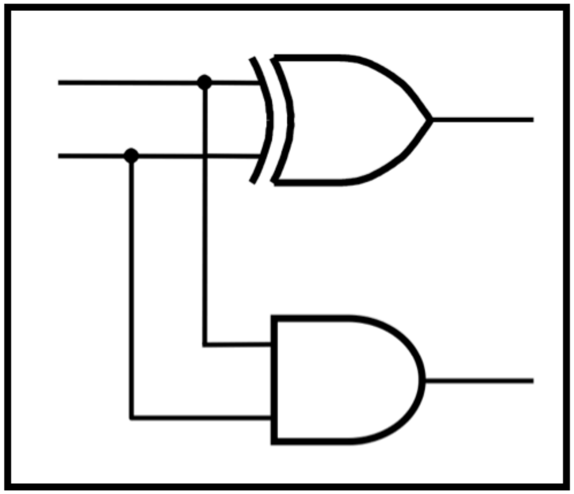


CprE 281: Digital Logic

Instructor: Alexander Stoytchev

http://www.ece.iastate.edu/~alexs/classes/

Examples of Solved Problems

CprE 281: Digital Logic
Iowa State University, Ames, IA

Administrative Stuff

- No homework is due next week.

Administrative Stuff

- Midterm Exam #1
- When: Friday Sep 25.
- Where: This classroom
- What: Chapter 1 and Chapter 2 plus number systems
- The exam will be open book and open notes (you can bring up to 3 pages of handwritten notes).

Topics for the Midterm Exam

- Binary Numbers
- Octal Numbers
- Hexadecimal Numbers
- Conversion between the different number systems
- Truth Tables
- Boolean Algebra
- Logic Gates
- Circuit Synthesis with AND, OR, NOT
- Circuit Synthesis with NAND, NOR
- Converting an AND/OR/NOT circuit to NAND circuit
- Converting an AND/OR/NOT circuit to NOR circuit
- SOP and POS expressions

Topics for the Midterm Exam

- Mapping a Circuit to Verilog code
- Mapping Verilog code to a circuit

- Multiplexers
- Venn Diagrams
- K-maps for 2, 3, and 4 variables

- Minimization of Boolean expressions using theorems
- Minimization of Boolean expressions with K-maps

- Incompletely specified functions (with don’t cares)
- Functions with multiple outputs

Example 1

Determine if the following equation is valid

$$\bar{x}_1\bar{x}_3 + x_2x_3 + x_1\bar{x}_2 = \bar{x}_1x_2 + x_1x_3 + \bar{x}_2\bar{x}_3$$

$$\bar{x}_1\bar{x}_3 + x_2x_3 + x_1\bar{x}_2 \stackrel{?}{=} \bar{x}_1x_2 + x_1x_3 + \bar{x}_2\bar{x}_3$$

$$\underbrace{\bar{x}_1\bar{x}_3 + x_2x_3 + x_1\bar{x}_2}_{\text{LHS}} \stackrel{?}{=} \underbrace{\bar{x}_1x_2 + x_1x_3 + \bar{x}_2\bar{x}_3}_{\text{RHS}}$$

Left-Hand Side (LHS)

Row number	x_1	x_2	x_3	$\overline{x_1}\,\overline{x_3}$	$x_2 x_3$	$x_1 \overline{x_2}$	f
0	0	0	0				
1	0	0	1				
2	0	1	0				
3	0	1	1				
4	1	0	0				
5	1	0	1				
6	1	1	0				
7	1	1	1				

Left-Hand Side (LHS)

Row number	x_1	x_2	x_3	$\overline{x_1}\,\overline{x_3}$	$x_2\,x_3$	$x_1\,\overline{x_2}$	f
0	0	0	0	1	0	0	
1	0	0	1	0	0	0	
2	0	1	0	1	0	0	
3	0	1	1	0	1	0	
4	1	0	0	0	0	1	
5	1	0	1	0	0	1	
6	1	1	0	0	0	0	
7	1	1	1	0	1	0	

Left-Hand Side (LHS)

Row number	x_1	x_2	x_3	$\overline{x}_1 \overline{x}_3$	$x_2 x_3$	$x_1 \overline{x}_2$	f
0	0	0	0	1	0	0	1
1	0	0	1	0	0	0	0
2	0	1	0	1	0	0	1
3	0	1	1	0	1	0	1
4	1	0	0	0	0	1	1
5	1	0	1	0	0	1	1
6	1	1	0	0	0	0	0
7	1	1	1	0	1	0	1

Right-Hand Side (RHS)

Row number	x_1	x_2	x_3	$\overline{x_1}x_2$	x_1x_3	$\overline{x_2}\ \overline{x_3}$	f
0	0	0	0				
1	0	0	1				
2	0	1	0				
3	0	1	1				
4	1	0	0				
5	1	0	1				
6	1	1	0				
7	1	1	1				

Right-Hand Side (RHS)

Row number	x_1	x_2	x_3	$\overline{x_1} x_2$	$x_1 x_3$	$\overline{x_2}\,\overline{x_3}$	f
0	0	0	0	0	0	1	
1	0	0	1	0	0	0	
2	0	1	0	1	0	0	
3	0	1	1	1	0	0	
4	1	0	0	0	0	1	
5	1	0	1	0	1	0	
6	1	1	0	0	0	0	
7	1	1	1	0	1	0	

Right-Hand Side (RHS)

Row number	x_1	x_2	x_3	$\overline{x}_1 x_2$	$x_1 x_3$	$\overline{x}_2\ \overline{x}_3$	f
0	0	0	0	0	0	1	1
1	0	0	1	0	0	0	0
2	0	1	0	1	0	0	1
3	0	1	1	1	0	0	1
4	1	0	0	0	0	1	1
5	1	0	1	0	1	0	1
6	1	1	0	0	0	0	0
7	1	1	1	0	1	0	1

$$\underbrace{\bar{x}_1\bar{x}_3 + x_2x_3 + x_1\bar{x}_2}_{\text{LHS}} \stackrel{?}{=} \underbrace{\bar{x}_1x_2 + x_1x_3 + \bar{x}_2\bar{x}_3}_{\text{RHS}}$$

f
1
0
1
1
1
1
0
1

f
1
0
1
1
1
1
0
1

Example 2

Design the minimum-cost product-of-sums expression for the function

$$f(x_1, x_2, x_3) = \Sigma\, m(0, 2, 4, 5, 6, 7)$$

Minterms and Maxterms (with three variables)

Row number	x_1	x_2	x_3	Minterm	Maxterm
0	0	0	0	$m_0 = \overline{x}_1\overline{x}_2\overline{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \overline{x}_1\overline{x}_2 x_3$	$M_1 = x_1 + x_2 + \overline{x}_3$
2	0	1	0	$m_2 = \overline{x}_1 x_2\overline{x}_3$	$M_2 = x_1 + \overline{x}_2 + x_3$
3	0	1	1	$m_3 = \overline{x}_1 x_2 x_3$	$M_3 = x_1 + \overline{x}_2 + \overline{x}_3$
4	1	0	0	$m_4 = x_1\overline{x}_2\overline{x}_3$	$M_4 = \overline{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1\overline{x}_2 x_3$	$M_5 = \overline{x}_1 + x_2 + \overline{x}_3$
6	1	1	0	$m_6 = x_1 x_2\overline{x}_3$	$M_6 = \overline{x}_1 + \overline{x}_2 + x_3$
7	1	1	1	$m_7 = x_1 x_2 x_3$	$M_7 = \overline{x}_1 + \overline{x}_2 + \overline{x}_3$

[Figure 2.22 from the textbook]

Minterms and Maxterms (with three variables)

Row number	x_1	x_2	x_3	Minterm	Maxterm
0	0	0	0	$m_0 = \overline{x}_1\overline{x}_2\overline{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \overline{x}_1\overline{x}_2 x_3$	$M_1 = x_1 + x_2 + \overline{x}_3$
2	0	1	0	$m_2 = \overline{x}_1 x_2\overline{x}_3$	$M_2 = x_1 + \overline{x}_2 + x_3$
3	0	1	1	$m_3 = \overline{x}_1 x_2 x_3$	$M_3 = x_1 + \overline{x}_2 + \overline{x}_3$
4	1	0	0	$m_4 = x_1\overline{x}_2\overline{x}_3$	$M_4 = \overline{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1\overline{x}_2 x_3$	$M_5 = \overline{x}_1 + x_2 + \overline{x}_3$
6	1	1	0	$m_6 = x_1 x_2\overline{x}_3$	$M_6 = \overline{x}_1 + \overline{x}_2 + x_3$
7	1	1	1	$m_7 = x_1 x_2 x_3$	$M_7 = \overline{x}_1 + \overline{x}_2 + \overline{x}_3$

The function is 1 for these rows

Minterms and Maxterms (with three variables)

