\
+\ (-2) \\
\hline
(-9)
\end{array}
\qquad
\begin{array}{rr}
 & 1\,0\,0\,1 \\
+ & 1\,1\,1\,0 \\
\hline
 & 1\,0\,1\,1\,1
\end{array}
$$

Examples of determination of overflow

$x_3 = 0$
$y_3 = 0$
$s_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{rr} & 0\,1\,1\,1 \\ + & 0\,0\,1\,0 \\ \hline & 1\,0\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{rr} & 1\,0\,0\,1 \\ + & 0\,0\,1\,0 \\ \hline & 1\,0\,1\,1 \end{array}$$

$x_3 = 1$
$y_3 = 0$
$s_3 = 1$

$x_3 = 0$
$y_3 = 1$
$s_3 = 0$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{rr} & 0\,1\,1\,1 \\ + & 1\,1\,1\,0 \\ \hline & 1\,0\,1\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{rr} & 1\,0\,0\,1 \\ + & 1\,1\,1\,0 \\ \hline & 1\,0\,1\,1\,1 \end{array}$$

$x_3 = 1$
$y_3 = 1$
$s_3 = 0$

Examples of determination of overflow

$$x_3 = 0 \quad y_3 = 0 \quad s_3 = 1$$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{r} 0\,1\,1\,1 \\ +\ 0\,0\,1\,0 \\ \hline 1\,0\,0\,1 \end{array}$$

$$x_3 = 1 \quad y_3 = 0 \quad s_3 = 1$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{r} 1\,0\,0\,1 \\ +\ 0\,0\,1\,0 \\ \hline 1\,0\,1\,1 \end{array}$$

$$x_3 = 0 \quad y_3 = 1 \quad s_3 = 0$$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{r} 0\,1\,1\,1 \\ +\ 1\,1\,1\,0 \\ \hline 1\,0\,1\,0\,1 \end{array}$$

$$x_3 = 1 \quad y_3 = 1 \quad s_3 = 0$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{r} 1\,0\,0\,1 \\ +\ 1\,1\,1\,0 \\ \hline 1\,0\,1\,1\,1 \end{array}$$

In 2's complement, both +9 and -9 are not representable with 4 bits.

Examples of determination of overflow

$$x_3 = 0 \quad y_3 = 0 \quad s_3 = 1$$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{r} 0\,1\,1\,1 \\ +\ 0\,0\,1\,0 \\ \hline 1\,0\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{r} 1\,0\,0\,1 \\ +\ 0\,0\,1\,0 \\ \hline 1\,0\,1\,1 \end{array}$$

$$x_3 = 1 \quad y_3 = 0 \quad s_3 = 1$$

$$x_3 = 0 \quad y_3 = 1 \quad s_3 = 0$$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{r} 0\,1\,1\,1 \\ +\ 1\,1\,1\,0 \\ \hline 1\,0\,1\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{r} 1\,0\,0\,1 \\ +\ 1\,1\,1\,0 \\ \hline 1\,0\,1\,1\,1 \end{array}$$

$$x_3 = 1 \quad y_3 = 1 \quad s_3 = 0$$

Overflow occurs only in these two cases.

Examples of determination of overflow

$x_3 = 0$
$y_3 = 0$
$s_3 = 1$

$$\begin{array}{r} (+7) \\ +\ (+2) \\ \hline (+9) \end{array} \qquad \begin{array}{r} 0\,1\,1\,1 \\ +\quad 0\,0\,1\,0 \\ \hline 1\,0\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (+2) \\ \hline (-5) \end{array} \qquad \begin{array}{r} 1\,0\,0\,1 \\ +\quad 0\,0\,1\,0 \\ \hline 1\,0\,1\,1 \end{array}$$

$x_3 = 1$
$y_3 = 0$
$s_3 = 1$

$x_3 = 0$
$y_3 = 1$
$s_3 = 0$

$$\begin{array}{r} (+7) \\ +\ (-2) \\ \hline (+5) \end{array} \qquad \begin{array}{r} 0\,1\,1\,1 \\ +\quad 1\,1\,1\,0 \\ \hline 1\,0\,1\,0\,1 \end{array}$$

$$\begin{array}{r} (-7) \\ +\ (-2) \\ \hline (-9) \end{array} \qquad \begin{array}{r} 1\,0\,0\,1 \\ +\quad 1\,1\,1\,0 \\ \hline 1\,0\,1\,1\,1 \end{array}$$

$x_3 = 1$
$y_3 = 1$
$s_3 = 0$

$$\text{Overflow} = \overline{x}_3\,\overline{y}_3\,s_3 + x_3\,y_3\,\overline{s}_3$$

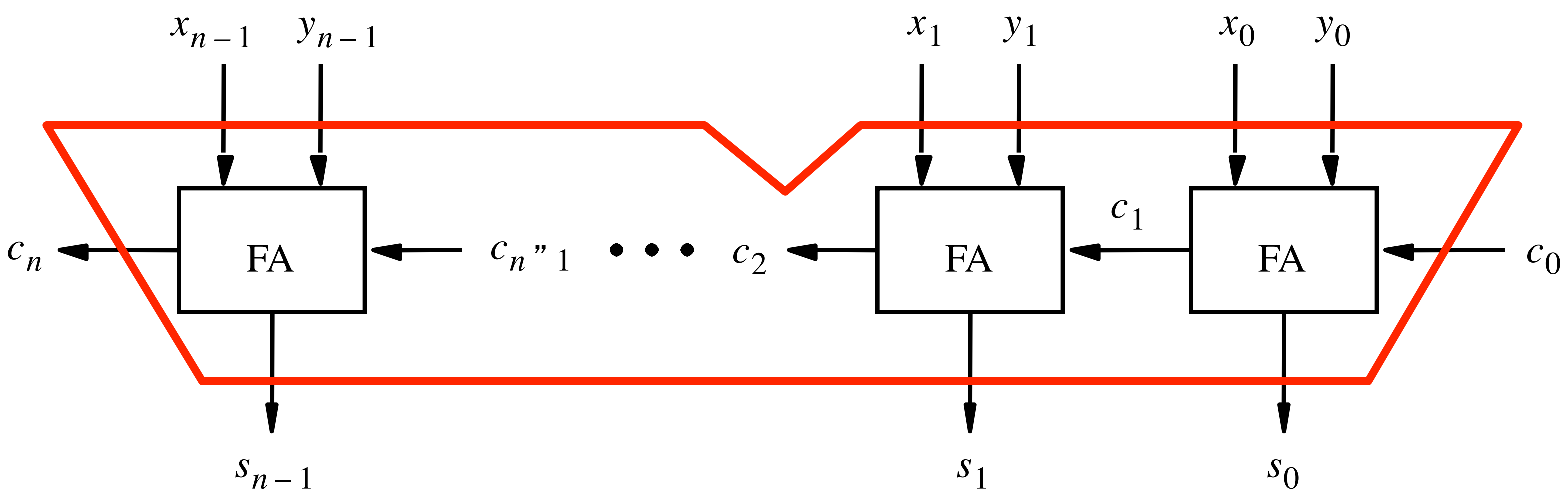
Another way to look at the overflow issue

$$
\begin{array}{rcccc}
 & X= & x_3 & x_2 & x_1 & x_0 \\
+ & Y= & y_3 & y_2 & y_1 & y_0 \\
\hline
 & S= & s_3 & s_2 & s_1 & s_0
\end{array}
$$

If both numbers that we are adding have the same sign but the sum does not, then we have an overflow.

$$\text{Overflow} = \overline{x}_3\,\overline{y}_3\,s_3 + x_3\,y_3\,\overline{s}_3$$

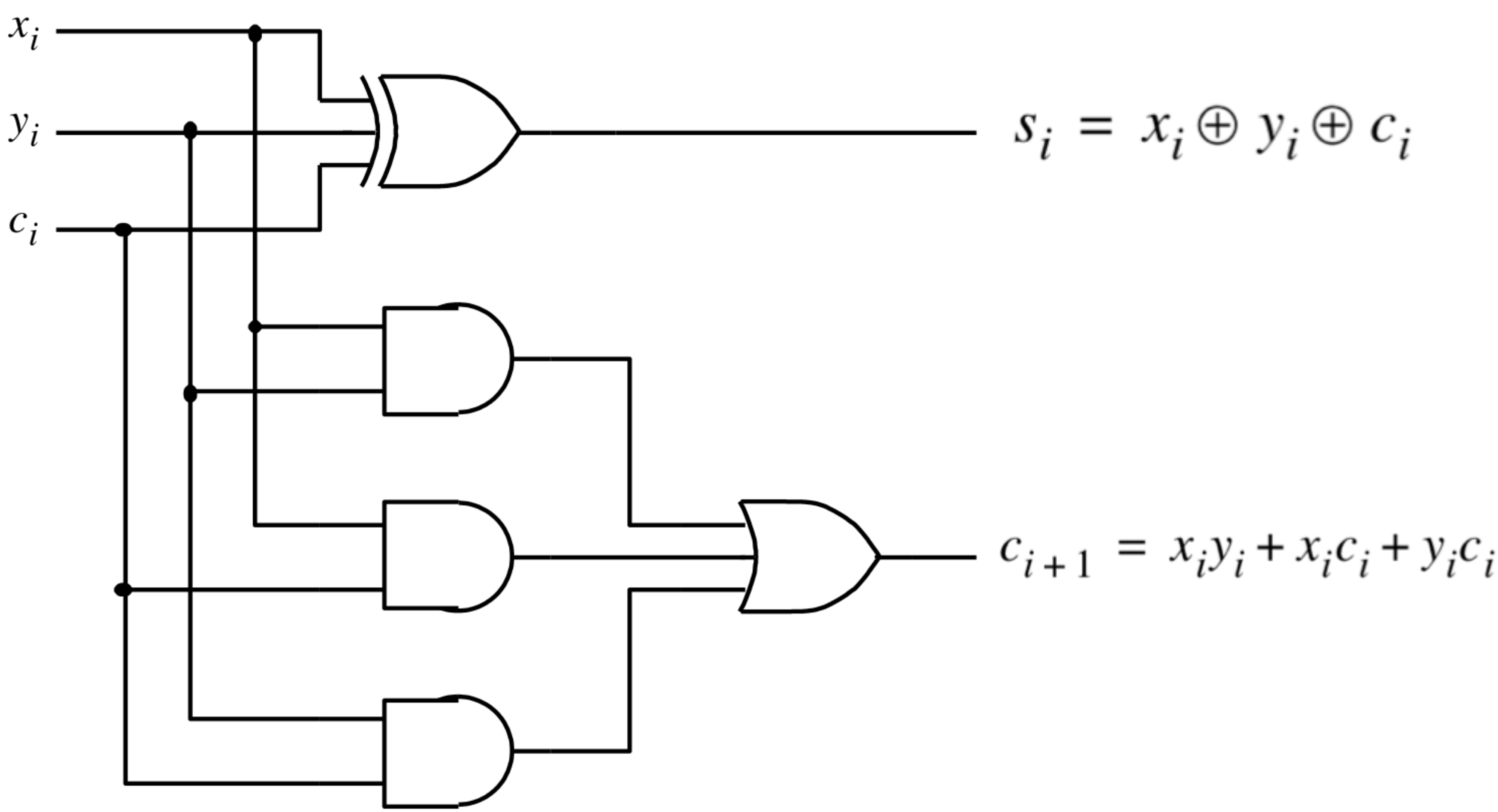
How long does it take to compute all sum bits and all carry bits?



Can we perform addition even faster?