Row number	x_1	x_2	x_3	Minterm	Maxterm
0	0	0	0	$m_0 = \overline{x}_1\overline{x}_2\overline{x}_3$	$M_0 = x_1 + x_2 + x_3$
1	0	0	1	$m_1 = \overline{x}_1\overline{x}_2x_3$	$M_1 = x_1 + x_2 + \overline{x}_3$
2	0	1	0	$m_2 = \overline{x}_1x_2\overline{x}_3$	$M_2 = x_1 + \overline{x}_2 + x_3$
3	0	1	1	$m_3 = \overline{x}_1x_2x_3$	$M_3 = x_1 + \overline{x}_2 + \overline{x}_3$
4	1	0	0	$m_4 = x_1\overline{x}_2\overline{x}_3$	$M_4 = \overline{x}_1 + x_2 + x_3$
5	1	0	1	$m_5 = x_1\overline{x}_2x_3$	$M_5 = \overline{x}_1 + x_2 + \overline{x}_3$
6	1	1	0	$m_6 = x_1x_2\overline{x}_3$	$M_6 = \overline{x}_1 + \overline{x}_2 + x_3$
7	1	1	1	$m_7 = x_1x_2x_3$	$M_7 = \overline{x}_1 + \overline{x}_2 + \overline{x}_3$

The function is 1 for these rows

The function is 0 for these rows

Two different ways to specify the same function f of three variables

$$f(x_1, x_2, x_3) = \Sigma\, m(0, 2, 4, 5, 6, 7)$$

$$f(x_1, x_2, x_3) = \Pi\, M(1, 3)$$

The POS Expression

$$M_1 = x_1 + x_2 + \overline{x}_3 \qquad M_3 = x_1 + \overline{x}_2 + \overline{x}_3$$

$$f(x_1, x_2, x_3) = \Pi\, M(1, 3)$$

$$= M_1 \bullet M_3$$

$$= (x_1 + x_2 + \overline{x_3}) \bullet (x_1 + \overline{x_2} + \overline{x_3})$$

The Minimum POS Expression

$$f(x_1, x_2, x_3) = (x_1 + x_2 + \overline{x_3}) \bullet (x_1 + \overline{x_2} + \overline{x_3})$$

$$= (x_1 + \overline{x_3} + x_2) \bullet (x_1 + \overline{x_3} + \overline{x_2})$$

$$= (x_1 + \overline{x_3})$$

Hint: Use the following Boolean Algebra theorem

14b. $(x + y) \bullet (x + \overline{y}) = x$

Alternative Solution Using K-Maps

x_1	x_2	x_3	
0	0	0	m_0
0	0	1	m_1
0	1	0	m_2
0	1	1	m_3
1	0	0	m_4
1	0	1	m_5
1	1	0	m_6
1	1	1	m_7

(a) Truth table

x_3 \ $x_1 x_2$	00	01	11	10
0	m_0	m_2	m_6	m_4
1	m_1	m_3	m_7	m_5

(b) Karnaugh map

Alternative Solution Using K-Maps

x_1	x_2	x_3	
0	0	0	m_0
0	0	1	m_1
0	1	0	m_2
0	1	1	m_3
1	0	0	m_4
1	0	1	m_5
1	1	0	m_6
1	1	1	m_7

(a) Truth table

x_3 \ x_1x_2	00	01	11	10
0	m_0	m_2	m_6	m_4
1	m_1	m_3	m_7	m_5

(b) Karnaugh map

Alternative Solution Using K-Maps

x_1	x_2	x_3	
0	0	0	m_0
0	0	1	m_1
0	1	0	m_2
0	1	1	m_3
1	0	0	m_4
1	0	1	m_5
1	1	0	m_6
1	1	1	m_7

(a) Truth table

x_3 \ x_1x_2	00	01	11	10
0	m_0	m_2	m_6	m_4
1	m_1	m_3	m_7	m_5

(b) Karnaugh map

Alternative Solution Using K-Maps

$x_3 \backslash x_1x_2$	00	01	11	10
0	m_0	m_2	m_6	m_4
1	m_1	m_3	m_7	m_5

Alternative Solution Using K-Maps

$x_3 \backslash x_1x_2$	00	01	11	10
0	1	1	1	1
1	0	0	1	1

Alternative Solution Using K-Maps

x_3 \ x_1x_2	00	01	11	10
0	1	1	1	1
1	0	0	1	1

$$(x_1 + \overline{x}_3)$$

Example 3

Problem: A circuit that controls a given digital system has three inputs: x_1, x_2, and x_3. It has to recognize three different conditions:

- Condition A is true if x_3 is true and either x_1 is true or x_2 is false
- Condition B is true if x_1 is true and either x_2 or x_3 is false
- Condition C is true if x_2 is true and either x_1 is true or x_3 is false

The control circuit must produce an output of 1 if at least two of the conditions A, B, and C are true. Design the simplest circuit that can be used for this purpose.

Condition A

Condition A is true if x_3 is true and either x_1 is true or x_2 is false

Condition A

Condition A is true if x_3 is true and either x_1 is true or x_2 is false

$$A = x_3(x_1 + \bar{x}_2) \;=\; x_3x_1 + x_3\bar{x}_2$$

Condition B

Condition B is true if x_1 is true and either x_2 or x_3 is false

Condition B

Condition B is true if x_1 is true and either x_2 or x_3 is false

$$B = x_1(\bar{x}_2 + \bar{x}_3) \;=\; x_1\bar{x}_2 + x_1\bar{x}_3$$

Condition C

Condition C is true if x_2 is true and either x_1 is true or x_3 is false

Condition C

Condition C is true if x_2 is true and either x_1 is true or x_3 is false

$$C = x_2(x_1 + \bar{x}_3) = x_2x_1 + x_2\bar{x}_3$$

The output of the circuit can be expressed as
$$f = \boxed{AB} + AC + BC$$

$$\begin{aligned}
AB &= (x_3x_1 + x_3\overline{x}_2)(x_1\overline{x}_2 + x_1\overline{x}_3) \\
&= x_3x_1x_1\overline{x}_2 + x_3x_1x_1\overline{x}_3 + x_3\overline{x}_2x_1\overline{x}_2 + x_3\overline{x}_2x_1\overline{x}_3 \\
&= x_3x_1\overline{x}_2 + 0 + x_3\overline{x}_2x_1 + 0 \\
&= x_1\overline{x}_2x_3
\end{aligned}$$

The output of the circuit can be expressed as
f = AB + AC + BC

$$
\begin{aligned}
AC &= (x_3x_1 + x_3\bar{x}_2)(x_2x_1 + x_2\bar{x}_3) \\
&= x_3x_1x_2x_1 + x_3x_1x_2\bar{x}_3 + x_3\bar{x}_2x_2x_1 + x_3\bar{x}_2x_2\bar{x}_3 \\
&= x_3x_1x_2 + 0 + 0 + 0 \\
&= x_1x_2x_3
\end{aligned}
$$

The output of the circuit can be expressed as

f = AB + AC + BC

$$
\begin{aligned}
BC &= (x_1\bar{x}_2 + x_1\bar{x}_3)(x_2x_1 + x_2\bar{x}_3) \\
&= x_1\bar{x}_2x_2x_1 + x_1\bar{x}_2x_2\bar{x}_3 + x_1\bar{x}_3x_2x_1 + x_1\bar{x}_3x_2\bar{x}_3 \\
&= 0 + 0 + x_1\bar{x}_3x_2 + x_1\bar{x}_3x_2 \\
&= x_1x_2\bar{x}_3
\end{aligned}
$$

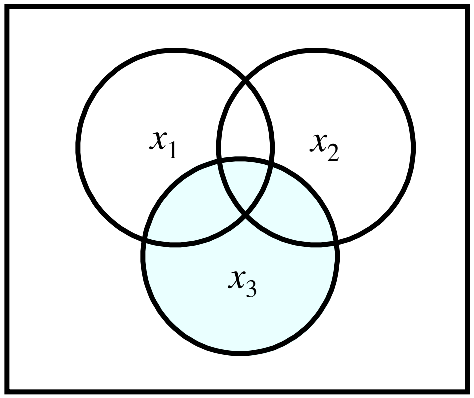
Finally, we get

$$\begin{aligned} f &= x_1\bar{x}_2x_3 + x_1x_2x_3 + x_1x_2\bar{x}_3 \\ &= x_1(\bar{x}_2 + x_2)x_3 + x_1x_2(x_3 + \bar{x}_3) \\ &= x_1x_3 + x_1x_2 \\ &= x_1(x_3 + x_2) \end{aligned}$$

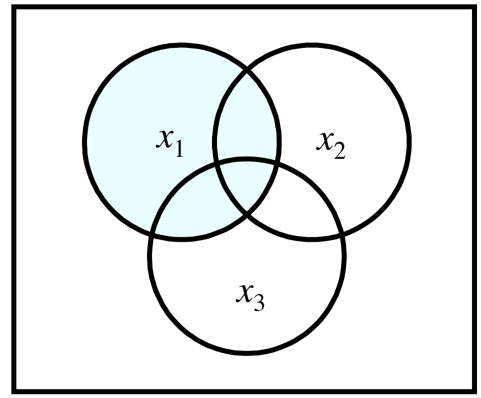
Example 4

Solve the previous problem using Venn diagrams.