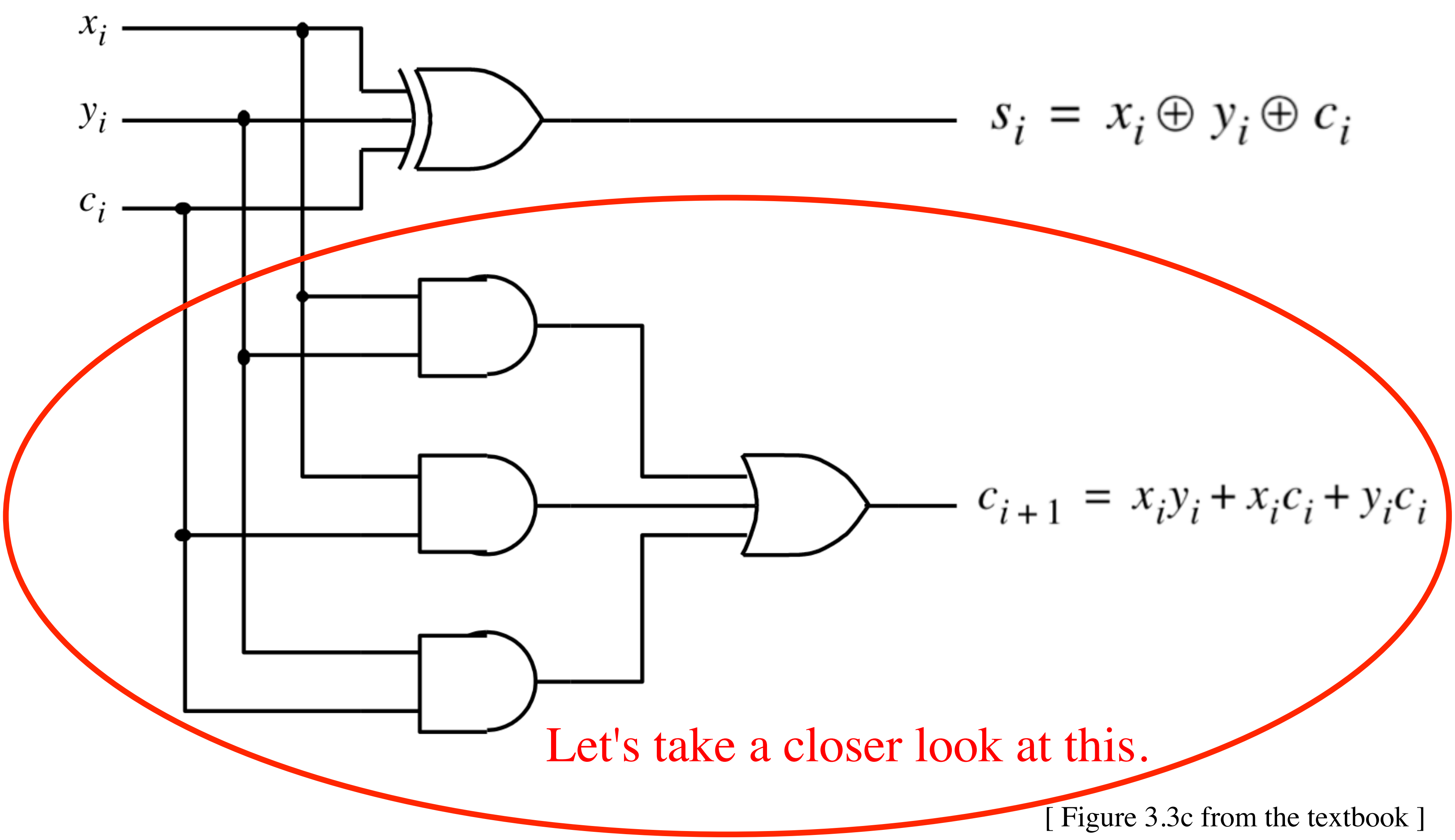
The goal is to evaluate very fast if the carry from the previous stage will be equal to 0 or 1.

The Full-Adder Circuit



[Figure 3.3c from the textbook]

The Full-Adder Circuit



[Figure 3.3c from the textbook]

Decomposing the Carry Expression

$$c_{i+1} = x_i\, y_i + x_i\, c_i + y_i\, c_i$$

Decomposing the Carry Expression

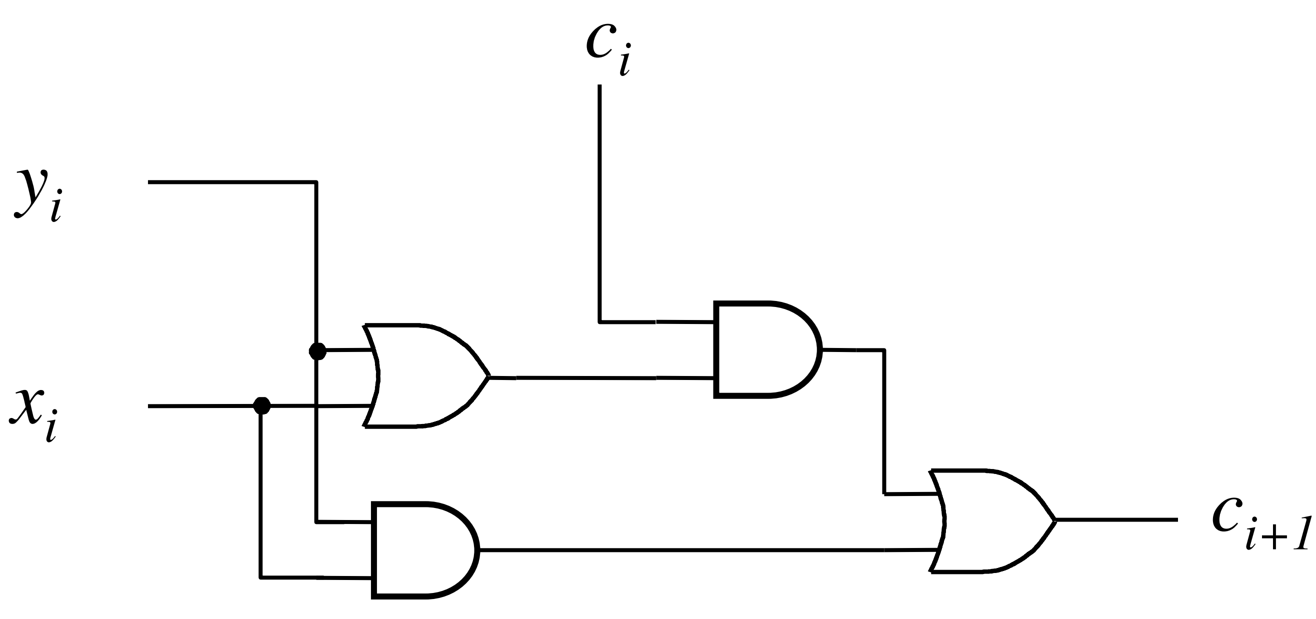
$$c_{i+1} = x_i y_i + x_i c_i + y_i c_i$$

$$c_{i+1} = x_i y_i + (x_i + y_i) c_i$$

Decomposing the Carry Expression

$$c_{i+1} = x_i y_i + x_i c_i + y_i c_i$$

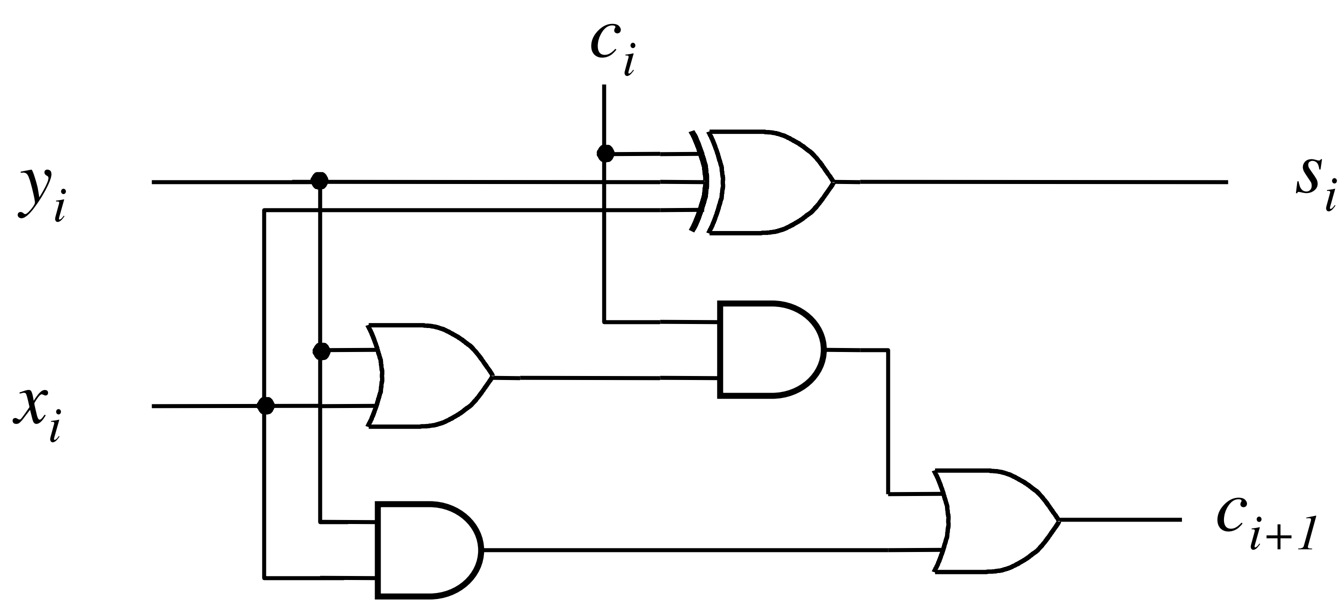
$$c_{i+1} = x_i y_i + (x_i + y_i) c_i$$



Another Way to Draw the Full-Adder Circuit

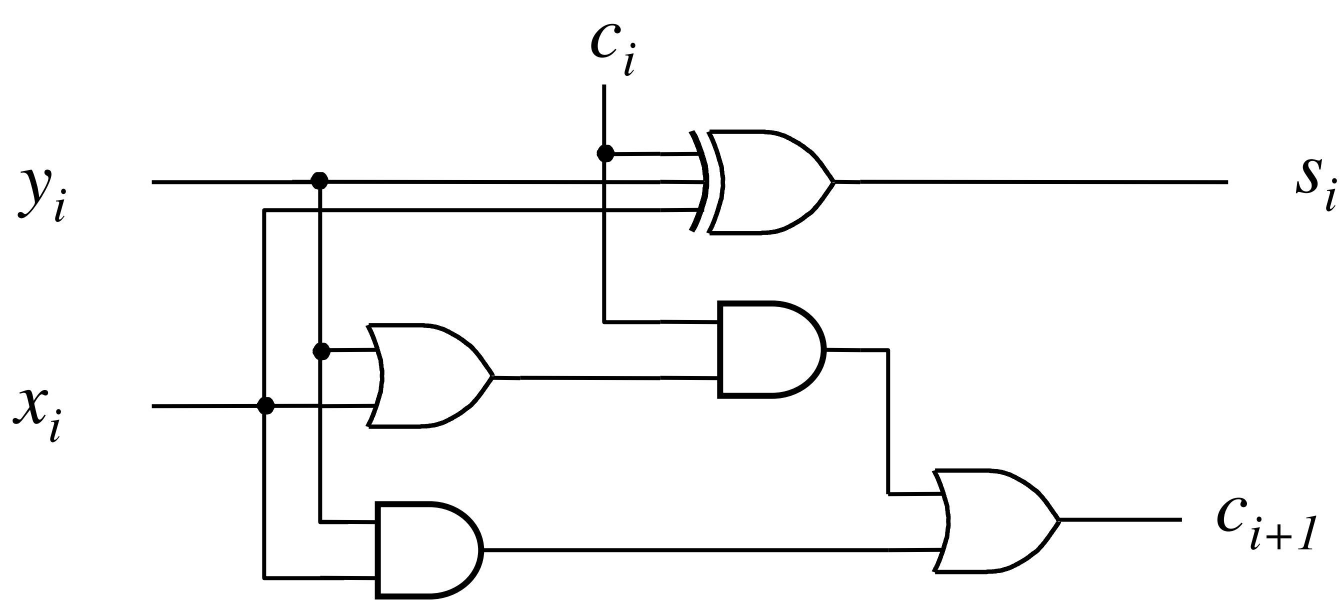
$$c_{i+1} = x_i y_i + x_i c_i + y_i c_i$$

$$c_{i+1} = x_i y_i + (x_i + y_i) c_i$$



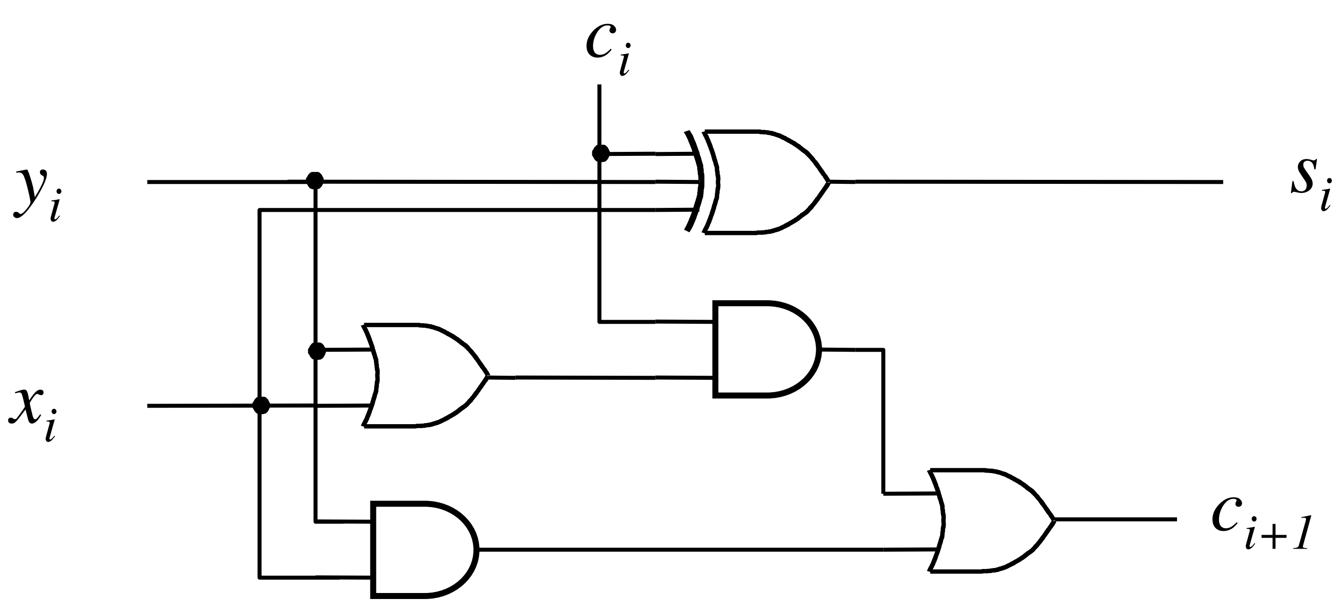
Another Way to Draw the Full-Adder Circuit

$$c_{i+1} = x_i y_i + (x_i + y_i)c_i$$



Another Way to Draw the Full-Adder Circuit

$$c_{i+1} = \underbrace{x_i y_i}_{g_i} + \underbrace{(x_i + y_i)}_{p_i} c_i$$

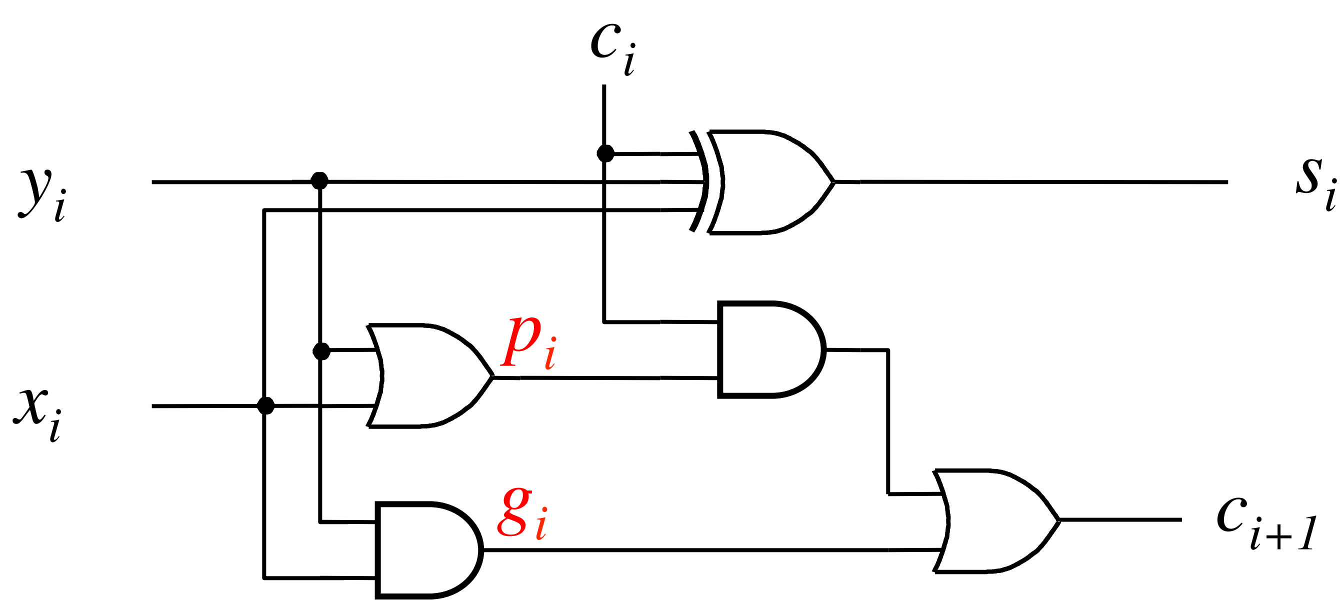


Another Way to Draw the Full-Adder Circuit

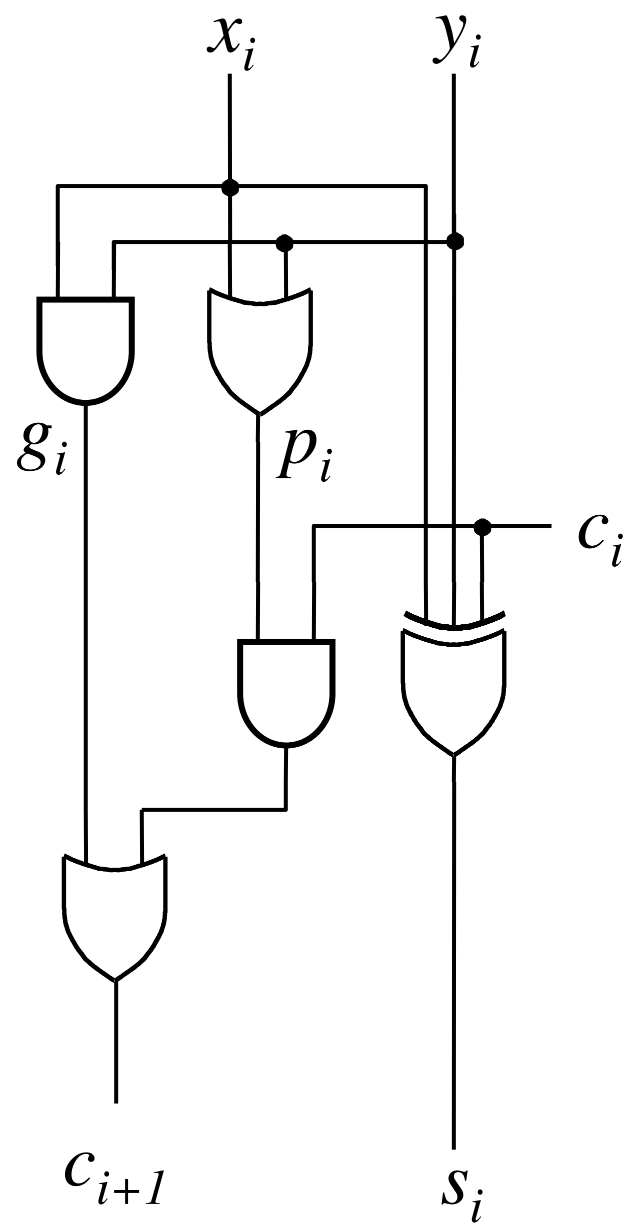
g - generate

p - propagate

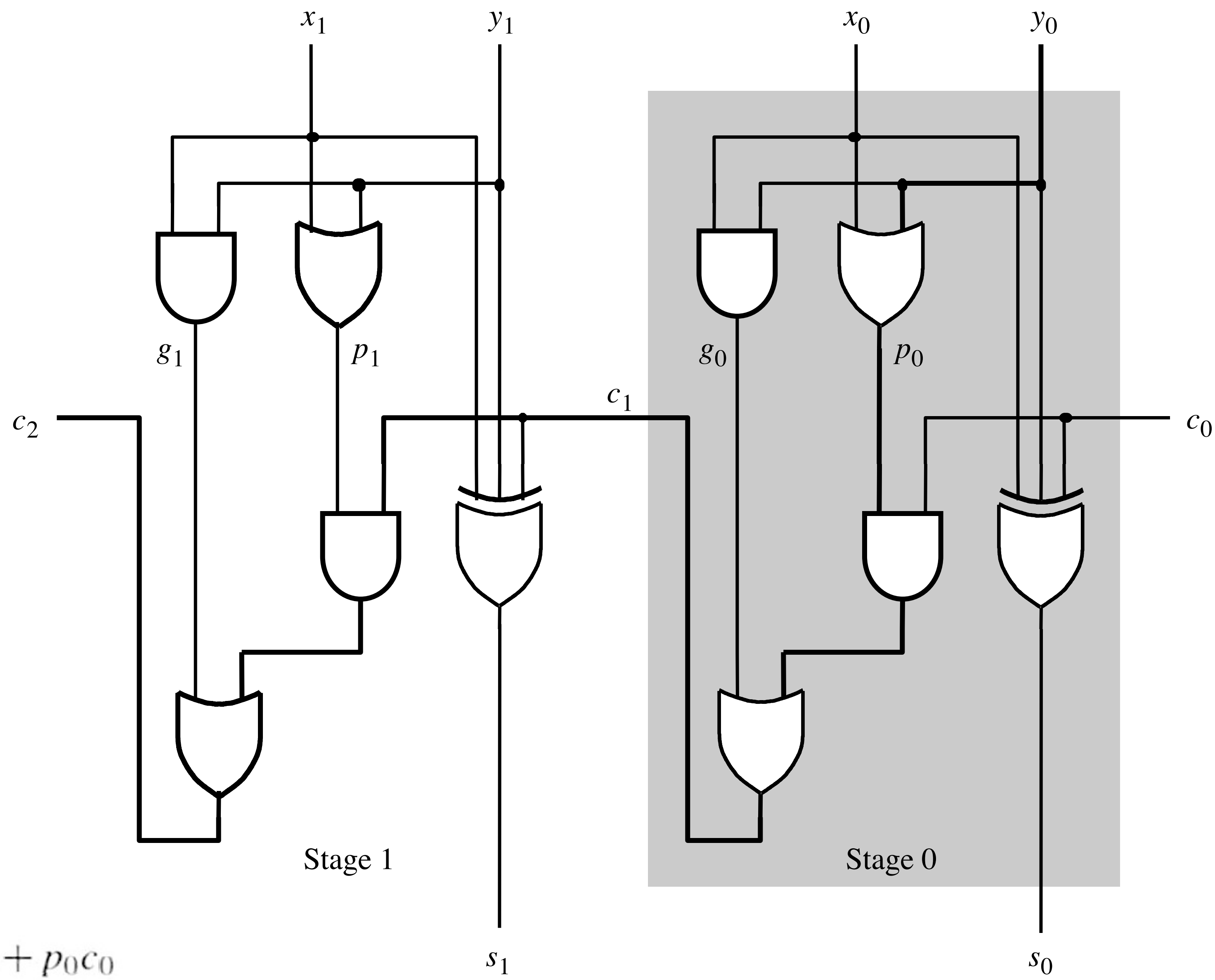
$$c_{i+1} = \underbrace{x_i y_i}_{g_i} + \underbrace{(x_i + y_i)}_{p_i} c_i$$