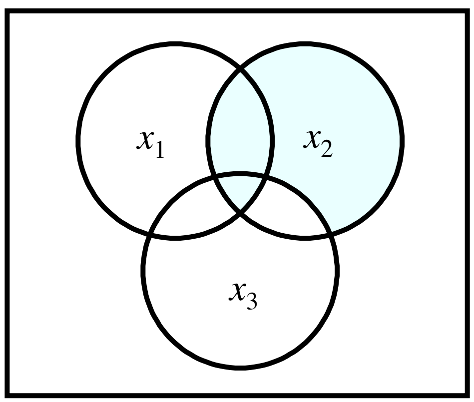
Venn Diagrams
(find the areas that are shaded at least two times)



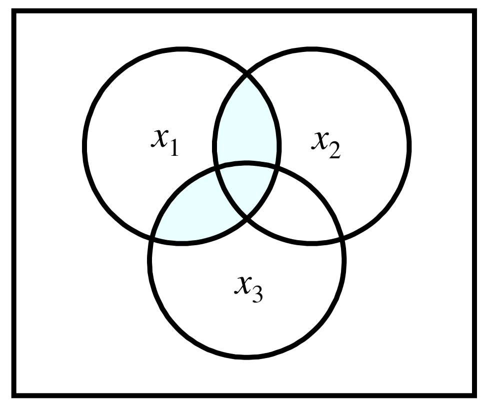
(a) Function A



(b) Function B



(c) Function C



(d) Function f

[Figure 2.66 from the textbook]

Example 5

Design the minimum-cost SOP and POS expression for the function

$$f(x_1, x_2, x_3, x_4) = \Sigma\, m(4, 6, 8, 10, 11, 12, 15) + D(3, 5, 7, 9)$$

Let's Use a K-Map

$$f(x_1, x_2, x_3, x_4) = \Sigma\, m(4, 6, 8, 10, 11, 12, 15) + D(3, 5, 7, 9)$$

$x_3x_4 \backslash x_1x_2$	00	01	11	10
00	m_0	m_4	m_{12}	m_8
01	m_1	m_5	m_{13}	m_9
11	m_3	m_7	m_{15}	m_{11}
10	m_2	m_6	m_{14}	m_{10}

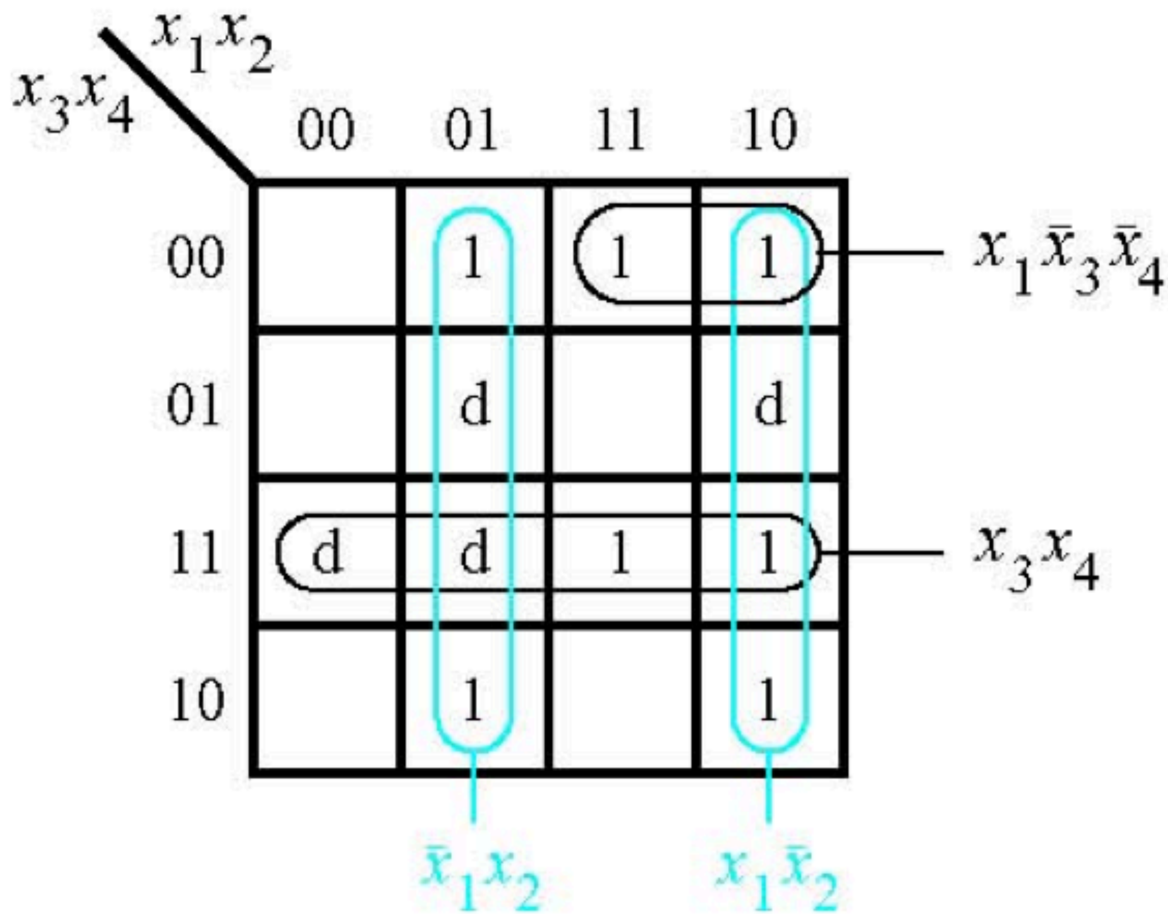
Let’s Use a K-Map

$$f(x_1, x_2, x_3, x_4) = \Sigma\, m(4, 6, 8, 10, 11, 12, 15) + D(3, 5, 7, 9)$$

x_3x_4 \ x_1x_2	00	01	11	10
00	0	1	1	1
01	0	d	0	d
11	d	d	1	1
10	0	1	0	1

x_1 x_2 x_3 x_4

The SOP Expression



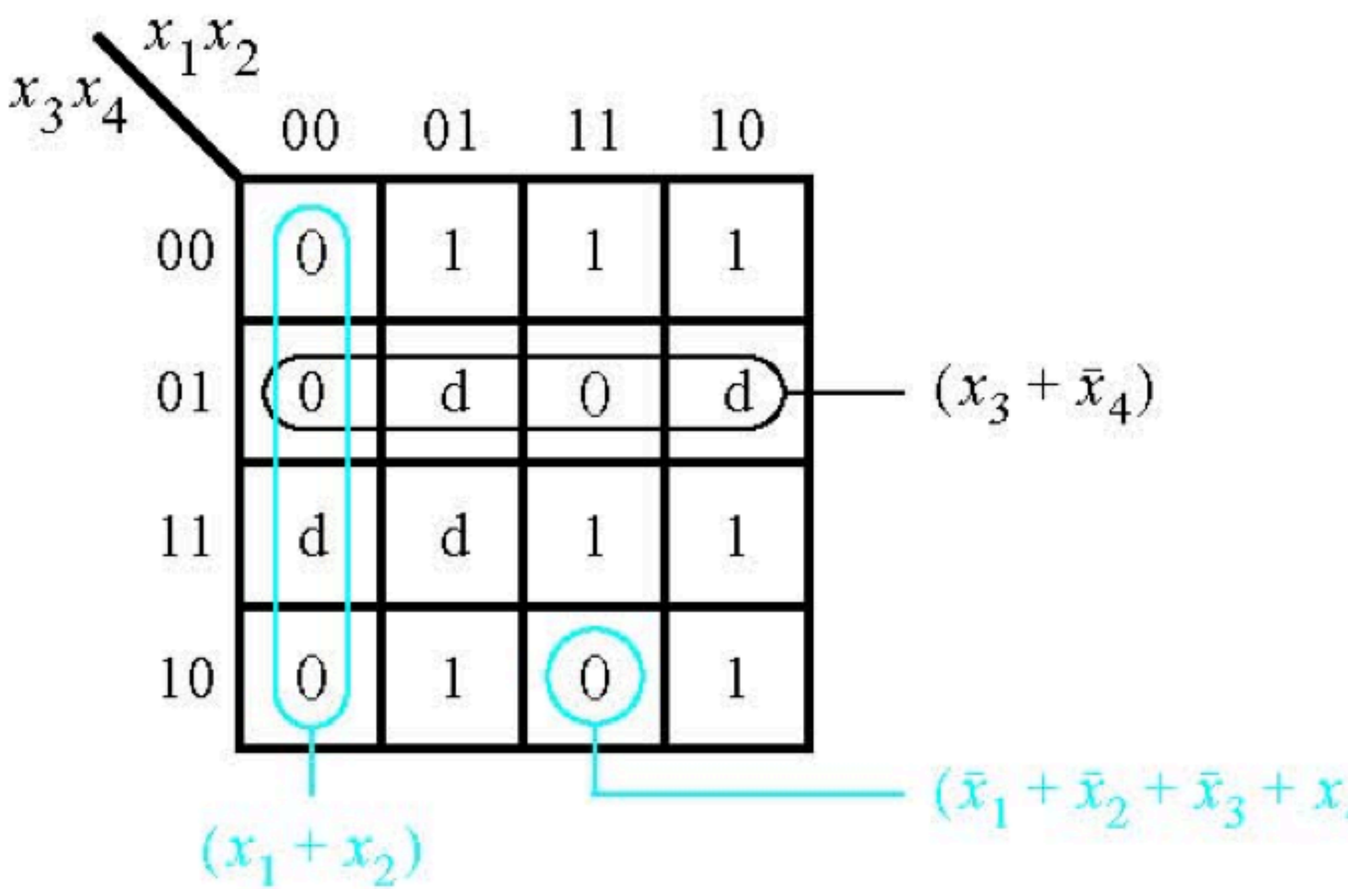
[Figure 2.67a from the textbook]

What about the POS Expression?

$$f(x_1, x_2, x_3, x_4) = \Sigma\, m(4, 6, 8, 10, 11, 12, 15) + D(3, 5, 7, 9)$$

x_3x_4 \ x_1x_2	00	01	11	10
00	0	1	1	1
01	0	d	0	d
11	d	d	1	1
10	0	1	0	1

The POS Expression



[Figure 2.67b from the textbook]

Example 6

Use K-maps to find the minimum-cost SOP and POS expression for the function

$$f(x_1, \ldots, x_4) = \overline{x}_1\overline{x}_3\overline{x}_4 + x_3x_4 + \overline{x}_1\overline{x}_2x_4 + x_1x_2\overline{x}_3x_4$$

assuming that there are also don't-cares defined as $D = \sum(9, 12, 14)$.