Yet Another Way to Draw It (Just Rotate It)



Now we can Build a Ripple-Carry Adder

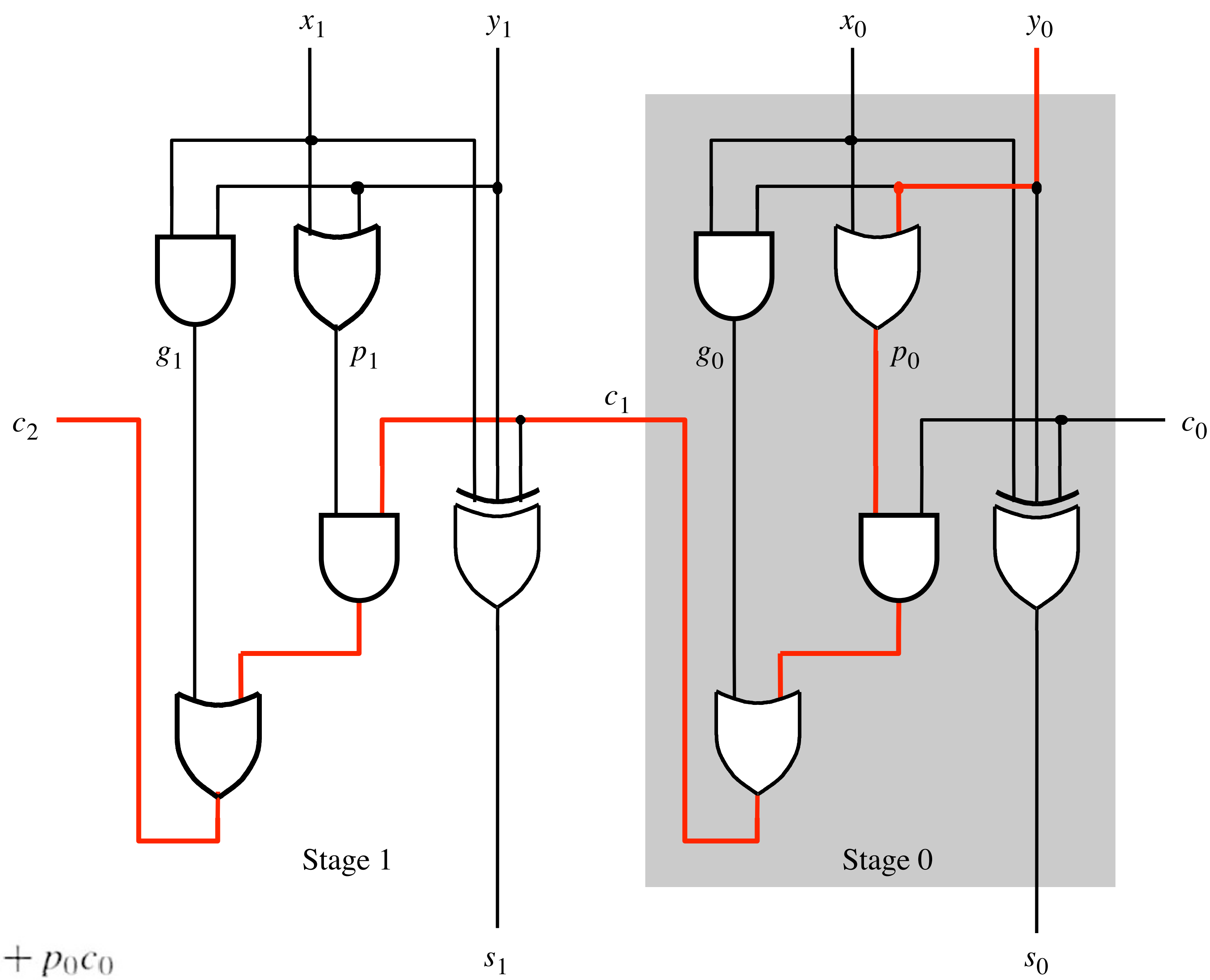


$$c_1 = g_0 + p_0 c_0$$

$$c_2 = g_1 + p_1 g_0 + p_1 p_0 c_0$$

[Figure 3.14 from the textbook]

Now we can Build a Ripple-Carry Adder

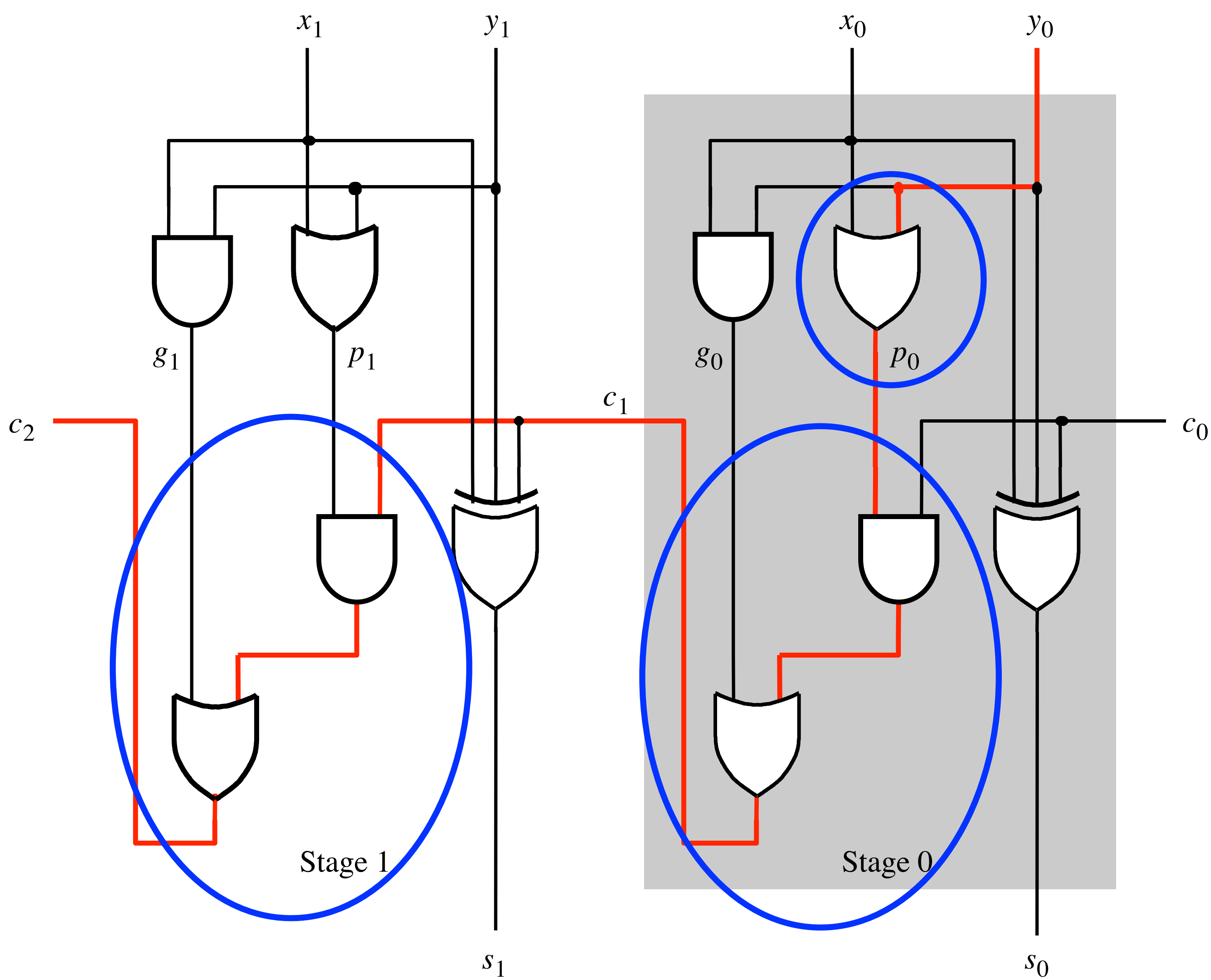


$$c_1 = g_0 + p_0 c_0$$

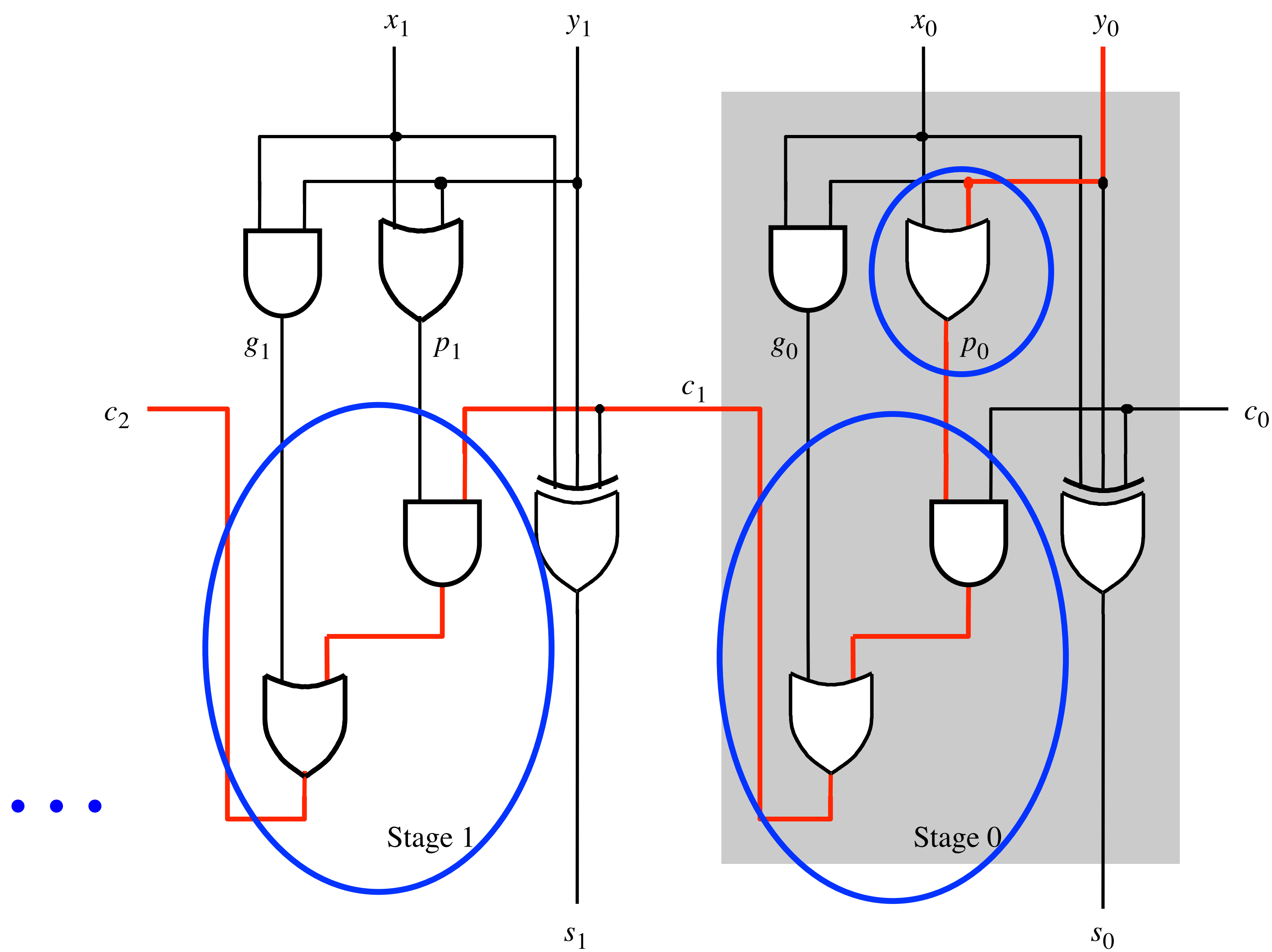
$$c_2 = g_1 + p_1 g_0 + p_1 p_0 c_0$$

[Figure 3.14 from the textbook]