Let’s map the expression to the K-Map

$$f(x_1, \ldots, x_4) = \bar{x}_1\bar{x}_3\bar{x}_4 + x_3x_4 + \bar{x}_1\bar{x}_2x_4 + x_1x_2\bar{x}_3x_4$$

$$D = \sum(9, 12, 14).$$

$x_3x_4 \backslash x_1x_2$	00	01	11	10
00	m_0	m_4	m_{12}	m_8
01	m_1	m_5	m_{13}	m_9
11	m_3	m_7	m_{15}	m_{11}
10	m_2	m_6	m_{14}	m_{10}

Let's map the expression to the K-Map

$$f(x_1, \ldots, x_4) = \bar{x}_1\bar{x}_3\bar{x}_4 + x_3x_4 + \bar{x}_1\bar{x}_2x_4 + x_1x_2\bar{x}_3x_4$$

$$D = \sum(9, 12, 14).$$

$x_3x_4 \backslash x_1x_2$	00	01	11	10
00	m_0	m_4	d	m_8
01	m_1	m_5	m_{13}	d
11	m_3	m_7	m_{15}	m_{11}
10	m_2	m_6	d	m_{10}

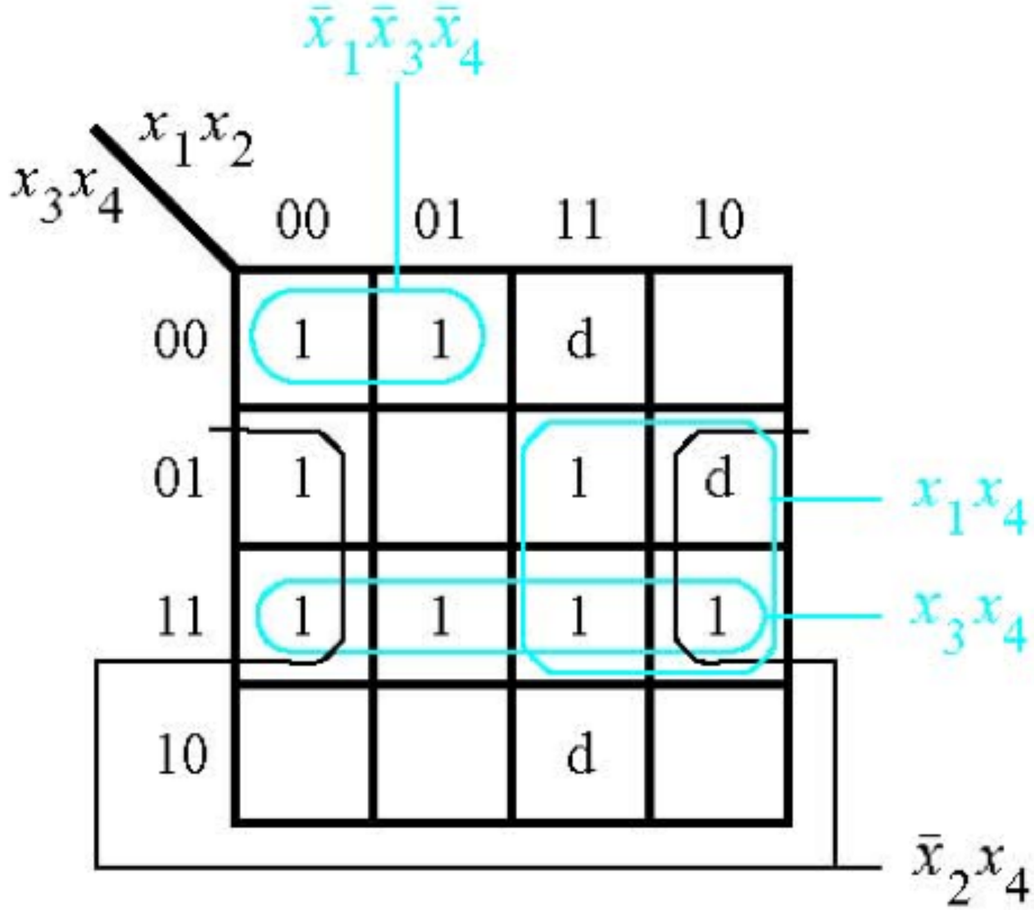
Let’s map the expression to the K-Map

$$f(x_1, \ldots, x_4) = \bar{x}_1\bar{x}_3\bar{x}_4 + x_3x_4 + \bar{x}_1\bar{x}_2x_4 + x_1x_2\bar{x}_3x_4$$

x_3x_4 \ x_1x_2	00	01	11	10
00	m_0	m_4	d	m_8
01	m_1	m_5	m_{13}	d
11	m_3	m_7	m_{15}	m_{11}
10	m_2	m_6	d	m_{10}

The SOP Expression

$$f(x_1, \ldots, x_4) = \bar{x}_1\bar{x}_3\bar{x}_4 + x_3x_4 + \bar{x}_1\bar{x}_2x_4 + x_1x_2\bar{x}_3x_4$$



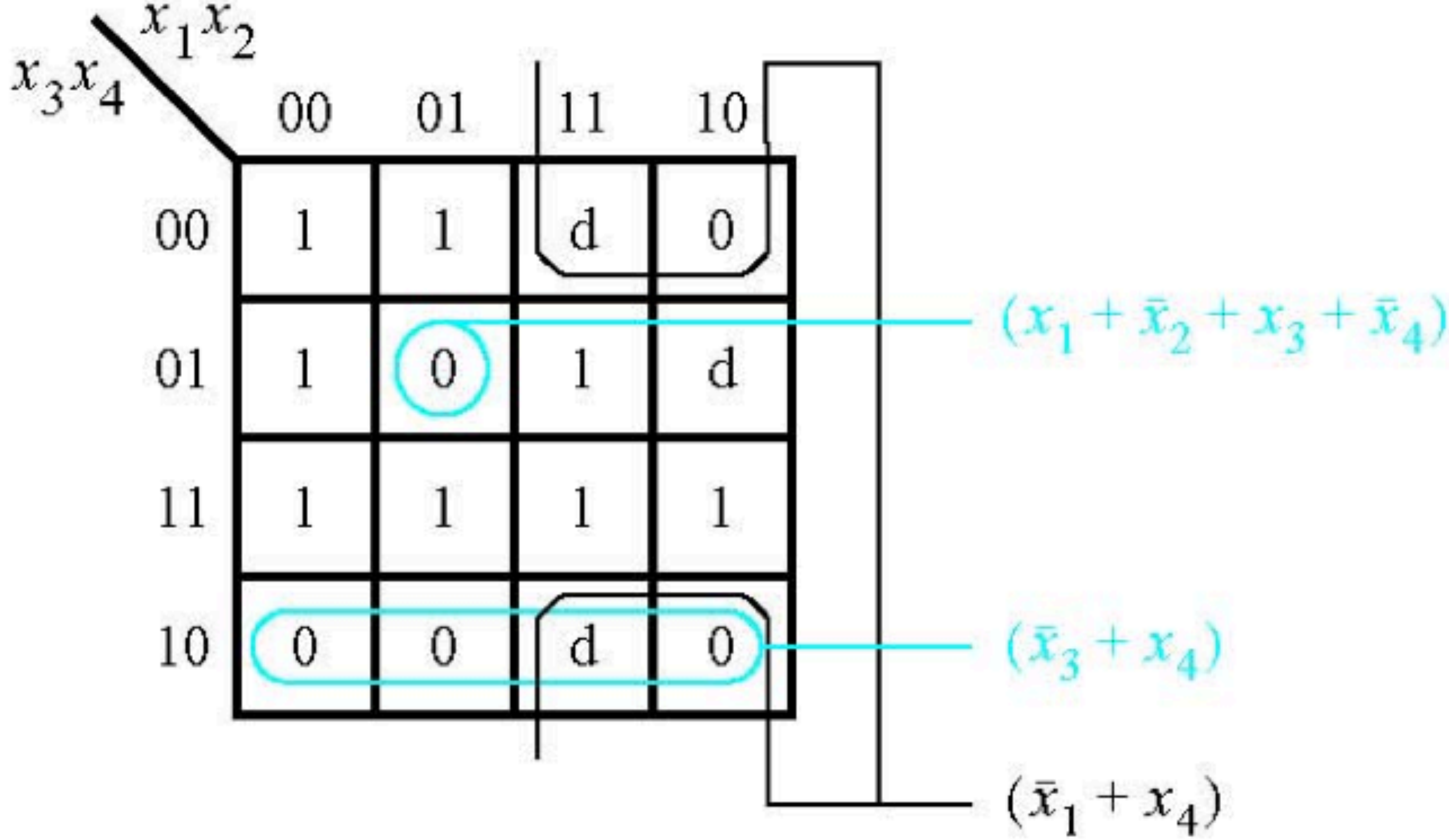
$$f = x_3x_4 + \bar{x}_1\bar{x}_3\bar{x}_4 + \bar{x}_2x_4 + x_1x_4$$