The delay is 5 gates (1+2+2)



n-bit ripple-carry adder: 2n+1 gate delays



Decomposing the Carry Expression

$$c_{i+1} = x_i y_i + x_i c_i + y_i c_i$$

$$c_{i+1} = \underbrace{x_i y_i}_{g_i} + \underbrace{(x_i + y_i)}_{p_i} c_i$$

$$c_{i+1} = g_i + p_i c_i$$

$$c_{i+1} = g_i + p_i (g_{i-1} + p_{i-1} c_{i-1})$$

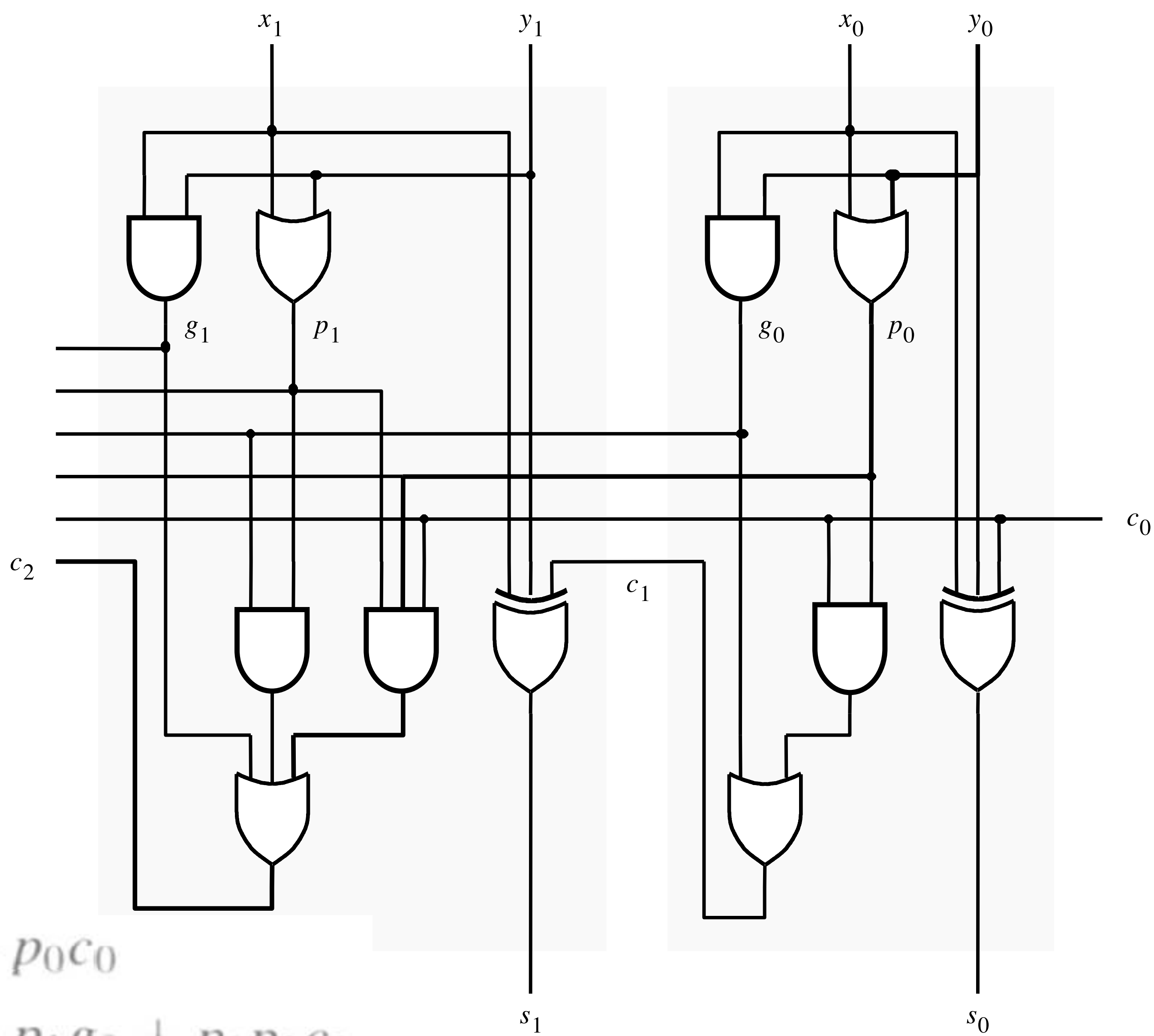
$$= g_i + p_i g_{i-1} + p_i p_{i-1} c_{i-1}$$

Carry for the first two stages

$$c_1 = g_0 + p_0 c_0$$

$$c_2 = g_1 + p_1 g_0 + p_1 p_0 c_0$$

The first two stages of a carry-lookahead adder

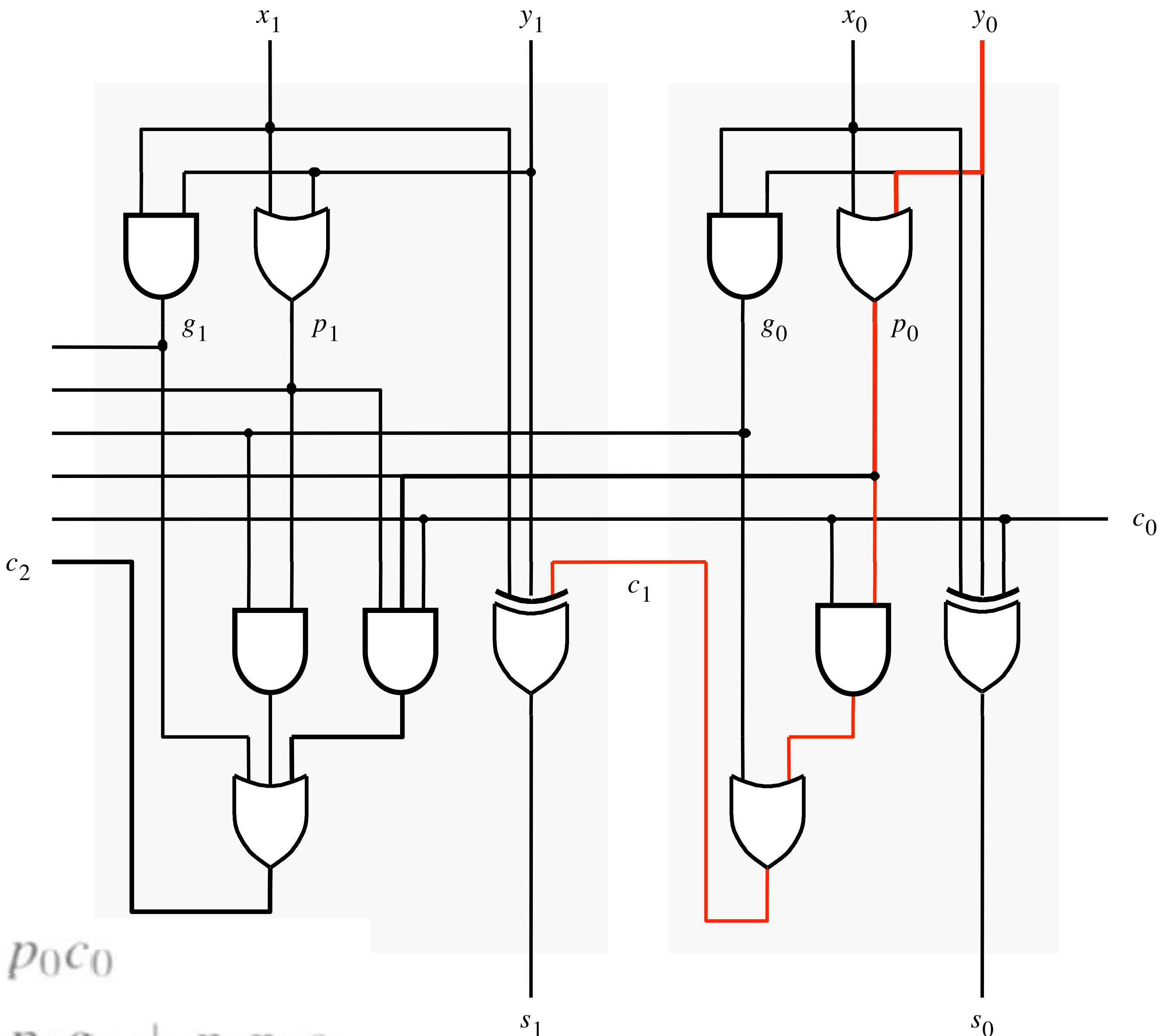


$$c_1 = g_0 + p_0c_0$$

$$c_2 = g_1 + p_1g_0 + p_1p_0c_0$$

[Figure 3.15 from the textbook]