[Figure 2.68a from the textbook]

What about the POS Expression?

x_3x_4 \ x_1x_2	00	01	11	10
00	1	1	d	0
01	1	0	1	d
11	1	1	1	1
10	0	0	d	0

The POS Expression



$$f = (\bar{x}_3 + x_4)(\bar{x}_1 + x_4)(x_1 + \bar{x}_2 + x_3 + \bar{x}_4)$$

[Figure 2.68b from the textbook]

Example 7

Derive the minimum-cost SOP expression for

$$f = s_3(\bar{s}_1 + \bar{s}_2) + s_1 s_2$$

First, expand the expression using property 12a

$$f = s_3(\bar{s}_1 + \bar{s}_2) + s_1 s_2$$

$$f = \bar{s}_1 s_3 + \bar{s}_2 s_3 + s_1 s_2$$

Construct the K-Map for this expression

$$f = \bar{s}_1 s_3 + \bar{s}_2 s_3 + s_1 s_2$$

s_1	s_2	s_3	
0	0	0	m_0
0	0	1	m_1
0	1	0	m_2
0	1	1	m_3
1	0	0	m_4
1	0	1	m_5
1	1	0	m_6
1	1	1	m_7

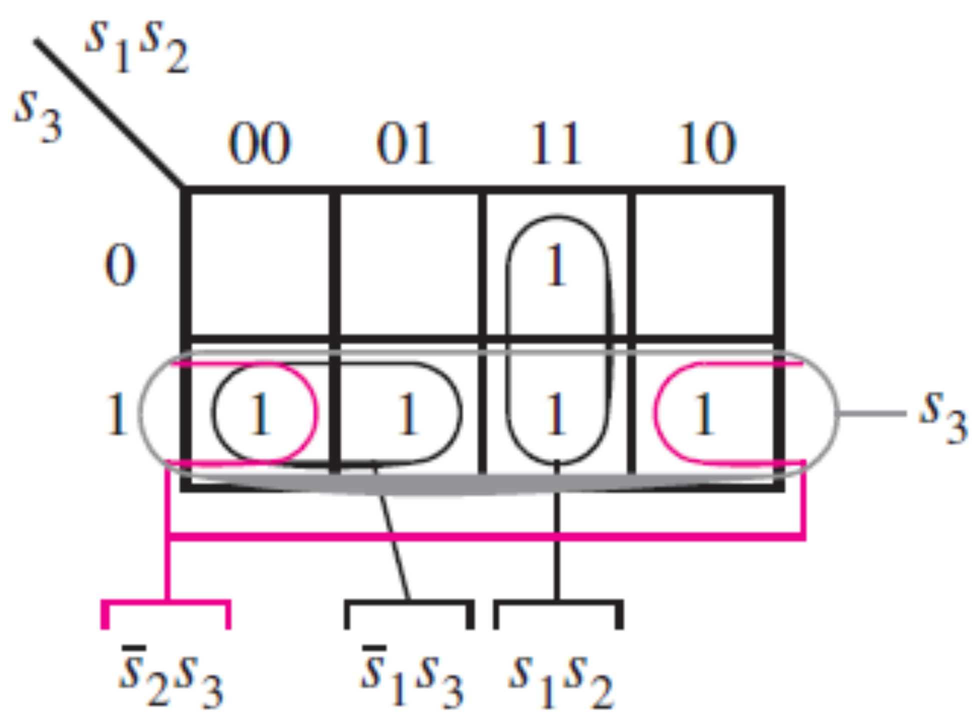
(a) Truth table

s_3 \ $s_1 s_2$	00	01	11	10
0	m_0	m_2	m_6	m_4
1	m_1	m_3	m_7	m_5

(b) Karnaugh map

Construct the K-Map for this expression

$$f = \bar{s}_1 s_3 + \bar{s}_2 s_3 + s_1 s_2$$

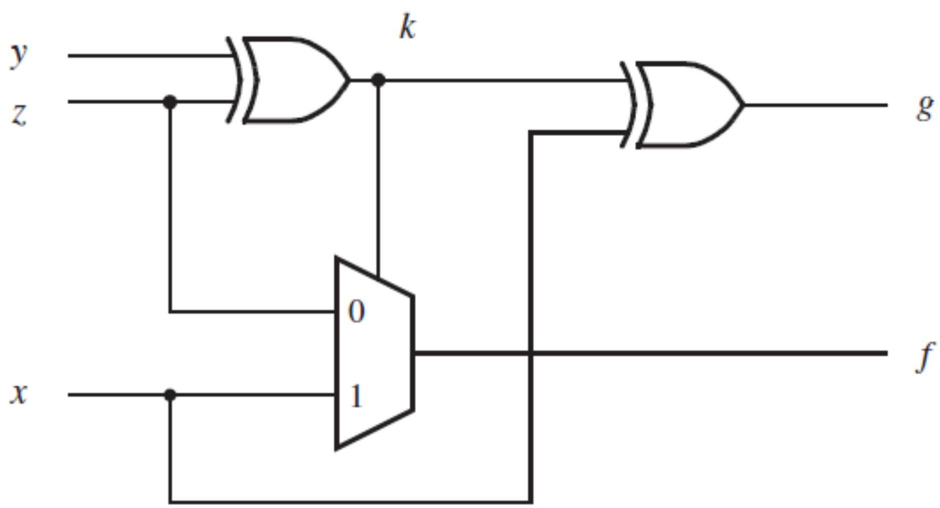


[Figure 2.69 from the textbook]

Example 8

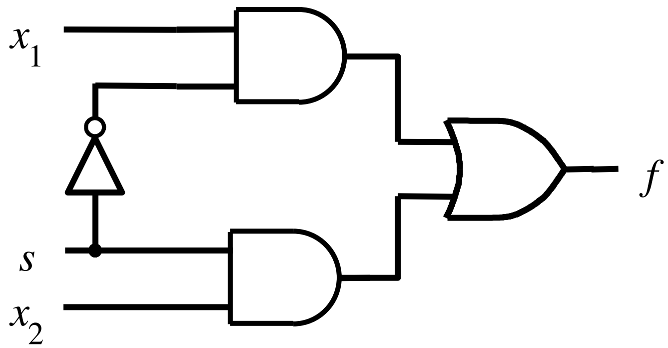
Write the Verilog code for the following circuit …

Logic Circuit

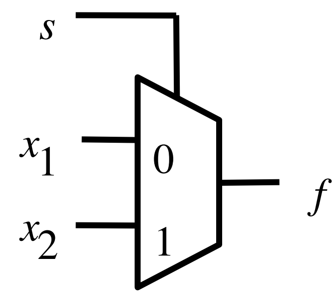


[Figure 2.70 from the textbook]

Circuit for 2-1 Multiplexer



(b) Circuit

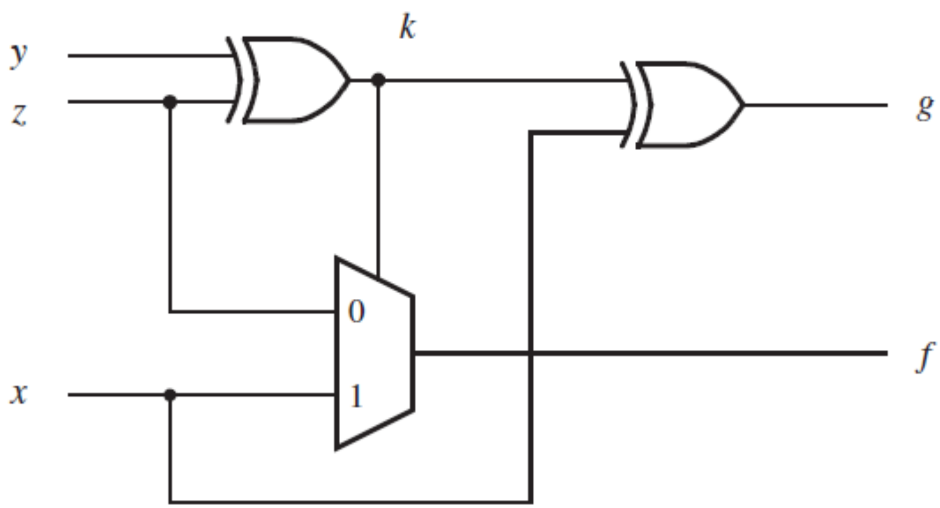


(c) Graphical symbol

$$f(s, x_1, x_2) = \overline{s}\, x_1 + s\, x_2$$

[Figure 2.33b-c from the textbook]

Logic Circuit vs Verilog Code




```
module f_g (x, y, z, f, g);
    input x, y, z;
    output f, g;
    wire k;

    assign k = y ^ z;
    assign g = k ^ x;
    assign f = (~k & z) | (k & x);

endmodule
```

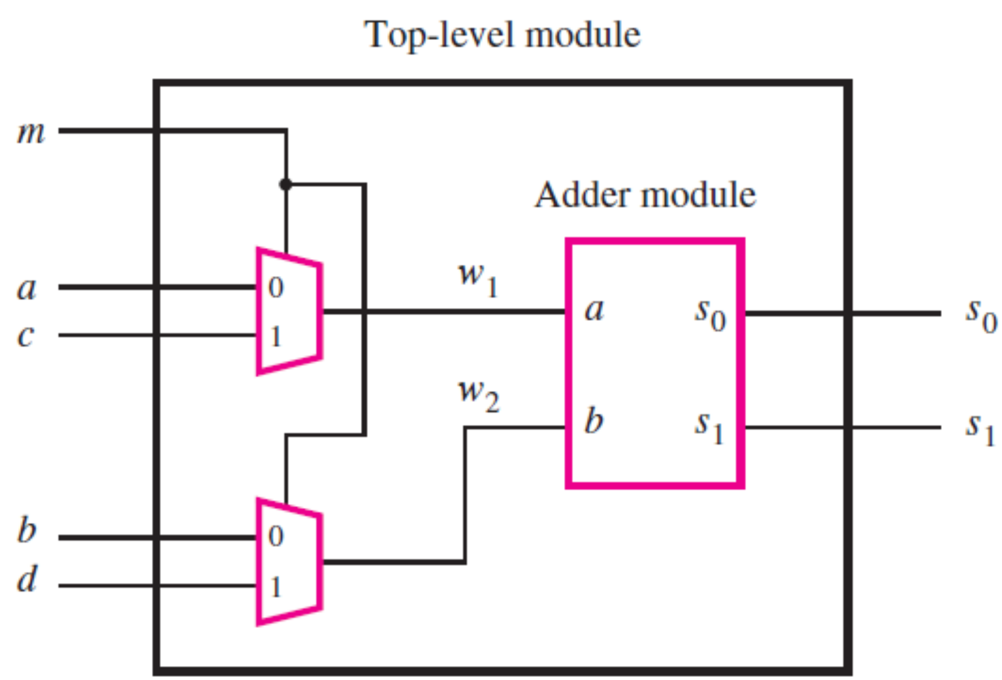
[Figure 2.70 from the textbook]

[Figure 2.71 from the textbook]

Example 9

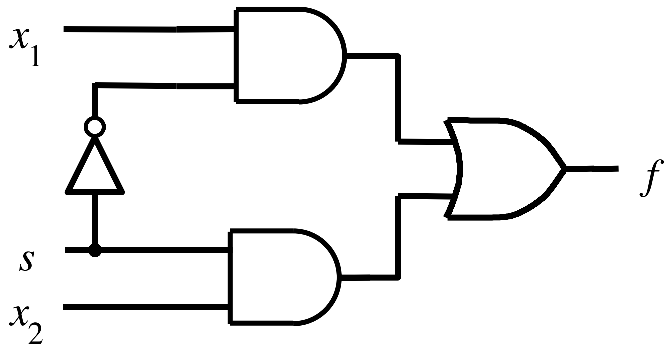
Write the Verilog code for the following circuit …

The Logic Circuit for this Example

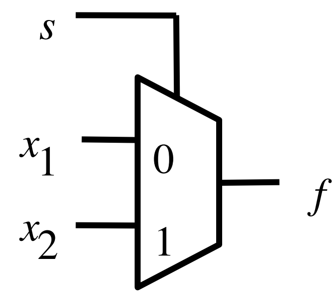


[Figure 2.72 from the textbook]

Circuit for 2-1 Multiplexer



(b) Circuit

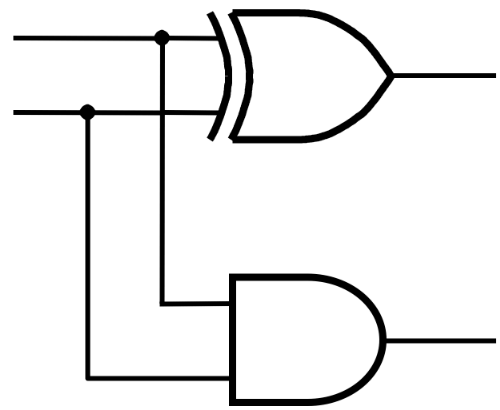


(c) Graphical symbol

$$f(s, x_1, x_2) = \overline{s}\, x_1 + s\, x_2$$

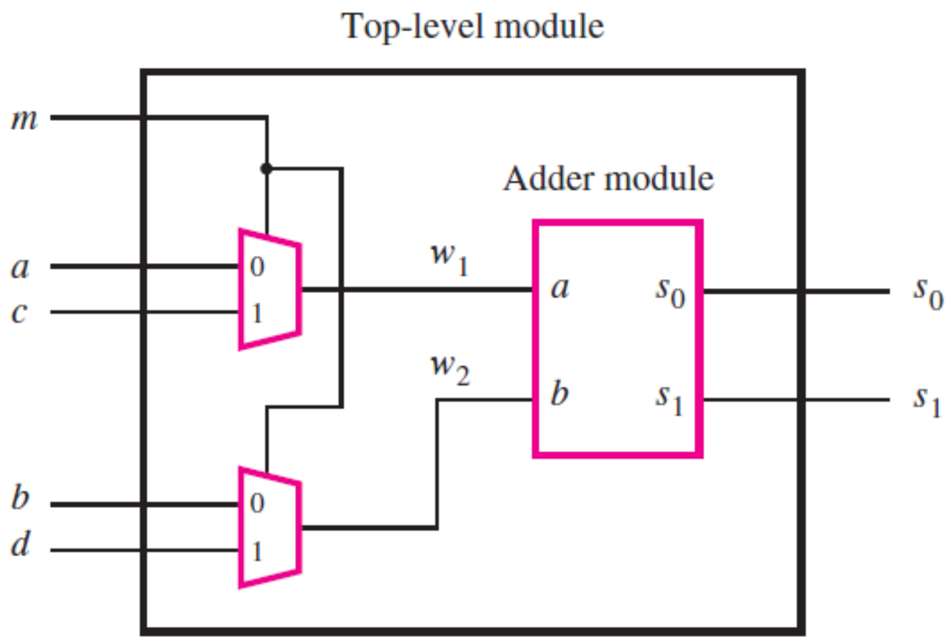
[Figure 2.33b-c from the textbook]