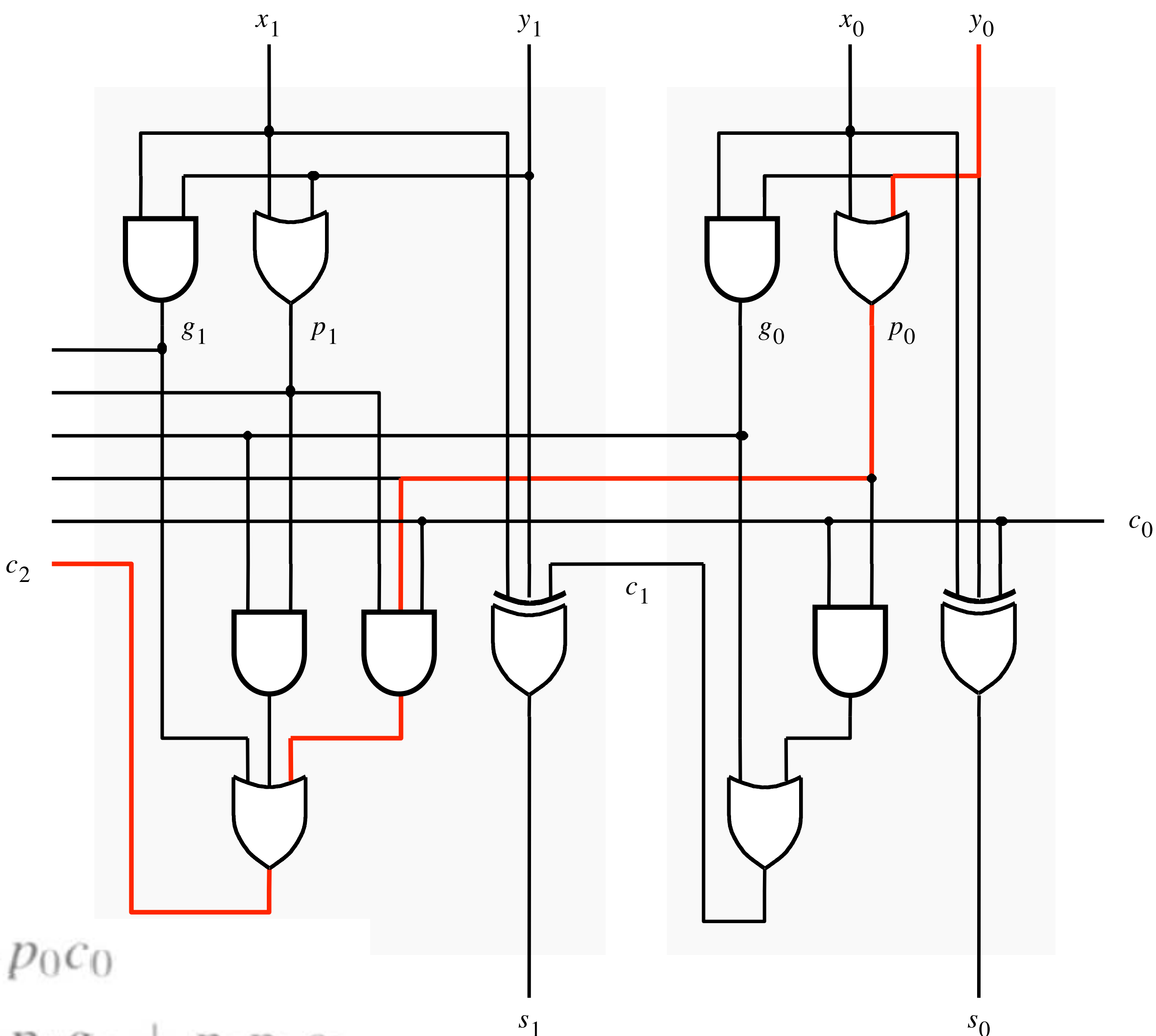
It takes 3 gate delays to generate c_1



$$c_1 = g_0 + p_0 c_0$$

$$c_2 = g_1 + p_1 g_0 + p_1 p_0 c_0$$

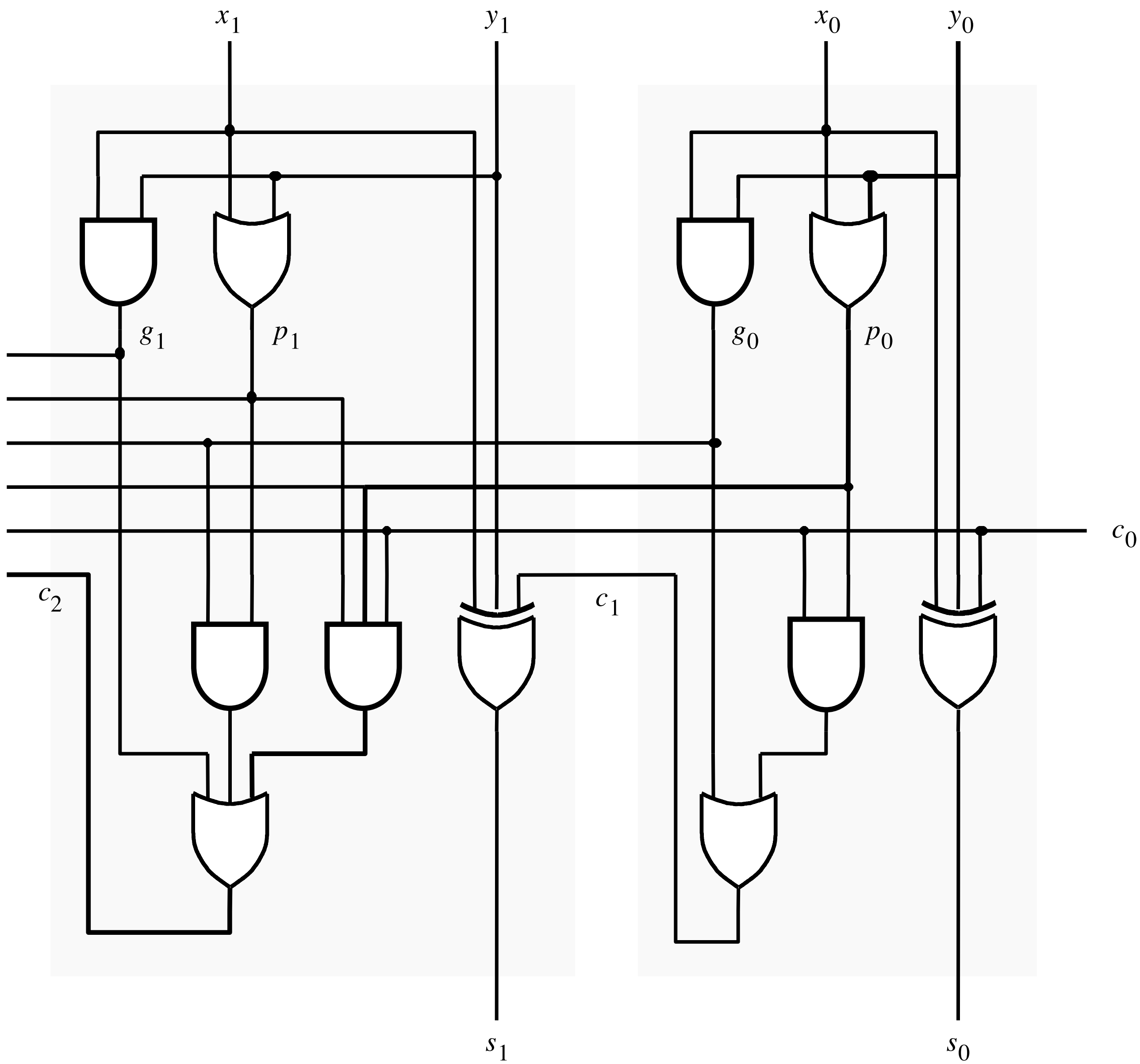
It takes 3 gate delays to generate c_2



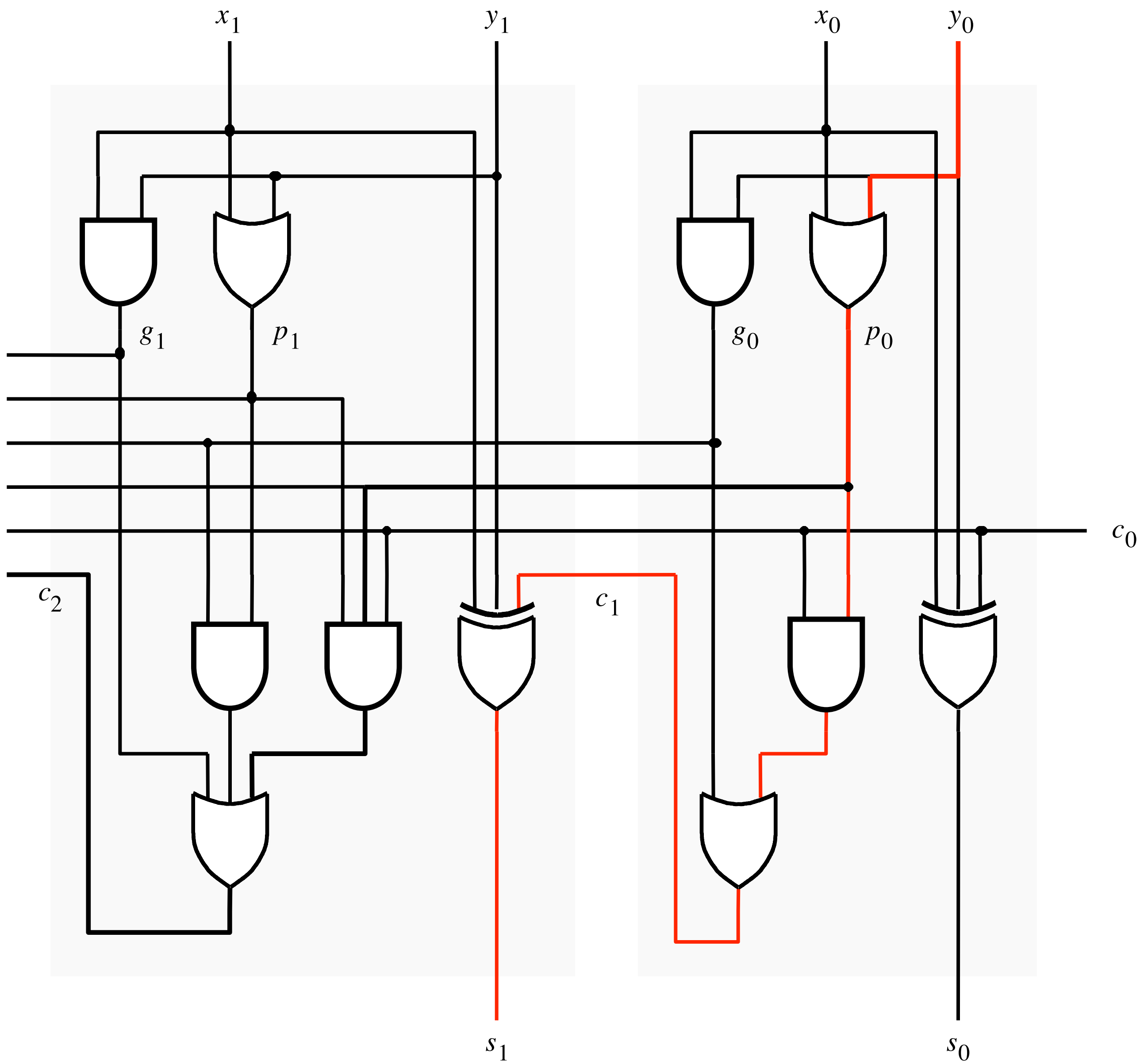
$$c_1 = g_0 + p_0 c_0$$

$$c_2 = g_1 + p_1 g_0 + p_1 p_0 c_0$$

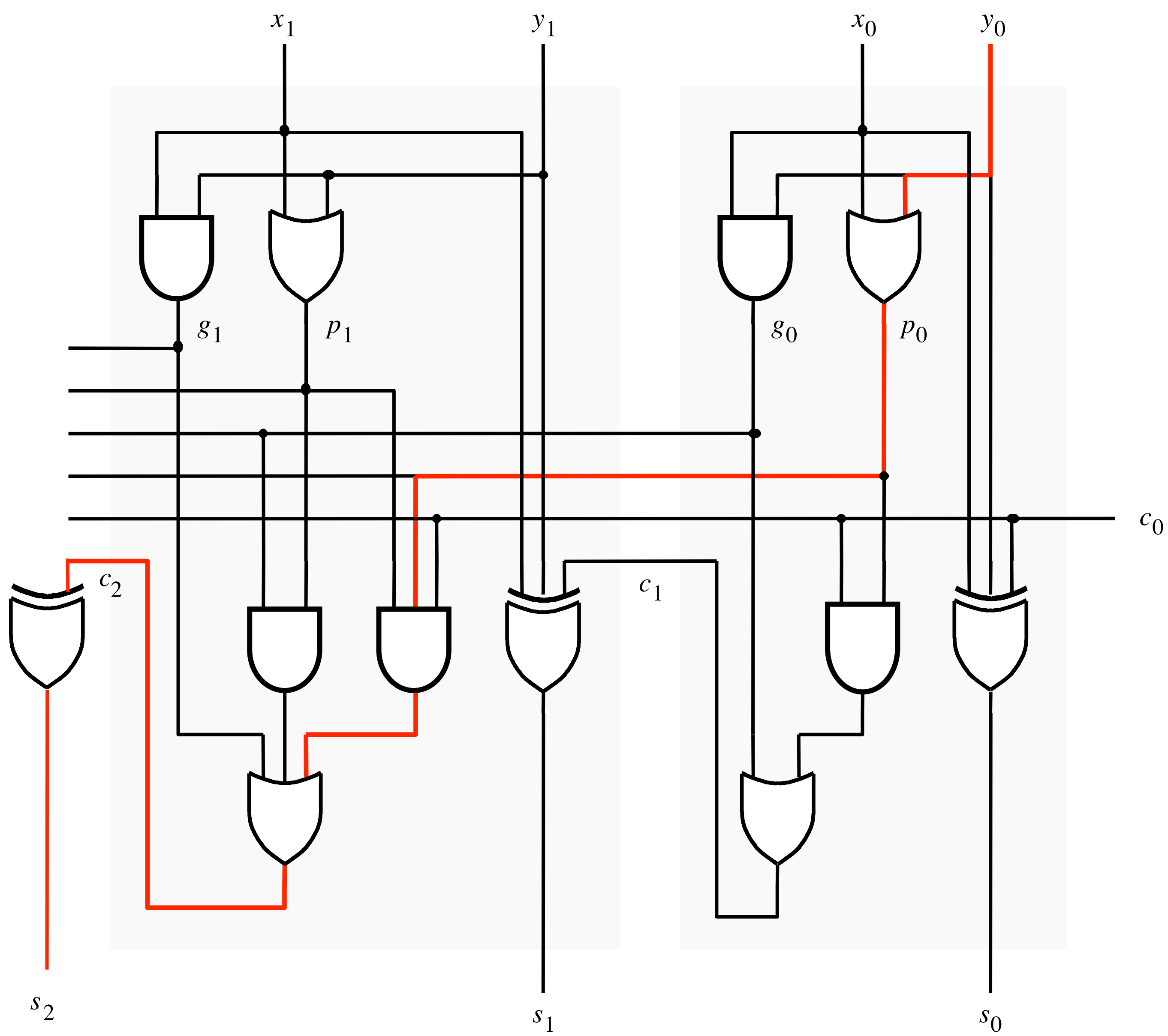
The first two stages of a carry-lookahead adder