Addition of Binary Numbers



a	b	s_1	s_0
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

Logic Circuit vs Verilog Code

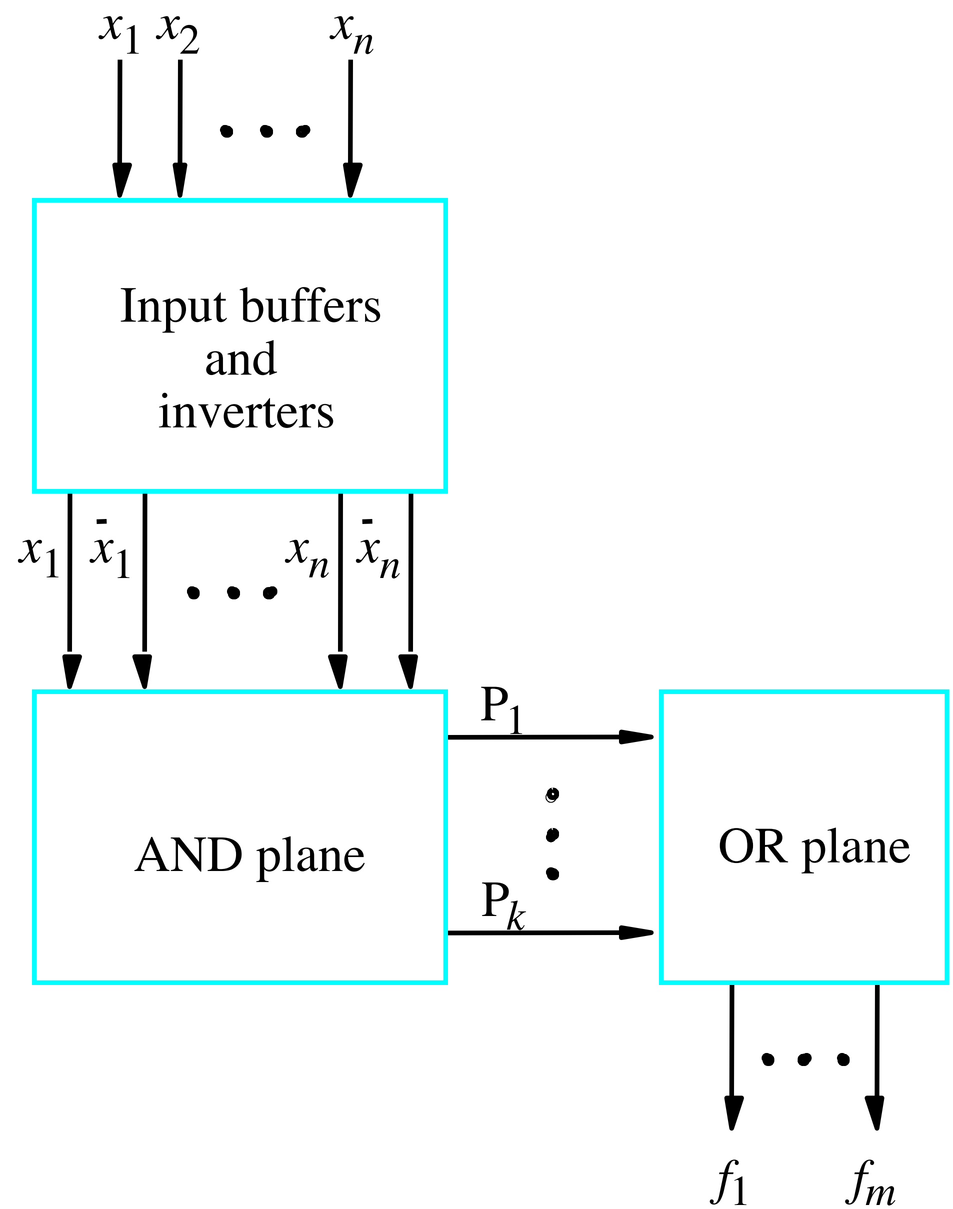



```
module shared (a, b, c, d, m, s1, s0);
    input a, b, c, d, m;
    output s1, s0;
    wire w1, w2;
    mux2to1 U1 (a, c, m, w1);
    mux2to1 U2 (b, d, m, w2);
    adder U3 (w1, w2, s1, s0);
endmodule

module mux2to1 (x1, x2, s, f);
    input x1, x2, s;
    output f;
    assign f = (~s & x1) | (s & x2);
endmodule

module adder (a, b, s1, s0);
    input a, b;
    output s1, s0;
    assign s1 = a & b;
    assign s0 = a ^ b;
endmodule
```

[Figure 2.73 from the textbook]

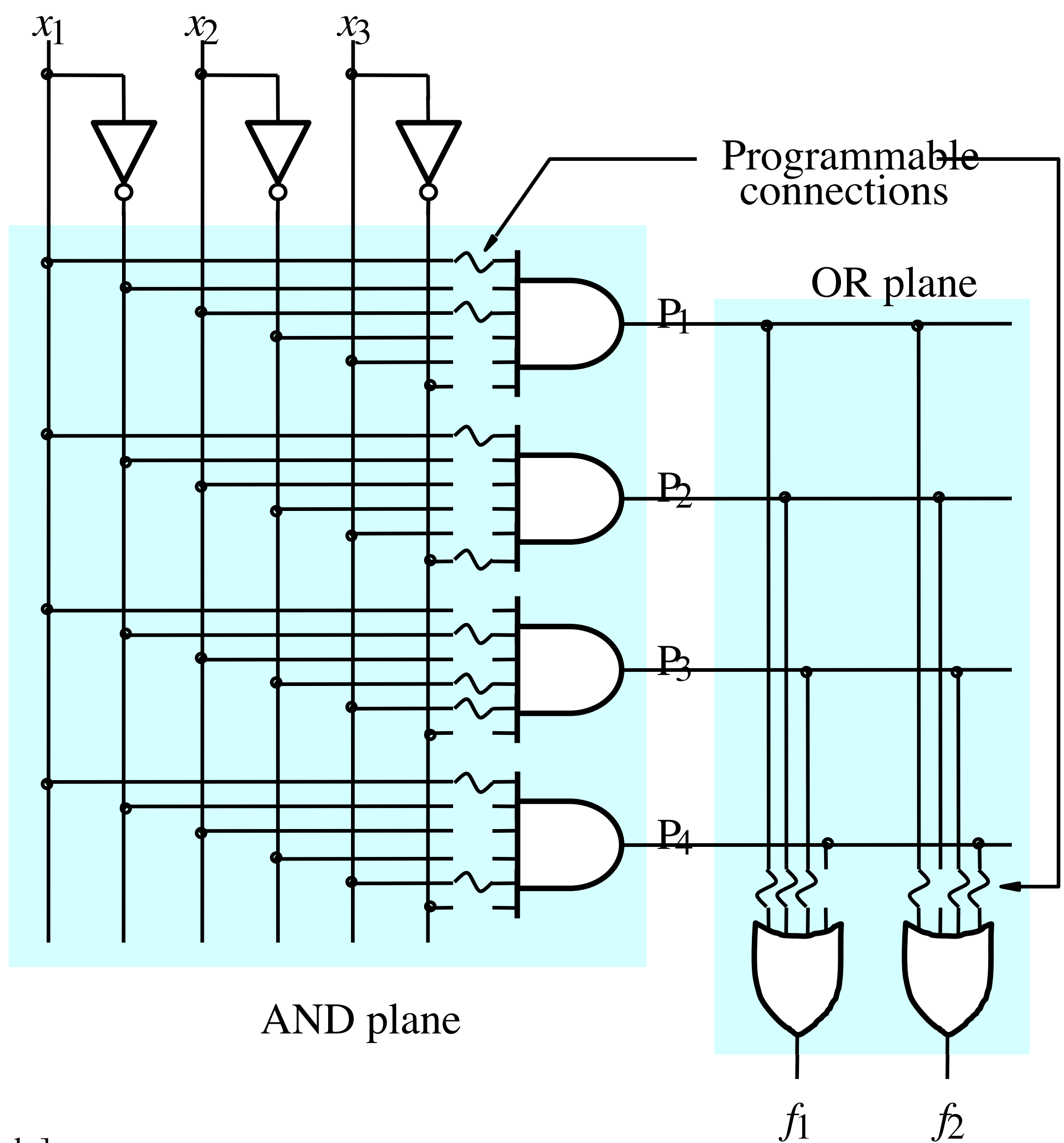
Some material form Appendix B
(that is needed for HW5)

Programmable Logic Array (PLA)



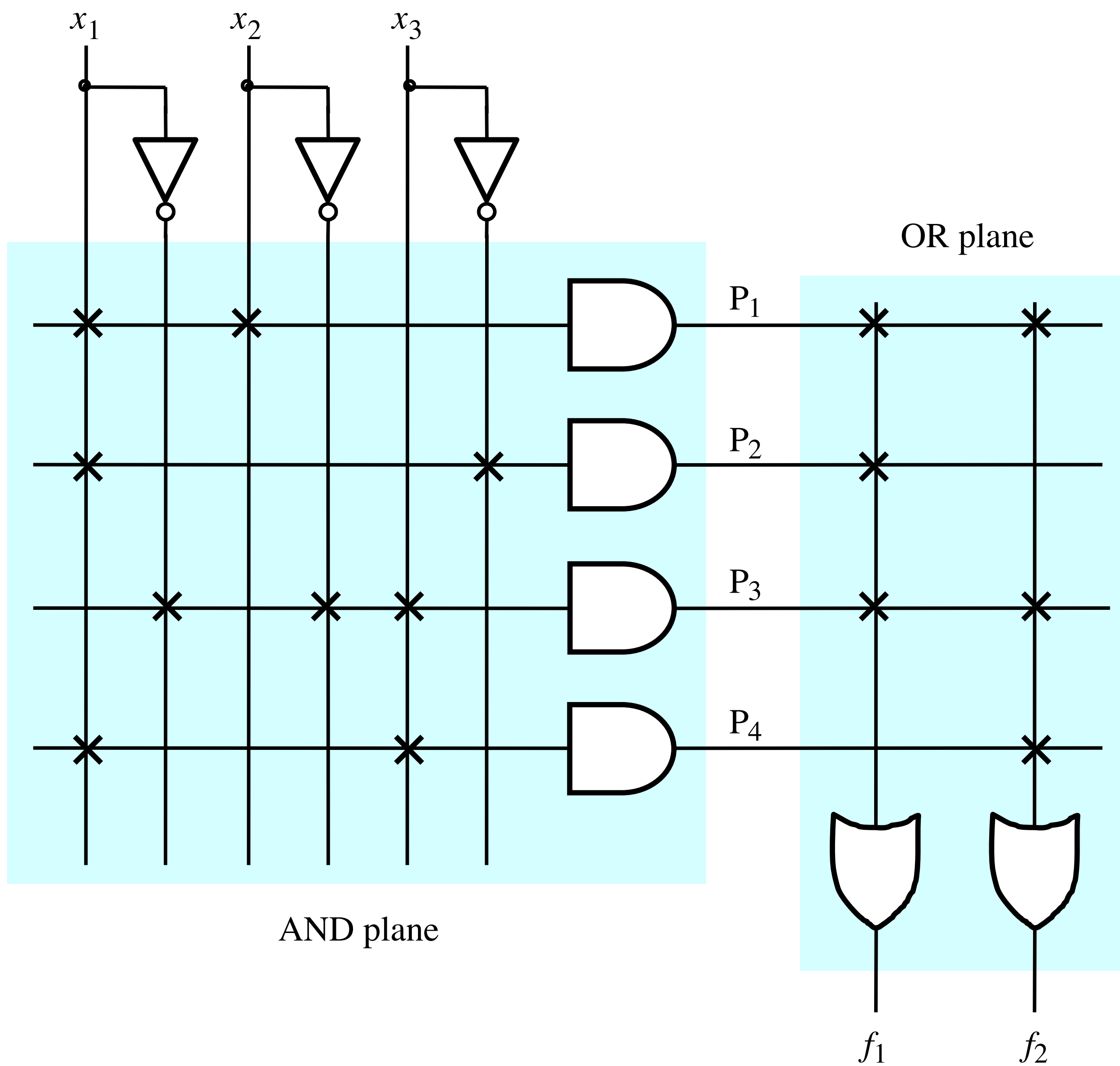
[Figure B.25 from textbook]

Gate-Level Diagram of a PLA



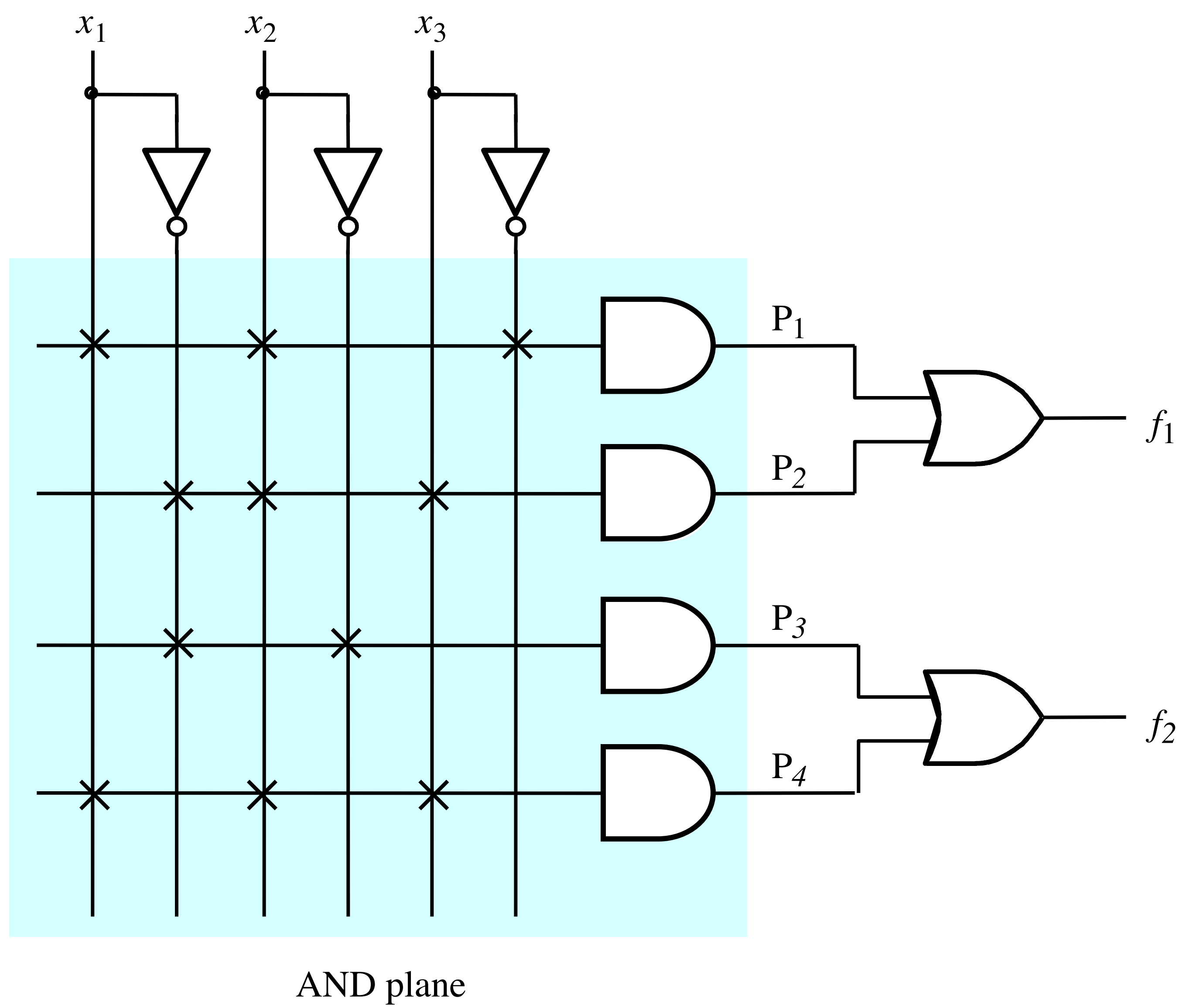
[Figure B.26 from textbook]

Customary Schematic for PLA



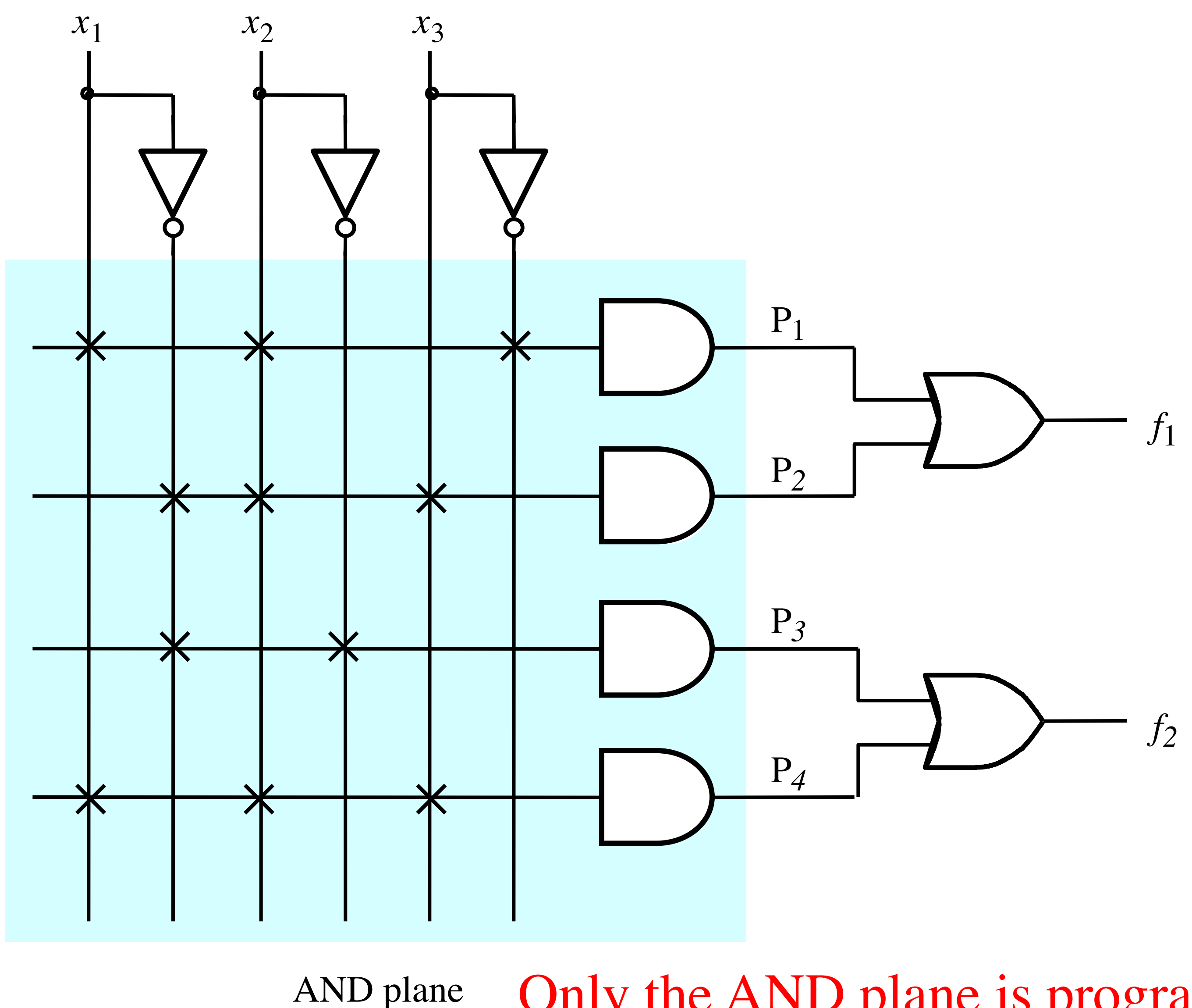
[Figure B.27 from textbook]

Programmable Array Logic (PAL)



[Figure B.28 from textbook]

Programmable Array Logic (PAL)



Only the AND plane is programmable.
The OR plane is fixed.

[Figure B.28 from textbook]

Questions?

THE END