It takes 4 gate delays to generate s_1



It takes 4 gate delays to generate s_2



N-bit Carry-Lookahead Adder

- **It takes 3 gate delays to generate all carry signals**
- **It takes 1 more gate delay to generate all sum bits**
- **Thus, the total delay through an n-bit carry-lookahead adder is only 4 gate delays!**

Expanding the Carry Expression

$$c_{i+1} = g_i + p_i c_i$$

$$c_1 = g_0 + p_0 c_0$$

$$c_2 = g_1 + p_1 g_0 + p_1 p_0 c_0$$

$$c_3 = g_2 + p_2 g_1 + p_2 p_1 g_0 + p_2 p_1 p_0 c_0$$

$$\cdots$$

$$\begin{aligned} c_8 = {} & g_7 + p_7 g_6 + p_7 p_6 g_5 + p_7 p_6 p_5 g_4 \\ & + p_7 p_6 p_5 p_4 g_3 + p_7 p_6 p_5 p_4 p_3 g_2 \\ & + p_7 p_6 p_5 p_4 p_3 p_2 g_1 + p_7 p_6 p_5 p_4 p_3 p_2 p_1 g_0 \\ & + p_7 p_6 p_5 p_4 p_3 p_2 p_1 p_0 c_0 \end{aligned}$$

Expanding the Carry Expression

$$c_{i+1} = g_i + p_i c_i$$

$$c_1 = g_0 + p_0 c_0$$

$$c_2 = g_1 + p_1 g_0 + p_1 p_0 c_0$$

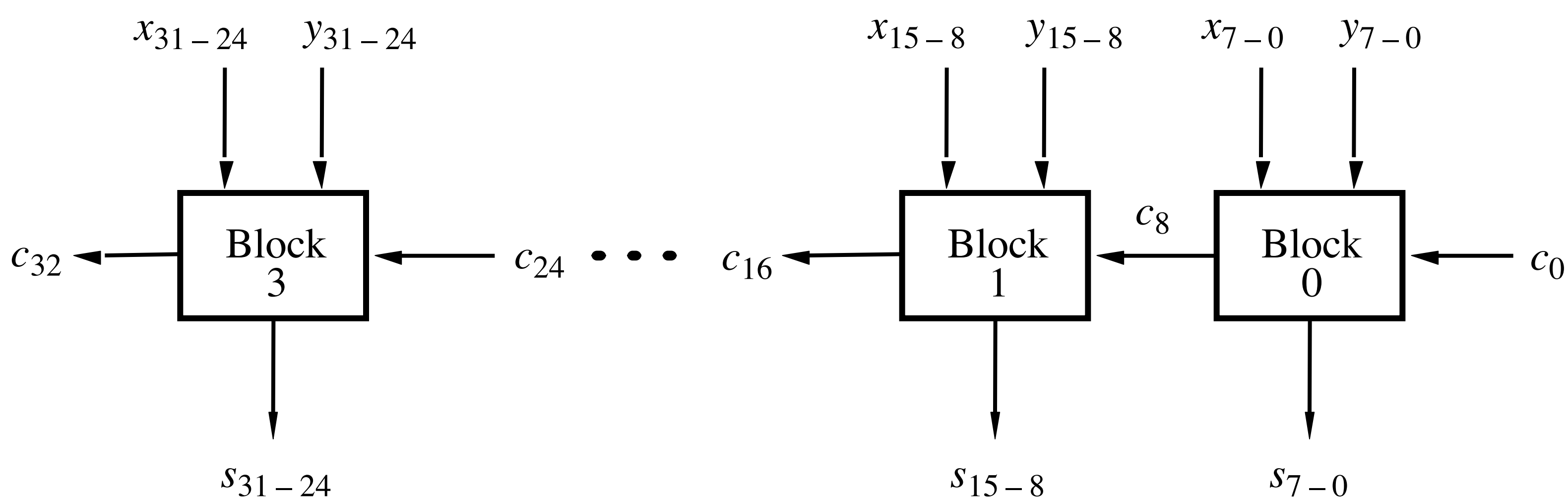
$$c_3 = g_2 + p_2 g_1 + p_2 p_1 g_0 + p_2 p_1 p_0 c_0$$

$$\cdots$$

$$c_8 = g_7 + p_7 g_6 + p_7 p_6 g_5 + p_7 p_6 p_5 g_4$$
$$+ p_7 p_6 p_5 p_4 g_3 + p_7 p_6 p_5 p_4 p_3 g_2$$
$$+ p_7 p_6 p_5 p_4 p_3 p_2 g_1 + p_7 p_6 p_5 p_4 p_3 p_2 p_1 g_0$$
$$+ p_7 p_6 p_5 p_4 p_3 p_2 p_1 p_0 c_0$$

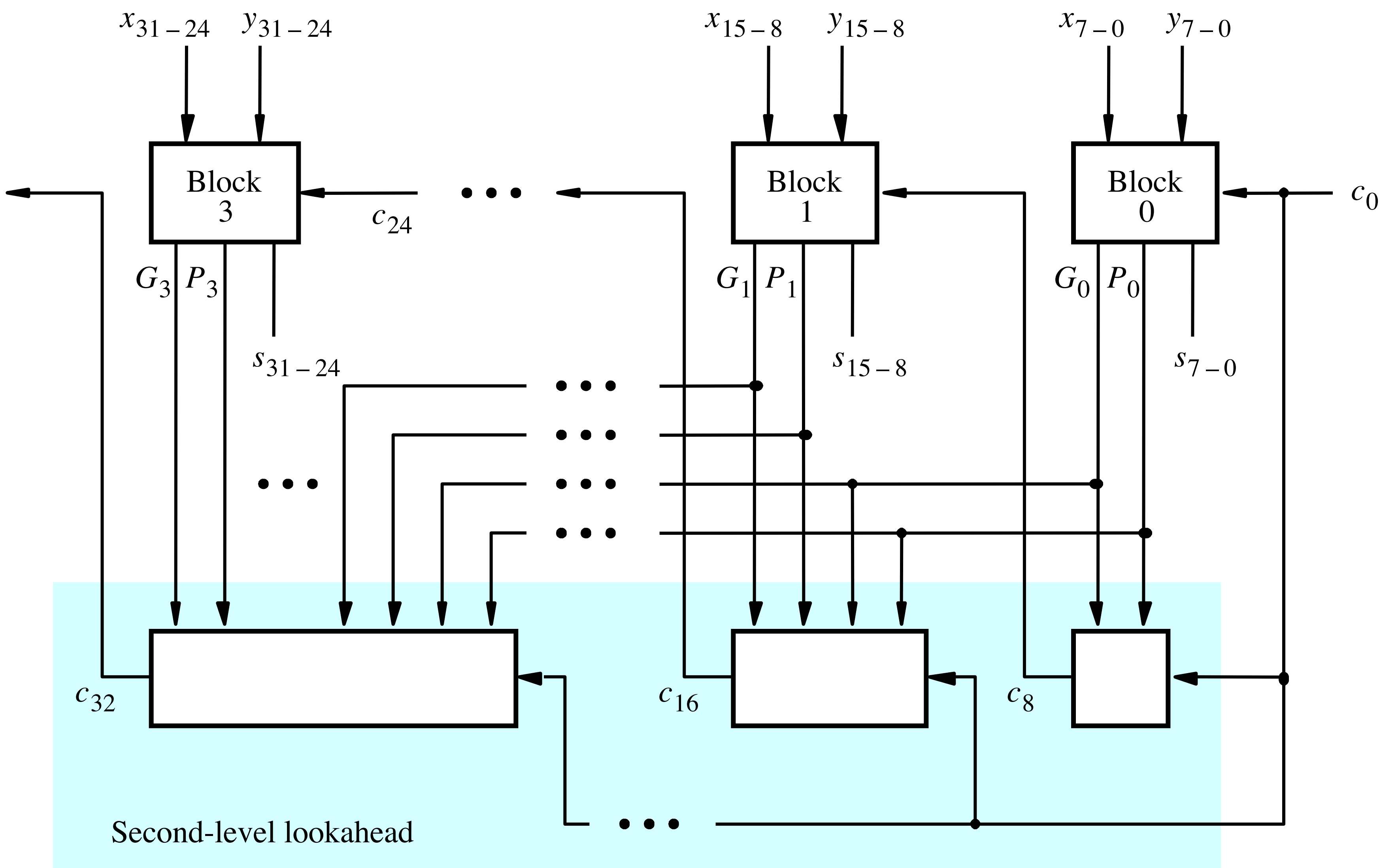
Even this takes only 3 gate delays

A hierarchical carry-lookahead adder with ripple-carry between blocks



[Figure 3.16 from the textbook]

A hierarchical carry-lookahead adder



[Figure 3.17 from the textbook]

The Hierarchical Carry Expression

$$\begin{aligned} c_8 = \; & g_7 + p_7 g_6 + p_7 p_6 g_5 + p_7 p_6 p_5 g_4 \\ & + p_7 p_6 p_5 p_4 g_3 + p_7 p_6 p_5 p_4 p_3 g_2 \\ & + p_7 p_6 p_5 p_4 p_3 p_2 g_1 + p_7 p_6 p_5 p_4 p_3 p_2 p_1 g_0 \\ & + p_7 p_6 p_5 p_4 p_3 p_2 p_1 p_0 c_0 \end{aligned}$$

The Hierarchical Carry Expression

$$c_8 = g_7 + p_7g_6 + p_7p_6g_5 + p_7p_6p_5g_4 + p_7p_6p_5p_4g_3 + p_7p_6p_5p_4p_3g_2 + p_7p_6p_5p_4p_3p_2g_1 + p_7p_6p_5p_4p_3p_2p_1g_0 + p_7p_6p_5p_4p_3p_2p_1p_0c_0$$

The Hierarchical Carry Expression

$$c_8 = g_7 + p_7g_6 + p_7p_6g_5 + p_7p_6p_5g_4$$
$$+ p_7p_6p_5p_4g_3 + p_7p_6p_5p_4p_3g_2$$
$$+ p_7p_6p_5p_4p_3p_2g_1 + p_7p_6p_5p_4p_3p_2p_1g_0$$
$$+ p_7p_6p_5p_4p_3p_2p_1p_0c_0$$

G_0 (blue box: $g_7 + \dots + p_7p_6p_5p_4p_3p_2p_1g_0$)

P_0 (red box: $p_7p_6p_5p_4p_3p_2p_1p_0$)

The Hierarchical Carry Expression

$$c_8 = \underbrace{g_7 + p_7g_6 + p_7p_6g_5 + p_7p_6p_5g_4 + p_7p_6p_5p_4g_3 + p_7p_6p_5p_4p_3g_2 + p_7p_6p_5p_4p_3p_2g_1 + p_7p_6p_5p_4p_3p_2p_1g_0}_{G_0} + \underbrace{p_7p_6p_5p_4p_3p_2p_1p_0}_{P_0}c_0$$

$$c_8 = G_0 + P_0c_0$$

The Hierarchical Carry Expression

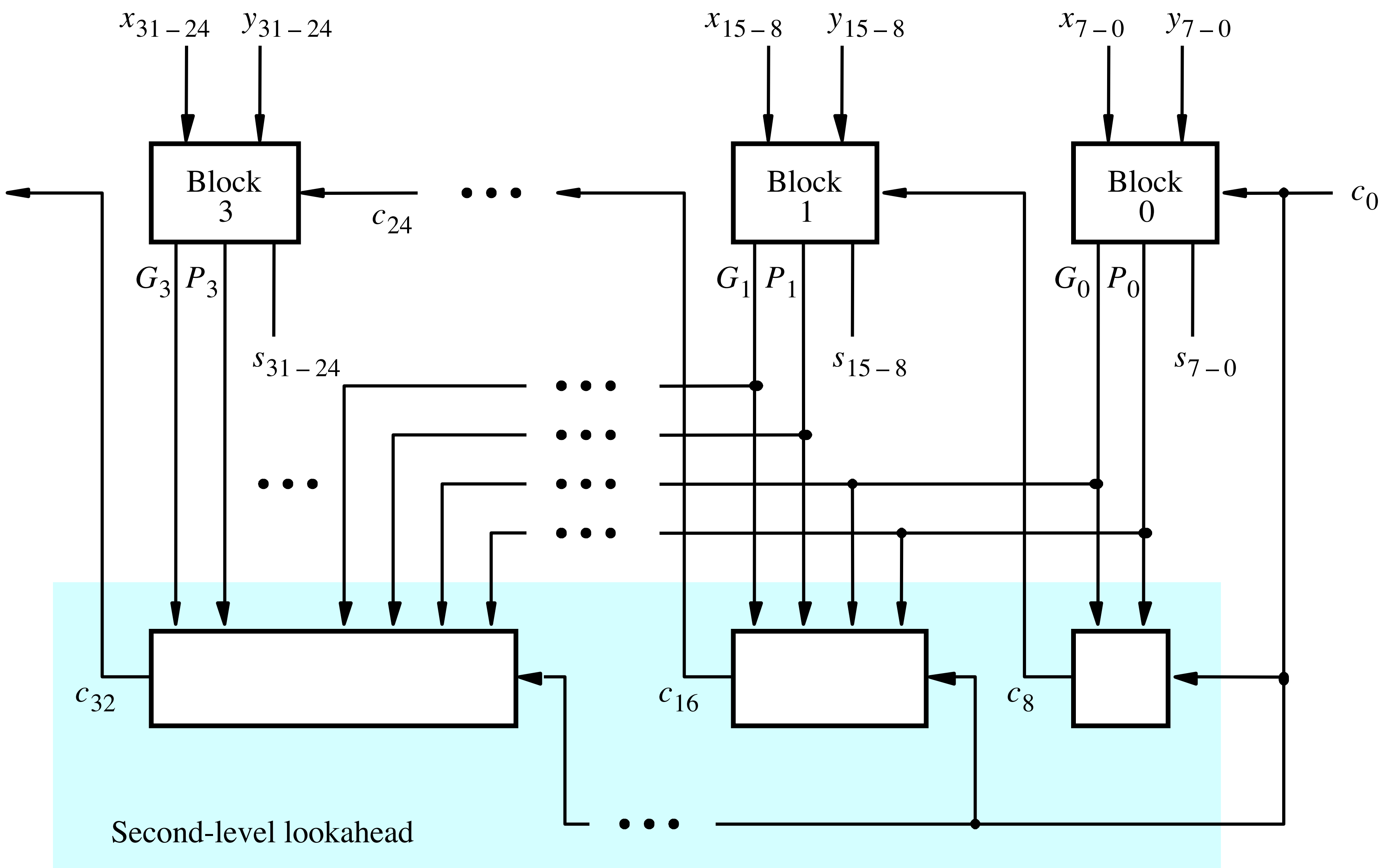
$$c_8 = G_0 + P_0 c_0$$

$$\begin{aligned} c_{16} &= G_1 + P_1 c_8 \\ &= G_1 + P_1 G_0 + P_1 P_0 c_0 \end{aligned}$$

$$c_{24} = G_2 + P_2 G_1 + P_2 P_1 G_0 + P_2 P_1 P_0 c_0$$

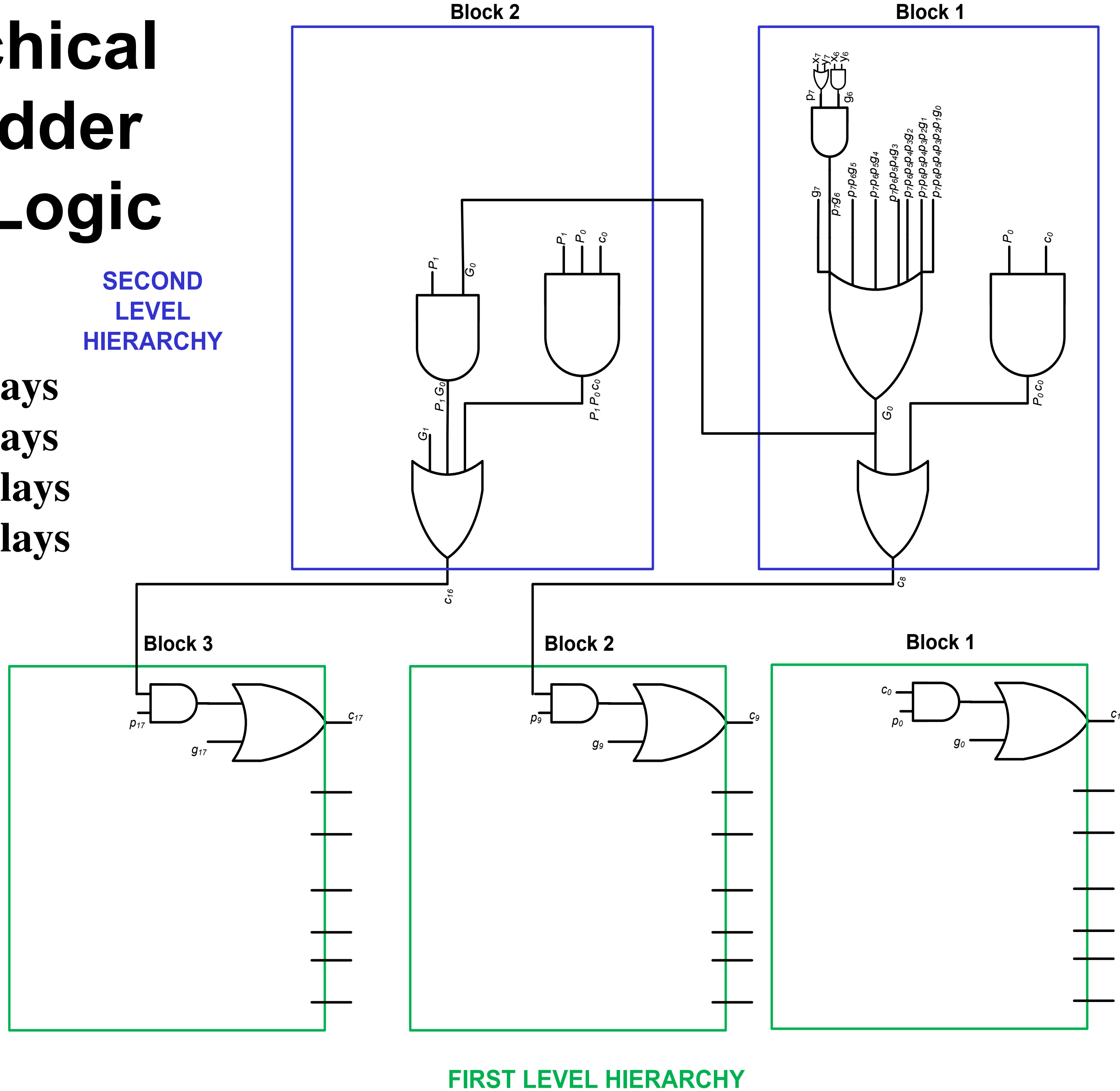
$$c_{32} = G_3 + P_3 G_2 + P_3 P_2 G_1 + P_3 P_2 P_1 G_0 + P_3 P_2 P_1 P_0 c_0$$

A hierarchical carry-lookahead adder



[Figure 3.17 from the textbook]

Hierarchical CLA Adder Carry Logic



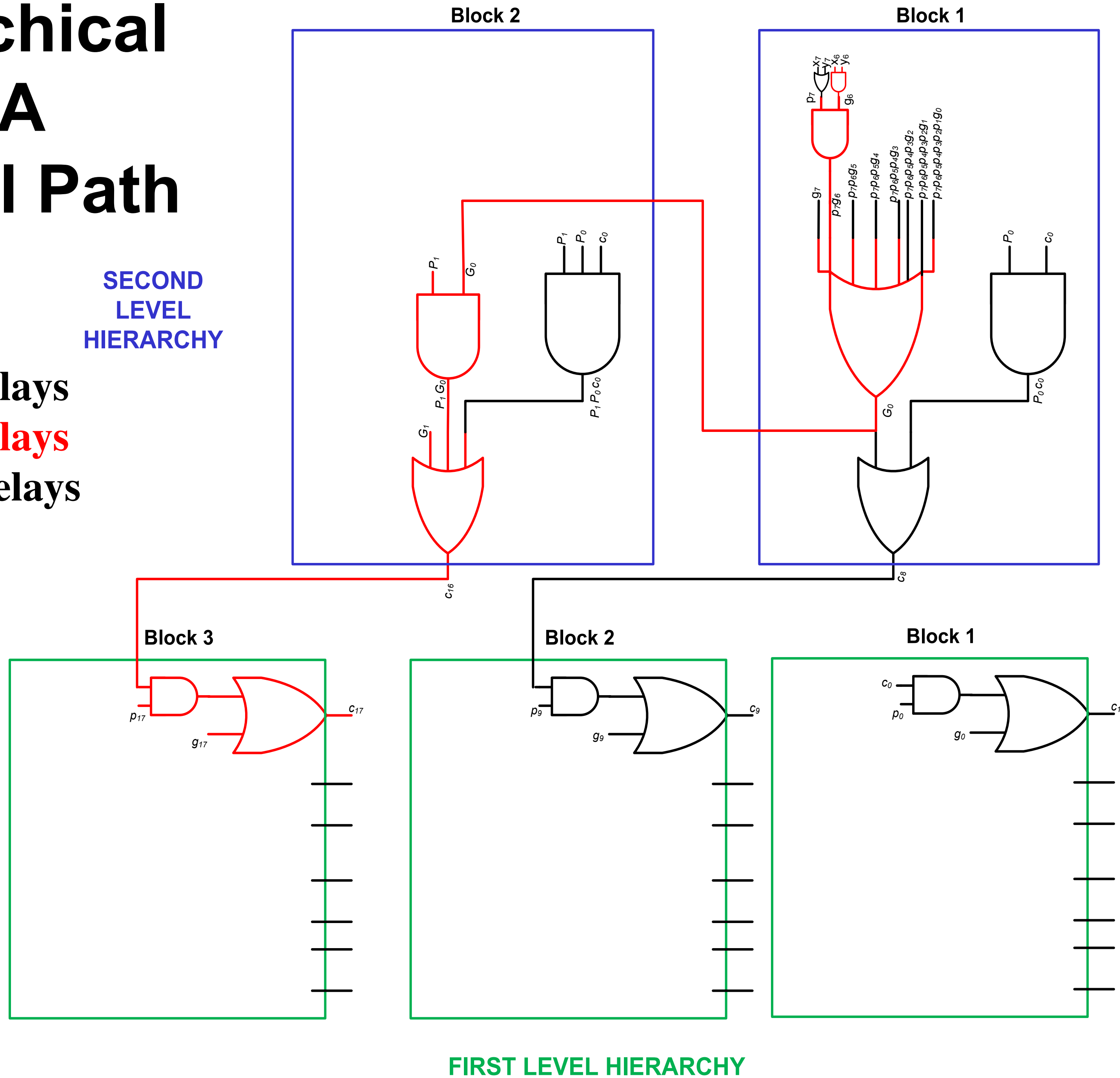
C8 – 5 gate delays
C16 – 5 gate delays
C24 – 5 Gate delays
C32 – 5 Gate delays

Hierarchical CLA Critical Path

C9 – 7 gate delays

C17 – 7 gate delays

C25 – 7 Gate delays



Total Gate Delay Through a Hierarchical Carry-Lookahead Adder

- **Is 8 gates**
  - **3 to generate all Gj and Pj**
  - **+2 to generate c8, c16, c24, and c32**
  - **+2 to generate internal carries in the blocks**
  - **+1 to generate the sum bits (one extra XOR)**

Questions?

THE END